

## Project 1: MIPS programming with MARS

In this project, you will write a MIPS assembly program and verify that it runs correctly with the simulator MARS. This program is a simple version of the “Perceptron Algorithm” widely used in AI, and will be used for subsequent projects of this class.

### Simplified Perceptron Algorithm

The job is to “guess” the hidden parameters (weights) of A and B, in a simple linear function  $F(x, y) = Ax + By$ , where all the numbers are assumed to be integers. A and B are not given, but whether the result  $F(x[i], y[i])$  is positive ( $\geq 0$ ) or negative ( $< 0$ ), is given for a large number of sample points  $(x[1], y[1])$ ,  $(x[2], y[2])$ , ... The algorithm goes as follows:

Use a “guess” version of  $f(x, y) = ax + by$ , initialized with  $a = b = 1$ , to go through all the sample points  $(x[i], y[i])$ , and update (improve) a and b accordingly:

If  $f(x[i], y[i]) < 0$  but  $F(x[i], y[i]) \geq 0$ :  
    update  $a = a + x[i]$ ;  
    update  $b = b + y[i]$ ;

If  $f(x[i], y[i]) \geq 0$  but  $F(x[i], y[i]) < 0$ :  
    update  $a = a - x[i]$ ;  
    update  $b = b - y[i]$ ;

Otherwise (f and F have the same sign for  $x[i], y[i]$ ):  
    keep a, b, as they are.

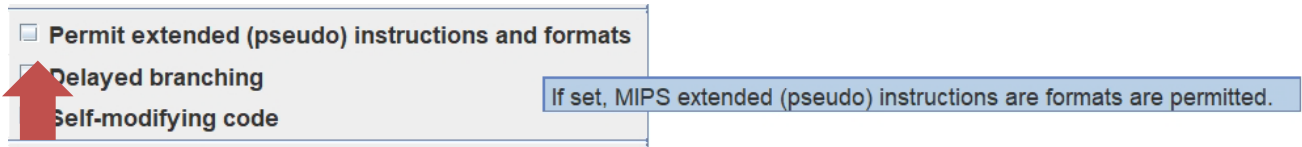
### Restrictions:

You are limited to use the instructions listed in the MIPS reference card (attached in Appendix), NOT including the pseudo-instructions.

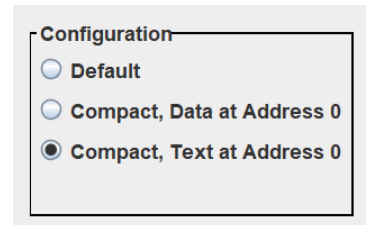
## MARS Setup:

when programming your MIPS assembly code, MARS needs to be set with the following:

- “Permit extended (pseudo) instructions and formats” option UNCHECKED such that extended (pseudo) instructions are NOT permitted.



- MIPS Memory Configuration should be set to the 3<sup>rd</sup> option: Compact, Text at Address 0



### (60%) Part A: generate the sample points via a “right-up” spiral.

Write MIPS code to generate n sample points from (x[1], y[1]) to (x[n], y[n]) in the following way: starting from (0, 0), always follow the sequence of going right, up, left, and down, in a “spiral” fashion that can reach any point exactly once, within enough steps.

Specifically:

first	1 step	right	(x1, y1) = (1, 0)		
then	1 step	up	(x2, y2) = (1, 1)		
then	2 steps	left	(x3, y3) = (0, 1)	(x4, y4) = (-1, 1)	
then	2 steps	down	(x5, y5) = (-1, 0)	(x6, y6) = (-1, -1)	
then	3 steps	right	(x7, y7) = (0, -1)	(x8, y8) = (1, -1)	(x9, y9) = (2, -1)
then	3 steps	up	(x10, y10) = (2, 0)	(x11, y11) = (2, 1)	(x12, y12) = (2, 2)
then	4 steps	left	...		
then	4 steps	down	...		

### Other Specifics:

In the end, your program should store the word arrays:

x[1] at M[0x2020], followed by x[2] at 0x2024, x[3] at 0x2028, etc  
y[1] at M[0x2420], followed by y[2] at 0x2424, y[3] at 0x2428, etc

Your program must begin with the following line of code which is used to initialize the number of sample points (n=100). Therefore, if we modify this instruction, we can change the number of points to generate. (This should work for Part A, B, and C)

```
addi $3, $0, 100    # n = 100
```

**(80%) Part B: evaluate the classification array for  $F(x, y)$**

Your program must at the beginning use the following 2 lines of code to specify A and B:

```
addi $1, $0, 5      # A = 5
addi $2, $0, -6     # B = -6
```

Again, if we modify these instructions, we can change A, B to other values.

Your program should use these given A, B values, to provide the “classification” results:

```
C[i] = 1      when  $F(x, y) = A * x[i] + B * y[i] \geq 0$ 
C[i] = -1     when  $F(x, y) = A * x[i] + B * y[i] < 0$ 
```

In the end, your program should store the word array:

`C[1]` at `M[0x2820]`, followed by `c[2]` at `0x2824`, `c[3]` at `0x2828`, etc.

**(100%) Part C: perceptron algorithm**

Begin at  $a = b = 1$ , use the arrays  $x$ ,  $y$ , and  $C$  in memory, to update  $a$  and  $b$ , according to the perceptron algorithm.

In the end, write the final  $a$  and  $b$  values in memory:

```
a at M[0x2000]
b at M[0x2004]
```

**(10% extra credit) NO multiplication instructions**

Your implementation should NOT include any multiplication instructions, but still achieve reasonable runtime.

## Report components:

1. General intro and Q&A (-2 points each if answer is missing)
  - a) Which parts (A, B, C, extra credit) does your program achieve?
  - b) How many hours have you spend in total on this project?
  - c) What was the most difficult (or time-consuming) part of this project?
2. Program functionality

Run your program for each of the following configurations:

- a) A = 5, B = -6, n = 100
- b) A = 0, B = 5, n = 50
- c) A = 7, B = 0, n = 150
- d) A = 2, B = 13, n = 200

Provide MARS screenshots of your program's results for each of the above configurations, including:

- final results for a, b, x, y, c
  - show MARS' data memory content at relevant regions of:
    - 0x2000, 0x2020, 0x2420, 0x2820.
- instruction count of completing your program
  - use MARS Tools -> Instruction Statistics

We will also test your program with other values .

3. Appendix

Include your MIPS assembly code here (-5 if missing).

**Deadline:** week 4 Sunday (end of the day) on Bb.

**What to submit:** upload the following two files on Bb

- **report.pdf** the report specified above.
- **p1.asm** your MIPS assembly code (so we can run your code on MARS - missing it will result in -5 points).

Appendix: allowed MIPS instructions – choose up to 15 different ones among them

<b>add</b>	rd, rs, rt	Add	$rd = rs + rt$	R 0 / 20
<b>sub</b>	rd, rs, rt	Subtract	$rd = rs - rt$	R 0 / 22
<b>addi</b>	rt, rs, imm	Add Imm.	$rt = rs + imm_{\pm}$	I 8
<b>addu</b>	rd, rs, rt	Add Unsigned	$rd = rs + rt$	R 0 / 21
<b>subu</b>	rd, rs, rt	Subtract Unsigned	$rd = rs - rt$	R 0 / 23
<b>addiu</b>	rt, rs, imm	Add Imm. Unsigned	$rt = rs + imm_{\pm}$	I 9
<b>mult</b>	rs, rt	Multiply	$\{hi, lo\} = rs * rt$	R 0 / 18
<b>div</b>	rs, rt	Divide	$lo = rs / rt; hi = rs \% rt$	R 0 / 1a
<b>multu</b>	rs, rt	Multiply Unsigned	$\{hi, lo\} = rs * rt$	R 0 / 19
<b>divu</b>	rs, rt	Divide Unsigned	$lo = rs / rt; hi = rs \% rt$	R 0 / 1b
<b>mfhi</b>	rd	Move From Hi	$rd = hi$	R 0 / 10
<b>mflo</b>	rd	Move From Lo	$rd = lo$	R 0 / 12
<b>and</b>	rd, rs, rt	And	$rd = rs \& rt$	R 0 / 24
<b>or</b>	rd, rs, rt	Or	$rd = rs   rt$	R 0 / 25
<b>nor</b>	rd, rs, rt	Nor	$rd = \sim(rs   rt)$	R 0 / 27
<b>xor</b>	rd, rs, rt	eXclusive Or	$rd = rs \wedge rt$	R 0 / 26
<b>andi</b>	rt, rs, imm	And Imm.	$rt = rs \& imm_0$	I c
<b>ori</b>	rt, rs, imm	Or Imm.	$rt = rs   imm_0$	I d
<b>xori</b>	rt, rs, imm	eXclusive Or Imm.	$rt = rs \wedge imm_0$	I e
<b>sll</b>	rd, rt, sh	Shift Left Logical	$rd = rt \ll sh$	R 0 / 0
<b>srl</b>	rd, rt, sh	Shift Right Logical	$rd = rt \gg sh$	R 0 / 2
<b>sra</b>	rd, rt, sh	Shift Right Arithmetic	$rd = rt \gg sh$	R 0 / 3
<b>sllv</b>	rd, rt, rs	Shift Left Logical Variable	$rd = rt \ll rs$	R 0 / 4
<b>srlv</b>	rd, rt, rs	Shift Right Logical Variable	$rd = rt \gg rs$	R 0 / 6
<b>srav</b>	rd, rt, rs	Shift Right Arithmetic Variable	$rd = rt \gg rs$	R 0 / 7
<b>slt</b>	rd, rs, rt	Set if Less Than	$rd = rs < rt ? 1 : 0$	R 0 / 2a
<b>sltu</b>	rd, rs, rt	Set if Less Than Unsigned	$rd = rs < rt ? 1 : 0$	R 0 / 2b
<b>slti</b>	rt, rs, imm	Set if Less Than Imm.	$rt = rs < imm_{\pm} ? 1 : 0$	I a
<b>sltiu</b>	rt, rs, imm	Set if Less Than Imm. Unsigned	$rt = rs < imm_{\pm} ? 1 : 0$	I b
<b>j</b>	addr	Jump	$PC = PC \& 0xF0000000   (addr_0 \ll 2)$	J 2
<b>jal</b>	addr	Jump And Link	$\$ra = PC + 8; PC = PC \& 0xF0000000   (addr_0 \ll 2)$	J 3
<b>jr</b>	rs	Jump Register	$PC = rs$	R 0 / 8
<b>jalr</b>	rs	Jump And Link Register	$\$ra = PC + 8; PC = rs$	R 0 / 9
<b>beq</b>	rt, rs, imm	Branch if Equal	$\text{if } (rs == rt) \text{ PC} += 4 + (imm_{\pm} \ll 2)$	I 4
<b>bne</b>	rt, rs, imm	Branch if Not Equal	$\text{if } (rs != rt) \text{ PC} += 4 + (imm_{\pm} \ll 2)$	I 5
<b>syscall</b>		System Call	$c0\_cause = 8 \ll 2; c0\_epc = PC; PC = 0x80000080$	R 0 / c
<b>lui</b>	rt, imm	Load Upper Imm.	$rt = imm \ll 16$	I f
<b>lb</b>	rt, imm(rs)	Load Byte	$rt = \text{SignExt}(M_1[rs + imm_{\pm}])$	I 20
<b>lbu</b>	rt, imm(rs)	Load Byte Unsigned	$rt = M_1[rs + imm_{\pm}] \& 0xFF$	I 24
<b>lh</b>	rt, imm(rs)	Load Half	$rt = \text{SignExt}(M_2[rs + imm_{\pm}])$	I 21
<b>lhu</b>	rt, imm(rs)	Load Half Unsigned	$rt = M_2[rs + imm_{\pm}] \& 0xFFFF$	I 25
<b>lw</b>	rt, imm(rs)	Load Word	$rt = M_4[rs + imm_{\pm}]$	I 23
<b>sb</b>	rt, imm(rs)	Store Byte	$M_1[rs + imm_{\pm}] = rt$	I 28
<b>sh</b>	rt, imm(rs)	Store Half	$M_2[rs + imm_{\pm}] = rt$	I 29
<b>sw</b>	rt, imm(rs)	Store Word	$M_4[rs + imm_{\pm}] = rt$	I 2b
<b>ll</b>	rt, imm(rs)	Load Linked	$rt = M_4[rs + imm_{\pm}]$	I 30
<b>sc</b>	rt, imm(rs)	Store Conditional	$M_4[rs + imm_{\pm}] = rt; rt = \text{atomic} ? 1 : 0$	I 38