Local equivalence of stabilizer states

a graphical characterisation

Nathan Claudet and Simon Perdrix

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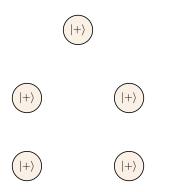






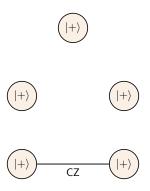


Graph states, local unitary equivalence, local Clifford equivalence & local complementation



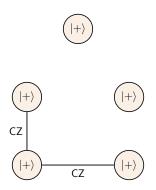
¹Edges do not have a direction.

²No multiples edges and no loops.



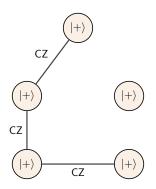
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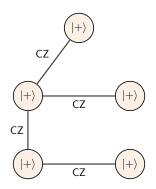
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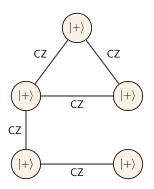
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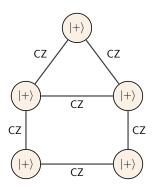
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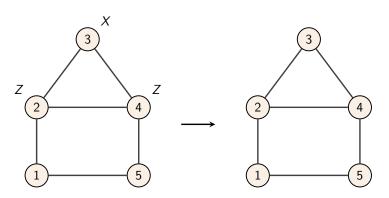


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Stabilizer states

Graph states are a subfamily of stabilizer states because for each vertex u, applying X on u and applying Z on the neighbours of u leaves the graph state invariant.



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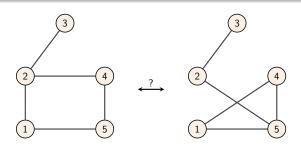
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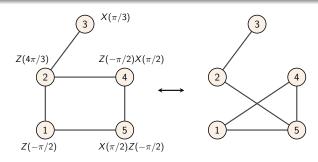


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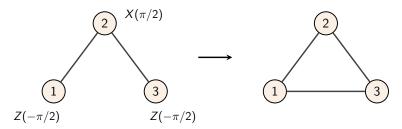
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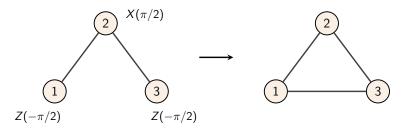
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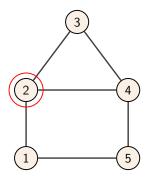
Theorem (Van den Nest, Dehaene, De Moor, 2004)

Two graph states are local Clifford equivalent iff the two corresponding graphs are related by **local complementations**.

Local complementation

Definition (Kotzig, 1966)

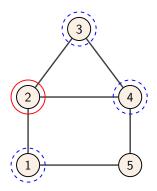
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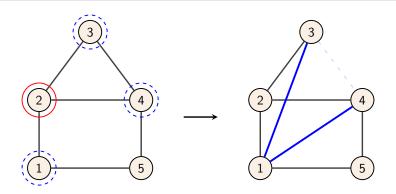
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Algorithmic aspect of local Clifford equivalence

There exists an efficient algorithm (Bouchet, 1991) to recognise whether two graphs are related by local complementations, implying an efficient algorithm to recognise whether two graph states are local Clifford equivalent.

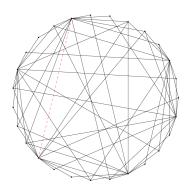
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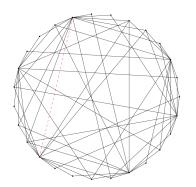
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Consequence: local complementation does **not** capture the local unitary equivalence of graph states.

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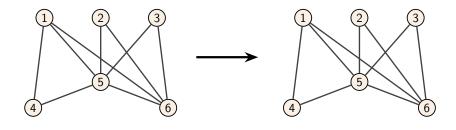
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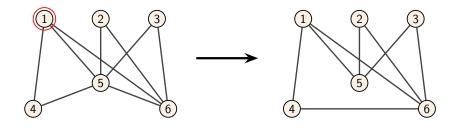
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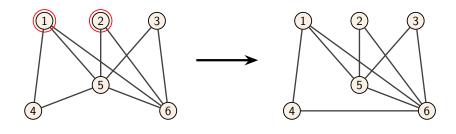
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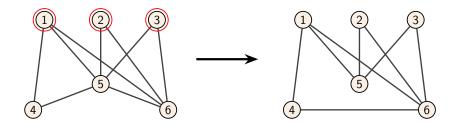
But what about local unitary equivalence for **any** graph? Can we construct a graphical characterisation?

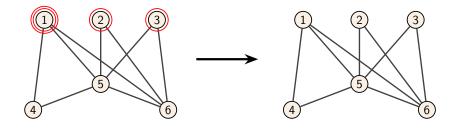
Generalising local complementation to capture local unitary equivalence

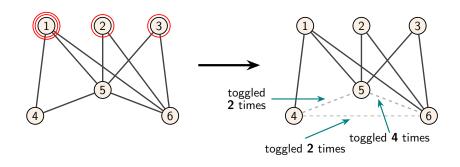






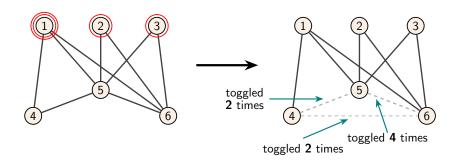






A refinement of idempotent local complementations

A sequence of local complementations may leave the graph invariant.

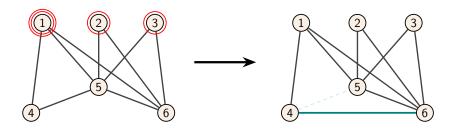


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r-local complementation

- 3-local complementation is a refinement of idempotent 2-local complementation, and so on...
- \rightarrow Infinite family of graphical operations parametrised by an integer r:

r-local complementations

1-local complementation = local complementation.

Main results

Theorem (this work)

Two graphs are related by r-local complementations iff the two corresponding graph states are related by local unitaries in the level r+1 of the Clifford hierarchy.

For r = 1, we recover local Clifford \Leftrightarrow local complementation.

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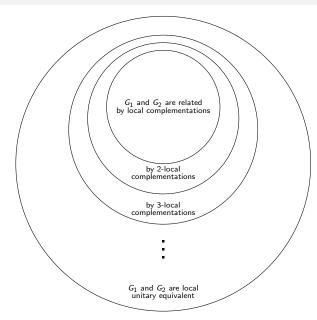
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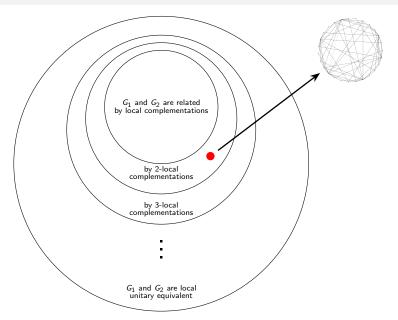
Corollary

If two graph states are local unitary equivalent, the local unitaries can be chosen to be in the Clifford hierarchy.

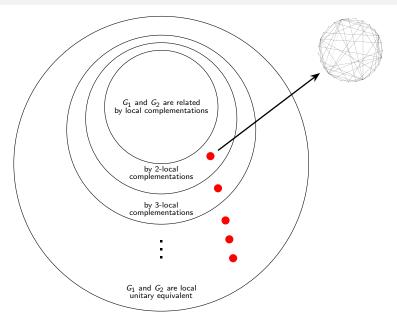
An infinite hierarchy of local equivalences



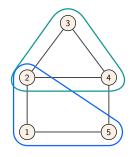
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Proof sketch: Minimal local set

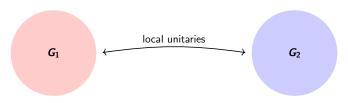


Minimal local sets are subsets of vertices that are invariant by local unitary equivalence and carry information on the possible local unitaries that maps graph states to other graph states.

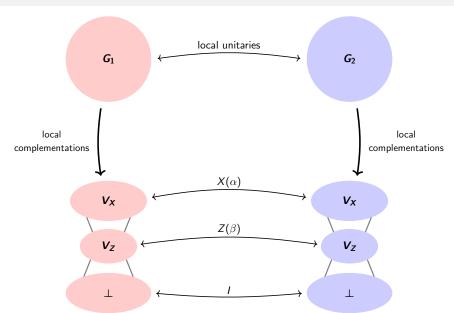
Theorem (C, Perdrix, 2024)

Each vertex of a graph is covered by at least one minimal local set.

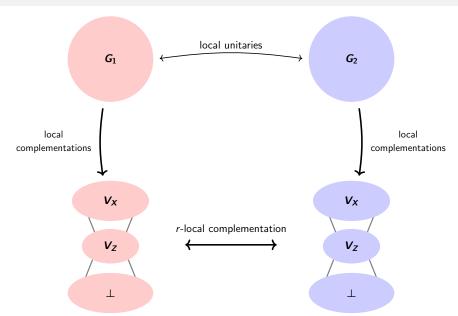
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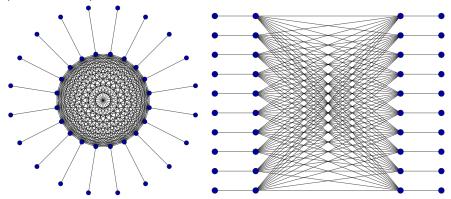


Proof sketch: Standard form for graph states



Application 1: local equivalence of repeater graph states

It was conjectured that LU=LC holds for some repeater graph states (Tzitrin, 2018). We showed that this is indeed the case.



Application 2: LU=LC for graph states up to 19 qubits

Previously: LU=LC for graph states up to 8 qubits, and there exists a 27-qubit pair for which LU \neq LC.

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Application 3: A quasi-polynomial algorithm to decide local unitary equivalence

Previously: exponential algorithm for deciding local unitary equivalence of graph states (Burchardt, de Jong, Vandré, 2024).

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Theorem (C, Perdrix, 2025)

There exists an algorithm that decides whether two graph states are local unitary equivalent with runtime $n^{\log_2(n)+O(1)}$.

The generalised local complementation is a graph rule that completely captures the local unitary equivalence of graph states.

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Open questions:

 Does there exist a counter-example to LU=LC between 20 and 26 qubits?

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Open questions:

- Does there exist a counter-example to LU=LC between 20 and 26 qubits?
- Does there exist a polynomial-time algorithm for local unitary equivalence?

Thanks



arXiv:2409.20183