

Project 1

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Abstract

The aim of this project is two-fold: primarily we aim to evaluate the limits of machine accuracy in the context of solving differential equations, and it functions as an introduction to problem solving in C++. We will be comparing the efficiency of vector and matrix methods of storage with respect to computation time while exercising basic C++ skills (i.e. dynamic memory allocation, compiling of cpp files, etc.).

1 Introduction

Before we begin, although C++ will be used for some parts of this project in order to learn the language, the majority of our coding will be done in Python. Our group has experience in using this language and we decided that it was best to use the coding language we are most comfortable with. We will do some work with C++ and a discussion of this will be provided in our report.

1.1 Mathematical Motivation

The motivation behind this project is simple: Many important differential equations in science can be written as linear second-order differential equations

$$\frac{d^2y}{dx^2} + k^2(x)y = f(x),$$

where f is normally called the inhomogeneous term and k^2 is a real function.

A classical equation from electromagnetism is Poisson's equation. The electrostatic potential Φ is generated by a localized charge distribution $\rho(\mathbf{r})$. In three dimensions it reads

$$\nabla^2\Phi = -4\pi\rho(\mathbf{r}).$$

With a spherically symmetric Φ and $\rho(\mathbf{r})$ the equations simplifies to a one-dimensional equation in r , namely

$$\frac{1}{r^2} \frac{d}{dr} \left(r^2 \frac{d\Phi}{dr} \right) = -4\pi\rho(r),$$

which can be rewritten via a substitution $\Phi(r) = \phi(r)/r$ as

$$\frac{d^2\phi}{dr^2} = -4\pi r\rho(r).$$

The inhomogeneous term f or source term is given by the charge distribution ρ multiplied by radius r and the constant -4π .

We will rewrite this equation by letting $\phi \rightarrow u$ and $r \rightarrow x$. The general one-dimensional Poisson equation reads then

$$-u''(x) = f(x).$$

2 Solution

2.1 Problem Set-Up

In this project we solved the one-dimensional Poisson equation with Dirichlet boundary conditions by rewriting it as a set of linear equations written as:

$$-u''(x) = f(x), \quad x \in (0, 1), \quad u(0) = u(1) = 0.$$

In our case we will assume that the source term is $f(x) = 100e^{-10x}$, and keep the same interval and boundary conditions. Then the above differential equation has a closed-form solution given by $u(x) = 1 - (1 - e^{-10})x - e^{-10x}$. We will compare our numerical solution with this result in the future exercise.

We fed the program an input value i , which was then used to determine x and the step size h . We defined these values as follows:

$$i = [input], \quad n = 10^i, \quad h = \frac{1.0}{n}$$

We approximate the second derivative of u with

$$-\frac{v_{i+1} + v_{i-1} - 2v_i}{h^2} = f_i \quad \text{for } i = 1, \dots, n,$$

where $f_i = f(x_i)$.

We began with a tridiagonal matrix, \mathbf{A} , and transformed the problem into the following:

$$\mathbf{A}\mathbf{v} = \tilde{\mathbf{b}},$$

where \mathbf{A} is an $n \times n$ tridiagonal matrix which we rewrite as

$$\mathbf{A} = \begin{bmatrix} 2 & -1 & 0 & \dots & \dots & 0 \\ -1 & 2 & -1 & 0 & \dots & \dots \\ 0 & -1 & 2 & -1 & 0 & \dots \\ & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & & -1 & 2 & -1 \\ 0 & \dots & & 0 & -1 & 2 \end{bmatrix},$$

and $\tilde{b}_i = h^2 f_i$. \mathbf{A} can be derived by looking at the formula for the approximation of the second derivative. Note how the v_{i+1} and the v_{i-1} terms have a coefficient of -1 . This is where the off diagonal elements come from in our matrix \mathbf{A} . Since the two off diagonal entries are equivalent, we can approximate them into one vector, e . However, since all of the entries are -1 , we decided to not allocate memory to this. Instead, we simplified our algorithm to account for these values. We will discuss our problem solving algorithms in the following section.

The reason for the conversion from a matrix to vectors comes about in our number of floating point operations (FLOPS). In using the matrix approximation, we were looking at $\frac{2}{3}n^3$ FLOPS in calculating our solutions^[2]. However, in being able to reduce this down to a vector problem, we managed to reduced our calculations to $9n$, and then further to $4n$, again to be discussed in the next section. The short of this discussion is that we were able to reduce our floating point operations by two orders of magnitude. In the sense of computational efficiency, this makes for much smarter problem solving and much more effective computation.

2.2 Algorithms

The algorithms that we implemented in the solving of this problem are as follows. We denote the diagonal elements as $d[i]$, the solution values as $u[i]$, and the values of the function as $f[i]$. The problem starts when look at a general case for a 4×4 matrix:

$$\mathbf{A} = \begin{bmatrix} d_1 & e_1 & 0 & 0 \\ e_1 & d_2 & e_2 & 0 \\ 0 & e_2 & d_3 & e_3 \\ 0 & 0 & e_3 & d_4 \end{bmatrix}, \quad \mathbf{u} = \begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ u_4 \end{bmatrix}, \quad \mathbf{f} = \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ f_4 \end{bmatrix},$$

And now we can look at solving $\mathbf{A}\mathbf{u} = \tilde{\mathbf{f}}$, through the process of Gaussian Elimination.^[2] We begin by multiplying the second row of A by the ratio of the terms in the first column of rows one and two. Explicitly, this looks like $\frac{e_1}{d_1}$ on both the right hand side and left hand side. This transforms the equation into the following:

$$\begin{bmatrix} d_1 & e_1 & 0 & 0 \\ 0 & d_2 - \frac{e_1^2}{d_1} & e_2 & 0 \\ 0 & e_2 & d_3 & e_3 \\ 0 & 0 & e_3 & d_4 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ u_4 \end{bmatrix} = \begin{bmatrix} f_1 \\ f_2 - f_1 \frac{e_1}{d_1} \\ f_3 \\ f_4 \end{bmatrix}$$

We now can redefine $\tilde{d}_2 = d_2 - \frac{e_1^2}{d_1}$ and $\tilde{f}_2 = f_2 - f_1 \frac{e_1}{d_1}$. This changes the matrix to a new form that we can use to repeat the process further.

$$\begin{bmatrix} d_1 & e_1 & 0 & 0 \\ 0 & \tilde{d}_2 & e_2 & 0 \\ 0 & e_2 & d_3 & e_3 \\ 0 & 0 & e_3 & d_4 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ u_4 \end{bmatrix} = \begin{bmatrix} f_1 \\ \tilde{f}_2 \\ f_3 \\ f_4 \end{bmatrix}$$

Now, we can perform the same operation and multiply the second row by $\frac{e_2}{\tilde{d}_2}$. As we show with Gaussian Elimination, the final result will be an upper triangular matrix of the form:

$$\begin{bmatrix} d_1 & e_1 & 0 & 0 \\ 0 & \tilde{d}_2 & e_2 & 0 \\ 0 & 0 & \tilde{d}_3 & e_3 \\ 0 & 0 & 0 & \tilde{d}_4 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ u_4 \end{bmatrix} = \begin{bmatrix} f_1 \\ \tilde{f}_2 \\ \tilde{f}_3 \\ \tilde{f}_4 \end{bmatrix}$$

At this point, we can begin the process of forward substitution and equate $\tilde{d}_4 * u_4 = \tilde{f}_4$. Using this value, we can solve for u_4 and then solve $\tilde{d}_3 * u_3 + \tilde{d}_4 * u_4 = \tilde{f}_3$ for u_3 , and so on and so forth. With pre-populated arrays for u and f , we now have our solution. We can now extend this process to a more general case for a square matrix of any size. This process yields the algorithms used in our program.

$$\tilde{d}_i = d_i - \frac{e_{i-1}^2}{d_{i-1}} \quad (1)$$

$$\tilde{f}_i = f_i - f_{i-1} \frac{e_{i-1}}{d_{i-1}} \quad (2)$$

$$u_i = \frac{\tilde{f}_i - e_i u_{i+1}}{d_i} \quad (3)$$

Now we can use the fact that for our specific case all of the e_i values are equal to -1 to simplify our algorithms. This also causes a drastic simplification for the d_i values.

$$\tilde{d}_i = \frac{i+1}{i} \quad (4)$$

$$\tilde{f}_i = f_i + f_{i-1} \frac{1}{d_{i-1}} \quad (5)$$

$$u_i = \frac{\tilde{f}_i + u_{i+1}}{d_i} \quad (6)$$

Thus, we see a reduced floating point operation count down to $\approx 4n$. This is where the efficiency of this approach really becomes evident.

Some examples of our code can be seen in a sample from our C++ code below.

```
u[0] = 0.0; u[n] = 0.0; //u[0]=u[n]=0;
d[0] = d[n] = 2;
```

```

for (int i = 0; i <= n; i++){
    x[i] = i*h;
    f[i] = hh*ff(i*h);
}

for (int i = 1; i<n; i++) d[i] = (i+1.0)/( (double) i );
//Forward substitution
for (int i = 2; i< n ; i++) f[i] = f[i] + f[i-1]/d[i-1];
//Backwards substitution
u [ n-1] = f[n-1]/d[n-1];
for (int i = n-2; i > 0; i--) u[i] = (f[i]+u[i+1])/d[i];

```

2.3 Relative Error

We calculated the relative error in the data set $i = 1, \dots, n$, by setting up

$$\epsilon_i = \log_{10} \left(\left| \frac{v_i - u_i}{u_i} \right| \right),$$

as function of $\log_{10}(h)$ for the function values u_i and v_i . Here, it is important to note that the i value corresponds to the step size, h , as discussed above as $h = \frac{1.0}{n}$. The relative error was calculated for each power of ten, corresponding to an input value, i . For each step length, we saw that the error was relatively constant, meaning that for each value of x , the relative error was approximately the same. Since the error was the exact same for all x values, we elected to sample only one error entry *specifically the first entry* for the sake of efficiency (i.e. $i = 1, x = 0.1$ $i = 2, x = 0.01$, etc.). This is a standard practice for looking at error trends. We were consistent with our choice, so whatever error that this may have imposed, we were at least consistent with this and can account it to systematic error. We then conducted multiple trials utilizing values of i ranging from $i = 0$ to $i = 8$. Or, step sizes ranging from $10^{-6} \rightarrow 10^{-8}$. The results we got are plotted in **Figure 2**.

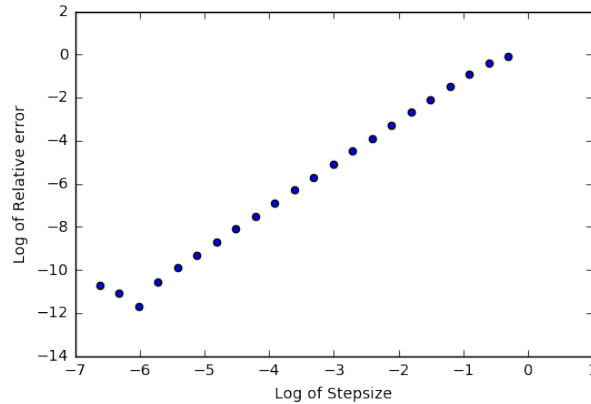


Figure 1: Plot of Log of Relative Error. On the x-axis, we see the log of the step-size. This is equivalent to looking at which i value we chose to use in order to approximate our solution. The y-axis shows the log of the relative error given in the equation above. What we see is that the log of error has a slope of approximately -2, as expected in our initial mathematic discussion. The error begins to increase again after $i = 6$ or $n = 10^6, h \approx 10^{-6}$. This is due to the truncation error that is getting overwritten by the number of bits. The increase in error can be accounted to machine error.

2.4 LU Decomposition Comparison

LU-Decomposition is a process by which a matrix A can be turned into an upper and lower diagonal matrix to make it easier to deal with. But there are ways to make a program to complete this process in an efficient manner. To set it up, say

$$A = LU$$

with row elements i and column elements j . So, the i th row and j th column of matrix A would be denoted by A_{ij} .

To start, there are a few tricks that we can use to easily get some values for U. In the first step, $U_{1j}=A_{1j}$. This will allow us to get the entire first row of the U matrix with little work, this just needs to be iterated over the dimensions of the matrix. We can use this simple equation because the matrix A is a multiplication of matrix L and U. The first row of L is $L = \begin{bmatrix} 1 & 0 & 0 & 0 \end{bmatrix}$ and the first row of L times the first column of U will always be equal to the first row of matrix U, or in other words,

$$A_{1j} = U_{1j}$$

In step two, we will derive some values for L. All diagonal elements of the L matrix is 1, due to the fact that L is a lower triangular matrix. When we multiply a row of matrix L with a column of matrix U, we will find the value of a value in matrix a with the row and column index equal to the row and column number used in the previously mentioned multiplication. We can find a pattern, which can be automated with the following equation:

$$L_{ij} = \frac{A_{ij} - \sum_{k=1}^{j-1} L_{ik}U_{kj}}{U_{ij}}.$$

Now we have all the values of matrix L

In step three, we will be finding the values for the diagonal elements of U. This is done in a similar method as above, where we use patterns from simple matrix multiplication and with the knowledge of what the result of this multiplication would be, we can predict the elements that were multiplied within matrix U. The following equation can be used for this purpose:

When $i=j$,

$$U_{ij} = A_{ij} - \sum_{k=1}^{i-1} L_{ik}U_{ki}$$

In the final step, we will be concerned with all elements within U, where $i < j$. This will get all the diagonal elements of U, which will complete our knowledge of the values of the elements of U

$$U_{ij} = A_{ij} - \sum_{k=1}^{i-1} L_{ik}U_{kj}$$

We then solved this problem again but with a different technique: LU decomposition. The process of LU decomposition is described in our algorithms section and reference 5.

The code that we used to set up the LU decomposition is described below: (Python Code)

```
def tridiag(a, b, c, k1=-1, k2=0, k3=1):
    #function taken from
    #http://stackoverflow.com/questions/5842903/block-tridiagonal-matrix-python
    return np.diag(a, k1) + np.diag(b, k2) + np.diag(c, k3)

def plot(x, y):
    pass

#def time_take(n):
#    #time_array = np.zeros(lim)
#    #n_array = np.zeros(lim)
#    a = [-1] * (n-1)
#    b = [2] * n
```

```

    c = [-1] * (n-1)
    tic = time.clock()
    A = tridiag(a, b, c)
    B=scipy.linalg.lu(A)
    toc=time.clock()
    delta_t=toc-tic

    return delta_t

def LU-Decomp(n):
    a = [-1] * (n-1)
    b = [2] * n
    c = [-1] * (n-1)
    #tic = time.clock()
    A = tridiag(a, b, c)
    B=scipy.linalg.lu(A)
    #toc=time.clock()
    #delta_t=toc-tic

    return B

#change exp to 4 at your own caution.
#It takes absurdly long for reasons I cant figure out.
#The computation should be the same as above
exp = 3
trials = 20
n_array = np.zeros(exp-1)
#time_array = np.zeros(exp-1)

colors = ['r', 'b', 'g', 'y']
patches = []
for i in range(1,exp):
    #loop for within exponent
    n = 10**exp
    n_array[i-1] = i
    #time_array[i-1] = time_take(n)
    time_array = np.zeros(trials)
    for j in range(trials):
        #loop for each trial
        tic = time.clock()
        B = LU-Decomp(n)
        toc = time.clock()
        delta_t = toc-tic
        time_array[j] = delta_t

```

To make this matrix, we took some code from online to initialize the tridiagonal matrix^[4]. The code doesn't make an $n \times n$ matrix as that would take a massive amount of memory. Instead it creates 3 arrays, which correspond to the 3 diagonals that have non-zero numbers within the tridiagonal matrix. Much of the code is to create the matrix, and there is a module in numpy which does the LU-decomposition, so we didn't have to hard code the process in. To see how long this process takes, we used the time.clock function from the time module and took the difference in times between the end of initializing the matrix and completing the LU-decomposition process. Then we ran the code 100 times and made a scatter plot of the time this took. We took the average time and found that to make a 1000×1000 matrix, it took an average of about .10 s to complete the code.

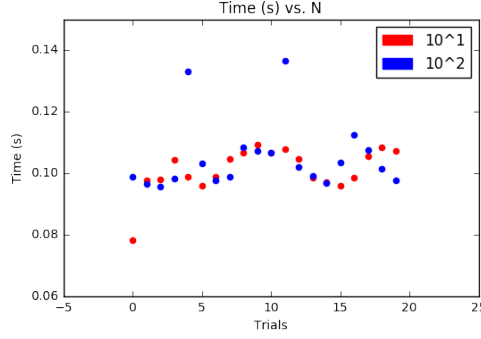


Figure 2: The plot above shows the trials for LU decomposition that we ran to calculate the time that it takes to achieve our numerical results. The x-axis denotes the trial number (just an integer associated with trial 1,2,3,... etc.) and the y-axis was the time that the program took to complete in seconds. The final result was an average time over 100 trials. Note, the size of the matrix is denoted by the color given in the key. The 10^3 sized matrix timed out and Python was unable to compile it.

The LU decomposition of the matrix was then used to solve for the solution to the differential equation as we did using our algorithms in our earlier code. While this may seem like an incredibly fast process, but the true difference became noticeable when we looked at larger matrices. For the $10^4 \times 10^4$ and $10^5 \times 10^5$ cases, the amount of memory required to complete the calculations and the time necessary to compile the code caused errors in our computer systems. Thus, we saw the fatal error of this approach. For our general case using simplified algorithms, we were able to effectively solve the differential equations for matrices up to the order of $10^8 \times 10^8$ case without issue. The amount of time and memory required to use LU decomposition for the sake of solving differential equations of this form make this an ineffective method.

Also, fun fact, it takes 42 minutes, roughly, to compile the 10^5 square matrix case.

In comparing to our C++ code, we included a time function to keep track of the amount of time required to perform the algorithms. The results were as followed:

Computation Time and Matrix Size

Matrix Size	Computational Time
10^1	2e-06 Seconds
10^2	3e-06 Seconds
10^3	2.6e-05 Seconds
10^4	0.000295 Seconds
10^5	0.002819 Seconds
10^6	0.02787 Seconds
10^7	0.291171 Seconds

Table 1. The table below summarizes the size of the square matrices and shows the time required to perform the numerical algorithms.

In the LU decomposition method, we saw that for 10^1 and 10^2 matrices, the computation time over 20 trials averaged out to roughly 0.1 seconds. The 10^3 square matrix could not compile without producing a timeout error (Python). In performing our original numerical algorithms method, we ran through 10^7 square matrices and the computational time reached roughly 0.3 seconds. This is on the same order of magnitude as the LU methods, but in matrix sizes that we couldn't even compile with the LU method. This shows us that our original approach was by far much more superior in terms of computational efficiency.

3 Concluding Remarks

The numerical ODE solver performed as expected, it showed increasing accuracy up until the relative error began to approach 10^{-12} . At that point the truncation errors of the machine began to accumulate,

decreasing the overall accuracy. By using the algorithmic tricks we discussed in class we could reduce the total computations to the order of N . this resulted in the vector method being much faster than the LU analogue, so much so as to give several orders of increased accuracy for comparable computation time.

4 References

1. Block tridiagonal matrix python. (n.d.). Retrieved January 27, 2017, from <http://stackoverflow.com/questions/5842903/block-tridiagonal-matrix-python>
2. <https://github.com/CompPhysics/ComputationalPhysicsMSU> , This is a github repository, which had some code in it that we used in our ipython notebooks.
3. Press, W. H., Vetterling, W. T., Teukolsky, S. A., and Flannery, B. P. (1992). *Numerical recipes in C. the art of scientific computing*. Cambridge University Press.