Goblint: Mixed Flow-Sensitive Abstract Interpretation

Simmo Saan¹ Julian Erhard^{2,3} **Michael Schwarz**² Karoliine Holter¹ Michael Petter² Vesal Vojdani¹ Helmut Seidl²

SV-COMP Community Meeting 2025

 $^1 \text{University}$ of Tartu $^2 \text{Technical University}$ of Munich $^3 \text{Ludwig-Maximilians}$ Universität München m.schwarz@tum.de

Goblint in 4 Bullet Points

- Static Analyzer for C programs
- Based on abstract interpretation sound!
- Overapproximating no violations!
- Specializes in multi-threaded programs

Goblint in 4 Bullet Points

- Static Analyzer for C programs
- Based on abstract interpretation sound!
- Overapproximating no violations!
- Specializes in multi-threaded programs

A basic example approach:

- Accumulate all possible values for globals
- Track local states per program point

```
int g = 0;
                                   t1:
                                       g = 10;
main:
                                       . . .
    x = 3;
                                       g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
```

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
     . . .
                        [g]:
```

```
int g = 0;
                                   t1:
                                       g = 10;
main:
                                        . . .
    x = 3;
                                       g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
                        [g]:
```

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
                         [g]:[0,0]
```

```
int g = 0;
                                   t1:
                                        g = 10;
main:
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    }
    . . .
                        [g]:[0,0]
```

```
int g = 0;
                                  t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
                        [g]:[0,0]
```

```
int g = 0;
                                  t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
   g = 8;
    if(g > 10) {
      error();
    . . .
                        [g]:[0,0]
```

```
int g = 0;
                                  t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
   g = 8;
    if(g > 10) {
      error();
    . . .
                        [g]:[0,8]
```

```
int g = 0;
                                  t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
   g = 8;
    if(g > 10) {
      error();
    . . .
                        [g]:[0,8]
```

```
int g = 0;
                                   t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
   g = 8; \{x \mapsto [3,3]\}
    if(g > 10) {
      error();
    }
    . . .
                        [g]:[0,8]
```

```
int g = 0;
                                   t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
    g = 8; \{x \mapsto [3,3]\}
    if(g > 10) {
      error();
    }
    . . .
```

[g]:[0,8]

```
int g = 0;
                                   t1:
                                       g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                    g = 42;
    g = 8; \{x \mapsto [3,3]\}
    if(g > 10) {
      error();
    }
    . . .
                        [g]:[0,8]
```

```
int g = 0;
                                     t1:
                                          g = 10;
main:
    x = 3; \{x \mapsto [3,3]\}
                                       g = 42;
    g = 8; \{x \mapsto [3,3]\}
    if(g > 10) {
      error();
    }
                    \{x \mapsto [3,3]\}
     . . .
                          [g]:[0,8]
```

```
int g = 0;
                                    t1:
                                         g = 10;
main:
                                         . . .
    x = 3;
                                         g = 42;
    g = 8;
    if(g > 10) {
      error();
    }
    . . .
                         [g]:[0,8]
```

```
int g = 0;
                                    t1:
                                         g = 10;
main:
                                         . . .
    x = 3;
                                         g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
                        [g]:[0,10]
```

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
```

[g]:[0,10]

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
```

[g]:[0,10]

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
```

[g]:[0,42]

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
```

[g]:[0,42]

```
int g = 0;
                                    t1:
                                        g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    . . .
```

[g]:[0,42]

```
int g = 0;
                                    t1:
                                         g = 10;
main:
                                         . . .
    x = 3;
                                        g = 42;
    g = 8;
    if(g > 10) {
      error();
    }
    . . .
                        [g]:[0,42]
```

Zooming out:

- Accumulate information for global variables
- Track local states per program point

,.... Thread-Modular Shape Analysis

Alexey Gotsman University of Cambridge Alexey.Gotsman@cl.cam.ac.uk

Josh Berdine Microsoft Research jib@microsoft.com

Byron Cook Microsoft Research bycook@microsoft.com

Mooly Sagiv* Tel-Aviv University msagiv@post.tau.ac.il

We present the first shope analysis for multithreaded programs that we present use man arouse amonym our minimum control postgrams that avoids the explicit commerciation of execution-interleavings. Our aparounts the expired enumeration or execution-interteavings, turn ap-proach is to automatically infer a resource invariant associated with present is to automatically under a resource uncarrant associated with each lock that describes the part of the beap peaceted by the lock each lock that describes the part of the heap protected by the lock. This allows us to use a sequential shape analysis on each threat. We show that resource invariants of a certain class can be characwer some that resource invariants or a corrain class can be characterized as least fixed points and computed via repeated applications

any given thread, the resource invariant restricts how other threads any given thread, the remember attractants need that a control attraction can interfere with it. If resource invariants are known, analyzing a can interfere with it. If resource invariants are known, analyzing a multithrended program does not require cusmerating uner-paying and on he. and on he. done using a sequential shape analysis. The challenge is to infer the resource invariants.

to infer the resource invariants.

A resource invariant describes two orthogonal kinds of inforis resource invarient describes care contentions among a more mation; it simultaneously curves out the part of the heap proceed mation: it simultaneously carves out the past of the heap protected by the lock and defines the possible shapes that this part can have during amountum structure. There information parameter invariance

-----Thread-Modular Shape Analysis

Josh Berdine

Microsoft Research

Wh@microsoft.com

Byron Cook

Mooly Sagiv* Tel-Aviv University

Alexey Gotsman University of Cambridge Alexey Gotsman@cl.cam.ac.uk

Microsoft Research bycook@microsoft.com msagiv@post.tau.ac.il

We present the first shape analysis for multithreaded programs that we present use may arouse anaryon to manufactured programs that avoids the explicit enumeration of execution-interleavings. Our aparounds the explicit entantertation or execution-interteavings, you ap-proach is to automatically infer a resource invariant associated with present is to automatically under a resource uncarrant associated with each lock that describes the part of the beap peaceted by the lock each tock that describes the part of the heap protected by the lock. This allows us to use a sequential shape analysis on each truck. We show that resource invariants of a certain class can be characwe show that ensuring interesting in a contain stand was no volume, terized as least fixed points and computed via repeated applications

any given thread, the resource invariant restricts how other threads any given thread, the remotive streament tentricis each often a stream can interfere with it. If resource invariants are known, analyzing a can interfere with it. If resource invariants are known, analyzing a maintifered program does not require cusmicrating interleaving as and can be then using a sequential shape analysis. The challenge is A resource invariant describes two orthogonal kinds of inforto infer the resource invariants.

as resource invariant structures was assumption a among matter in simultaneously curves out the part of the heap protected usation: it simultaneously carves out the part of the healt pronected by the lock and defines the possible shapes that this part can have by the lock and defines the possible shapes that this part can have during parameter aparenties. Hence independible parameter incurrence

STATIC ANALYSIS OF RUN-TIME ERRORS IN EMBEDDED REAL-TIME PARALLEL C PROGRAMS

ANTOINE MINÉ

CNRS & École Normale Sundrieure, 45 rue d'Ulm, 75005 Paris, France e-mail address: minefirli ens fr

ABSTRACT. We present a static analysis by Abstract Interpretation to check for run-time errors in parallel and multi-threaded C programs. Following our work on Astrée, we focus on embedded critical programs without recursion nor dynamic memory allocation, but extend the analysis to a static set of threads communicating implicitly through a shared memory and explicitly using a finite set of mutual exclusion locks, and scheduled according to a real-time scheduling policy and fixed priorities. Our method is thread-modular. It is

..... Thread-Modular Shape Analysis

Byron Cook

Microsoft Research

bycook@microsoft.com

Josh Berdine Microsoft Research Wh@microsoft.com

Mooly Sagiv Tel-Aviv University msagiv@post.tau.ac.il

Alexey Gotsman

University of Cambridge

Alexey Gotsman@cl.cam.ac.uk

We present the first shape analysis for multithreaded programs that we present the new sample analysis for manufactories programs that avoids the explicit connectation of execution-interleavings. Our apavoids the exprict enumeration or execution-intersectings. Our ap-proach is to automatically infer a resource invariant associated with proach is to automaticany uner a resource unsariant associated with each lock that describes the part of the heap protected by the lock. each tock that describes me part or me map protected of all closes. This allows us to use a sequential shape analysis on each thread. This answer us to use a sequential snape manyon on costs smeller. We show that resource invariants of a certain class can be characwe show that ensuring investment in a constant cases were to columns trived as least fixed points and compared via repeated applications

any given thread, the resource invariant restricts how other threads can interfere with it. If resource invariants are known, analyzing a can interfere with it. It resource uncartaints are autown, amoryzing a multithreaded program does not require enumerating interfeavings minimeraded program does not require enimerating unerviewings and can be done using a sequential shape analysis. The challenge is A resource invariant describes two orthogonal kinds of inforto infer the resource invariants.

anation: it simultaneously curves out the part of the heap protected to the lock and defines the possible shapes that this part can have by me took and defines the possente unages that this part can have desiren namezons aractricon. Manyos informable, passances incurioses

STATIC ANALYSIS OF RUN-TIME ERRORS IN EMBEDDED REAL-TIME PARALLEL C PROGRAMS

ANTOINE MINÉ

CNRS & Écolo Normalo Survérieure, 45 rms d'Ulm, 75005 Paris, France, e-mail address: minefirli ens fr

ABSTRACT. We present a static analysis by Abstract Interpretation to check for run-time errors in parallel and multi-threaded C programs. Following our work on Astrée, we focus on embedded critical programs without recursion nor dynamic memory allocation, but extend the analysis to a static set of threads communicating implicitly through a shared memory and explicitly using a finite set of mutual exclusion locks, and scheduled according to a real-time scheduling policy and fixed priorities. Our method is thread-modular. It is

Relational Thread-Modular Static Value Analysis by Abstract Interpretation*

Antoine Miné

CNRS & École Normale Supérieure 45, rue d'Ulm 75005 Paris. France mine@di.ens.fr

Abstract. We study thread-modular static analysis by abstract interpretation to infer the values of variables in concurrent programs. We show how to go beyond the state of the art and increase an analysis precision by adding the ability to infer some relational and history-sensitive properties of thread interferences. The fundamental basis of this work is the formalization by abstract interpretation of a rely-guarantee concrete semantics which is thread-modular, constructive, and complete for safety

..... Thread-Modular Shape Analysis

Alexey Gotsman University of Cambridge Alexey Gotsman@cl.cam.ac.uk

Josh Berdine Microsoft Research Wh@microsoft.com

Byron Cook Microsoft Research bycook@microsoft.com

> any given thread, the resource invariant can interfere with it. If resource invarian multithreaded program does not require and can be done using a sequential shap to infer the resource invariants. A resource invariant describes two

We present the first shape analysis for multithreaded programs that we present the new sample analysis for manufactories programs that avoids the explicit connectation of execution-interleavings. Our apavoids the expirit enumeration of execution-instreavings. Dif ap-proach is to automatically infer a resource invariant associated with mation: it simultaneously curves out the proach is to automaticany uner a resource unsariant associated with each lock that describes the part of the heap protected by the lock. by the lock and defines the possible is each tock that describes me part or me map protected of all closes. This allows us to use a sequential shape analysis on each thread. This antows us to use a sequential shape analysis on each tireda.

We show that resource invariants of a certain class can be characduring paperson against the in we show that ensuring investment in a constant cases were to columns trived as least fixed points and compared via repeated applications

Improving Thread-Modular Abstract Interpretation

Michael Schwarz¹, Simmo Saan², Helmut Seidl¹, Kalmer Apinis², Julian Erhard¹, and Vesal Voidani²

> ¹ Technische Universität München, Garching, Germany (n.schwarz, helmut.seidl, julian.erhard)@tum.de ² University of Tartu, Tartu, Estonia. {simmo.saan, kalmer.apinis, vesal.vojdani}@ut.ee

Abstract. We give thread-modular non-relational value analyses as abstructions of a local trace semantics. The semantics as well as the analyses are formulated by means of global invariants and side-effecting constraint systems. We show that a generalization of the analysis provided by the static analyzer Gonlint as well as a natural improvement of Antoine Miné's approach can be obtained as instances of this general scheme. We show that these two analyses are incomparable w.r.t. precision and BEDDED

or run-time ée, we focus cation, but gh a shared d according

> Static Value rpretation*

CNRS & École Normale Supérieure 45, rue d'Ulm 75005 Paris. France mine@di.ens.fr

Abstract. We study thread-modular static analysis by abstract interpretation to infer the values of variables in concurrent programs. We show how to go beyond the state of the art and increase an analysis precision by adding the ability to infer some relational and history-sensitive properties of thread interferences. The fundamental basis of this work is the formalization by abstract interpretation of a rely-guarantee concrete semantics which is thread-modular, constructive, and complete for safety

Thread-Modular Shape Analysis

Alexey Gotsman University of Cambridge Alexey Gotsman@cl.cam.ac.uk

Josh Berdine Microsoft Research Wh@microsoft.com

prosent is to automatically more a remounce dissinating management and cash lock that describes the part of the heap protected by the lock.

This allows us to use a sequential shape analysis on each threat.

THE MINOR IS IN USE A SEQUENCES SHOPE MANAGERY IN CASE DESCRIPTION OF A COPTAIN CLASS CAN be Charac-

we show that remarks increases on a various value was no values.

Serviced as least fixed points and computed via repeated applications.

Byron Cook Microsoft Research bycook@microsoft.com

to infer the resource invariants. A resource invariant describes two mation: it simultaneously curves out the by the lock and defines the possible is

any given thread, the resource invariant can interfere with it. If resource invarian multithreaded program does not require and can be done using a sequential shap We present the first shape analysis for multithreaded programs that we pronoun on new sample analysis for manufactured programs that avoids the explicit commercian of execution-interleavings. Our apavoids the expirit enumeration of execution-unerseavings. Dat ap-proach is to automatically infer a resource invariant associated with

during paneron aracution Mayon b

Improving Thread-Modular

Abstract Interpretation Michael Schwarz¹, Simmo Saan², Helmut Seidl¹, Kalmer Apinis², Julian Erhard¹, and Vesal Voidani²

> 1 Technische Universität München, Garching, Germany (n.schwarz, helmut.seidl, julian.erhard)@tum.de ² University of Tartu, Tartu, Estonia {simmo.saan, kalmer.apinis, vesal.vojdani}@ut.ee

Abstract. We give thread-modular non-relational value analyses as abstructions of a local trace semantics. The semantics as well as the analyses are formulated by means of global invariants and side-effecting constraint systems. We show that a generalization of the analysis provided by the static analyses Gourgest as well as a natural improvement of Antoine line's approach can be obtained as instances of this general scheme. to show that these two analyses are incomparable w.r.t. precision and REDDED

or run-time ie, we focus cation, but gh a shared d according

> Static Value rpretation*

CNRS & École Normale Supérieure 45, rue d'Ulm 75005 Paris. France mine@di.ens.fr

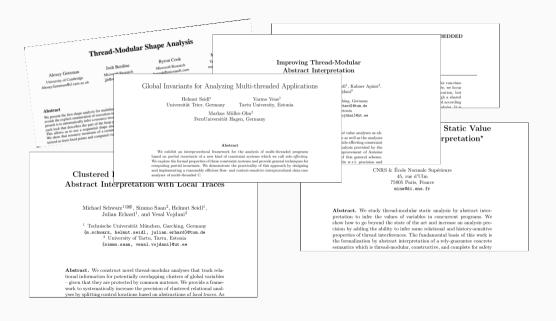
Abstract. We study thread-modular static analysis by abstract interpretation to infer the values of variables in concurrent programs. We show how to go beyond the state of the art and increase an analysis precision by adding the ability to infer some relational and history-sensitive properties of thread interferences. The fundamental basis of this work is the formalization by abstract interpretation of a rely-guarantee concrete semantics which is thread-modular, constructive, and complete for safety

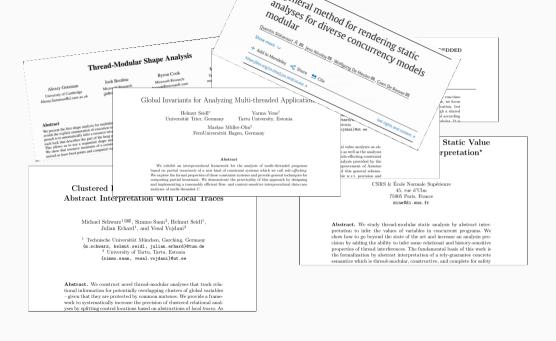
Clustered Relational Thread-Modular Abstract Interpretation with Local Traces

Michael Schwarz¹(⊠), Simmo Saan², Helmut Seidl¹, Julian Erhard¹, and Vesal Voidani²

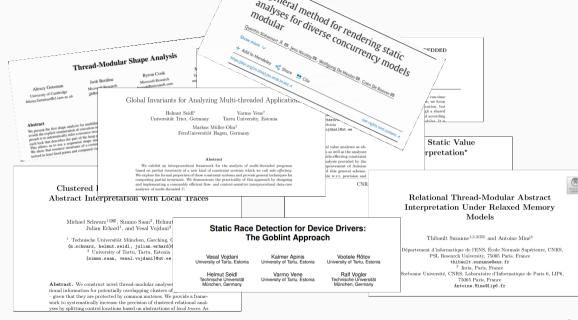
¹ Technische Universität München, Garching, Germany (m.schwarz, helmut.seidl, julian.erhard)@tum.de ² University of Tartu, Tartu, Estonia {simmo.saan, vesal.voidani}@ut.ee

Abstract. We construct novel thread-modular analyses that track relational information for potentially overlapping clusters of global variables given that they are protected by common mutexes. We provide a framework to systematically increase the precision of clustered relational analvses by splitting control locations based on abstractions of local traces. As









Zooming out further

Differentiate:

- Flow-insensitive for some information:
- Flow-sensitive for other information:

Zooming out further

Differentiate:

- Flow-insensitive for some information: Globals
- Flow-sensitive for other information: Locals

Zooming out further

Differentiate:

- Flow-insensitive for some information: Globals
- Flow-sensitive for other information: Locals

Mixed Flow-Sensitivity

Example: Thread-Modular Analysis of Multi-Threaded Programs

• Globals: program globals

• Locals: program points of threads

Further Examples within Goblint

Partial Contexts / Context Lifters

- Globals: Start points of procedures
- Locals: Other program points of procedures

Further Examples within Goblint

Non-Local Control Flow via setjmp/longjmp

- Globals: Targets of longjumps
- Locals: Other program points of procedures

Mixed Flow-Sensitivity

Differentiate:

- Flow-insensitive for some information: Globals
- Flow-sensitive for other information: Locals

Generic concept arising in many settings

Mixed Flow-Sensitivity

Differentiate:

- Flow-insensitive for some information: Globals
- Flow-sensitive for other information: Locals

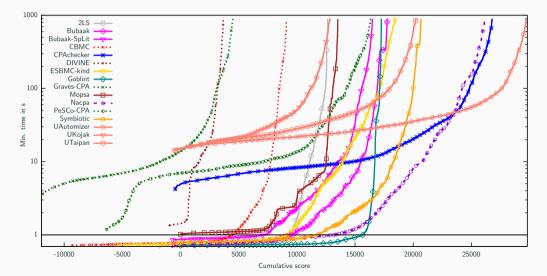
Generic concept arising in many settings

An implementation using side-effecting constraint systems lies at the heart of ${\tt GOBLINT}$

Results in this year's competition

- Best in class for data race freedom
 - \bullet behind metaverifier CoOperace which incorporates Goblint
- Only tool to support all properties without producing any incorrect verdicts
 - (Second year in a row, this needs to change!)

Overall



Ockham Criterion

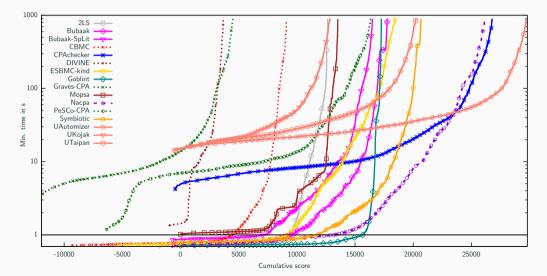
Proposed by [Black and Ribeiro, '16] for SATE V:

- (1) The analyzer's findings are claimed to always be correct.
- (2) It produces findings for most of a program.
- (3) Even one incorrect finding disqualifies an analyzer.

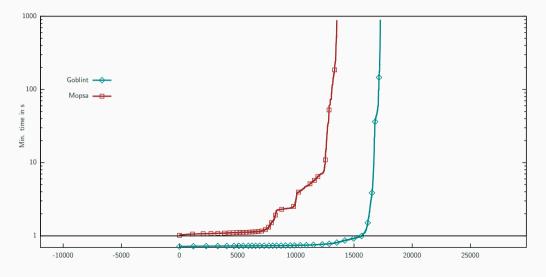


William of Ockham (c. 1287-1347 CE)

Overall



Overall (Ockham)



Thank you!

- Mixed Flow-Sensitivity
 - Flow-sensitive for some unknowns
 - Flow-insensitive for others
- Best contestant for data races
- Best overall score without incorrect verdicts



O/goblint/analyzer



This work was supported in part by Deutsche Forschungsgemeinschaft (DFG) – 378803395/2428 CONVEY.