Longest Common Subsequence

subsequence

 An ordered combination of each member of the sequence

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• Sequence = (w,a,l,k,i,n,g)
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- Subsequence Ex1 = (w,a,l,k) >> (w,a,l,k,i,n,g)
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- Subsequence
$$Ex2 = (k,i,n,g) >> (w,a,l,k,i,m,g)$$

- Subsequence Ex3 =
$$(w,g)$$
 >> (w,a,l,k,i,n,g)

- Subsequence
$$Ex4 = (w,l,n,g) >> (w,a,l,k,i,m,g)$$

The problem

- Given two sequences A,B
 - Find a subsequence s of both A and B such that the length of s is longest

- Example
 - -A = (w,a,l,k,i,n,g)
 - -B = (a,l,i,e,n)
 - Longest Common Subsequence = (a,l,i,n)
 - (a,l,i,n) is a subsequence of A (w,a,l,k,i,m,g)
 - (a,l,i,n) is a subsequence of B (a,l,i,e,m)

Notation

- Let the first index of A and B be 1
 - E.g., A[1] = 'w', A[2] = 'a', A[3] = 'l', ...
- Let |A| = n
- Let |B| = m
- Let A_i be the substring from position 1 to i of A
 - $E.g. A_1 = 'w'$
 - E.g. $A_2 = 'wa'$
 - E.g. $A_5 =$ 'walki'
 - $-A_0 = "$

The sub-problem

- If we wish to know LCS(A,B)
 - Does LCS of (A_x, B_y) helps us?

– What sub problem shall we use?

Think Backward (or forward?)

- If we know LCS(A,B)
- How does it help?
 - i.e., where LCS(A,B) contribute to?
 - Try the very obvious case...
 - Does it help us solve
 - LCS(A + 'c',B + 'c')?
 - Sure!
 - LCS(A + 'c', B + 'c') = LCS(A, B) + 'c'
 - Because they both ends with 'c'
 - E.g. A = 'walking', B = 'alien'
 - What is LCS('walkingC', 'alienC')?
 - alinC

Think Backward (or forward?)

- Any more case to consider?
- If we know LCS(A,B)
 - Does it help us solve
 - LCS(A,B + 'c')?
 - Yes
 - Adding 'c' would have only two outcomes
 - it does not change the LCS
 - \rightarrow So LCS(A,B + 'c') = LCS(A,B)
 - It does change the LCS
 - » So LCS(A,B + 'c') = something ending with 'c'
 - » To be continue...

Using LCS(A,B + 'c')

The case that LCS is changed

Notice that, in this case, what comes after both B and A_{k-1} is 'c'

- Is that possible?
- Yes, when there are 'c' in A that comes after LCS(A,B)
 - Assume that that point is at A[k] (hence, A[k] = c)
 - LCS(A,B + 'c') would be LCS(A_{k-1} ,B) + 'c'
- Check that $LCS(A_{k-1},B)$ + 'c' is actually LCS(A,B + 'c')
 - LCS(A_{k-1} ,B +'c') will be the same as
 - $LCS(A_k, B+'c')$
 - $LCS(A_{k+1},B+'c')$
 - $-LCS(A_{k+2},B+c)$
 - **–** ...
 - LCS(A,B+'c')

This means that $LCS(A,B + 'c') = LCS(A_{n-1},B+'c')$.. So, LCS(A,B) **does not contribute** to LCS(A,B + 'c')

Think Backward (or forward?)

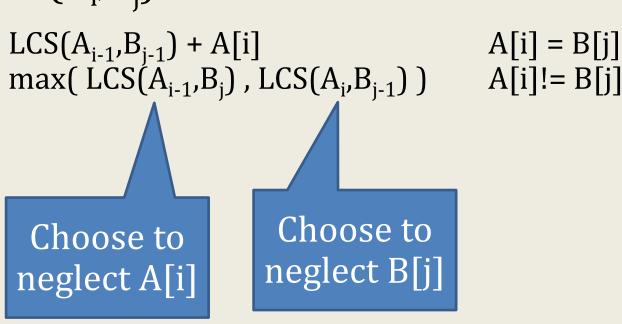
- The remaining case
- If we know LCS(A,B)
 - Does it help us solve
 - LCS(A + 'c',B)?
 - Yes
 - Similar to the case of A,B + 'c'
- In conclusion, this means that
 - $LCS(A + 'c',B) = LCS(A+'c',B_{n-1})$

conclusion

- LCS(A,B) will be LCS(A,B + 'c') when 'c' does not constitute the longer common subsequence
 - If 'c' is in the longer common subsequence
 - LCS(A,B + 'c') will be $LCS(A_{n-1},B + 'c')$ instead!!!!
 - Not our case
- So, backwardly,
 - LCS(A,B) is either
 - LCS(A_{n-1},B) - Or
 - LCS(A,B_{m-1})
 - Just select the longer one!!!!

Recurrence

• $LCS(A_i, B_j) =$

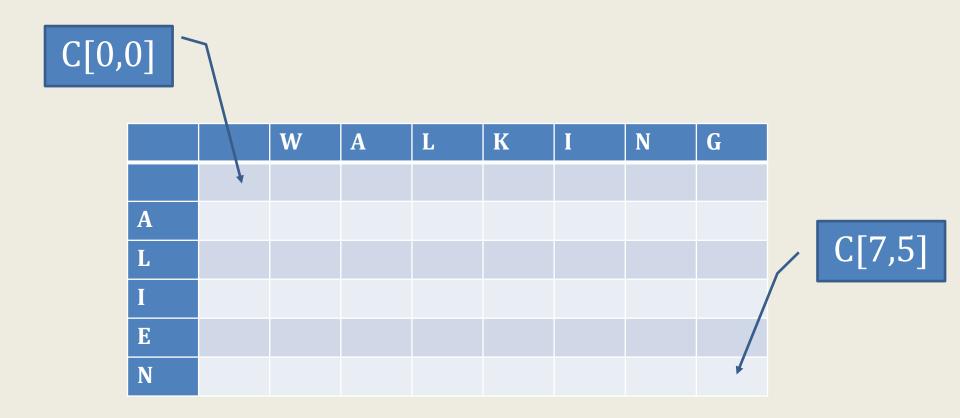


Solution to the LCS

- Simplify problem
 - To find the length of LCS

Let c(i,j) be the length of LCS(A_i,B_j)

$$c(i,j) = \begin{cases} 0 & \text{if } i=0 \text{ or } j=0 \\ c(i-1,j-1)+1 & \text{if } i>0, j>0 \text{ and } A[i]=B[j] \\ \max(c(i-1,j),c(i,j-1)) & \text{if } i>0, j>0 \text{ and } A[i]!=B[j] \end{cases}$$



$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

if i=0 or j=0 if i>0, j>0 and A[i] = B[j] if i>0, j>0 and A[i]!= B[j]

Fill the trivial case

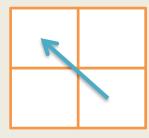
		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0							
L	0							
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$



$$A[i] != b[j]$$



$$A[i] = b[j]$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0						
L	0							
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	Q	0	0	0	0	0	0
A	0	0						
L	0							
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1					
L	0							
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1				
L	0							
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1			
L	0							
Ι	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1		
L	0							
Ι	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	Ι	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	ل ا
L	0							
Ι	0							
E	0							
N	0							

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0							
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0						
I	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1					
Ι	0							
E	0							
N	0							

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1				
Ι	0							
E	0							
N	0							

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	ل ا		
Ι	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2		
Ι	0							
E	0							
N	0							

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

if i=0 or j=0 if i>0, j>0 and A[i] = B[j] if i>0, j>0 and A[i]!= B[j]

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	← Ĭ
Ι	0							
E	0							
N	0							

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
Ι	0							
E	0							
N	0							

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
Ι	0						2222	
E	0	-						-
N	0	-						>

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
N	0	0	1	2	2	3	3+1	4

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

if i=0 or j=0 if i>0, j>0 and A[i] = B[j] if i>0, j>0 and A[i]!= B[j]

Recovering the Actual Solution

 We know particulality which case c(i, j) is from

$$c(i,j) = \begin{cases} 0 & \text{if } I > 0, j > 0 \\ c(i-1,j-1) + 1 & \text{if } I > 0, j > 0 \text{ and } A[i] = B[j] \\ \max(c(i-1,j), c(i,j-1)) & \text{if } I > 0, j > 0 \text{ and } A[i]! = B[j] \end{cases}$$

• If it is the second case, it simply means that A[i] is the last member in LCS

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
N	0	0	1	2	2	3	3+1	4

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
N	0	0	1	2	2	3	3+1	4

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
N	0	0	1	2	2	3	3+1	4

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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
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$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
Ι	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
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		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
Ι	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
N	0	0	1	2	2	3	3+1	4

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$

		W	A	L	K	I	N	G
	0	0	0	0	0	0	0	0
A	0	0	0+1	1	1	1	1	1
L	0	0	1	1+1	2	2	2	2
I	0	0	1	2	2	2+1	3	3
E	0	0	1	2	2	3	3	3
N	0	0	1	2	2	3	3+1	4

$$c(i,j) = \begin{cases} 0 \\ c(i-1,j-1) + 1 \\ max(c(i-1,j),c(i,j-1)) \end{cases}$$