

ICCS310: Assignment 3  
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**1: NFA vs DFA Expressiveness**

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**(1)**

Let construct an NFA,  $M = (Q, \Sigma, \delta, q_0, F)$  where  $Q = \{0, 1, \dots, k\}$ . Let  $\delta(0, b) = 0, \delta(0, 1) = \{0, a\} = \{0, 1\}$  and  $\delta(i-1, a) = i$ , for  $2 \leq i \leq k$ . Then set  $q_0 = 0$  and  $F = \{k\}$ . We know that the machine will start at state 0 (starting state). When the machine locate an  $a$  it wil guess that it is a  $k$ th character to the right and will move to state 1. When it reaches state  $k$ , it will only accept if there are exactly  $k-1$  bits following the one that move from  $b$  to  $a$ .

**(2)**

Let the input have  $k$  character. We know that the characters can be either  $a$  or  $b$ . Let  $x$  and  $y$  be a string with  $k$  bit such that  $x, y \in \Sigma^*$  and that  $|x| = |y| = k$ . Let  $i$  be a position such that  $x_i \neq y_i$ . Hence either  $x$  or  $y$  contain an  $a$  at the  $i$ th position. Let  $z = b^{i-1}$ , then  $z$  distinguish  $x$  and  $y$  as one of the  $xz$  and  $yz$  has an  $a$  at  $k$ th postion from the right. Since there are  $2^k$  string of length  $k$  that are all distinguishable from the above prove, a DFA that accept this language need to have  $2^k$  states.

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**4: HackerRank Challenge**

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