## Assignment 4

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1. Answer each part TRUE or FALSE

a. 
$$2^{\log^2 n} \in \mathcal{O}(n^k)$$

d. 
$$n^k \in o(2^{\log n})$$

b. 
$$n! \in 2^{\mathcal{O}(n \log n)}$$

e. 
$$2^n \in o(n!)$$

c. 
$$3^n \in \mathcal{O}(2^{n \log n})$$

f. 
$$\frac{1}{n} \in o(1)$$

## Solution:

a. 
$$2^{\log^2 n} \in \mathcal{O}(n^k)$$

True

Proof:  $2^{\log^2 n} = n$  and  $n \in \mathcal{O}(n^k)$ 

b. 
$$n! \in 2^{\mathcal{O}(n \log n)}$$

c. 
$$3^n \in \mathcal{O}(2^{n \log n})$$

False

As  $3^n$  grows much faster than  $2^{n \log n}$ 

d. 
$$n^k \in o(2^{\log n})$$
 True

- e.  $2^n \in o(n!)$
- f.  $\frac{1}{n} \in o(1)$  True as 1/n is strictly smaller than 1 after n>1
- 2. Show that P is closed under union, complement, concatenation, and star.
- 3. Show that NP is closed under union, concatenation, and star
- 4. We normally assume natural numbers are represented in binary, such that a number  $n \in \mathcal{N}$  is represented by the string  $b_{\lfloor \log n \rfloor} b_{\lfloor \log n \rfloor 1} \dots b_0, b_i \in \{0, 1\}$ , and  $n = \sum_{i=0}^{\lfloor \log n \rfloor} b_i 2^i$ . We could also write a number in unary, where a number  $n \in \mathcal{N}$  is represented by n consecutive 1s. The problem of factoring a number in binary is not known to be in P, but what if the number is given in unary? Prove your answer.
- 5. Suppose the number k and the graph G are given, and we want to know if there exists a subset S of size k from the vertices of G such that there is no edge between them in G. Prove that this problem is NP-Complete.

6. Show the following language is NP-Complete:

$$DOUBLE - SAT = \{ \langle \phi \rangle | \phi \text{ has at least 2 satisfying assignments } \}$$

7. Let S be a set and let C be a collection of subsets of S. A set  $S' \subseteq S$  is called a set hitting set for C if every subset in C contains at least an element in S'. Let

$$HITSET = \{\langle C, k \rangle | C \text{ has a hitting set of size } k \}$$

Prove that HITSET is NP-Complete.

- 8. a. Prove that NP = coNP iff there is an NP-Complete problem in coNP.
  - b. Show that if  $coNP \neq NP$  then  $P \neq NP$ .
- 9. Show that P is closed under homomorphism iff P = NP
- 10. Let  $CNF_k = \{\langle \phi \rangle | \phi \text{ is a satisfiable cnf-formula where each variable appears in at most } k \text{ places } \}.$ 
  - a. Show that  $CNF_2 \in P$
  - b. Show that  $CNF_3$  is NP-complete