ASSIGNMENT-1 CH. Venkata Navameells.

Describe the concept of Abstract data type (ADT) and how they differ from concrete data structures. Design on ADT for a clack and implement it using arrays and linked lest in c. Include operations leke purh, pop , peck, isempty is fall and part.

Abstract data type (ADT):

An abstract data type (ADT) is a theoretical model that defines a set of operations and the semantice (behaviour) of those Operations on a data structure, without specifing how the data structure should be implemented. It provides a high level description of what operations can be performed on the data and what constraints apply to those operations.

Characteristics of ADT :

operations: - Defines a set of operations - that can be performed on the data structure.

Semantics :- spenfies the behaviour of each operation.

Encapsulation: - Hødes the implementation details.

focuseing on the interface provided to user.

ADT for stack:

A stack is a fundamental data structure that follows the last In, first out (LIFO) principle.

It supports the following operations:

> push: - Adds an element to the top of the stack
> pop - Removes and returns the element from the
top of stack.

Peck: - Returns the element from the top of the Stack without removing it. is empty: - Checks if the stack is empty. isfull: - Checks of the stack is full. Concrete data structures: The implementations using arrays and linked list ADT focusses on the operations and their behaviour course concrete data structures focus on how those operations are realized using specific programming control. Advantages of ADT: By separating the ADT from Its implementation, you ackieve modulanty, encapsulation and fleribility en designing and using dular structures in program. This Separation allows foir easier maintenance, code reuse, and abstraction of the complex operations. Implementation in c using linked list # molude 2 std: 0. h> =# include 2 std 90. b> tope det struct Node of ent dater ;
struct Node* nent; 3 Node 3 ent main () E Node top = Null; Node * new node = (node *) molloc (size of (node));

of (now now = + cold) } of " " money withing failed to " There . James rock - State - De me me - mit - mi to - morale mens note - (note #) mostore (high of (nous 1) \$ et (new made + + Nall) print + (premary allocation fooled 1). faction 13 area rade - tata = 30 3 new note -> rest = top; top - new rocks of (top ! + 10 in (1) } prat (Top charit of & la top addit)) t else { point ("stack is empty : 10"); of (top == Noll) } note temp = top; print (paped charat: "tow", temp-lite);

Implementation in cusing array's # molude < stdio.b> # define max - size 100 type of Struct & int items [MAX - size] 3 int top s I stack array; ent main () qu Stack orray Stack; stack. top =-1; Stack : "tems [++ stack . top] = 10; Stack: items [++ stack. top]=20; Stack. items [++ stack -top] = 30; Ef (stack . top ! = -1) { print f (" top element: 1/2 d/m", stack. items (stack. top)); I cheef printf ("stack is empty ! 10"); if (stock top! = -1) } prønt (" popped element: /od/n", stack. Items [stack top-]); 4 else of printf ("stack underflow ! | n"); of (stack top! = -1) { printf ("popped element: "/. d/or", stack: "tems [stack.top--]); 4 else f

top = top = mext : It free (temp) s 3 else 2 printf (" stack underflow: \n")) cokele (top! = Null) { node* temp = top;

top = top > ment;

free (temp); I return on the first of the and properties of the conceptor of the The university amnounced the selected candiates register number for placement training. The student xxx, reg mo. 20142010 coishes to Check wheather his name is listed or not. The list is not sorted in any order. Identify the searching techniq that can be applied and explain the searching steps with the suitable procedure. Linear Search : 1000 Linear Search works by decreasing each element in the last one by one until the desired element is found or the end of the list is reached. It's a Simple Searching technique that doesn't require only prior sorting of the data. steps for linear Search; 1. Starts from the first element. a. Check if the current element is equal to the target element 3. If the current element is not target, move to next.

A. Continue this pracess until either the target element is found or you reach the end. 5. If the target is found, return its position. If the end of the lest is reached. Procedure :-Given the lest: 2014 2015, 20142033, 20142011, 20142017, 2014 2010, 20142056, 20142003 1. Start at the first element of the lest. 2. Compare 20142010" with "20142015" 3- compare "2014 2010" with "20142010" They one equal. The element "2014 2010" is found at the fifth element C code for linear Search: # include 2 std io. h> ent main () of int reg numbers [] = {2014 2015, 2014 2033, 2014 2011, 2014 2017 20142056, 201420037 int target = 2014 2010; ent n= size of (reg numbers)/-size of (reg numbers) int found = 0; for (9=0;9<7) \$ if (regnumbers [] = target) of prent (" Regestration number of de found at ender found = 1; 3 y break 3

Explanation of the code?—

Explanation of the code?—

The reg numbers array contains—the list of registration 2. "target" is the registration number use are sparching 3. Iterate through each element matches the target. 4. If the loop completes with out finding the target. 5. The program will print the index of the found regestration number or entialize that the registration ce not present. Write pseudo code for stack operations. 1. Intéalize stack Operations () Intéalize necessary variable or structures to respond the stack. 2. push (elements): if stack is full? print " stack overflow" add element to the top of the stack increment top pointer. 3. pop (): if stack is empty: print ("stack under flow") return null (or appropriate error value) remove and return element from the top of the

4) peck (): if "stack is empty", print a stack is empty". return mull (or appropriate error value) return element at the top of the stack 5) is empty () return true if top 95 -1 (stack the empty) return true, ef top is each to man nige -1 6) if full: Other wisk return false. Pseudo code: - Institulize the necessary variables of data structures to resprésent à stack. -> Adds am element to the top of the stack check of the stack is full. -> Removes and returns the element from the top of the stack. checks of the stack. -> Checks of the stack of empty by inspecting the top pointer or equipment variable. -> Chees if the stack is full by comparing the top pointer or equivalent variable.