1.4 ARCHITECTURAL CLASSIFICATION SCHEMES

Three computer architectural classification schemes are presented in this section. Flynn's classification (1966) is based on the multiplicity of instruction streams and data streams in a computer system. Feng's scheme (1972) is based on serial versus parallel processing. Händler's classification (1977) is determined by the degree of parallelism and pipelining in various subsystem levels.

1.4.1 Multiplicity of Instruction-Data Streams

In general, digital computers may be classified into four categories, according to the multiplicity of instruction and data streams. This scheme for classifying computer organizations was introduced by Michael J. Flynn. The essential computing process is the execution of a sequence of instructions on a set of data. The term *stream* is used here to denote a sequence of items (instructions or data) as executed or operated upon by a single processor. *Instructions* or data are defined with respect to a referenced machine. An *instruction stream* is a sequence of instructions as executed by the machine; a data stream is a sequence of data including input, partial, or temporary results, called for by the instruction stream.

Computer organizations are characterized by the multiplicity of the hardware provided to service the instruction and data streams. Listed below are Flynn's four machine organizations:

- Single instruction stream-single data stream (SISD)
- Single instruction stream-multiple data stream (SIMD)
- Multiple instruction stream-single data stream (MISD)
- Multiple instruction stream-multiple data stream (MIMD)

These organizational classes are illustrated by the block diagrams in Figure 1.16. The categorization depends on the multiplicity of simultaneous events in the system components. Conceptually, only three types of system components are needed in the illustration. Both instructions and data are fetched from the memory modules. Instructions are decoded by the control unit, which sends the decoded instruction stream to the processor units for execution. Data streams flow between the processors and the memory bidirectionally. Multiple memory modules may be used in the shared memory subsystem. Each instruction stream is generated by an independent control unit. Multiple data streams originate from the subsystem of shared memory modules. I/O facilities are not shown in these simplified block diagrams.

SISD computer organization This organization, shown in Figure 1.16a, represents most serial computers available today. Instructions are executed sequentially but may be overlapped in their execution stages (pipelining). Most SISD uniprocessor systems are pipelined. An SISD computer may have more than one functional unit in it. All the functional units are under the supervision of one control unit.