



**Mid semester examination – introduction To AI**

**Name – Naveen kumar**

**Branch- CSE (AIMI)**

**Section – B**

**Roll no.- 48**

**University roll no.- 202401100400121**

**Problem statement :- N-Queens Problem.**

## **INTRODUCTION :-**

### **Problem Definition:**

*The N-Queens Problem is a classic problem where you are asked to place **N queens** on a **N x N chessboard** in such a way that no two queens can attack each other.*

### **Key Points:**

#### **1. Chessboard:**

- *The chessboard is a square grid of size **N x N**. This means the board has **N rows** and **N columns**.*

#### **2. Queens:**

- *A queen in chess can attack another queen in three possible ways:*
  - **Same row:** *Queens in the same row can attack each other.*
  - **Same column:** *Queens in the same column can attack each other.*

#### **3. Objective:**

- *The goal is to place **exactly N queens** on the chessboard, where each queen is placed in a different row and a different column, such that no two queens threaten each*

*other. In other words, there should be no queens on the same row, column, or diagonal.*



**Constraints:**

- **Row and Column Constraint:** Each queen must be placed in a different row and a different column.
- **Diagonal Constraint:** No two queens can share the same diagonal. The two main diagonals on a chessboard are:
  - The **main diagonal** (top-left to bottom-right), where the difference between the row and column indices is constant.
  - The **anti-diagonal** (top-right to bottom-left), where the sum of the row and column indices is constant.

**Solution Approach:**

*To solve this problem, a common strategy is **backtracking**, where you explore all possible placements of queens row by row and backtrack when you hit a dead-end (i.e., a situation where no valid placement can be made for the next queen).*

- **Backtracking:** You place a queen in a valid position, then move on to the next row. If you can't place a queen in any column of the next row (because of conflicts), you backtrack and move the previous queen to a new position.

***Steps in the Methodology:***

### **1. Starting with an empty board:**

- *Imagine an empty  $N \times N$  chessboard. Initially, no queens are placed.*

### **2. Placing queens row by row:**

- *Start by placing a queen in the first row.*
- *For each row, attempt to place a queen in every column, one by one. For each column:*
  - *Check whether placing the queen in the current column will result in a **valid configuration***

### **3. Checking for conflicts:**

- *A queen can attack another queen if they are:*
  - *In the **same column**.*
  - *On the **same diagonal***
  - *In the **same row***

### **4. Recursive step:**

- *If placing a queen in the current column of the current row is **safe**, move to the next row and try placing a queen there.*
- *Repeat the process for each row until all queens are placed on the board.*

### **5. Backtracking step:**

- *If you reach a row where no valid column for the queen exists, **backtrack**:*
  - *Remove the queen from the previous row*

- *Try placing the queen in the next column in the previous row.*
- *This ensures that you explore all possible configurations without skipping any valid solutions.*

## **6. Termination condition:**

- *If all queens are placed successfully (i.e., you've filled all rows), then you've found a valid solution.*
- *If you backtrack all the way to the first row and cannot find a valid placement, then there is no solution for that configuration.*

## **Typed code :**

```
import random

# Count conflicts for each queen's position
def count_ conflicts(board, n):

    return [sum(board[i] == board[j] or abs(board[i] - board[j]) == abs(i - j) for j in range(n) if i != j) for i in range(n)]

# Find optimal column for a queen in a given row
def best_ column(board, row, n):

    options = [(col, sum(count_ conflicts(board[:row] + [col] + board[row+1:], n))) for col in range(n) if col != board[row]]

    return min(options, key=lambda x: x[1])[0] if options else board[row]

# Solve N-Queens using local search
def solve _n_ queens(n, retries=10, steps=1000):

    for _ in range(retries):

        # Start with a random board

        board = [random . randint (0, n - 1) for _ in range(n)]
```

```

for _ in range(steps):
    conflicts = count_conflicts(board, n)

    # If no conflicts, return the solution
    if not sum(conflicts):
        return board

    # Select a row with the maximum conflicts
    row = random.choice([i for i in range(n) if conflicts[i] == max(conflicts)])

    # Move the queen in that row to the best column
    board[row] = best_column(board, row, n)

return None

# Display the chessboard
def render_board(board):
    print("No solution" if not board else "\n".join(" ".join('Q' if board[r] == c else '.' for c in range(len(board))) for
r in range(len(board))))

# Main execution
if __name__ == "__main__": # Corrected __name__ condition
    for n in map(int, input("Enter board sizes: ").split()):
        print(f"\n Solving for N={n}")
        render_board(solve_n_queens(n))

```

○ ***Output code images :-***

The screenshot shows a Jupyter Notebook titled "Naveen\_202401100400121.ipynb". The interface includes a top menu bar with "File", "Edit", "View", "Insert", "Runtime", "Tools", and "Help". Below the menu is a toolbar with "Commands", "+ Code", and "+ Text" buttons. The notebook content consists of a code cell that has been executed, indicated by a green checkmark and "3s" in the left margin. The code cell contains the following text:

```
Enter board sizes: 5 6

Solving for N=5
. . . . Q
. . Q . .
Q . . . .
. . . Q .
. Q . . .

Solving for N=6
. . Q . . .
. . . . . Q
. Q . . . .
. . . . Q .
Q . . . . .
. . . Q . .
```

At the bottom of the notebook, a status bar shows a green checkmark, "3s", and "completed at 3:05PM".

## ***References :***

***1: Chat GPT.***

***2:CLASS NOTES .***

***3:WIKIPEDIA.***