

# Assignment 3

● Graded

## Group

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## Total Points

65 / 70 pts

## Question 1

Team name

0 / 0 pts

✓ + 0 pts Correct

+ 0 pts Incorrect

## Question 2

Commands

10 / 10 pts

✓ + 10 pts go/enter, climb/enter, pluck/pick,back/climb,give,back,back, thrnxtzy, read or correct combination of the above ("c" and "put" can be ignored).

+ 0 pts Incorrect

## Question 3

Analysis

 45 / 50 pts

✓ + 15 pts Finding at least two distinct powers of  $g$ .

✓ + 25 pts Finding the values of  $g$  by repeated division or Extended Euclid's algorithm or any other method.

✓ + 5 pts The value of  $g$  is 192847283928500239481729

✓ + 5 pts Finding *password* using the information of  $g$ .

+ 0 pts Wrong answer or NA.

+ 50 pts Solving the assignment using an entirely different approach.

💬 - 5 pts you should have compiled it carefully, it is not readable

## Question 4

Password

10 / 10 pts

+ 0 pts Incorrect

✓ + 10 pts 3608528850368400786036725

Question 5

Codes

0 / 0 pts

✓ + 0 pts Correct

### Q1 Team name

0 Points

NAA

### Q2 Commands

10 Points

List the commands used in the game to reach the ciphertext.

enter - enter - pluck -climb - give -  
(take magic words thrnxtzy) - back -  
back - ( thrnxtzy) - read

### Q3 Analysis

50 Points

Give a detailed analysis of how you figured out the password? (Explain in less than 500 words)

```
\documentclass[a4paper]{article}

\usepackage[english]{babel}
\usepackage[utf8]{inputenc}
\usepackage{algorithm}
\usepackage[noend]{algpseudocode}

\title{Assignment 3}

\author{NAA}

\date{\today}
\def\changemargin#1#2{\list{}{\rightmargin#2\leftmargin#1}\item[]}
\let\endchangemargin=\endlist
\begin{document}
\maketitle
\section{ Theory}
\begin{itemize}
\item  $Z_{p^*}$  is a multiplicative group as  $p$  is an prime number.
\item the group action on  $Z_{p^*}$  is denoted by  $*_p$  which is defined by:
\begin{center}
\hspace{0cm}  $*_p \colon G \times G \rightarrow G$  \\
\hspace{4cm}  $(x,y) \mapsto x*y \pmod p$ 
\end{center}
\item We are going to use the fact that every element in the group has a
\textbf{unique} inverse.
\item Also as  $Z_{p^*}$  is a group  $*_p$  is \textbf{associative} .
\item We use \textbf{Fermat's little theorem} to calculate the inverse which
states that if  $p$  is a prime number, then for any integer  $a$ , the number  $a^p - a$ 
is an integer multiple of  $p$ . If  $a$  is not divisible by  $p$ , Fermat's little theorem is
equivalent to the statement that  $a^{p-1} - 1$  is an integer multiple of  $p$ .
\begin{center}
 $a^{p-1} \equiv 1 \pmod p$  \\
 $a*a^{p-2} \equiv 1 \pmod p$ 
\end{center}
\end{itemize}
```

```

\end{center}
therefore  $a^{-1} = a^{p-2} \pmod p$ 
\item Next we are going to use an recursive method to find g.
\end{itemize}
\section{Initial calculation}

\begin{center}
\begin{tabular}{|c|c|}
\hline
a & password  $*_p g^a \pmod p$  \\
\hline
324 & 11226815350263531814963336315 \\
\hline
2345 & 9190548667900274300830391220 \\
\hline
9513 & 4138652629655613570819000497 \\
\hline
\end{tabular}
\end{center}

\newline Let password  $*_p g^{3^2^4} = x_1$ 
\newline Therefore we can substitute this variable in other two equation. (We
know the inverse of  $(x_1)^{-1} = 11676372716222599136085566753$ )
\begin{enumerate}
\item  $(x_1) *_p g^{2^0^2^1} = 7021284369301638640577066679$ 
\item  $(x_1) *_p g^{9^1^8^9} = 9190548667900274300830391220$ 
\end{enumerate}
\newline Multiplying by inverse of  $(x_1)$  both sides give:
\begin{enumerate}
\item  $g^{2^0^2^1} = 3426347385144995225825016781$ 
\item  $g^{9^1^8^9} = 4138652629655613570819000497$ 
\end{enumerate}
\section{Algorithm}
\heading Used this algorithm while solving the equation  $g^a = x_1$  and
 $g^b = x_2$  st  $a > b$ 
\begin{itemize}
\item Both  $x_1, x_2 \in G$  where  $G$  is the multiplicative group.
\item  $*_p$  is the group operation that is  $*_p(x_1, x_2) \rightarrow x_1 *_p x_2 \pmod p$ 
\newline If  $a \in G$  then  $a^x$  implies  $a *_p a *_p a \dots$  x times a also
 $a^x \in G$ 
\end{itemize}

```

```

\begin{algorithm}
\caption{Algorithm}
\begin{algorithmic}
\Procedure{}{$a,b$}\Comment{Use recursive method to solve }
\State $rem$ \gets $a\bmod b$
\State $Q$ \gets [$a/b$]\Comment{[x] the greatest integer less than x}
\While{$rem\neq 0$}\Comment{Check if remainder is not zero}
\State $g^{r^e^m}$ \gets $g^{a*_p((g^b)^{-1})^Q}$ \Comment{ $*_p$ is the group
operation }
\State $a$ \gets $b$
\State $b$ \gets $rem$
\State $rem$ \gets $a\bmod b$
\State $Q$ \gets [$a/b$]
\EndWhile\label{euclidendwhile}
\State \textbf{return} $(b,g^b)$\Comment{the function will return
$(rem,g^{r^e^m})$}
\EndProcedure
\end{algorithmic}
\end{algorithm}

```

Let's do the first step of this algorithm explicitly to get the idea

```

\begin{center}
\hspace{-4.7cm}We have:\newline
$g^{9^{1^8^9}}=3426347385144995225825016781$
\newline
$g^{2^{0^2^1}}=3426347385144995225825016781$
\newline
\newline We can write {9189= 4*2021 + 1105} [the rem $\gets$ 9189 $\bmod$
2021 = 1105]
\begin{center}
\item $\rightarrow g^{9^{1^8^9}}=(g^{2^{0^2^1}})^{4*_pg^{1^1^0^5}}$
\item $\rightarrow ((g^{2^{0^2^1}})^4)^{-1}*_pg^{9^{1^8^9}}=g^{1^1^0^5}$
\end{center}

```

```

\end{center}
\hspace{.5cm}Therefore we will be able to get the value of $g^{1^1^0^5}$
\section{ Important notes on algorithm}
\begin{itemize}
\item The multiplication used in the above algorithm is the Group operation
($*_p$)
\item If we know any element the group we can easily take out it's inverse.

```

Hence all the values can be computed for the steps of the algorithm (We know  $x_1, x_2$ )

\end{itemize}

\section{Implementing the Algorithm in our case }

\begin{center}

\begin{tabular}{|c c c c|}

\hline

a & b & Q & rem \\ [0.5ex]

\hline\hline

9189 & 2021 & 4 & 1105 \\

\hline

2021 & 1105 & 1 & 916 \\

\hline

1105 & 916 & 1 & 189 \\

\hline

916 & 189 & 4 & 160 \\

\hline

189 & 160 & 1 & 29 \\

\hline

160 & 29 & 5 & 15 \\

\hline

29 & 15 & 1 & 14 \\

\hline

15 & 14 & 1 & 1 \\ [1ex]

\hline

\end{tabular} \hspace{-3.5cm}

\newline

\begin{tabular}{|c c|}

\hline

a &  $g^a$  \\

\hline\hline

9189 & 3426347385144995225825016781 \\

\hline

2021 & 7021284369301638640577066679 \\

\hline

1105 & 1332524359715193692493602650 \\

\hline

916 & 16928329349929603757418032233 \\

\hline

```

189 & 7233340894988383169873081319 \\
\hline
160 & 15480832131739101784049259744 \\
\hline
29 & 14409628835368808838382787765\\
\hline
15 & 9862566087568179051837025782\\
\hline
14 & 11662011900497299711580345247\\
\hline
1 & 192847283928500239481729\\
[1ex]
\hline
\end{tabular}

\end{center}

\section{ Finally solving the equation }
We have got the value of g now we will use : \vspace{.125cm}
\newline
password $_p$ $ g^3^2^4$=11226815350263531814963336315
\vspace{.125cm}
\newline password= 11226815350263531814963336315 $_p$ $ (g^3^2^4)^-
^1$ \vspace{.125cm}
\newline password= 11226815350263531814963336315 $_p$
726117032386935245054894092 \vspace{.125cm}
\newline password= 3608528850368400786036725

\end{document}

```

#### Q4 Password

10 Points

What was the final command used to clear this level?

3608528850368400786036725



## Q5 Codes

0 Points

Upload any code that you have used to solve this level.

```
1  def M(a,b,mod):
2      res=1
3      while(b>0):
4          if(b%2==1):
5              res=(res*a)%mod
6              a=((a%mod) *(a%mod))
7              b=b//2
8      return res
9
10 def I(a,mod):
11     return M(a, mod-2, mod)
12
13
14 # print(I(324,19807040628566084398385987581))
15 ans = 324*I(324,19807040628566084398385987581)
16 # print(ans//19807040628566084398385987581 ,
17     ans%19807040628566084398385987581)
18
19 mod=19807040628566084398385987581
20 print()
21
22
23 # 11226815350263531814963336315
24 g9189=4138652629655613570819000497*I(11226815350263531814963336315,1980704
25 print("g9189      :",g9189%mod)
26 print()
27
28
29 g2021=I(11226815350263531814963336315, mod)*9190548667900274300830391220
30 print("g2021      :",g2021%mod)
31 print()
32
33 g2021_4_num = (((g2021)**4,mod)) * (g9189))%mod
34 # print("g2021_4_num :",g2021_4_num%mod)
35 # print()
36
37 g1105=g2021_4_num
38 print("g1105      :",g1105%mod)
39 print()
40
41 g916= (g2021*(I(g1105,mod)))%mod
42 print("g916       :",g916%mod)
43 print()
44
45
```

```
46 g189 = ((g1105)*(I(g916,mod)))%mod
47 print("g189      :",g189)
48 print()
49
50
51 g160= (I((g189**4),mod) * g916 )%mod
52 print("g160      :",g160%mod)
53 print()
54
55
56 g29= (I(g160, mod) * g189 )%mod
57 print("g29       :",g29%mod)
58 print()
59
60
61 g15 = (I(g29**5, mod)*g160) %mod
62 print("g15       :",g15%mod)
63 print()
64
65
66 g14 = (I(g15, mod)*g29) %mod
67 print("g14       :",g14%mod)
68 print()
69
70
71 g1 = (I(g14, mod)*g15) % mod
72 print("g1        :",g1%mod)
73 print()
74
75 inverse_g = I(g1, mod)
76 print("g : ",inverse_g)
77 print()
78
79 # password =11226815350263531814963336315//(I(g1**324,mod))
80
81 tmp= inverse_g
82
83
84 for i in range(1,324):
85     inverse_g = (inverse_g * tmp) % mod
86
87 password = (inverse_g * 11226815350263531814963336315 )% mod
88
89
90 print("password  :",password)
91
92
93
94
```

95	# enter - enter - pluck -climb - give - (take magic words ) - back - back - (magic word ) - read - pasword ( 3608528850368400786036725)
96	
97	# (324, 11226815350263531814963336315)
98	# (2345,9190548667900274300830391220)
99	# (9513, 4138652629655613570819000497)

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