

$$1a) 5n^3 + 2n^2 + 3n \leq 5n^3 + 2n^2 + 3n^2 \text{ for } n \geq 1$$

$$5n^3 + 2n^2 + 3n^2 = 5n^3 + 5n^2$$

$$5n^3 + 5n^2 \leq 5n^3 + 5n^3 \text{ for } n \geq 1$$

$$5n^3 + 5n^3 = 10n^3$$

$$5n^3 + 2n^2 + 3n \leq 10n^3 \text{ for } n \geq 1$$

Thus, we have found a c and n_0 where $cn^3 \geq 5n^3 + 2n^2 + 3n$ for all $n \geq n_0$

$$1b) \sqrt{7n^2 + 2n - 8} \leq \sqrt{7n^2 + 2n}$$

$$\sqrt{7n^2 + 2n} \leq \sqrt{7n^2 + 2n^2} \text{ for } n \geq 1$$

$$\sqrt{7n^2 + 2n^2} = 3n$$

$$\sqrt{7n^2 + 2n - 8} \leq 3n \text{ for } n \geq 1$$

$$\sqrt{7n^2 + 2n - 8} \geq \sqrt{7n^2 + 2n - 8n} \text{ for } n \geq 1$$

$$\sqrt{7n^2 + 2n - 8n} = \sqrt{7n^2 - 6n}$$

$$\sqrt{7n^2 - 6n} \geq \sqrt{7n^2 - 6n^2} \text{ for } n \geq 1$$

$$\sqrt{7n^2 - 6n^2} = n$$

$$\sqrt{7n^2 + 2n - 8} \geq n \text{ for } n \geq 1$$

Thus, $\sqrt{7n^2 + 2n - 8}$ is both $O(n)$ and $\Omega(n)$, making it $\Theta(n)$

1c) We have

$$d(n) = O(f(n)) \quad \text{and} \quad e(n) = O(g(n))$$

We can rewrite as

$$\begin{array}{ll} d(n) \leq c_1(f(n)) & \text{and} \quad e(n) \leq c_2(g(n)) \\ \text{for all } n \geq n_{01} & \text{for all } n \geq n_{02} \end{array}$$

Since we know all 4 terms are nonnegative, we can do:

$$\begin{array}{l} d(n) \cdot e(n) \leq c_1 \cdot f(n) \cdot c_2 \cdot g(n) \\ \text{for all } n \geq \max(n_{01}, n_{02}) \end{array}$$

Rewriting, we get

$$\begin{array}{l} d(n)e(n) \leq (c_1 \cdot c_2) f(n) g(n) \\ \text{for all } n \geq \max(n_{01}, n_{02}) \end{array}$$

Since c_1 and c_2 are just arbitrary constants, we can rephrase as

$$d(n)e(n) \leq C(f(n)g(n)) \quad \text{for } n \geq n_0$$

$$\text{where } C = c_1 \cdot c_2 \quad \text{and} \quad n_0 = \max(n_{01}, n_{02})$$

Thus, we have found a valid C and n_0 , to satisfy the definition of big- O . Therefore...

$$d(n)e(n) = O(f(n)g(n))$$

2)

example 1: $\Theta(n^2)$

example 2: $\Theta(n)$

example 3: $\Theta(\log(n))$

example 4: $\Theta(n)$

Justification for example 4:



$n + \frac{n}{2} + \dots + 2 + 1 \approx 2n$, which is $\Theta(n)$, not $\Theta(n \log n)$

Justification for example 3:

$\log_2(n^2) = 2 \cdot \log_2(n)$, which is just $\Theta(\log(n))$

Justification for example 2:

only one loop, whose body iterates in constant time, n times

Justification for example 1:

$1 + 2 + 3 + 4 \dots (n-2) + (n-1) + n$ is $\Theta(n^2)$