

Q) Explain Caesar Cipher technique with example.

The earliest known, and the simplest, use of substitution cipher was by Julius Caesar. The Caesar cipher involves replacing each letter of the alphabet with the letter standing three places further down the alphabet.

For example:

Plaintext : meet me after the foggy party.

Ciphertext : PHHW PH DIWHU PKH WRJD SDUWB.

Note that the alphabet is wrapped around, so that the letter following Z is A. We can define the transformation by listing all possibilities, as follows:

Let us, assign a numerical equivalent to each letter:

a	b	c	d	e	f	g	h	i	j	k	l	m
0	1	2	3	4	5	6	7	8	9	10	11	12
n	o	p	q	r	s	t	u	v	w	x	y	z
13	14	15	16	17	18	19	20	21	22	23	24	25

Plaintext: abcdefghijklmnopqrstuvwxyz

Ciphertext: defghijklmnopqrstuvwxyz abc

Then the algorithm can expressed as follows:

For each plaintext letter, p, substitute the ciphertext letter c

$$\therefore C = E(P, P) = (P+3) \bmod 26$$

A shift can be any amount. so that the general Caesar algorithm is.

$$C = E(K, P) = (P+K) \bmod 26$$

where takes on a value in the range 1 to 25. The decryption algorithm follow.

$$P = D(C, K) = C - K$$

$$P = D(K, C) = (C-K) \bmod 26$$

In this case, the plaintext leaps out as occupying the third line. The three important characteristics of this alphabet problem enabled us to use a brute-force on cryptanalysis:

1. The encryption and decryption algorithm is known.
2. There are only 25 keys to try.
3. The language of the plaintext is known and easily recognizable.

Q6 Encrypt the following message using Caesar cipher technique and use encryption key 134: "meet me after the toga party".

Ans: To encrypt the following message "meet me after the toga party" using Caesar cipher technique with a key 134. First we reduce the key modulo 26.

$$134 \bmod 26 = 4$$

Then, the effective encryption key is 4. Now, we shift the letter in the message by 4 positions in the alphabets.

Plaintext: meet me after the toga party.

Ciphertext: qiiix qij ejxiv xlis xske tevxe

EMMAU : festina

HEFTOM : festina

TADDA : festina

VHDEFO : festina

YDZSTI

Gibrigexy

Q6 Explain Mono-alphabetic cipher with an example.

= Monoalphabetic Cipher: A mono-alphabetic cipher is a type of substitution cipher where each letter in the plaintext is replaced by a fixed, different letter from the alphabet.

In mono-alphabetic Cipher, the mapping is randomly done and the difference between the letters is not uniform.

Let's consider an example, we define a substitution pattern by shifting the letters of the alphabet randomly.

Plaintext Alphabet : A B C D E F G H I J K L M N O P Q R S T U V W X Y Z  
Ciphertext Alphabet : Q W E R T Y U I O P A S D F G H J K L Z C V B N M

Now let's use this substitution to encrypt a message.

Plaintext : HELLOW	Plaintext : UAMME
Ciphertext : <del>O T F D D H V</del>	Ciphertext : G Q D D T
Ciphertext : ITSSGV	

5. Explain Playfair Cipher with Example.

= The best-known multiple-letter cipher is the playfair cipher, which treats the digraphs in the ~~key~~ plain text as single units and translates these units into the plaintext digraphs. The playfair cipher is more secure than monoalphabetic cipher because it uses the matrix of the alphabets to perform the substitution, creating greater complexity by working with pairs of letters.

The encryption algorithm consists of two steps:

Generate the key square ( $5 \times 5$ ). The key square is  $5 \times 5$  grids of the alphabets that acts as the key for encrypting the plaintext. Each of the 25 alphabets must be unique and one letter of the alphabet (usually J) is omitted from the table (each table can hold only 25 alphabets). If the plaintext is J, then it replaced by I.

Encrypted Digraphs: The message to be encrypted is divided into pairs of letters. If a pair contains the same letters like (LL), X is inserted between them. If the message has an odd number of letters, a trailing Z is added to the final plain text.

Encryption: There are three main rules for encryption of playfair cipher for each pair of letters.

1. Same Row: If both letters of the digraph appear in the same row of the matrix, each letter is replaced by the letter immediately to its right. (each letter to be start wrapping around to the start of the row if needed)

Same Column: If both letters are in the same column of the matrix, each letter is replaced by the letter immediately below it. (wrapping around to the top if necessary)

Same Rectangle: If the letters form the corners of the same rectangle, each letter is replaced by the one on the same row but in the other corner of the rectangle.

Example of the playfair cipher:

Suppose, the key is Keyword. Then the matrix.

K	E	Y	W	O
R	D	A	B	C
F	G	H	IJ	L
M	N	P	Q	S
T	U	V	X	Z

Suppose let the plaintext message : HELLO WORLD

Then the final plaintext : HE LX LO WO RL DZ

Then the ciphertext : GY IZ QSC OK CF CU

Q5 Explain the Hill cipher with an example.

Ans: Another interesting multiple letter cipher is a Hill cipher, developed by the mathematician Lester Hill in 1929. Hill cipher is a polygraphic substitution cipher based on the linear algebra. Each letter is represented by a number modulo 26. Often simple scheme is  $A=0, B=1, \dots, Z=25$  is used, but this is not an essential features of a cipher. To encrypt a message, each block of  $n$  letters (considered as an  $n$ -component vector) is multiplied by an invertible  $n \times n$  matrix, against modulo 26. To decrypted a message, each block of letters multiplied by inverse matrix for encryption.

Example: Encryption.

Suppose, A plaintext is "Hello World" and the key is : DDCF

Suppose we choose a  $2 \times 2$  as our key matrix,

$$K = \begin{bmatrix} 3 & 3 \\ 2 & 5 \end{bmatrix}$$

Let, encrypt the plaintext : Hi World  
the final plaintext : Hi Wo rl dx

2. Convert plaintext to the numbers:

Hi using mapping  $A=0, B=1 \dots Z=25$

$$\bullet H \rightarrow 7 \quad R \rightarrow 17$$

$$\bullet I \rightarrow 8 \quad L \rightarrow 11$$

$$\bullet W \rightarrow 22$$

$$\bullet O \rightarrow 14 \quad d \rightarrow 3 \\ X \rightarrow 23$$

So the plaintext vector is :

$$P_1 = \begin{bmatrix} 7 \\ 8 \end{bmatrix} \quad P_2 = \begin{bmatrix} 22 \\ 14 \end{bmatrix} \quad P_3 = \begin{bmatrix} 17 \\ 11 \end{bmatrix} \quad P_4 = \begin{bmatrix} 3 \\ 23 \end{bmatrix}$$

We will now encrypt each pair of letter using the formula:

$$C = K \cdot P \bmod 25$$

First pair:

$$C_1 = K \cdot P_1 \bmod 25$$

$$= \begin{bmatrix} 3 & 3 \\ 2 & 5 \end{bmatrix} \begin{bmatrix} 7 \\ 8 \end{bmatrix} \bmod 25$$

$$= \begin{bmatrix} 21+24 \\ 14+40 \end{bmatrix} \bmod 25$$

$$= \begin{bmatrix} 45 \\ 54 \end{bmatrix} \bmod 25$$

$$= \begin{bmatrix} 19 \\ 2 \end{bmatrix}$$

So the ciphered <sup>for</sup> Hill is  $c_2$

2nd Pair:

$$c_2 = kp_2 \bmod 26$$
$$= \begin{bmatrix} 3 & 3 \\ 25 & \end{bmatrix} \begin{bmatrix} 22 \\ 14 \end{bmatrix} \bmod 26$$
$$= \begin{bmatrix} 108 \\ 119 \end{bmatrix} \bmod 26 \quad \begin{bmatrix} 8 & 6 \\ 2 & 5 \end{bmatrix}^{-1} \cdot 4$$
$$\Rightarrow \begin{bmatrix} 4 \\ 10 \end{bmatrix}$$

So the ciphertext for [two is] ek

3rd Pair:

$$c_3 = kp_3 \bmod 26 = \begin{bmatrix} 3 & 3 \\ 25 & \end{bmatrix} \begin{bmatrix} 17 \\ 11 \end{bmatrix} \bmod 26$$
$$= \begin{bmatrix} 84 \\ 89 \end{bmatrix} \bmod 26$$
$$= \begin{bmatrix} 6 \\ 11 \end{bmatrix}$$

So, the ciphertext for nl is gl

4th Pair:

$$c_4 = \begin{bmatrix} 3 & 3 \\ 25 & \end{bmatrix} \begin{bmatrix} 3 \\ 23 \end{bmatrix} \bmod 26$$
$$= \begin{bmatrix} 78 \\ 121 \end{bmatrix} \bmod 26$$
$$= \begin{bmatrix} 0 \\ 17 \end{bmatrix}$$

So, the ciphertext for dn is an

Then the final ciphertext for Hi world is te ekglan

Decryption:

Given the ciphertext "TC EKGILAR" and the key matrix.

$$K = \begin{bmatrix} 3 & 3 \\ 2 & 5 \end{bmatrix}$$

then the inverse matrix of K is

$$K^{-1} = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix}$$

The ciphertext vector is

$$C_1 = \begin{bmatrix} 19 \\ 2 \end{bmatrix} \quad C_2 = \begin{bmatrix} 4 \\ 10 \end{bmatrix} \quad C_3 = \begin{bmatrix} 6 \\ 11 \end{bmatrix} \quad C_4 = \begin{bmatrix} 8 \\ 17 \end{bmatrix}$$

We will now decrypt each letter pair using formula

$$P = K^{-1} C \bmod 26$$

$$\therefore P_1 = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix} \begin{bmatrix} 19 \\ 2 \end{bmatrix} \bmod 26$$

$$= \begin{bmatrix} 319 \\ 398 \end{bmatrix} \bmod 26 = \begin{bmatrix} 7 \\ 8 \end{bmatrix} = \begin{bmatrix} 11 \\ 15 \end{bmatrix} = P_1$$

$$P_2 = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix} \begin{bmatrix} 4 \\ 10 \end{bmatrix} \bmod 26$$

$$= \begin{bmatrix} 280 \\ 170 \end{bmatrix} \bmod 26 = \begin{bmatrix} 22 \\ 14 \end{bmatrix} = \begin{bmatrix} 8 \\ 11 \end{bmatrix} = P_2$$

$$P_3 = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix} \begin{bmatrix} 6 \\ 11 \end{bmatrix} \bmod 26$$

$$= \begin{bmatrix} 17 \\ 11 \end{bmatrix} = rL$$

$$\text{for } P_4 = \begin{bmatrix} 15 & 17 \\ 20 & 9 \end{bmatrix} \begin{bmatrix} 0 \\ 17 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 3 \\ 23 \end{bmatrix}$$
  
$$= dx$$

->  $\bar{a} \cdot \bar{c} = 1$  or has fast inverses with respect to mod 26

**Exa**

-> initial diff to compute with previous pt is 1-09  
to compute diff. from previous point we do

$(1-09) + (1-09) = 2$  digital fast inverses with  
respect to multiplication

$$\text{and } -1 \cdot (-1) \equiv (4-1) \pmod{26} = 5 \pmod{26} = 5$$
  
$$\text{as base } (1-09) = 9-9$$

$$\text{as base } (1-09+1-09) = \text{as base } (11-11) \text{ as base } (11-11) =$$

$$(1-09+1-09) = \text{as base } (11-11) \text{ as base } (11-11)$$

to fast invert base so you just do regular invert all counts  
base are noted by base with fast inverses with respect to mod 26  
then we just to switch in term all digits will now be base  
noted for with fast inverses with respect to mod 26 diff and

Q6 Explain Vigenere Cipher with an example.

Ans: The Vigenere Cipher is a type of polyalphabetic substitution cipher, was developed by in the 16th century by the France cryptographer Blaise de Vigenere. We can express the vigenere cipher in the following manner:

Assume a sequence of the plaintext letter  $P = P_0, P_1, P_2 \dots P_{n-1}$ , a key consisting the sequence of the letter  $K = k_0, k_1, \dots, k_{m-1}$ , typically  $m < n$ . The sequence of the ciphertext letter  $C = C_0, C_1, C_2 \dots C_{n-1}$  is calculated as follows.

$$\begin{aligned} C = C_0, C_1, C_2 \dots C_{n-1} &= E(K, P) = E(k_0, k_1, k_2 \dots k_{n-1} \\ &\quad P_0, P_1, \dots, P_{m-1}) \bmod 26 \\ &= (P_0 + k_0) \bmod 26, (P_1 + k_1) \bmod 26 \dots (P_{m-1} + k_{m-1}) \bmod 26, \\ &\quad (P_m + k_0) \bmod 26, (P_{m+1} + k_1) \bmod 26 \dots (P_{2m-1} + k_{m-1}) \\ &\quad \bmod 26, \dots. \end{aligned}$$

Thus, the first letter of the key is added to the first letter of the plaintext, the second letter are added and so on through the first  $m$  letters of the plaintext. For the next  $m$  letter of the plaintext, the key letters

are repeated. This process is continued until all the sequence of plaintext is encrypted. A general-equation of the encryption process is -

$$C_i = (P_i + K_i \text{ mod } m) \text{ mod } 26$$

Then, the generalization of the decryption equation is

$$P_i = (C_i - K_i \text{ mod } m) \text{ mod } 26$$

Example: If the key is 'deceptive' and the plaintext message is "we are discovered save yourself" is encrypted as -

Key : deceptiive.deceptived.eceptive

Plaintext: wearediscovered.save.yourself

Key	3	4	2	4	15	19	8	21	4	3	4	2	4	15	19	8	21	4	3	4	2	4	15	19	8	21	4
plain text	22	4	0	17	4	3	8	18	2	14	21	4	17	4	3	18	0	21	4	24	14	20	17	18	1	5	
cipher	25	8	2	21	19	22	16	13	6	17	25	6	21	19	22	0	21	25	7	2	16	24	6	11	12	6	9

Ciphertext: zicvtwqngnzgvtwazheqyglmgj

Q: Explain the Vernam Cipher with an example.

The Vernam Cipher is a method of encrypting alphabetic text. It is one of the substitution techniques are converting the plaintext into ciphertext.

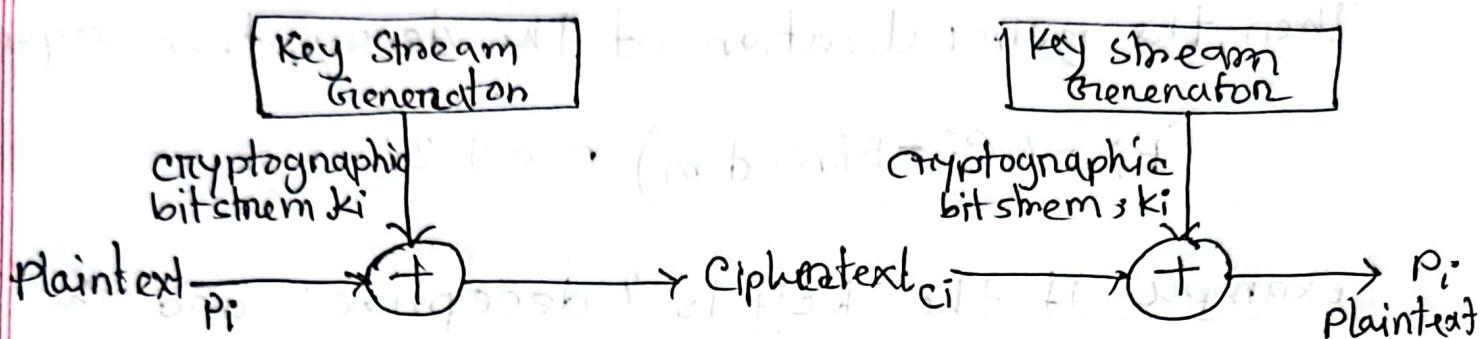


Fig: Vernam Cipher

This system works on binary data (bits) rather than the letters. This system can be express succinctly as follows:

$$c_i = p_i \oplus k_i$$

where, ~~p<sub>i</sub> = i<sup>th</sup> the number of di~~

~~p<sub>i</sub> = i<sup>th</sup> binary digit of plaintext.~~

~~c<sub>i</sub> = i<sup>th</sup> the binary digit of Ciphertext.~~

~~k<sub>i</sub> = i<sup>th</sup> the binary digit of key.~~

$\oplus$  = Exclusive OR operation.

Thus, the cipher ciphertext is generated by performing the bitwise XOR of the plaintext and the key. Because of the properties of XOR, decryption simply involves the same bitwise operations.

$$P_i = C_i \oplus K_i$$

### Example: Encryption:

Suppose, the key is FIRST and the plaintext message is UAMME

Plaintext position (P) :

Key position (K) :

(P+K)

Ciphertext position (C)

Ciphertext : ZIDEX

20	0	12	12	4
5	8	17	18	19
25	8	29	30	23
25	8	3	4	23

### Decryption:

Ciphertext position (C) :

Key position (K)

(C-K)

Plaintext position (P)

Plaintext : UAMME

25	8	3	4	23
5	8	17	18	19
20	0	14	-14	4
20	0	12	12	4