CSE 204: Data Structures and Algorithms

Assignment - 8: Divide and Conquer

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Course No: CSE 204

Complexity analysis:

The recursive function is:

```
public static List<Distance> secondClosestPoint(List<Point> pointsX, int
start, int end)
```

Here $\mathtt{List} < \mathtt{Point} > \mathtt{point} \times \mathtt{x}$ represents the points sorted by x. The cost is T(n)

The base cases:

```
if (size == 2) {
           List<Distance> distances = new ArrayList<>();
           distances.add(new Distance(pointsX.get(start),
pointsX.get(end)));
           Point temp = pointsX.get(start);
           if (pointsX.get(end).y < pointsX.get(start).y) {</pre>
               pointsX.set(start, pointsX.get(end));
               pointsX.set(end, temp);
           return distances;
       if (size == 3) {
           List<Distance> distances = new ArrayList<>();
           distances.add(new Distance(pointsX.get(start),
pointsX.get(start + 1)));
           distances.add(new Distance(pointsX.get(start + 1),
pointsX.get(end)));
           distances.add(new Distance(pointsX.get(start),
pointsX.get(end)));
           Collections.sort(distances, (d1, d2) -> {
               double dist = d1.distance - d2.distance;
               if (dist < 0)
                   return -1;
               if (dist > 0)
                   return 1;
               return 0:
           });
```

```
List<Point> temp = new ArrayList<>();
temp.add(pointsX.get(start));
temp.add(pointsX.get(start + 1));
temp.add(pointsX.get(end));
Collections.sort(temp, (d1, d2) -> {
    double dist = d1.y - d2.y;
    if (dist < 0)
       return -1;
    if (dist > 0)
       return 1;
    return 0;
});
for (int i = 0; i < 3; i++) {
    pointsX.set(start + i, temp.get(i));
if (distances.get(0).distance == distances.get(1).distance) {
    distances.remove(0);
} else
    distances.remove(2);
return distances;
```

In the base cases, the cost is O(1)

Divide step:

```
int mid = (start + end) / 2;
List<Distance> distances1 = secondClosestPoint(pointsX, start, mid);
List<Distance> distances2 = secondClosestPoint(pointsX, mid + 1,end);
```

The total cost of the divide step is 2T(n/2)

Conquer step:

Step 1:

```
List<Distance> minDistance = new ArrayList<> (distances1);
for (int i = 0; i < 2; i++) {</pre>
```

```
if (distances1.size() >= 1) {
               if (distances1.get(0).distance ==
distances2.get(0).distance)
                   distances1.remove(0);
               if (distances2.size() == 2 && distances1.size() >= 1) {
                   if (distances1.get(0).distance ==
distances2.get(1).distance)
                       distances1.remove(0);
       minDistance.addAll(distances2);
       Collections.sort(minDistance, (d1, d2) -> {
           double dist = d1.distance - d2.distance;
           if (dist < 0)
              return -1;
           if (dist > 0)
              return 1;
           return 0;
       });
       int minSize = minDistance.size();
       if (minSize == 4) {
          minDistance.remove(3);
          minDistance.remove(2);
       if (minSize == 3)
          minDistance.remove(2);
```

The cost of step-1 is O(1)

Step 2:

```
merge (pointsX, start, mid, end);
Here's the merge() function,
private static void merge (List<Point> points, int low, int mid, int high)
{
    int i = low, j = mid + 1;
    for (int k = low; k <= high; k++)
        Main.temp[k - low] = points.get(k);
    for (int k = low; k <= high; k++) {
        Point temp1;
        if (i > mid)
```

```
temp1 = Main.temp[j++ - low];
else if (j > high)
    temp1 = Main.temp[i++ - low];
else if (Main.temp[j - low].y < Main.temp[i - low].y)
    temp1 = Main.temp[j++ - low];
else
    temp1 = Main.temp[i++ - low];
points.set(k, temp1);
}</pre>
```

So, the cost of step -2 is O(n)

Step-3:

```
List<Point> midPoints = new ArrayList<>();

for (int i = start; i <= end; i++) {
    if (pointsX.get(i).x >= pointsX.get(mid).x -
    minDistance.get(0).distance
        && pointsX.get(i).x <= pointsX.get(mid).x +
    minDistance.get(0).distance) {
        midPoints.add(pointsX.get(i));
    }
}</pre>
```

The cost of step-3 is O(n)

Step-4:

The cost of step-4 is O(n)

So the total cost of conquer step is 3*O(n)+O(1)=O(n)

So,
$$T(n) = 2*T(n/2) + O(n)$$

Acording to Master theorem, T(n)=a*T(n/b) + f(n)

Here, a=2,b=2,f(n)=O(n)

Now, $log_22=1$

So, f(n) is polynomially equal and follows the 2nd condition of Master's Theorem.

So, T(n)=Theta(n¹*logn)=Theta(n*logn)