## Algorithms and Data Structures 2010 May Exam

## Nathan Sharp | s1869292 | B130263

## 15-12-2020

## Question 1

(a). State the definitions of a 'spanning subgraph', 'weight' of a subgraph and 'minimum spanning tree'.

**spanning subgraph** of a graph,  $\mathcal{G} = VE$ , is a subset of  $\mathcal{G}$  containing all vertices, V and a subset of the edges E.

weight of a subgraph the weight of a subgraph is the summation of the individual weight of all its edges.

**minimum spanning tree** of a graph  $\mathcal{G}$  is a connected spanning subgraph  $\mathcal{T}$  of  $\mathcal{G}$  of minimum weight.

(b). State prim' algorithm in pseudocode. For an input graph with n vertices and m edges what is its running time?

**Algorithm**  $Prim(\mathcal{G}: (VE), w: Weight Matrix) \to \mathcal{T}: MST$ 

- 1.  $\mathcal{T} \leftarrow \text{arbitrary } v \in G \quad \# v \text{ is an vertex in } V$
- 2. while fringe edge exists in  $\mathcal{T}$ :
- 3. add fringe edge of minimum weight to  $\mathcal{T}$
- 4. return  $\mathcal{T}$

where a **fringe edge** is an edge of the subtree  $\mathcal{T}$  with exactly one endpoint in  $\mathcal{T}$ .

Runtime is dependent on how the data structure is implemented, specifically, the *priority queue* used to find the minimum weight fringe edge in  $\mathcal{T}$  with runtimes as follows,

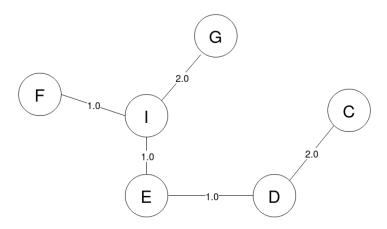
Data Structure	Runtime
Array Heap Fibonacci Heap	$T(n,m) = \Theta(n^2)$ $T(n,m) = \Theta((n+m)\log(n))$ $T(n,m) = \Theta(n\log(n) + m)$

- (c). Illustrate how Prim's algorithm processes the following input by indicating for each step of the main loop what edges Prim's algorithm has selected.
  - 1. {AB}
  - 2. {AB, BC}
  - 3. {AB, BC, CD}
  - 4. {AB, BC, CD, DE}
  - 5. {AB, BC, CD, DE, DH}
  - 6. {AB, BC, CD, DE, DH, EI}
  - 7. {AB, BC, CD, DE, DH, EI, IF}
  - 8. {AB, BC, CD, DE, DH, EI, IF, IG}
  - 9. {AB, BC, CD, DE, DH, EI, IF, IG}
- (d). Assume that  $\mathcal{G} = VE, w$ ) is a connected edge-weighted graph and that  $T \subset V$  is a non-empty set of vertices. A 'Steiner tree' for T is a connected (but not necessarily spanning) subgraph  $H = (V_H, E_H)$  of G that is a tree and whose vertex set  $V_H$  contains T as a subset. Its weight is

$$w(H) = \sum_{e \in E_H} w(e)$$

In the example graph from (c), find a steiner tree of minimum weight for  $T = \{C, E, F, G\}$ .

Solution,



(e). Suppose that an edge weighted connected graph G = (V, E, w) with n vertices

and m edges and a set  $T \subset V$  are given. Devise an algorithm that computer an optimal Steiner tree in time  $(2 + o(1))^n$ . Also provide an informl argument as to why your algorithm produces the correct output and why it has the desired running time.

```
def steiner(G,w,T) -> H:
    H = T[0] # any edge in T will do
    # initialise a priority queue of edges in T
# sorted by path distance to H
    Q = priorityQueue(T, dist_to_H)
    while not Q.isEmpty():
        t = Q.pop()
        H += min_path(t, H)

return H
```

The algorithm produces the correct output