

Operating Systems (INFR10079) 2020/2021 Semester 2

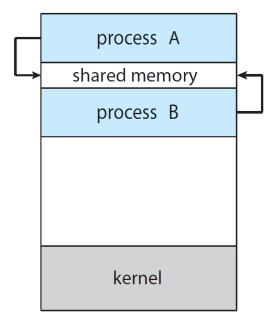
Processes (Interprocess Communication)

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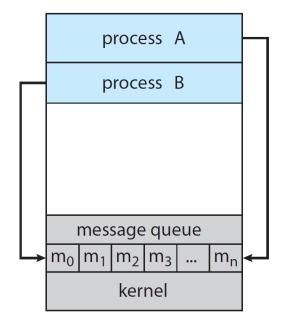
Interprocess Communication

- Independent processes
- Cooperating processes
 - Information sharing
 - More applications interested in the same information
 - Computation speedup
 - To exploit parallel hardware
 - Modularity
 - Reusability of components
- Require interprocess communication (IPC) mechanism to send and receive data
 - Shared memory
 - Message passing

Interprocess Communication



Shared memory



Message passing

Shared Memory

- Allow processes to communicate and synchronize
 - Sharing part of address space
 - OS doesn't mediates communication (no overhead)
 - Usually, OS prevents processes to access each other memory
 - Processes should agree (have permission) to void this restriction

Data

- Format decided by application
- Direct access (not mediated by the OS) very fast
- Application programmer fully manages the data transfer not trivial
- Possible use cases
 - Passing of large (single) objects (C&P, image, table, etc.)
 - Notification variable

Example of Shared Memory Code (POSIX)

```
#include <stdio.h>
                                                                      #include <stdio.h>
#include <stdlib.h>
                                                                      #include <stdlib.h>
#include <string.h>
                                                                      #include <fcntl.h>
#include <fcntl.h>
                                                                      #include <sys/shm.h>
#include <sys/shm.h>
                                                                      #include <sys/stat.h>
#include <sys/stat.h>
                                                                      #include <sys/mman.h>
#include <sys/mman.h>
int main() {
                                                                      int main() {
                                                                       const int SIZE = 4096; /* the size (bytes) of shared memory object */
 const int SIZE = 4096; /* the size (bytes) of shared memory object */
 const char *name = "OS"; /* name of the shared memory object */
                                                                       const char *name = "OS"; /* name of the shared memory object */
 const char *message 0 = "Hello"; /* written to shared memory */
 const char *message 1 = "World!"; /* written to shared memory */
 int fd; /* shared memory file descriptor */
                                                                       int fd; /* shared memory file descriptor */
 char *ptr; /* pointer to shared memory obect */
                                                                       char *ptr; /* pointer to shared memory obect */
 /* create the shared memory object */
                                                                       /* open the shared memory object */
 fd = shm open(name,O CREAT | O RDWR,0666);
                                                                       fd = shm open(name, O RDONLY, 0666);
 /* configure the size of the shared memory object */
                                                                       /* memory map the shared memory object */
 ftruncate(fd, SIZE):
                                                                       ptr = (char *)
 /* memory map the shared memory object */
                                                                        mmap(0, SIZE, PROT READ | PROT WRITE, MAP SHARED, fd, 0);
 ptr = (char *)
                                                                       /* read from the shared memory object */
  mmap(0, SIZE, PROT READ | PROT WRITE, MAP SHARED, fd, 0);
                                                                       printf("%s",(char *)ptr);
 /* write to the shared memory object */
                                                                       /* remove the shared memory object */
 sprintf(ptr,"%s",message 0);
                                                                       shm_unlink(name);
 ptr += strlen(message 0);
 sprintf(ptr,"%s",message 1);
                                                                       return 0:
 ptr += strlen(message 1);
 return 0;
```

Message Passing

- Allow processes to communicate and synchronize
 - Without sharing part of address space
 - OS mediates communication (may introduce overhead)
- Processes on the same machine and among different internetworked machines
 - Not possible with shared memory
- Message-passing facility provides at least two operations
 - send(message)
 - receive(message)
- Communication link
 - Several implementation tradeoffs, e.g., messages size
 - Fixed
 - Variable

Message Passing - Naming

- Communicating processes must refer to each other
- Direct communication
 - Symmetric: Explicit name of sender and receiver
 - send (P, message) send message to process P
 - receive (Q, message) receive message from process Q
 - Asymmetric: Explicit at least one end
 - send(p, message) send to p
 - receive (id, message) receive from any process, sender saved in id
- Indirect communication
 - No need to know/explicit in advance sender and/or receiver
 - Mailboxes (e.g., POSIX Mailbox)
 - send(A, message) into mailbox A
 - receive (A, message) from mailbox A
 - A mailbox can be accessed by more than two processes
 - Between processes multiple mailboxes may exist

Naming is a problem also in shared memory

Message Passing – Synchronization

- Different design options to implement send()/receive()
 - Blocking or synchronous
 - Nonblocking or asynchronous
- Different combinations may be offered
 - Blocking send
 - Nonblocking send
 - Blocking receive
 - Nonblocking receive

Rendezvous

- When both send() and receive() are blocking
- Simple solution to the consumer-producer problem***

Message Passing – Buffering

- Message exchanged reside in temporary buffers/queues
 - Zero capacity (no buffering)
 - No message waiting in it
 - Sender must block until the recipient receives the message
 - Bounded capacity
 - Finite length n; at most n messages can reside in it
 - If the queue is not full when a new message is sent, that is placed in the queue and the sender can continue execution
 - If the link is full, the sender has to wait
 - Unbounded capacity
 - Inifinite queue
 - The sender never blocks

Example of message passing: Pipes

- A pipe acts as a conduit allowing two processes to communicate
 - One way data flow
- Anonymous (ordinary)
 - Between parent and child
- Named
 - between any pair of processes)
 - In UNIX, FIFO

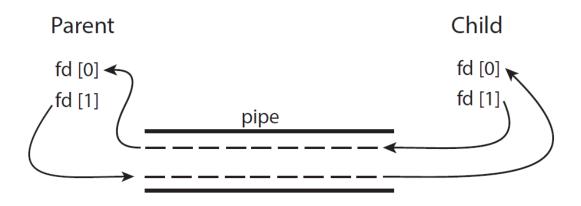


Figure 3.20 File descriptors for an ordinary pipe.

Example of Pipe Code (POSIX)

```
#include <sys/types.h>
#include <stdio.h>
#include <string.h>
#include <unistd.h>
#define READ END 0
#define WRITE END 1
int main(void) {
  char write msg[256] = "Greetings";
  char read msg[256];
  int fd[2];
  pid t pid;
  if (pipe(fd) == -1) {/* create the pipe */
    fprintf(stderr,"Pipe failed");
    return 1;
  pid = fork(); /* fork a child process */
  if (pid < 0) { /* error occurred */
    fprintf(stderr, "Fork Failed");
    return 1;
  if (pid > 0) { /* parent process */
    close(fd[READ_END]); /* close the unused end of the pipe */
    write(fd[WRITE END], write msg, strlen(write msg)+1); /* write to the pipe */
    close(fd[WRITE_END]); /* close the write end of the pipe */
  else { /* child process */
    close(fd[WRITE_END]); /* close the unused end of the pipe */
    read(fd[READ_END], read msg, 256); /* read from the pipe */
    printf("read %s",read msg);
    close(fd[READ_END]); /* close the read end of the pipe */
  return 0;
```

Client-Server Communication

Sockets Abstraction

- Endpoint for communication
- Identified by an IP address concatenated with a port number
- Servers implementing specific services (SSH, FTP, and HTTP) listen to well-known ports
 - an SSH server listens to port 22; an FTP server listens to port 21; and a web, or HTTP, server listens to port 80

Remote Procedure Call (RPC)

- Abstract the procedure-call mechanism
- For use between systems with network connections
- Similar in many respects to the IPC
 - Uses message-based communication to provide remote service

Signals

- OS mechanism to notify a process (one way)
 - Doesn't carry information
- From the OS
 - Can be thought as a software-generated interrupt/exception
 - Synchronous e.g., division by zero error
 - Asynchronous e.g., timer/alarm
 - Syscall: process -> OS; Signal: OS -> process
- From other processes (including the process itself)
 - Only notification, no data communication method
 - Management e.g., kill a process
 - Synchronization e.g., POSIX RT Signals
 - etc.

Handling Signals (Linux Example)

Signal handler

- Code to process a signal
- No return value (void)
- Handler must be registered
 - Or OS default action
- Send a signal
 - kill()
 - raise()

https://www.cs.auckland.ac.nz/references/unix/digital/APS33DTE/DOCU_006.HTM#realtime-handler-sec

```
#include <unistd.h>
#include <signal.h>
#include <stdio.h>
#include <sys/types.h>
#include <sys/wait.h>
#define SIG STOP CHILD SIGUSR1
main() {
 pid t pid;
 sigset t newmask, oldmask;
 if ((pid = fork()) == 0) { /* Child */
   struct sigaction action;
   void catchit();
   sigemptyset(&newmask);
   sigaddset(&newmask, SIG_STOP_CHILD);
   sigprocmask(SIG BLOCK, &newmask, &oldmask);
   action.sa flags = 0;
   action.sa_handler = catchit;
   if (sigaction(SIG STOP CHILD, &action, NULL) == -1) {
      perror("sigusr: sigaction");
      exit(1):
   sigsuspend(&oldmask);
 else { /* Parent */
   int stat:
   sleep(10):
   kill(pid, SIG_STOP_CHILD);
   pid = wait(&stat);
   printf("Child exit status = %d\n", WEXITSTATUS(stat));
   _exit(0);
void catchit(int signo) { /* Signal Handler */
   printf("Signal %d received from parent\n", signo);
   _exit(0);
```