

Introduction to Databases

Tutorial 5

Dr Paolo Guagliardo

Fall 2020 (week 7)

Problem 1 (mandatory). Consider the following set of FDs:

$$D \rightarrow AC, \quad AB \rightarrow DE, \quad FD \rightarrow E, \quad C \rightarrow F$$

(a) Indicate which of the above FDs hold in the following relation:

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>
1	1	2	3	0	4
2	1	2	3	1	4
1	1	3	3	0	5

Moreover, for each FD that does not hold provide two tuples (from the above relation) constituting a violation.

(b) Determine whether each of the following FDs is implied by the FDs above:

$$\begin{array}{llll} AC \rightarrow E & BD \rightarrow EF & EF \rightarrow BC & BC \rightarrow BF \\ AD \rightarrow CF & ABC \rightarrow DF & DEF \rightarrow AB & DF \rightarrow AE \\ CD \rightarrow DE & BE \rightarrow AC & CD \rightarrow ED & DE \rightarrow AF \end{array}$$

(c) For each FD in point (b) that is implied, write a derivation using Armstrong's axioms (including union and decomposition). **Requirement for tutorial marks:** attempt at least one derivation.

Solution.

(a) Let r_1 , r_2 and r_3 denote the rows of the given table (in the order in which they are shown, from top to bottom). Then, we have the following:

- $D \rightarrow AC$ does not hold because r_1 and r_2 agree on the value of D but they disagree on the values they assign to the pair of attributes A, C .
- $AB \rightarrow DE$ holds because all rows that agree on the values they assign to the pair of attributes A, B , namely r_1 and r_3 , also agree on the values they assign to the pair of attributes D, E .
- $FD \rightarrow E$ does not hold because r_1 and r_2 agree on the values they assign to the pair of attributes F, D but they disagree on the value of E .
- $C \rightarrow F$ holds because all rows that agree on the value of C , namely r_1 and r_2 , also agree on the value of F .

(b) Call Σ the given set of FDs. Recall that, for any FD $X \rightarrow Y$, we have $\Sigma \models X \rightarrow Y$ if and only if $Y \subseteq C_\Sigma(X)$.

- $\{E\} \not\subseteq C_\Sigma(AC) = ACF$, so $\Sigma \not\models AC \rightarrow E$
- $CF \subseteq C_\Sigma(AD) = ADCFE$, so $\Sigma \models AD \rightarrow CF$
- $DE \subseteq C_\Sigma(CD) = CDAFE$, so $\Sigma \models CD \rightarrow DE$
- $EF \subseteq C_\Sigma(BD) = BDACEF$, so $\Sigma \models BD \rightarrow EF$
- $DF \subseteq C_\Sigma(ABC) = ABCDEF$, so $\Sigma \models ABC \rightarrow DF$
- $AC \not\subseteq C_\Sigma(BE) = BE$, so $\Sigma \not\models BE \rightarrow AC$
- $BC \not\subseteq C_\Sigma(EF) = EF$, so $\Sigma \not\models EF \rightarrow BC$
- $AB \not\subseteq C_\Sigma(DEF) = DEFAC$, so $\Sigma \not\models DEF \rightarrow AB$
- $CD \rightarrow ED$ is the same as $CD \rightarrow DE$ because the l.h.s. and r.h.s. of an FD are *sets* of attributes
- $BF \subseteq C_\Sigma(BC) = BCF$, so $\Sigma \models BC \rightarrow BF$
- $AE \subseteq C_\Sigma(DF) = DFACE$, so $\Sigma \models DF \rightarrow AE$
- $AF \subseteq C_\Sigma(DE) = DEACF$, so $\Sigma \models DE \rightarrow AF$

(c) • Derivation of $AD \rightarrow CF$

- | | |
|--------------------------|---------------------------------------|
| (1) $D \rightarrow AC$ | [given in Σ] |
| (2) $AD \rightarrow AC$ | [from (1) by augmentation with A] |
| (3) $C \rightarrow F$ | [given in Σ] |
| (4) $AC \rightarrow ACF$ | [from (3) by augmentation with AC] |
| (5) $AD \rightarrow ACF$ | [from (2) and (4) by transitivity] |
| (6) $AD \rightarrow CF$ | [from (5) by decomposition] |

Step (6) above uses decomposition, which is a non-essential axiom. It could be replaced as follows:

- | | |
|---------------------------|-------------------------------------|
| (6a) $ACF \rightarrow CF$ | [by reflexivity] |
| (6b) $AD \rightarrow CF$ | [from (5) and (6a) by transitivity] |

In general, the decomposition axiom:

$$\frac{X \rightarrow YZ}{X \rightarrow Y} \text{ (decomposition)}$$

can be replaced by the following sequence that uses only essential axioms:

$$\frac{X \rightarrow YZ \quad \overline{YZ \rightarrow Y} \text{ (reflexivity)}}{X \rightarrow Y} \text{ (transitivity)}$$

• Derivation of $CD \rightarrow DE$

- | | |
|-------------------------|--------------------------------------|
| (1) $C \rightarrow F$ | [given in Σ] |
| (2) $CD \rightarrow FD$ | [from (1) by augmentation with D] |
| (3) $FD \rightarrow E$ | [given in Σ] |
| (4) $CD \rightarrow E$ | [from (2) and (3) by transitivity] |
| (5) $CD \rightarrow DE$ | [from (4) by augmentation with D] |

- Derivation of $BD \rightarrow EF$

- | | |
|--------------------------|--------------------------------------|
| (1) $D \rightarrow AC$ | [given in Σ] |
| (2) $DB \rightarrow ACB$ | [from (1) by augmentation with B] |
| (3) $DB \rightarrow AB$ | [from (2) by decomposition] |
| (4) $DB \rightarrow C$ | [from (2) by decomposition] |
| (5) $AB \rightarrow DE$ | [given in Σ] |
| (6) $DB \rightarrow DE$ | [from (3) and (5) by transitivity] |
| (7) $C \rightarrow F$ | [given in Σ] |
| (8) $DB \rightarrow F$ | [from (4) and (7) by transitivity] |
| (9) $DB \rightarrow DEF$ | [from (6) and (8) by union] |
| (10) $DB \rightarrow EF$ | [from (9) by decomposition] |

Step (9) above uses union, which is a non-essential axiom. It could be replaced as follows:

- | | |
|----------------------------|---------------------------------------|
| (9a) $DB \rightarrow DBE$ | [from (6) by augmentation with DB] |
| (9b) $DBE \rightarrow DEF$ | [from (8) by augmentation with DE] |
| (9c) $DB \rightarrow DEF$ | [from (9a) and (9b) by transitivity] |

In general, the union axiom:

$$\frac{X \rightarrow Y \quad X \rightarrow Z}{X \rightarrow YZ} \text{ (union)}$$

can be replaced by the following sequence that uses only essential axioms:

$$\frac{\frac{X \rightarrow Y}{X \rightarrow XY} \text{ (augmentation with } X) \quad \frac{X \rightarrow Z}{XY \rightarrow YZ} \text{ (augmentation with } Y)}{X \rightarrow YZ} \text{ (transitivity)}$$

- Derivation of $ABC \rightarrow DF$

- | | |
|---------------------------|---------------------------------------|
| (1) $AB \rightarrow DE$ | [given in Σ] |
| (2) $ABC \rightarrow DEC$ | [from (1) by augmentation with C] |
| (3) $C \rightarrow F$ | [given in Σ] |
| (4) $DEC \rightarrow DEF$ | [from (3) by augmentation with DE] |
| (5) $ABC \rightarrow DEF$ | [from (2) and (4) by transitivity] |
| (6) $ABC \rightarrow DF$ | [from (5) by decomposition] |

- Derivation of $CD \rightarrow ED$

- | | |
|-------------------------|--------------------------------------|
| (1) $C \rightarrow F$ | [given in Σ] |
| (2) $CD \rightarrow FD$ | [from (1) by augmentation with D] |
| (3) $FD \rightarrow E$ | [given in Σ] |
| (4) $FD \rightarrow ED$ | [from (3) by augmentation with D] |
| (5) $CD \rightarrow ED$ | [from (2) and (4) by transitivity] |

- Derivation of $BC \rightarrow BF$

- | | |
|-------------------------|----------------------------|
| (1) $C \rightarrow F$ | [given in Σ] |
| (2) $BC \rightarrow BF$ | [from (1) by augmentation] |

- Derivation of $DF \rightarrow AE$

- | | |
|--------------------------|--------------------------------------|
| (1) $DF \rightarrow E$ | [given in Σ] |
| (2) $DF \rightarrow DE$ | [from (1) by augmentation with D] |
| (3) $D \rightarrow AC$ | [given in Σ] |
| (4) $DE \rightarrow ACE$ | [from (3) by augmentation with E] |
| (5) $DF \rightarrow ACE$ | [from (2) and (4) by transitivity] |
| (6) $DF \rightarrow AE$ | [from (5) by decomposition] |

- Derivation of $DE \rightarrow AF$

- | | |
|--------------------------|--------------------------------------|
| (1) $D \rightarrow AC$ | [given in Σ] |
| (2) $C \rightarrow F$ | [given in Σ] |
| (3) $AC \rightarrow AF$ | [from (2) by augmentation with A] |
| (4) $D \rightarrow AF$ | [from (1) and (3) by transitivity] |
| (5) $DE \rightarrow AFE$ | [from (4) by augmentation with E] |
| (6) $DE \rightarrow AF$ | [from (5) by decomposition] |

Problem 2 (optional). Let R , S and T be relations on attributes A, B, C . Given the following set of INDs:

$$R[A, B] \subseteq S[B, C] \qquad S[C, B] \subseteq T[C, A]$$

determine which of the following INDs are implied:

$R[A] \subseteq T[A]$	$R[B] \subseteq T[B]$	$R[C] \subseteq T[C]$
$R[A] \subseteq T[B]$	$R[B] \subseteq T[A]$	$R[B] \subseteq T[C]$
$R[C] \subseteq T[B]$	$R[A] \subseteq T[C]$	$R[C] \subseteq T[A]$

Solution. Call Σ the given set of INDs. Using the axioms for INDs we can compute the closure Σ^+ of Σ :

- | | |
|-----------------------------------------------|------------------------------------|
| (1) $R[A, B] \subseteq S[B, C]$ | [given in Σ] |
| (2) $S[C, B] \subseteq T[C, A]$ | [given in Σ] |
| (3) $R[A] \subseteq S[B]$ | [from (1) by projection] |
| (4) $S[C] \subseteq T[C]$ | [from (2) by projection] |
| (5) $R[B, A] \subseteq S[C, B]$ | [from (1) by permutation] |
| (6) $S[B, C] \subseteq T[A, C]$ | [from (2) by permutation] |
| (7) $R[B] \subseteq S[C]$ | [from (5) by projection] |
| (8) $S[B] \subseteq T[A]$ | [from (6) by projection] |
| (9) $R[A, B] \subseteq T[A, C]$ | [from (1) and (6) by transitivity] |
| (10) $R[A] \subseteq T[A]$ | [from (9) by projection] |
| (11) $R[B, A] \subseteq T[C, A]$ | [from (9) by permutation] |
| (12) $R[B] \subseteq T[C]$ | [from (11) by projection] |
| (13) All trivial INDs obtained by reflexivity | |

So $R[A] \subseteq T[A]$ and $R[B] \subseteq T[C]$ are implied by Σ (because they are in Σ^+).