IDS Lecture 6: Predicate Logic

Free variables variables that are not in the scope of any quantifier. A variable that is not free is bound.

Interpretations

A formula may be true or false w.r.t a given interpretation.

Interpretation defines the semantics of the language; an assignment of variables that gives meaning to a statement.

Semantics of FOL: Interpretations

First Order Structure $\mathcal{I} = \langle \Delta, \cdot^{\mathcal{I}} \rangle$

 Δ non empty domain of objects (universe)

 $a^{\mathcal{I}}$ function which gives meaning to constant & predicate symbols

- $a^{\mathcal{I}} \in \Delta$ gives meaning to *constants*, "object a by means of interpretation function \mathcal{I} ".
- $R^{\mathcal{I}} \subseteq \Delta^1 \times ... \times \Delta^n$ gives meaning to *predicates*, "mapping it to an element in our domain (objects in the unverse)"

Variable Assignment (v) maps each variable to an object in Δ

• Notation: v[x/d] is v with $x \to d$

Semantics of FOL: Terms

Interpretation of terms under (\mathcal{I}, v)

$$x^{\mathcal{I},v} = v(x)$$

$$a^{\mathcal{I},v} = a^x$$

Formulas

 $(\mathcal{I}, v) \models \phi$ means interpretation (I, v) satisfies formula ϕ

$$I, v \models P(t_1, ..., t_n) \iff (t_1^{\mathcal{I}, v}, ..., t_n^{\mathcal{I}, v}) \in P$$

$$I, v \models \neg \phi \iff \mathcal{I}, v \nvDash \phi$$

$$I, v \models \phi \land \psi \iff \mathcal{I}, v \models \phi \text{ and } \mathcal{I}, v \models \psi$$

$$I, v \models \phi \lor \psi \iff \mathcal{I}, v \models \phi \text{ or } \mathcal{I}, v \models \psi$$

$$I, v \models \phi \to \psi \iff \mathcal{I}, v \models \phi \text{ then } \mathcal{I}, v \models \psi$$

$$I, v \models \forall x \phi \iff \text{for every } d \in \Delta : \mathcal{I}, v[x/d] \models \phi$$

$$I, v \models \exists x \phi \iff \text{there exists } d \in \Delta \text{ s.t } \mathcal{I}, v[x/d] \models \phi$$