

# Homework Assignment 3

$$1. \binom{2n}{n} = \frac{2n!}{n!(2n-n)!}$$

$$= \frac{2n!}{(n!)^2}$$

$$2n! = \sqrt{2\pi n} \cdot \left(\frac{2n}{e}\right)^{2n} \cdot (1 + \theta(\frac{1}{2n}))$$

$$(n!)^2 = (\sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot (1 + \theta(\frac{1}{2n})))^2$$

$$= \frac{1}{\sqrt{n\pi}} \cdot 4^n \cdot \frac{(1 + \theta(\frac{1}{2n}))}{(1 + \theta(\frac{1}{2n}))^2}$$

$$= \frac{1}{\sqrt{n\pi}} \cdot \frac{4^n}{n} \cdot \frac{(1 + \theta(\frac{1}{2n}))}{(1 + \theta(\frac{1}{2n}))^2}$$

$$\frac{4^n}{n} = \frac{1}{\sqrt{n\pi}} \cdot \frac{(1 + \theta(\frac{1}{2n}))}{(1 + \theta(\frac{1}{2n}))^2} \leftarrow \lim_{n \rightarrow \infty}$$

$$= \frac{1}{\sqrt{n\pi}} \cdot 1$$

$$= \frac{1}{\sqrt{n\pi}}$$

$$\frac{1}{\sqrt{n\pi}} \in (0, \infty)$$

2. a. Prove  $\forall n \geq n_0 : P(n) \rightarrow P(n+1)$

$$\text{Let } n_0 = 1 \rightarrow n \geq 1$$

$$\sum_{i=1}^1 i^3 = 1 = \left(\frac{1(1+1)}{2}\right)^2 \quad \checkmark$$

Assume  $P(n)$  is true, prove  $P(n+1)$

$$\begin{aligned} \sum_{i=1}^{n+1} i^3 &= \sum_{i=1}^n i^3 + (n+1)^3 \\ &= \left(\frac{n(n+1)}{2}\right)^2 + (n+1)^3 \quad * \text{ Inductive step} \\ &= \frac{n^2(n+1)^2}{4} + (n+1)^3 \\ &= \frac{n^2(n+1)^2 + 4(n+1)^3}{4} \end{aligned}$$

$$= \frac{n^4 + 6n^3 + 13n^2 + 10n + 4}{4}$$

$$= \frac{(n+1)^2(n+2)^2}{4}$$

$$= \left(\frac{(n+1)(n+2)}{2}\right)^2 \quad \checkmark$$

b. Prove  $\forall n > n_0 : P(n-1) \rightarrow P(n)$

$$\text{Let } n_0 = 1 \rightarrow n > 1$$

$$n=2 \quad \sum_{i=1}^1 i^3 = 1 = \left(\frac{(2-1)(2)}{2}\right)^2 \quad \checkmark$$

Assume  $P(n-1)$  is true, prove  $P(n)$

$$\begin{aligned} \sum_{i=1}^n i^3 &= \sum_{i=1}^{n-1} i^3 + n^3 \\ &= \left(\frac{(n-1)n}{2}\right)^2 + n^3 \quad * \text{ Inductive Step} \\ &= \frac{(n-1)^2 n^2 + 4n^3}{4} \end{aligned}$$



$$\begin{aligned}
 &= \frac{n^4 + 2n^3 + n^2}{4} \\
 &= \frac{n^2(n+1)^2}{4} \\
 &= \left(\frac{n(n+1)}{2}\right)^2 \quad \checkmark
 \end{aligned}$$

3. I. S'CD says

$$\begin{aligned}
 S(n) &\geq \lg(n) \\
 n &\geq 0
 \end{aligned} \quad \checkmark$$

II d. Let  $n \geq 1$ : Assume  $\forall k$   $1 \leq k < n$ :  $S(k) \geq \lg(k)$

prove:  $S(n) \geq \lg(n)$

$$S(n) \geq S(\lceil \frac{n}{2} \rceil) + 1$$

$$\geq \lg(\lceil \frac{n}{2} \rceil) + 1$$

\* induction step

$$\geq \lg(\frac{n}{2}) + 1$$

$$\geq \lg(n) - \lg(2) + 1$$

$$\geq \lg(n) - 1 + 1$$

$$\geq \lg(n)$$

$$\text{So } S(n) \geq \lg(n)$$

4. I. Base case  $S(n)$

$$T(n) \leq \frac{4}{3}n^2$$

$$1 \leq \frac{4}{3} \quad \checkmark$$

II d. Let  $n > 1$  be arbitrary

Assume  $\forall k$  in the range  $1 \leq k < n$

$$T(k) \leq \frac{4}{3}k^2$$

We must show:  $T(n) \leq \frac{4}{3}n^2$

$$T(n) = T(\lfloor \frac{n}{2} \rfloor) + n^2$$

$$\leq \frac{4}{3}(\lfloor \frac{n}{2} \rfloor)^2 + n^2$$

\* inductive step

$$\leq \frac{4}{3}(\frac{n}{2})^2 + n^2$$

$$\leq \frac{n^2}{3} + n^2$$

$$\leq \frac{n^2 + 3n^2}{3}$$

$$\leq \frac{4}{3}n^2$$

$$\text{So } T(n) \leq \frac{4}{3}n^2 \quad \text{hence } T(n) = O(n^2)$$



5. I. Base case  $T(1), T(2)$

$$T(1) \leq 3n^2 - 1$$

$$2 \leq 2 \quad \checkmark$$

$$T(2) \leq 3n^2 - 1$$

$$2 \leq 11 \quad \checkmark$$

II. Let  $n > 2$  be arbitrary

Assume  $\forall k \quad 1 \leq k < n$

$$T(k) \leq 3k^2 - 1$$

We must show  $T(n) \leq 3n^2 - 1$

$$T(n) = 9T\left(\left\lfloor \frac{n}{3} \right\rfloor\right) + 1$$

$$\leq 9(3\left(\left\lfloor \frac{n}{3} \right\rfloor\right)^2 - 1) + 1 \quad \text{* inductive step}$$

$$\leq 9 \cdot 3\left(\left\lfloor \frac{n}{3} \right\rfloor\right)^2 - 9 + 1$$

$$\leq 9 \cdot \frac{3n^2}{27} - 8$$

$$\leq 3n^2 - 8$$

$$\leq 3n^2 - 1$$

\*  $(-8 < -1)$  So we can

So:  $T(n) \leq 3n^2 - 1$  hence  $T(n) = O(n^2)$  replace the constant)

6. I. Base case  $T(1), T(2)$

$$T(1) \leq 6n$$

$$6 \leq 6 \quad \checkmark$$

$$T(2) \leq 6n$$

$$6 \leq 12 \quad \checkmark$$

II. Let  $n > 2$  be arbitrary

Assume  $\forall k \quad 1 \leq k < n$

$$T(k) \leq 6n$$

We must show:  $T(n) \leq 6n$

$$T(n) = 2T\left(\left\lfloor \frac{n}{3} \right\rfloor\right) + n$$

$$\leq 2 \cdot 6\left(\left\lfloor \frac{n}{3} \right\rfloor\right) + n \quad \text{* inductive step}$$

$$\leq 12\left(\frac{n}{3}\right) + n$$

$$\leq 4n + n$$

$$\leq 5n$$

$$\leq 6n$$

\*  $(5 < 6)$

So  $T(n) \leq 6n$  hence  $T(n) = O(n)$