

## Division with Remainder

In  $\mathbb{Z}$ , we had the extremely important Division Algorithm. This theorem states that if  $a$  and  $b$  are integers with  $b \neq 0$ , then there exists a “quotient”  $q$  and a “remainder”  $r$  such that  $a = qb + r$ , and, moreover, the remainder is not too large –  $0 \leq r < |b|$ . This is the result from which most of the interesting results and algorithms in  $\mathbb{Z}$  spring.

We’d like to generalize this property to integral domains. Notice that one problem is the appearance of absolute value in the bound on  $r$ : in general, rings do not have anything like absolute value, or a way to compare the “sizes” of two elements. However we did describe such a gadget for some rings: multiplicative norms. Recall that  $N : R \rightarrow \mathbb{N}$  is a multiplicative norm if (1)  $N(x) = 0$  iff  $x = 0$ , (2)  $N(xy) = N(x)N(y)$ , and (3) if  $N(x) = 1$  then  $x$  is a unit. These properties do generalize the absolute value.

**Definition 1** (Euclidean Norm). *Let  $R$  be a domain.*

- *We say that a multiplicative norm  $N : R \rightarrow \mathbb{N}$  is a Euclidean norm if for all  $a, b \in R$  with  $b \neq 0$ , there exist  $q, r \in R$  such that  $a = qb + r$  and  $0 \leq N(r) < N(b)$ .*
- *If there is a Euclidean norm on  $R$ , we say that  $R$  is a Euclidean Domain.*

Of course  $\mathbb{Z}$  is a Euclidean Domain with norm  $N(a) = |a|$ . The existence of a Euclidean norm on  $R$  is very powerful. For instance, many of the nice properties of  $\mathbb{Z}$  which we derived from the Division Algorithm have analogues in any Euclidean Domain. More generally, the norm allows us to recover some of the benefits of mathematical induction.

**Proposition 1.** *Every Euclidean Domain is also a GCD Domain.*

*Proof.* Let  $R$  be a Euclidean domain with norm  $N$ . We want to show that for all  $a \in R$ , for all  $b \in R$ , the set  $\gcd(a, b)$  is not empty. We proceed by strong induction on  $N(a)$ .

**Base case.** Suppose  $N(a) = 0$ . Then  $a = 0$ , and so  $b \in \gcd(a, b)$  for all  $b$ .

**Inductive Step.** Let  $a \in R$  and suppose that the result holds for all  $a'$  with  $1 \leq N(a') < N(a)$ . In particular, note that  $a \neq 0$ . Now let  $b \in R$ . By the division algorithm we may decompose  $b$  as  $b = qa + r$ , where  $0 \leq N(r) < N(a)$ . If  $r = 0$  then  $a|b$  and we have  $a \in \gcd(a, b)$ . If  $r \neq 0$ , then by the inductive hypothesis  $\emptyset \neq \gcd(r, a) = \gcd(b - qa, a) = \gcd(b, a)$  as needed.  $\square$

**Proposition 2.** *Every Euclidean domain is a Unique Factorization domain.*

The proof for  $\mathbb{Z}$  generalizes.

**Proposition 3.** *Every field is a Euclidean domain.*

*Proof.* Define a mapping  $N : F \rightarrow \mathbb{N}$  by  $N(x) = 0$  if  $x = 0$  and 1 if  $x \neq 0$ . We can see that  $N$  is a Euclidean norm.  $\square$

## Example: The Gaussian Integers

**Proposition 4.**  $\mathbb{Z}[i]$  is a Euclidean domain under the norm  $N(a+bi) = a^2 + b^2$ .

*Proof.* Let  $\alpha = a_1 + a_2i$  and  $\beta = b_1 + b_2i$  be Gaussian integers, with  $\beta \neq 0$ . Thinking of  $\alpha$  and  $\beta$  as elements of  $\mathbb{Q}(i)$ , we have

$$\frac{\alpha}{\beta} = t_1 + t_2i = \frac{a_1b_1 + a_2b_2}{b_1^2 + b_2^2} + \frac{a_2b_1 - a_1b_2}{b_1^2 + b_2^2}i.$$

Choose integers  $q_1$  and  $q_2$  such that  $|q_1 - t_1| \leq \frac{1}{2}$  and  $|q_2 - t_2| \leq \frac{1}{2}$ . (Note that this is always possible.) Let  $\gamma = q_1 + q_2i$ , and let  $\delta = \alpha - \gamma\beta$ . Note that by construction,  $\gamma$  and  $\delta$  are in  $\mathbb{Z}[i]$ .

We now have

$$\begin{aligned} N(\delta) &= N(\alpha - \gamma\beta) = N\left(\left(\frac{\alpha}{\beta} - \gamma\right)\beta\right) = N\left(\frac{\alpha}{\beta} - \gamma\right)N(\beta) \\ &= ((q_1 - t_1)^2 + (q_2 - t_2)^2)N(\beta) \leq \frac{1}{2}N(\beta) < N(\beta), \end{aligned}$$

as needed. □

**Corollary 5.**  $\mathbb{Z}[i]$  is a GCD domain and a UFD.

Here is a worked example of the division algorithm in the Gaussian integers. Let  $\alpha = 10 + 7i$  and  $\beta = 3 + 2i$ . Now

$$\frac{\alpha}{\beta} = \frac{44}{13} + \frac{1}{13}i = \left(3 + \frac{5}{13}\right) + \left(0 + \frac{1}{13}\right)i.$$

Let  $t_1 = 3$  and  $t_2 = 0$ , so that  $\gamma = 3$ . Now  $\delta = \alpha - \gamma\beta = 1 + i$ . We then have  $10 + 7i = 3(3 + 2i) + (1 + i)$  and  $N(1 + i) < N(3 + 2i)$ .

## Exercises

1. ( $k$ -stage Euclidean)
2. (Factorization in  $\mathbb{Z}[i]$ )