

Redes de Computadores II

Universidade do Algarve

Semana I I

https://github.com/ncatanoc/redes_algarve

Néstor Cataño

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Networking outlook

Goal:

to understand the basic underpinnings of TLS and HTTPS.

Roadmap

1. Transport Layer Security (TLS)

2. Certificate Authority (CA)

3. Man-in-the-Middle attacks

TLS - introduction

The Transport Layer Security Protocol (TLS)

Why a new protocol?

Protection Goals

- ▶ Confidentiality
- ▶ Integrity
- ▶ Availability

| |
|-------------------------------|
| Application HTTP, DNS, ... |
| Transport TCP, UDP |
| Internetwork IP |
| Link Ethernet |

The 3 first layers provide **Availability** but not **Confidentiality** and **Integrity**
⇒ TLS ensures confidentiality and integrity of message contents.

TLS - introduction

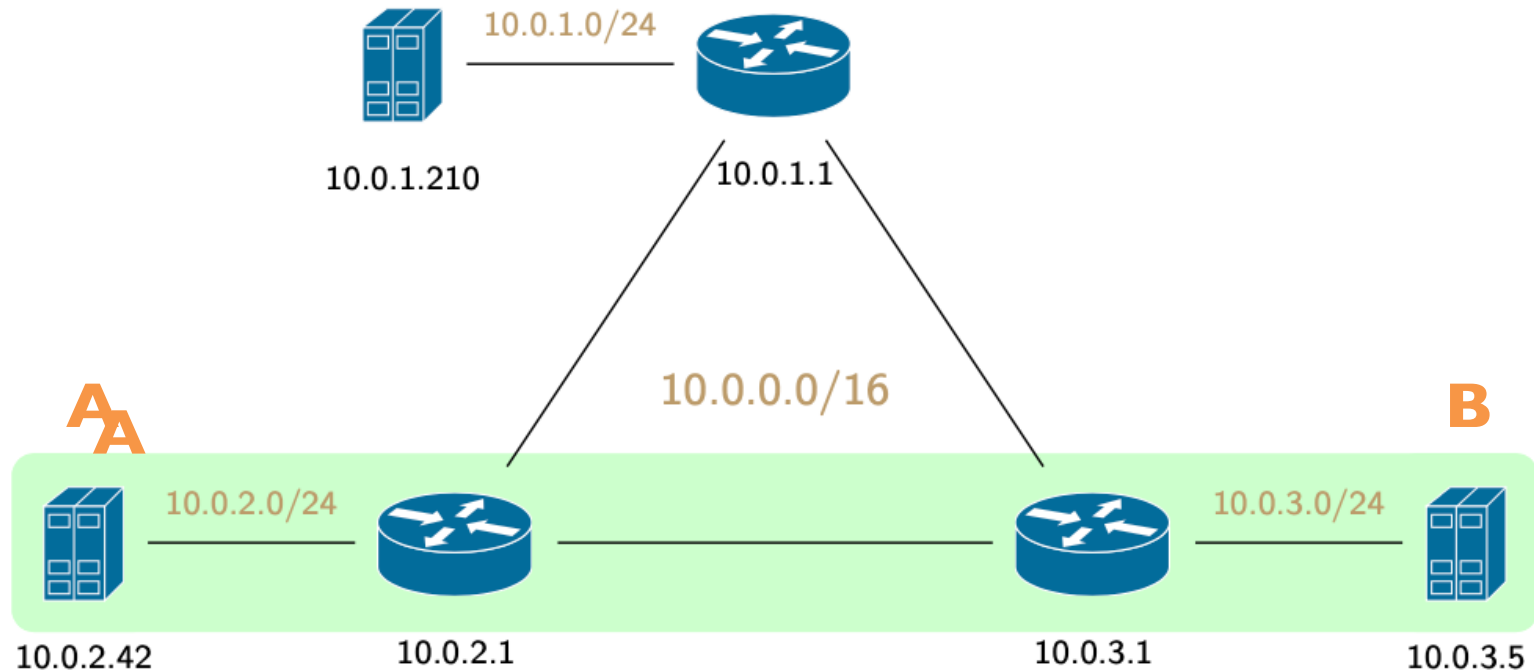
| |
|----------------|
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- The first 3 layers:
 - provide Availability
 - do not provide Confidentiality or Integrity

What do we need?

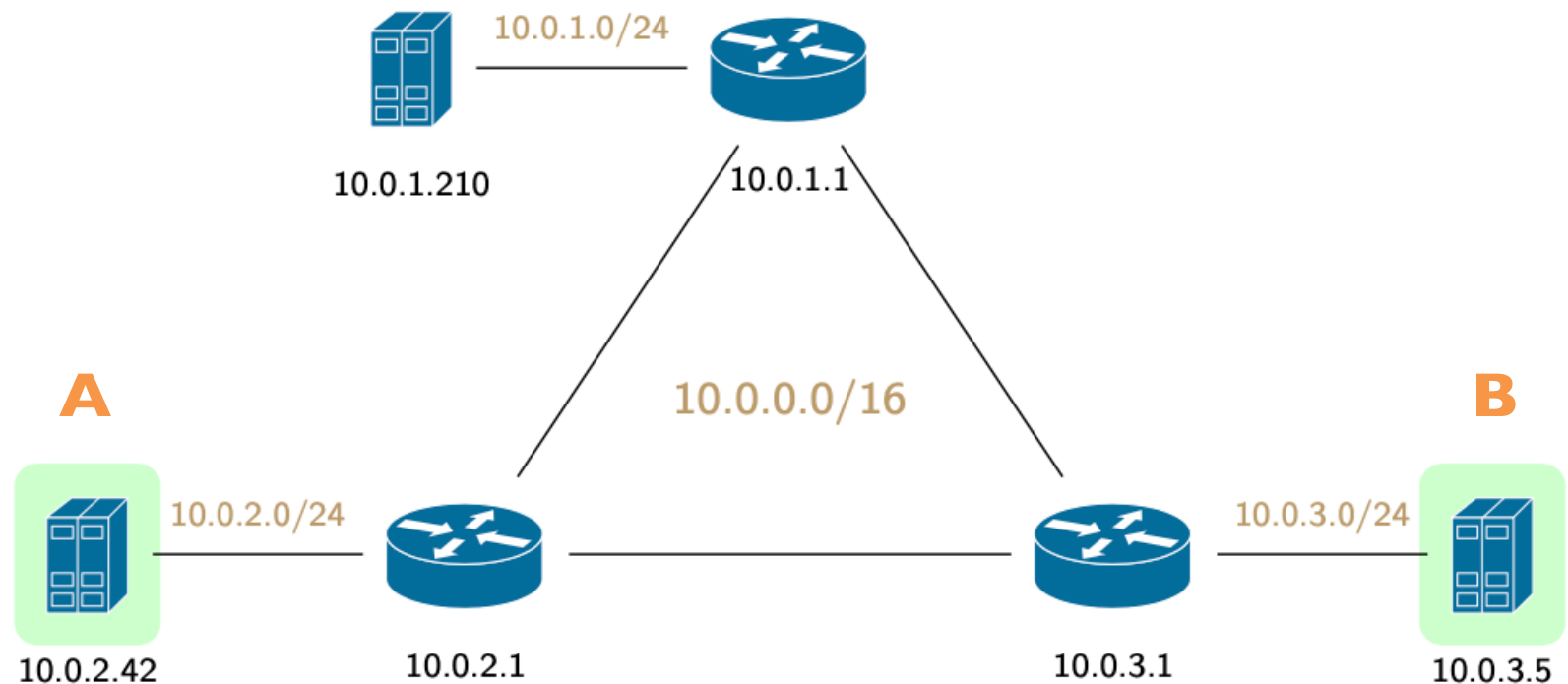
- Confidentiality: we use cryptography
 - We encrypt messages before passing them to TCP and UDP
- Integrity: we need to detect message modifications
- we need to make sure we are talking to the right person
 - e.g. Client browser talking to a server

Dolev-Yao: Trusted Areas (No TLS)



Anyone between **A** and **B** can modify messages
We need to trust 4 individual parts

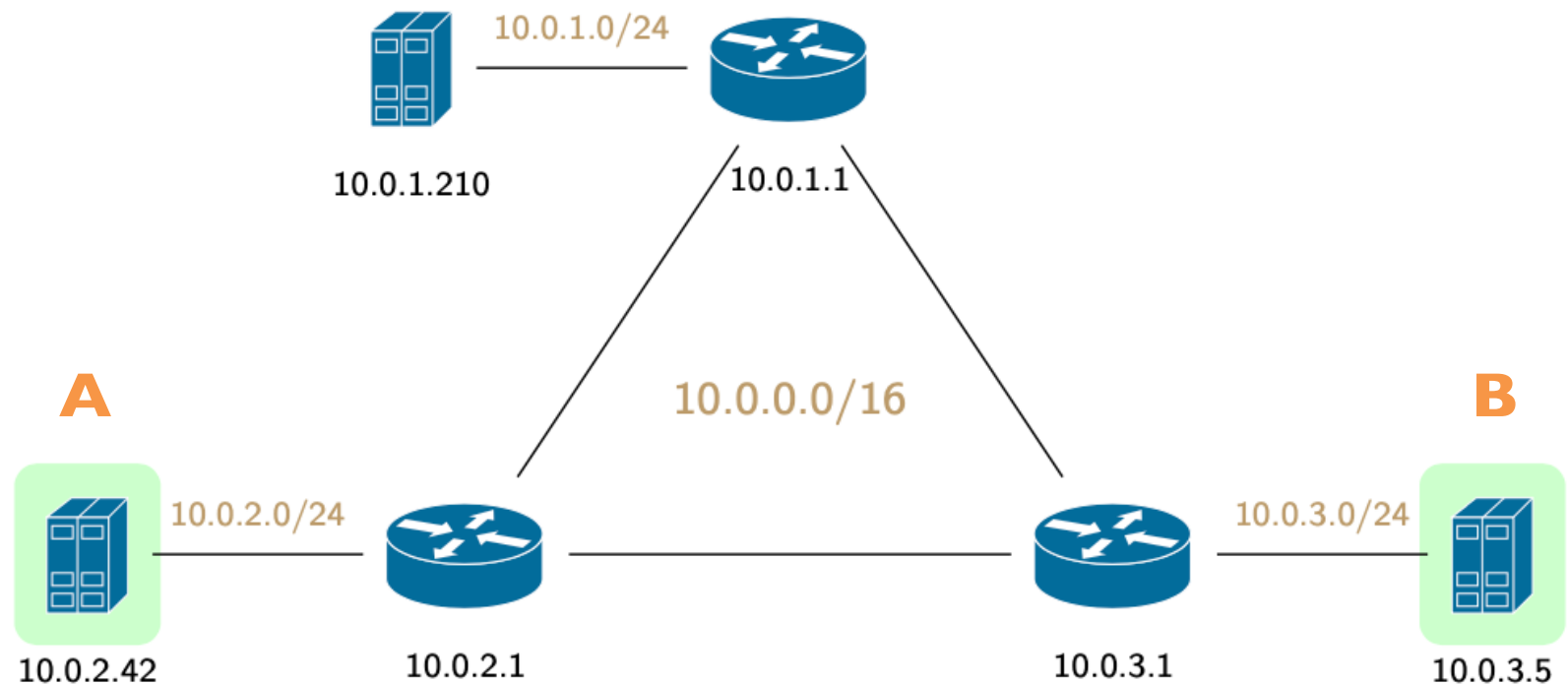
Dolev-Yao: Trusted Areas (TLS)



TLS' end-to-end principle:

We only need to trust **A** and **B**
routers between A and B cannot see their communication

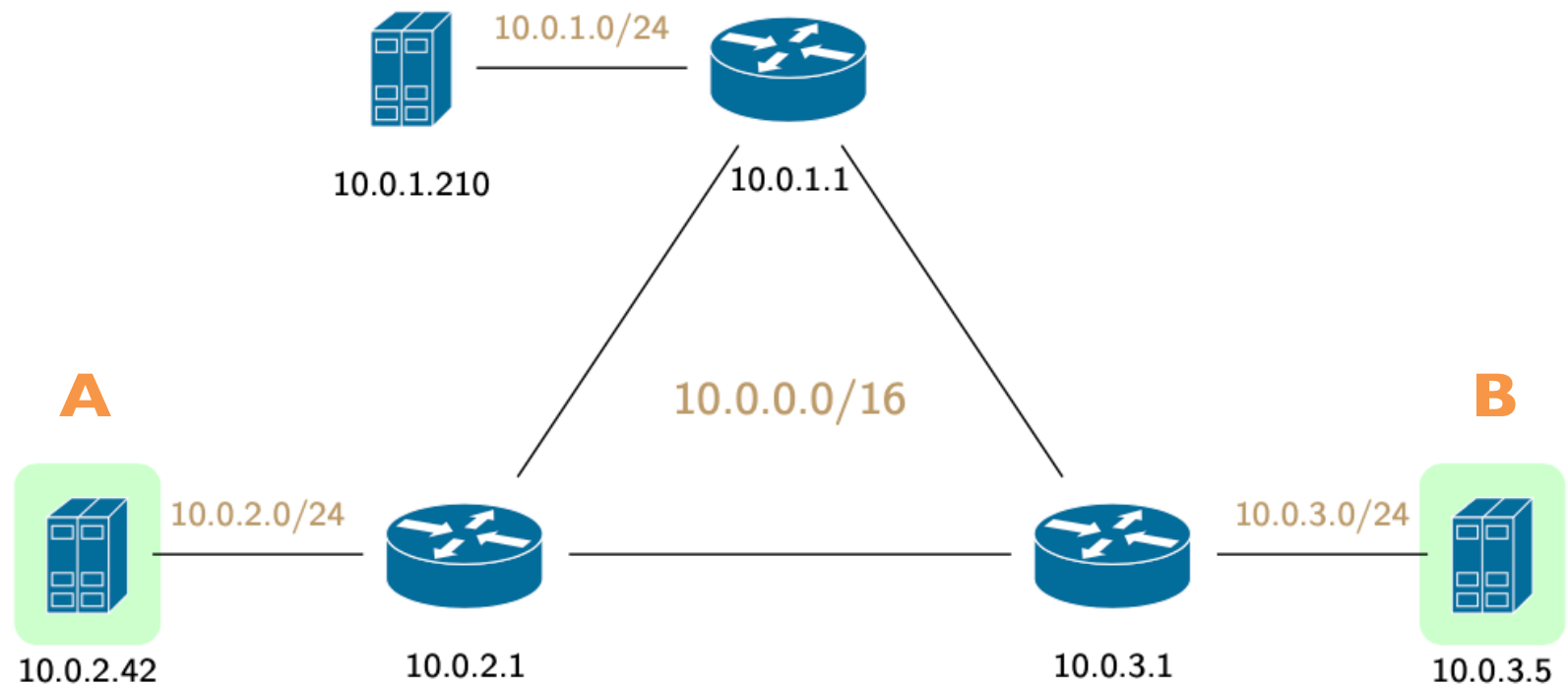
Dolev-Yao: Trusted Areas (TLS)



TLS' end-to-end principle:

Client **A**: If an attacker gains access to **A**'s machine, and can run arbitrary code then they can read and modify messages that A sends to **B**

Dolev-Yao: Trusted Areas (TLS)



TLS' end-to-end principle:

server **B**: if an attacker steals the server's certificate then the attacker can impersonate the server and bad things can happen

SSL vs. TLS



Ruins of Pompeii, Italy

SSL

The “S” in HTTPS.

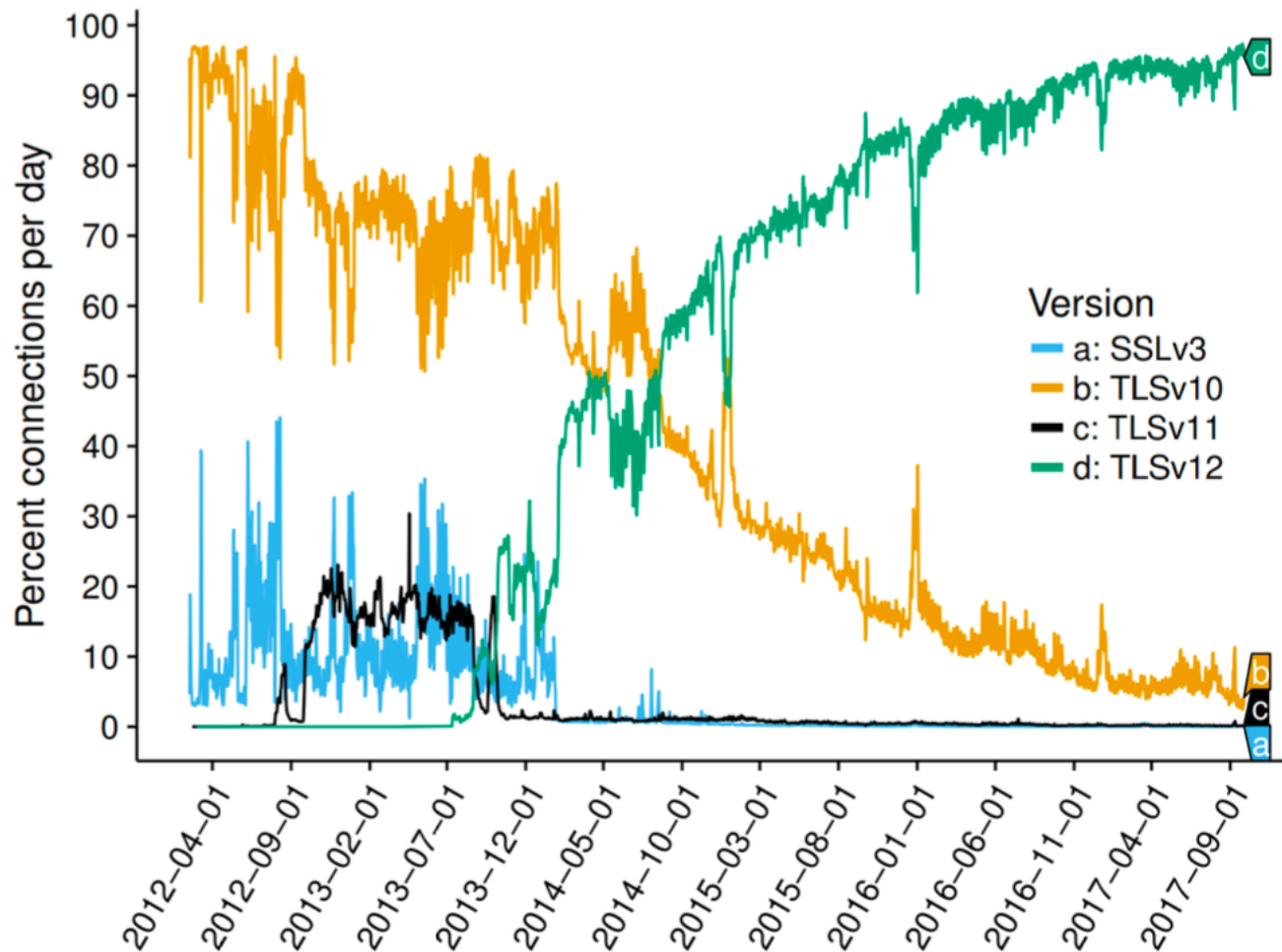


Castle Neuschwanstein, Germany

TLS

HTTPS means HTTP over TLS

TLS Versions



Johanna Amann, Oliver Gasser, Quirin Scheitle, Lexi Brent, Georg Carle, and Ralph Holz. 2017.

TLS Handshake

Alice

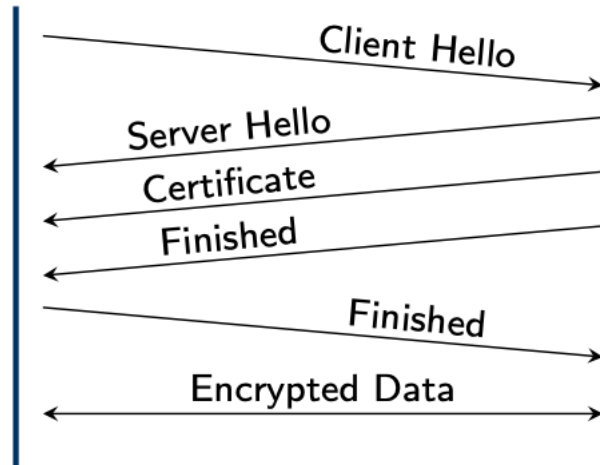


A

Bob



B



Basic TLS 1.3 Handshake

Confidentiality: Client and server negotiate “master key” used to encrypt data.

Integrity: Encryption primitives ensure that content modifications can be detected.

Authenticity: Server proves its identity by presenting certificate. TLS allows for clients to do this, too.

Client's Hello

- TLS version
- cryptographic primitives supported by the client
- domain name of the web site
- everything required by the key-exchange

TLS Handshake

Alice

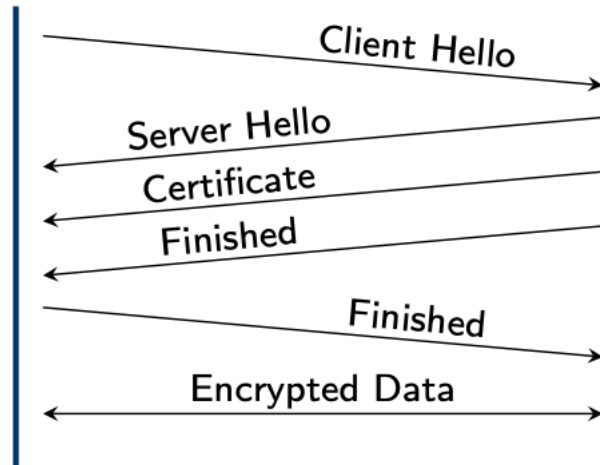


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TLS allows for clients to do this, too.

Server's certificate

- The client uses the certificate to prove that the server is the right entity
- otherwise, the client can end up communicating with an attacker

TLS Handshake

Alice

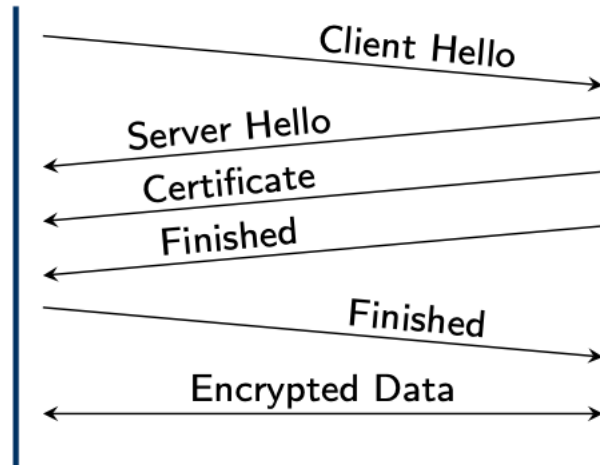


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Basic TLS 1.3 Handshake

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Client

- It does not prove its identity, so the server does not know to whom it's talking.
- There is **no TLS client's authentication**, but client authenticates on the **application layer**, e.g. servers send client a **cookie** over a **secure connection**

Question

Assume Alice and Bob securely communicate over TLS. Mallory controls the network. What can she do?

- ☐ See who is communicating with who.
- ☐ Secretly modify messages between Alice and Bob.
- ☐ Block communication.
- ☐ Impersonate Bob.

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Roadmap

1. Transport Layer Security (TLS)
- 2. Certificate Authority (CA)**
3. Man-in-the-Middle attacks

Establishing Trust

How does Alice know if she is talking to Bob?



Bob presents a certificate!

Establishing Trust

How does Alice know if she is talking to Bob?



Illustration: Olivia Lasting

Establishing Trust

How does Alice know if she is talking to Bob?

Website: example.com

Valid not before: May 4th, 2018

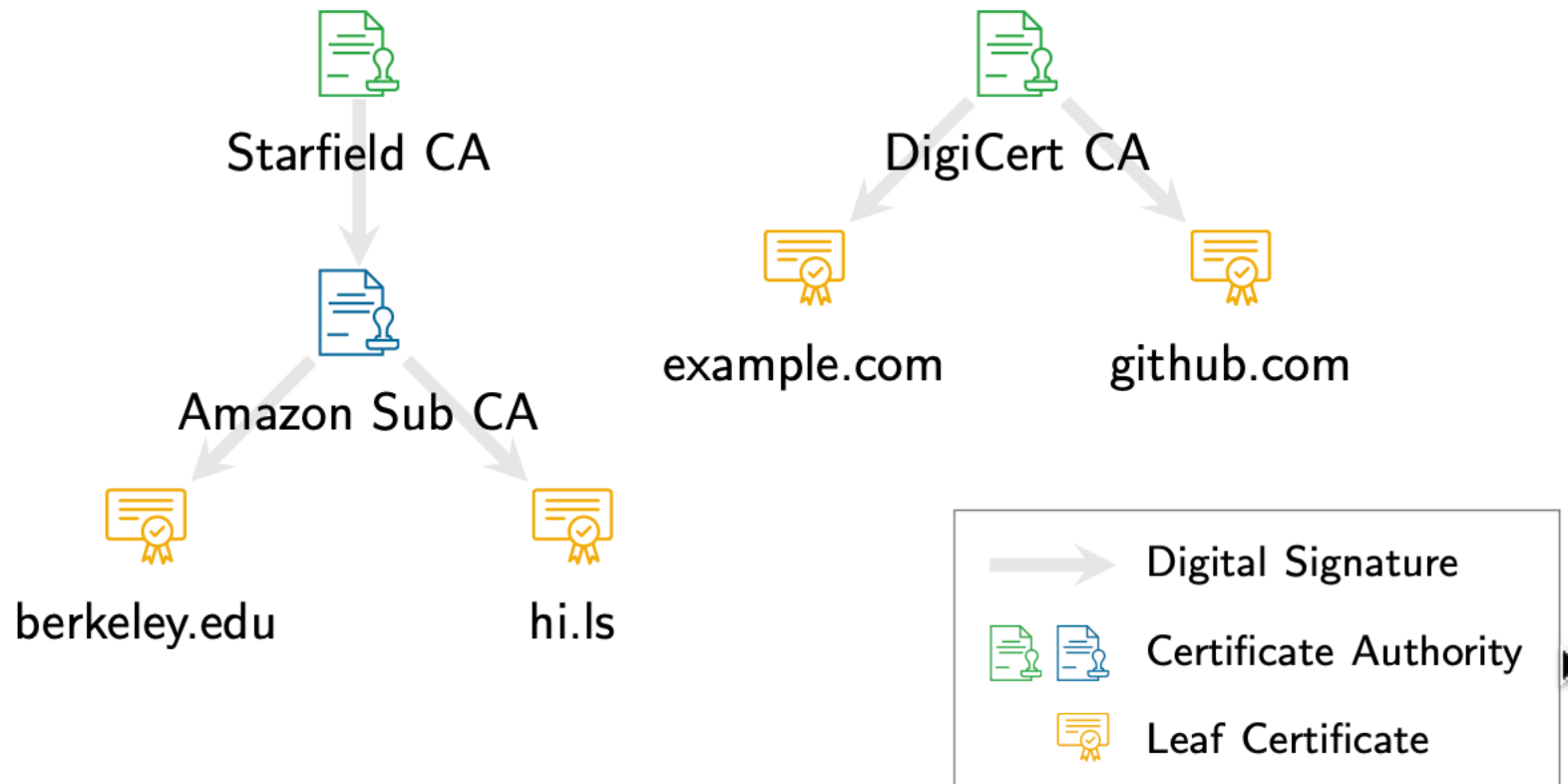
Valid not after: May 4th, 2019

Website Public Key: c5 55 45 d6..

Verified by: DigiCert Inc

Verifier Signature: 60 5f b3 66..

Internet Public Key Infrastructure



CA ~ Certification Authority

Your cell phone has a set of CAs that it trusts!!

TLS Cryptography

- ▶ Client and server agree on the cryptographic primitives during the handshake.
- ▶ Cryptography is hard:
 - ▶ Weak/broken ciphersuites
 - ▶ Implementation flaws
 - ▶ Cryptographic oracles
 - ▶ Downgrade attacks

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The complexity of TLS cryptography is mind-blowing
Open SSL contains 27 000 lines of Perl scripts
to generate assembly code for AES

Question

Check all statements that apply:

- ☐ TLS makes it possible to reduce the trusted areas while maintaining protection goals.
- ☐ Large parts of the internet are using SSLv3.
- ☐ When visiting Facebook, a browser proves your identity by presenting a TLS client certificate.
- ☐ TLS cryptography is a solved problem, hence it is not covered here.
- ☐ A website is trustworthy if it has a valid certificate.
- ☐ TLS does not protect against attacks on availability.

Question

Check all statements that apply:

- ☒ TLS makes it possible to reduce the trusted areas while maintaining protection goals.
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TLS Trust Model Issues



CAs



Certificates



Users

CA Compromise

Any CA can issue certificates for any domain and will be trusted.



 The compromise of a CA puts the whole ecosystem at risk.

CA Compromise

Any CA can issue certificates for any domain and will be trusted.



 The compromise of a CA puts the whole ecosystem at risk.

If an attacker wants a certificate for **Google.com**, the attacker can target the weakest CA in the ecosystem

Certificate Authority Selection

Who can sign certificates?

- ▶ Trusted certificate authorities (CAs) are selected by operating system vendors and some browsers, most notably Firefox.
- ▶ CAs apply for inclusion.
- ▶ This gives operating systems and browsers some negotiation power on behalf of users.
- ▶ On the other side: CAs might be too big to fail.
- ▶ Industry self-regulation:
CA/Browser Forum¹ defines baseline requirements.

Certificate Authority Selection

- ▶ Every operating system and some browsers ship with their own set of trusted certificate authorities (CAs).
 - ▶ Windows: \approx 130 trusted organizations.
- ▶ Root CAs can delegate trust to intermediary authorities.
 - ▶ > 650 organizations
 - ▶ > 50 jurisdictions
- ▶ Attacker can exploit the weakest link.



Certificate Authority Selection

Compelled Certificates



CA may be forced to issue certificates.

Roadmap

1. Transport Layer Security (TLS)

2. Certificate Authority (CA)

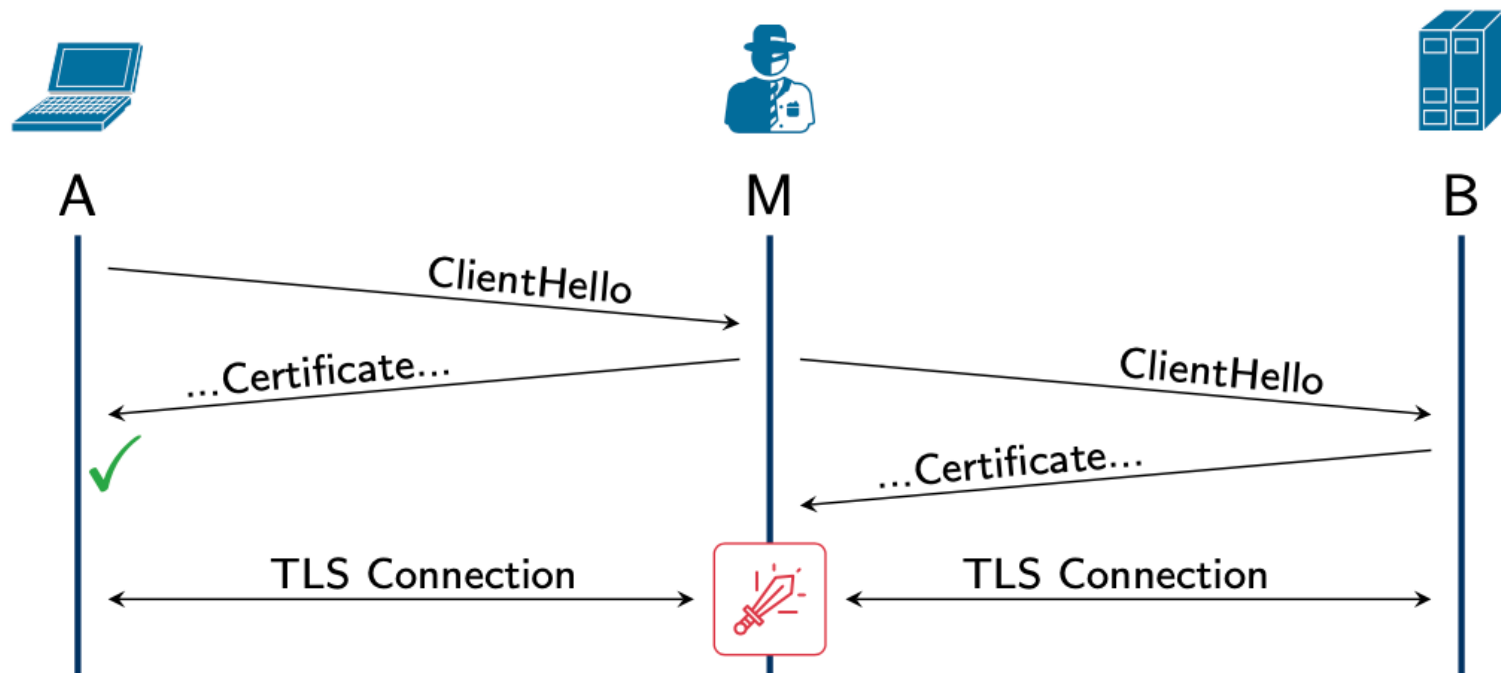
3. Man-in-the-Middle attacks

Man-in-the-middle attacks

How can we compromise protection goals?



Violate end-to-end principle with man-in-the-middle attack.



Attacker's Bucket List

1. Interception software
 - ▶ mitmproxy: <https://mitmproxy.org/>
2. A trusted certificate
 - ▶ Consent of client or server.
 - ▶ Trust model weakness.
 - ▶ Vulnerability in certificate verification.



The Most Dangerous Code in the World: Validating SSL Certificates in Non-Browser Software

Martin Georgiev
The University of Texas
at Austin

Subodh Iyengar
Stanford University

Suman Jana
The University of Texas
at Austin

Rishita Anubhai
Stanford University

Dan Boneh
Stanford University

Vitaly Shmatikov
The University of Texas

We demonstrate that SSL certificate validation is completely broken in many security-critical applications and libraries.

based on it; Amazon's and PayPal's merchant SDKs responsible for transmitting payment details from e-commerce sites to payment gateways; integrated shopping carts such as osCommerce, ZenCart,

shake, when the server presents its public-key certificate. In order for the SSL connection to be secure, the client must carefully verify that the certificate has been issued by a valid certificate authority,

The root causes of these vulnerabilities are badly designed APIs of SSL implementations (such as JSSE, OpenSSL, and GnuTLS) [...]

perns and pitfalls of SSL certificate validation in software based on these APIs and present our recommendations.

cessors such as PayPal and Amazon, (3) logging instant messenger clients into online services, and (4) authenticating servers to mobile

TLS in Non-Browser Software (2018)

Current State of OpenSSL:

- ▶ Certificate hostname is not verified by default.
- ▶ No functionality to use OS root certificate store.
- ▶ No automated revocation checks.
- ▶ TLS version specification is prone to error.



TLS outside of browsers is in a very sad state of affairs.