## Introduction to Software Verification 236342, Homework 1

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## 1

- A. Correct. Since the precondition is false, the postcondition is 'always' satisfied (since it is never tested).
- B. Incorrect. Counterexample x = -100, y = -99:
  - $l_0, -100, -99$
  - $l_1, -100, -99$
  - $l_2, -100, -99$
  - $l_3, -1, -99$
  - $l_*, -1, -99$

As can be seen, precondition is satisfied and postcondition is not.

- C. Correct. The postcondition is true, so regardless of anything else, for every input selection it will be evaluated as true (the program does not even have to terminate either).
- D. Incorrect. Counterexample x = 1, y = 9:
  - $l_0, 1, 9$
  - $l_1, 1, 9$
  - $l_2, -8, 9$
  - $l_3, -8, 9$
  - $l_*, -8, 9$

Postcondition is false, so it is not satisfied.

- E. Incorrect. Counterexample x = 1, y = 3:
  - $l_0, 1, 3$
  - $l_1, 1, 3$

- $l_2, 1, 3$
- $l_3, -2, 3$
- $l_4, -2, 3$
- $l_1, -2, -3$
- $l_2, -2, -3$
- $l_3, -2, -3$
- $l_4, -2, -3$
- $l_1, -2, 3$
- $l_2, -2, 3$
- $l_3, -5, 3$
- $l_4, -5, 3$
- F. Incorrect. Counterexample x = 1, y = 2, By running this example we see that the program gets stuck in a loop at labels  $l_1, l_2, l_3, l_4$ . And each 4 iterations result with the state being the same as the initial state at  $l_1$ . Since this is a total correcness condition on the specification, the correctness is contredicted by the program failing to terminate.

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3

A. To enforce the program to not finish on a spesific set on inputs, we can require that if said inputs have been given and the program finishes - the postcondition fails:

$$\{\forall p \in \mathbf{P}, x = p^2\}P\{false\}$$

B. To require that for a set of inputs a program finishes. we can use the precondition to apply the condition only to the relivant set and use total correctness to require the program to halt on these inputs:

$$< gcd(x,y) = 1 > P < true >$$