



RISC-V

Interrupts

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Interrupt Uses in Different Applications

- High-performance Unix-like systems
 - Interrupt handling small fraction of processing time
 - Fast cores, smart devices
 - Minimal interrupt handler
 - Scheduling in software
- Low/mid embedded systems
 - Interrupt handling significant fraction of processor time
 - Slow cores, dumb devices
 - Significant fraction of code in handlers
 - Interrupt controller acts as task scheduler
- High-performance real-time systems
 - Can't waste time on interrupt overhead
 - Handlers poll I/O devices with regular heartbeat
- And everything inbetween



RISC-V Interrupt Design Goals

- Simplicity
- Support all kinds of platforms from microcontrollers to virtualized servers
- Enable tradeoffs between performance and implementation cost
- Flexibility to support specialized needs



Categorizing Sources of RISC-V Interrupts

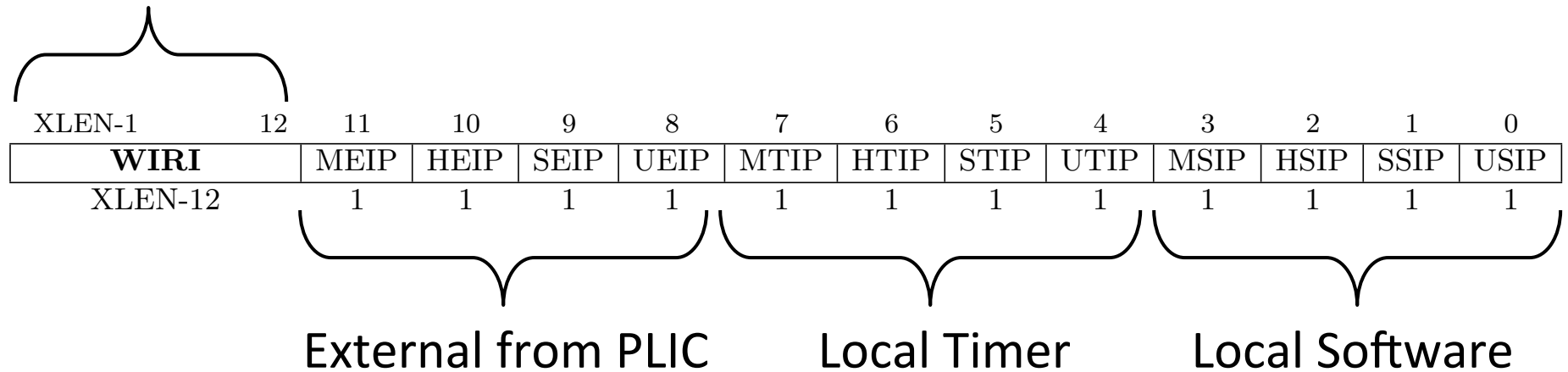
- Local Interrupts
 - Directly connected to one hart
 - No arbitration between harts to service
 - Determine source directly through xcause CSR
 - Only two standard local interrupts (software, timer)

- Global (External) Interrupts
 - Routed via Platform-Level Interrupt Controller (PLIC)
 - PLIC arbitrates between multiple harts claiming interrupt
 - Read of memory-mapped register returns source



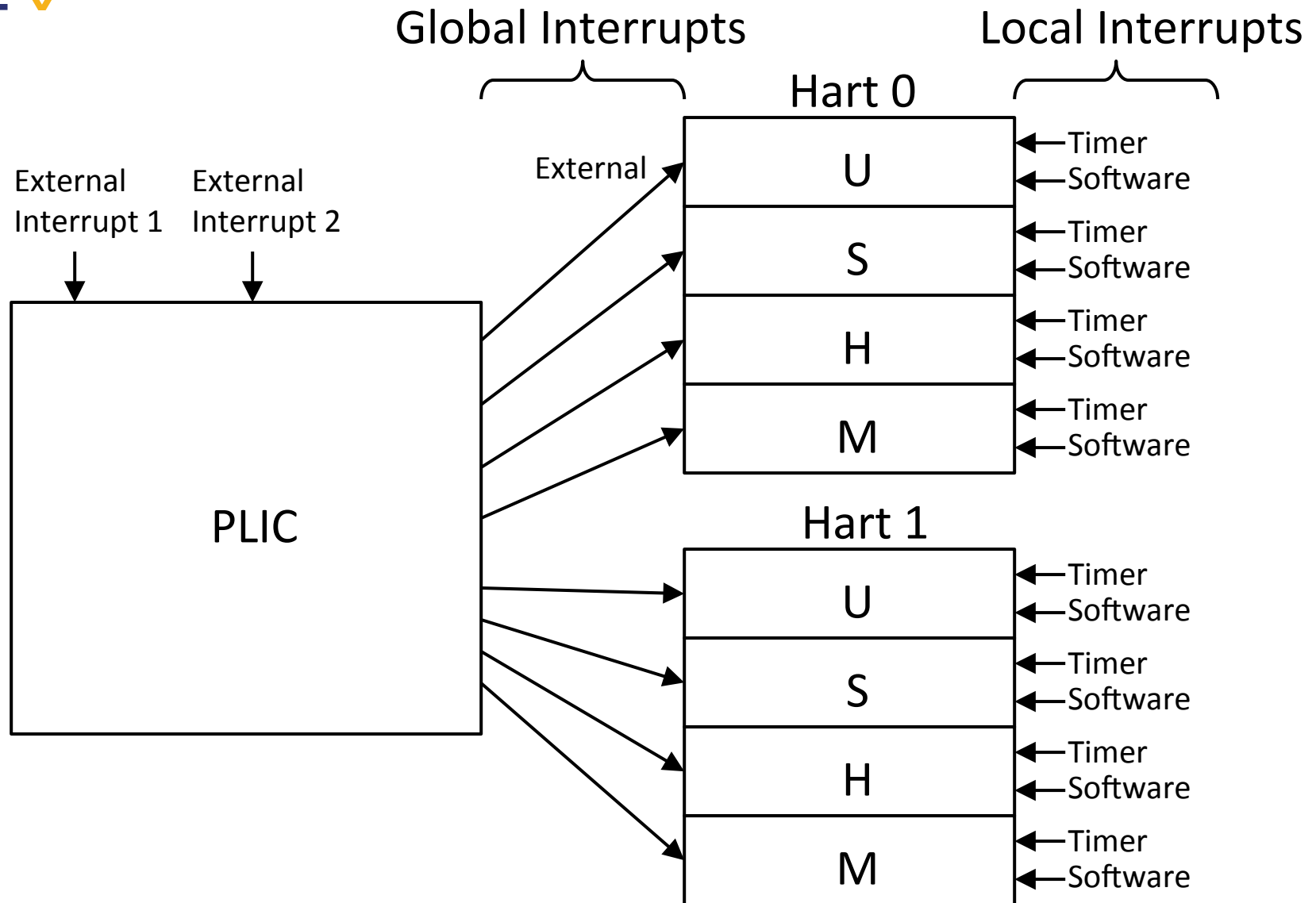
Machine Interrupt Pending CSR (**mip**)

*(Add Non-Standard
Local Interrupts Here)*



- **mip** reflects pending status of interrupts for hart
- Separate interrupts for each supported privilege level (M/H/S/U)
- User-level interrupt handling (“N”) optional feature when U-mode present (discussed later)

Platform-Level Interrupt Controller (PLIC)





Software Interrupts

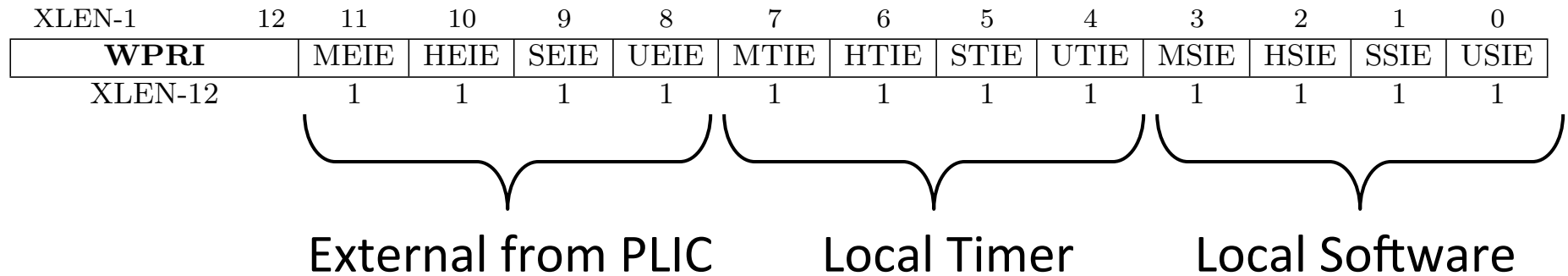
- MSIP
 - Only writeable in machine-mode via memory-mapped control register (mapping is platform-specific)
 - One hart can write to different hart's MSIP register
 - Mechanism for inter-hart interrupts
- HSIP, SSIP, USIP
 - Hart can only write bit xSIP in own **mip** register when running at privilege mode x or greater
- App/OS/Hypervisor can only perform inter-hart interrupts via ABI/SBI/HBI calls
 - Destination virtual hart might be descheduled
 - Interrupts virtualized by M-mode software using MSIP



Timer Interrupts

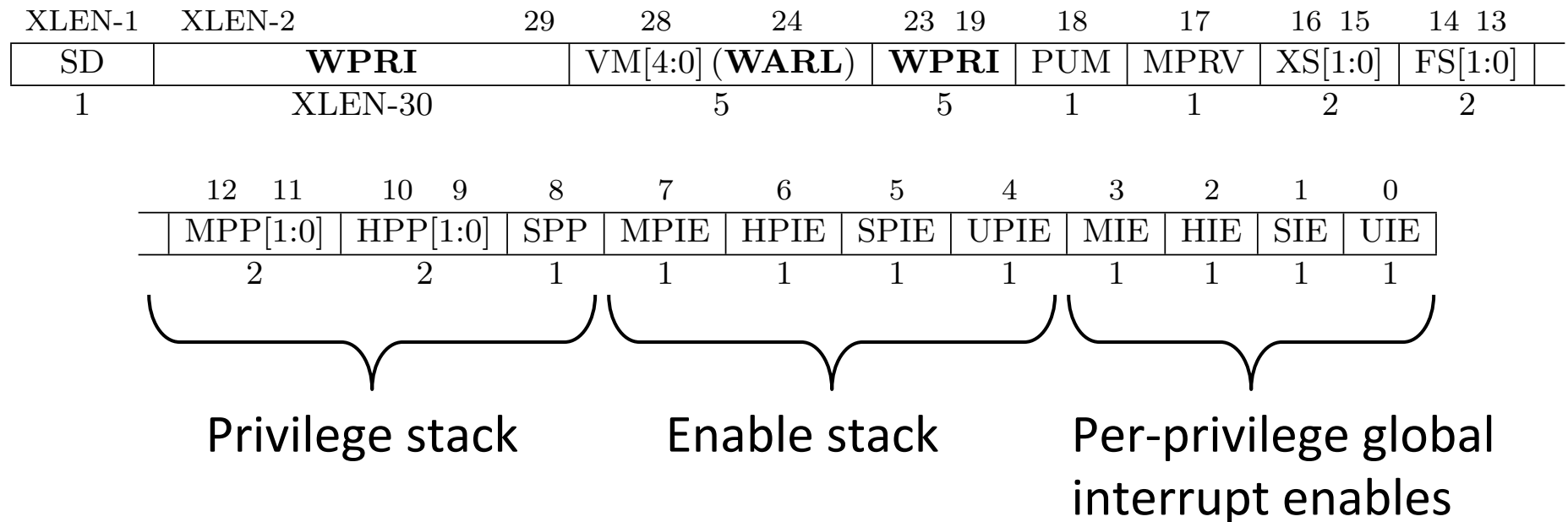
- MTIP
 - Single 64-bit real-time hardware timer and comparator in M-mode
 - MTIP set when **mtime** \geq **mtimecmp**
 - MTIP cleared by writing new **mtimecmp** value
- HTIP, STIP, UTIP
 - M-mode multiplexes single hardware timer and comparator for lower-privilege modes on same hart
 - ABI/SBI/HBI calls to set up timer
 - M-mode software writes/clears HTIP/STIP/UTIP
- Most systems will also have other hardware timers attached via PLIC etc.

Machine Interrupt Enable CSR (**mie**)



- **mie** mirrors layout of **mip**
- provides per-interrupt enables

Interrupts in mstatus



- Only take a pending interrupt for privilege mode x if $xIE=1$ and running in mode x or greater
- Interrupts always disabled for privileges less than current level



All interrupts trap to M-mode by default

- **mcause** CSR indicates which interrupt occurred
- M-mode can redirect to other privilege level by:
 - set up target interrupt and privilege stack
 - copy **mepc** to **hepc/sepc/uepc** respectively
 - copy **mcause** to **hcause/scause/ucause**
 - set **mepc** to target trap vector
 - set MPP to target privilege level, MPIE to false
 - execute **mret**

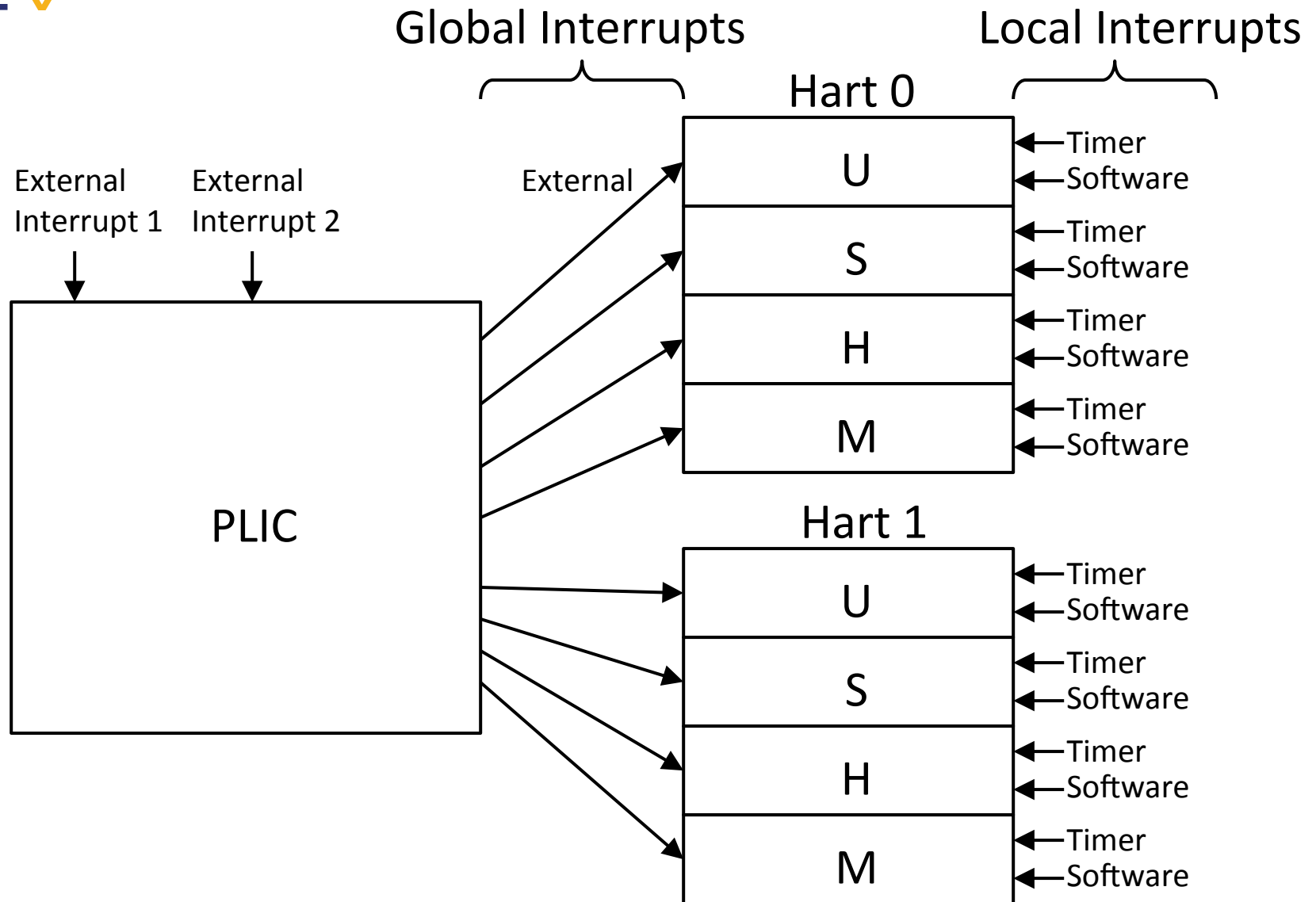
Interrupt	Exception Code	Description
1	0	User software interrupt
1	1	Supervisor software interrupt
1	2	Hypervisor software interrupt
1	3	Machine software interrupt
1	4	User timer interrupt
1	5	Supervisor timer interrupt
1	6	Hypervisor timer interrupt
1	7	Machine timer interrupt
1	8	User external interrupt
1	9	Supervisor external interrupt
1	10	Hypervisor external interrupt
1	11	Machine external interrupt
1	≥ 12	<i>Reserved</i>



Optional Interrupt Handler Delegation

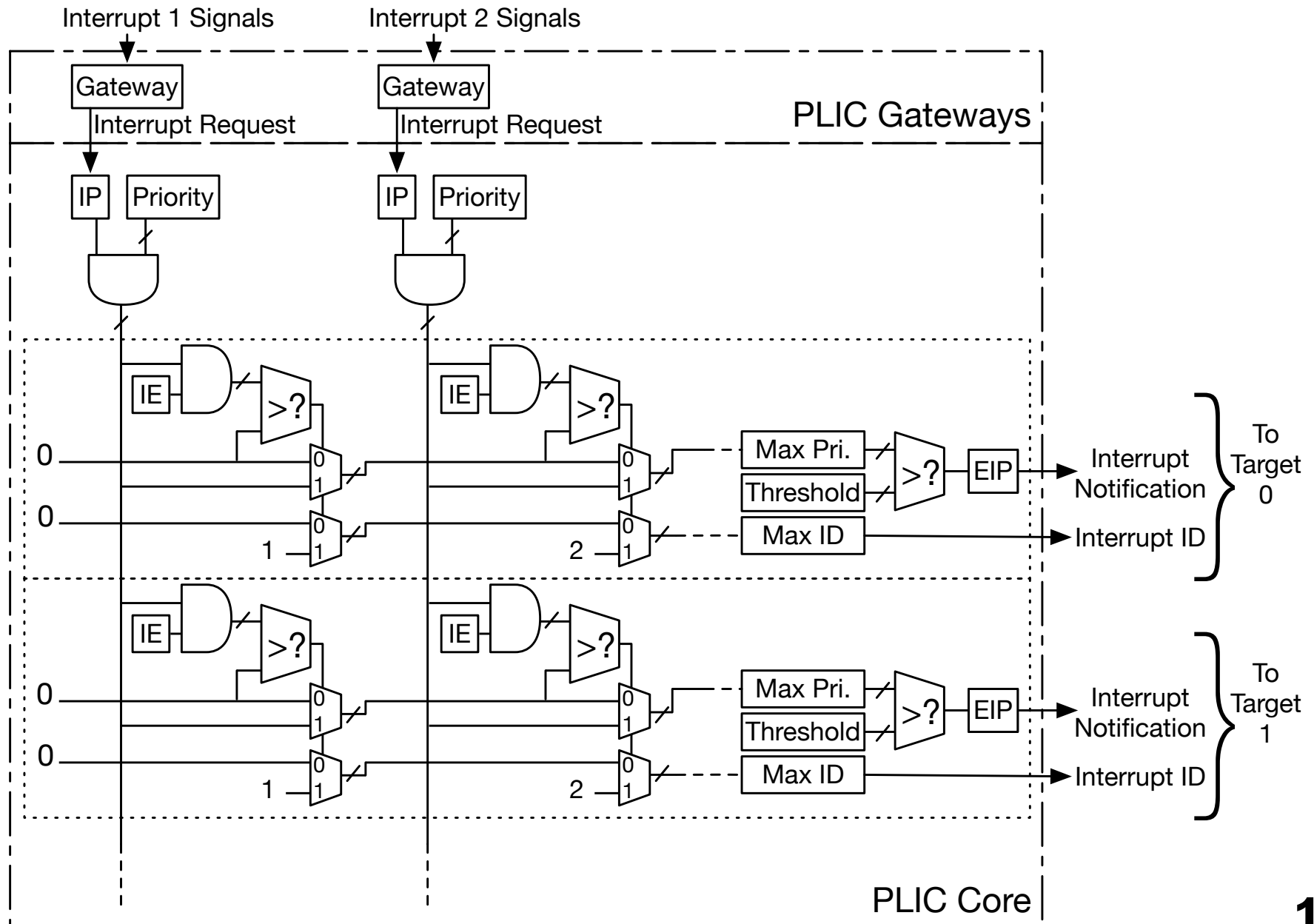
- Can delegate interrupt (and exception) handling to lower privilege level to reduce overhead
- **mideleg** has same layout as **mip**
- If a bit is set in **mideleg** then corresponding interrupt delegated to next lowest privilege level (H, S, or U)
- Can be delegated again using **hideleg/sideleg**
- Once delegated, the interrupt will not affect current privilege level (MIE setting ignored)

Platform-Level Interrupt Controller (PLIC)





PLIC Conceptual Block Diagram





PLIC Interrupt Gateways

Convert from external interrupt signal/message encoding to internal PLIC interrupt request, e.g.,

- Level-triggered gateways
- Edge-triggered gateways
- Message-signaled gateways
- XXX gateways in future

Will not forward a new request to PLIC core unless previous request's handler has signaled completion

- Level-triggered will issue new PLIC interrupt request if level still asserted after completion signaled
- Edge-triggered/message-signaled could queue requests



PLIC Per-Interrupt ID and Priority

- Each interrupt has ID and priority
- Interrupt IDs are integers from 1...N
- ID of zero means “no interrupt”
- Priorities are integers, larger number is higher priority
- Priority zero means “never interrupt”
- Priorities can be fixed or variable
 - Degenerate case, all are fixed at “1”.
- Ties broken by ID (lower ID is higher priority)

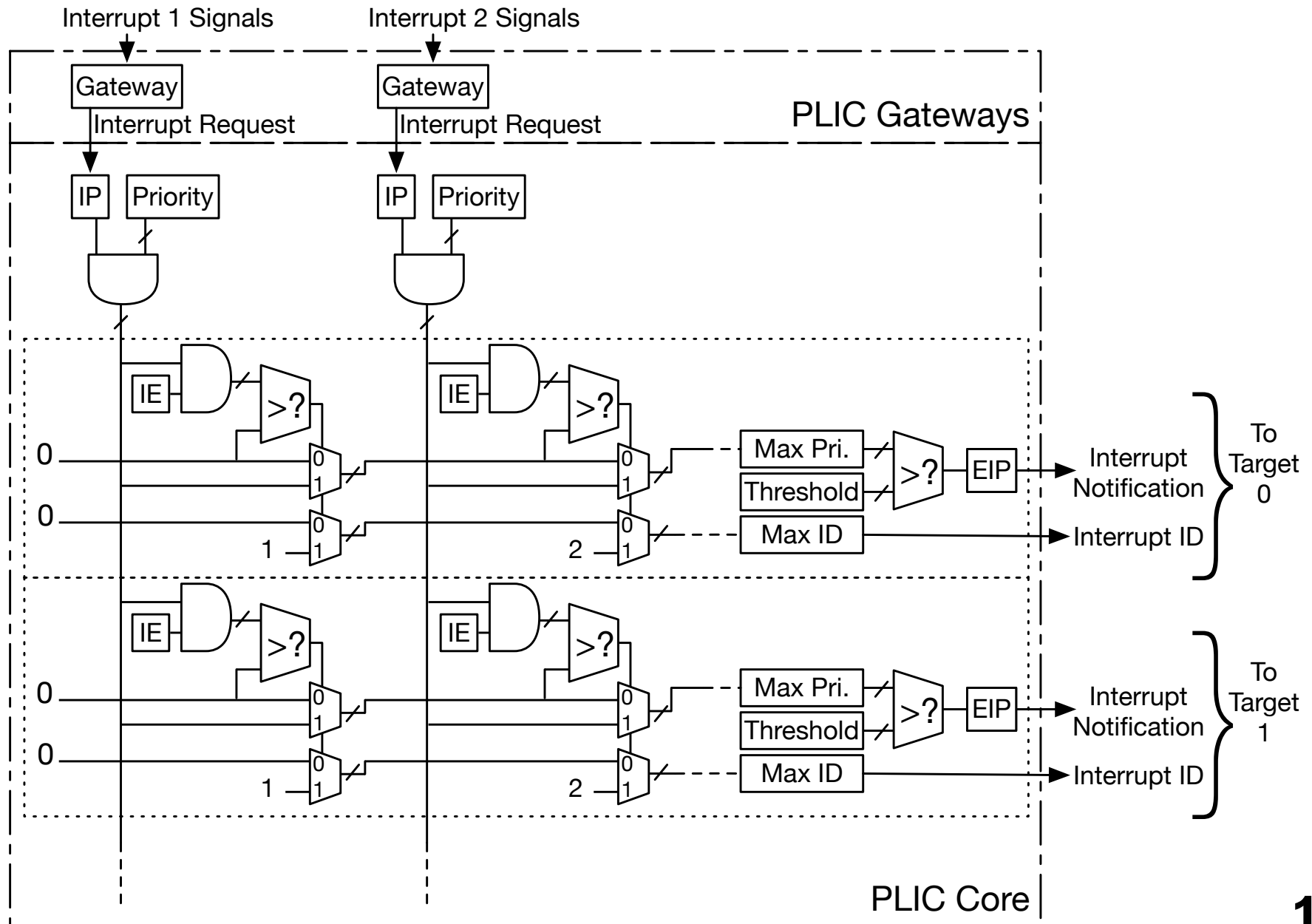


PLIC Per-Target Information

- Each target has vector of interrupt enables
- Each target has priority threshold
- Only interrupts with priority above threshold will cause interrupt
- Set threshold to 0, has no effect
 - Minimal implementation, hardwire threshold to zero
- Set threshold to MAX_PRI, then all interrupts masked
- Interrupt notifications asserted at target if enabled interrupt is above threshold
 - Notifications can take arbitrary time to arrive at target



PLIC Conceptual Block Diagram





PLIC Claim/Completion

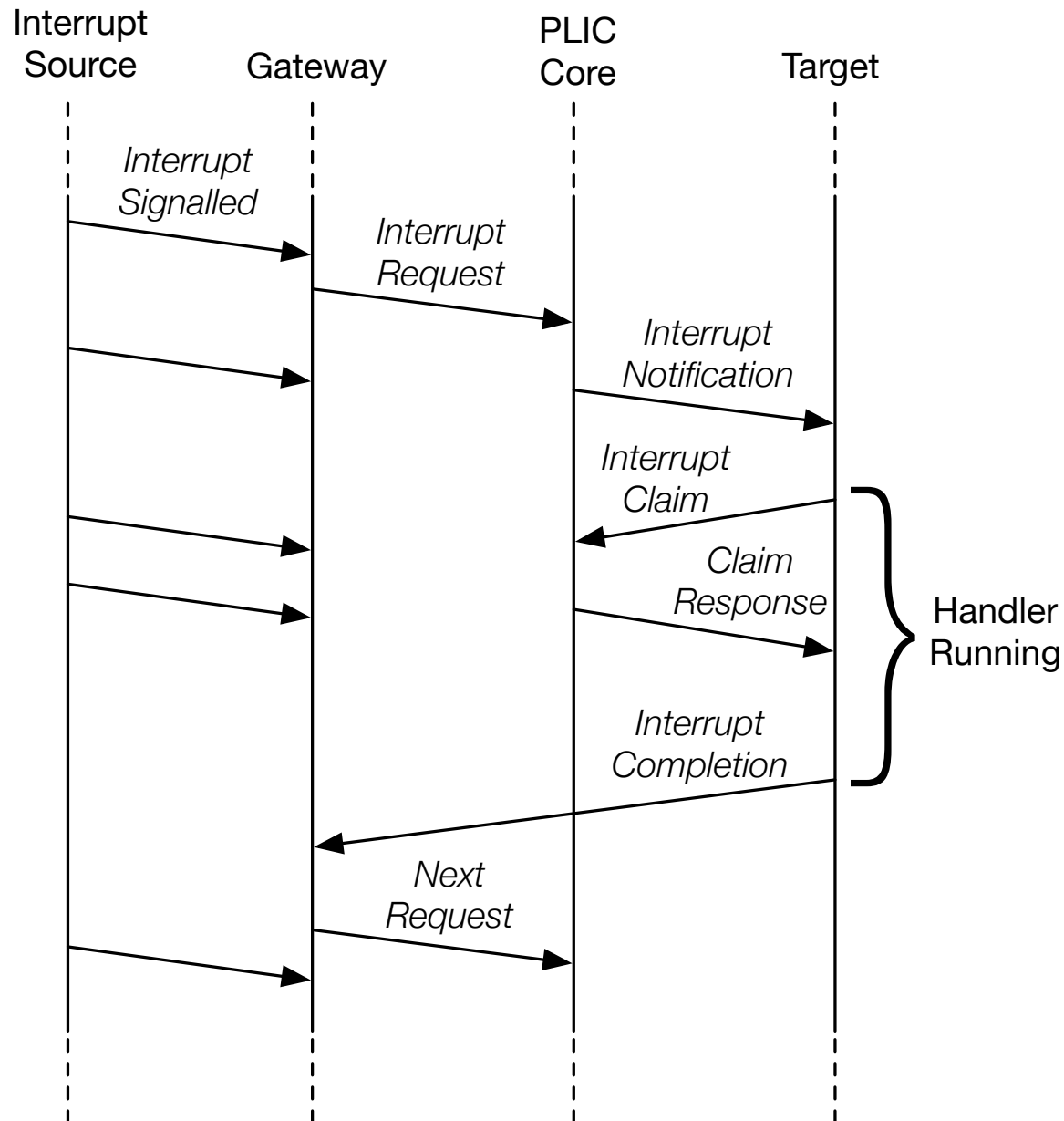
- Interrupted hart context *claims* interrupt from PLIC with read of memory-mapped register
- PLIC returns highest priority active interrupt for that hart
- Can return 0 if no active interrupts remain
 - Other hart might have claimed interrupt earlier
- Hart signals completion to gateway after handler finishes

PLIC Core Atomic Actions

- **Write Register:** A message containing a register write request is dequeued. One of the internal registers is written, where an internal register can be a priority, an interrupt-enable (IE), or a threshold.
- **Accept Request:** If the IP bit corresponding to the interrupt source is clear, a message containing an interrupt request from a gateway is dequeued and the IP bit is set.
- **Process Claim:** An interrupt claim message is dequeued. A claim-response message is enqueued to the requester with the ID of the highest-priority active interrupt for that target, and the IP bit corresponding to this interrupt source is cleared.

- Implementations can perform one action over many cycles, or many actions per cycle, provided behavior agrees with some sequence of these actions

PLIC Interrupt Flow





PLIC Interrupt Preemption/Nesting

- Preemption and nesting are function of the target core, not the PLIC
- Need a different hart context to receive nested interrupt
- Each standard RISC-V privilege level can provide one level of preemption/nesting
 - M-mode interrupt will preempt S-mode handler on hart
- Can add additional hart contexts to core to support nested interrupt handling, with per-cores rules on preemption/priority



PLIC Access Control

- PLIC registers are memory mapped, platform-specific
- M-mode-only access to interrupt enables and priorities
- Lower privilege modes only access claim, completion, and threshold registers
 - can only signal completion for inputs for which they're enabled



SiFive Freedom Platform PLIC Mapping

Address	Description
0x4000_0000	<i>Reserved</i>
0x4000_0004	source 1 priority
0x4000_0008	source 2 priority
...	
0x4000_0FFC	source 1023 priority
0x4000_1000	Start of pending array
...	(read-only)
0x4000_107C	End of pending array
0x4000_1800	
...	<i>Reserved</i>
0x4000_1FFF	
0x4000_2000	target 0 enables
0x4000_2080	target 1 enables
...	
0x401E_FF80	target 15871 enables
0x401F_0000	
...	<i>Reserved</i>
0x401F_FFFC	
0x4020_0000	target 0 priority threshold
0x4020_0004	target 0 claim/complete
0x4020_1000	target 1 priority threshold
0x4020_1004	target 1 claim/complete
...	
0x43FF_F000	target 15871 priority threshold
0x43FF_F004	target 15871 claim/complete

Machine-mode
only

Target per
page to
simplify
protection

Interrupt/Trap Vectors

- By default, single entry point per privilege level:
 - **mtvec/htvec/stvec/utvec**
- Useful in many systems where common handling code used, with bulk of work scheduled later
 - “Interrupt is data”
- Can optionally add differentiated entry points per trap type for embedded applications
 - “Interrupt is control”

User-Level Interrupts “N”

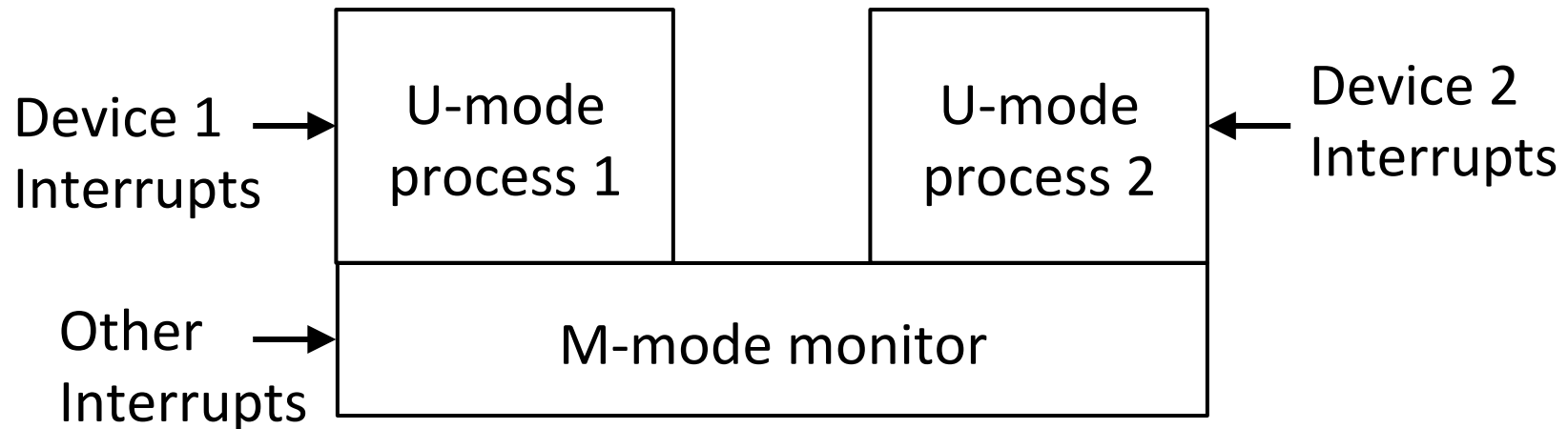
- Natural extension of interrupt model into user permissions
- Adds user CSRs, and **uret** instruction

Number	Privilege	Name	Description
User Trap Setup			
0x000	URW	ustatus	User status register.
0x004	URW	uie	User interrupt-enable register.
0x005	URW	utvec	User trap handler base address.
User Trap Handling			
0x040	URW	uscratch	Scratch register for user trap handlers.
0x041	URW	uepc	User exception program counter.
0x042	URW	ucause	User trap cause.
0x043	URW	ubadaddr	User bad address.
0x044	URW	uip	User interrupt pending.



Interrupts in Secure Embedded Systems (M, U modes)

- M-mode runs secure boot and runtime monitor
- Embedded code runs in U-mode
- Physical memory protection on U-mode accesses
- Interrupt handling can be delegated to U-mode code
- Provides arbitrary number of isolated subsystems





Questions?