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# Interpreting Pumping Lemma for Regular languages

Here is one way to interpret the Pumping Lemma for regular languages. Credit is due to Shai Halevi, who taught this in a recitation I was in at MIT in 1995.

## **Pumping Lemma**

```
If L is regular
                                                       then it has a DFA D recognizing it
                                                       where p is the number of states of D
then there exists p \geq 0
so that for all s \in L
                                                       which defines a path in D
   with |s| \geq p
                                                          so this path contains a cycle
there is a parse of s = xyz
                                                       x is everything before the first cycle
  with |y| \ge 1 and |xy| \le p
                                                       y is everything on the first cycle
                                                       z is everything after the first cycle
such that for any i > 0
                                                       no matter how many times we repeat the cycle
xy^iz \in L
                                                       we end up in a final state of D
```

## Showing that L is not regular

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For any integer p \ge 0 \Longrightarrow show that L is not recognizable by any DFA with p states suppose toward a contradiction that there was a DFA D having p states and recognizing L there exists a string s \in L such that |s| \ge p \Longrightarrow so in processing s we must visit more than p states so it defines a path with a cycle so that for any parse s = xyz \Longrightarrow no matter where this cycle is \Longrightarrow there exists some i \ge 0 \Longrightarrow we can repeat the cycle i times so that xy^iz \notin L \Longrightarrow and end up in a non-final state of D
```

### Example

Let the alphabet be  $\{a,b,c\}$ . Show that  $L=\{a^n w \mid w \in \{b,c\}^* \text{ and } |w|=2n\}$  is not regular.

#### Proof

```
Let p \geq 0. Consider s = a^p b^p c^p \in L.

Let s = xyz with |y| \geq 1 and |xy| \leq p.

Let i,j be integers such that i \geq 0 and j \geq 1 and i+j \leq p and x = a^i \quad \text{and} \quad y = a^j \quad \text{and} \quad z = a^{p-i-j} b^p c^p \ .
```

Consider  $xy^2z$ . We have

$$xy^{2}z = a^{i}a^{2j}a^{p-i-j}b^{p}c^{p}$$
$$= a^{i+2j+p-i-j}b^{p}c^{p}$$
$$= a^{p+j}b^{p}c^{p}$$

Since  $j \ge 1$ , we have that  $2(p+j) \ne p+p$ . So,  $xy^2z \notin L$ . Thus, L is not regular.