Authenticated Encryption

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Agenda: Authenticated Encryption

- 1. Authenticity is often necessary even when all you care about is privacy.
 - 1.1 Attack against CBC (TCP/IP)
 - 1.2 Attack against CTR with checksum (a remote terminal application)
 - 1.3 Side-channel attack against non-atomic decryption (SSH BPP)
- 2. Authenticated Encryption is the answer.
- 3. Generic composition: GCM, CCM, EAX
- 4. Handcrafted construction: OCB
- 5. Generic composition: caveats

TCP/IP uses CBC with random IV.

- 1. Usually webserver listens on port 80.
- 2. Suppose an attacker A listens on port 25.
- 3. Suppose client uses IPsec to encrypt data for webserver.
- 4. If the client only uses encryption to send M, the packet can be modified by an attacker so that it can get M on port 25.

Client uses CBC with random IV to send M to port 80:

$$C \leftarrow (IV, CBC_K(IV, dest = 80 \parallel M))$$

A can change C to

$$C' \leftarrow (IV', CBC_K(IV, dest = 80 \parallel M))$$

where

$$IV' = IV \oplus (\dots 80 \dots) \oplus (\dots 25 \dots)$$

To see that this attack works, try computing $M[0] = D_K(C[0]) \oplus IV$.

You can change the decryption to be whatever you want by manipulating ${\it IV}$.

CBC with random IV is IND-CPA but still isn't enough to ensure privacy!

Adding checksum to CTR mode is not enough to get authenticity.

Consider a remote terminal application. Suppose each keystroke is encrypted with CTR mode.

TCP/IP packet:

$$\mathsf{IP}\;\mathsf{hdr}\;\|\;\mathsf{TCP}\;\mathsf{hdr}\;\|\;\;T\;\;\|\;\;M$$

where T is a 16-bit checksum and M is a 1-byte keystroke.

TCP/IP packet: T is a 16-bit checksum and M is a 1-byte keystroke.

IP hdr
$$\parallel$$
 TCP hdr \parallel T \parallel M

Attack:

- 1. Try for many t and s,
 - 1.1 submit to the server

| IP hdr || TCP hdr ||
$$\oplus t$$
 || $\oplus s$

- 1.2 Watch the server whether it says the packet was valid.
- 2. Collect many equations of the form:

$$checksum(hdr, M) \oplus t = checksum(hdr, M \oplus s)$$

3. For some checksum, it is easy to compute M.

Attacking SSH Binary Packet Protocol

Weaknesses

- 1. Packet length is encrypted but is used immediately after decryption but before the MAC is verified.
- 2. Decryption is non-atomic.

Authenticated Encryption is Needed! Integrity of Ciphertext

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Subroutines

Subroutine Initialize K \overset{\$}{\leftarrow} KG \; ; \; S \leftarrow \emptyset \; ; \; \text{win} \leftarrow \text{false}

Subroutine \text{Enc}(M)
C \overset{\$}{\leftarrow} \text{Enc}(K,M) \; ; \; S \leftarrow S \cup \{C\}
Return C

Subroutine \text{Vf}(C)
M \leftarrow \text{Dec}(K,C)
If M \neq \bot and C \not\in S then win = true If M = \bot then return 0 else return 1

Subroutine Finalize Return win
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Experiment \operatorname{Exp}^{\operatorname{int-ctxt}}_{\operatorname{AE}}(A)

Initialize d \stackrel{5}{\sim} A^{\operatorname{Enc}, \operatorname{Vf}}
Return Finalize
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We define the **int-ctxt advantage** of an adversary *A* mounting an attack against integrity of ciphertexts of AE as

$$\textbf{Adv}^{\mathrm{int\text{-}ctxt}}_{\mathsf{AE}}(A) = \mathsf{Pr}\left[\; \textbf{Exp}^{\mathrm{int\text{-}ctxt}}_{\mathsf{AE}}(A) \Rightarrow \mathsf{true}\;\right] \;.$$

Authenticated Encryption

Authenticated Encryption

 $AE = IND-CPA \wedge INT-CTXT$

Syntax is slight different from that of encryption also.

Generic composition

Generic composition methods

Encrypt-and-MAC, Encrypt-then-MAC, MAC-then-encrypt

In practice,

- ► SSL uses MtE.
- ► IPsec uses EtM.
- ► SSH uses E&M.

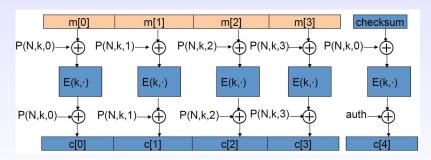
AE schemes in practice

- ► GCM (EtM): CTR\$ then CW-MAC
- ► CCM (MtE): CBC-MAC then CTR (802.11i)
- ► EAX (EtM): CTR then CMAC

All are nonce-based and supports AEAD. All are endorsed by NIST.

AE schemes in practice (cont.)

OCB is a dedicated AE scheme constructed from PRP. One application of PRP per block rather than two.



Generic composition: caveats

It is not as simple as Encrypt-then-MAC being the "best" composition method!