You are using countermodel generation for classical predicate logic.

Tableau for
$\exists y \forall x R(x,y)$
$\not\models \forall x \exists y R(x,y)$ :

3. $\forall x R(x,a)$ $(\exists,\ 2,\ [y/a]*)$

| | | |
|---|---|---|
| 4. | $\neg\exists y R(a,y)$ | $(\neg\forall,\ 1,\ [x/a])$ |
| 5. | $\neg R(a,a)$ | $(\neg\exists,\ 4,\ [y/a])$ |
| 6. | $R(a,a)$ | $(\forall,\ 3,\ [x/a])$ |
| | $\times$ | (5, 6) |

| | | |
|---|---|---|
| 11. | $\neg\exists y R(b,y)$ | $(\neg\forall,\ 1,\ [x/b]*)$ |
| 12. | $\neg R(b,a)$ | $(\neg\exists,\ 11,\ [y/a])$ |
| 13. | $R(a,a)$ | $(\forall,\ 3,\ [x/a])$ |
| 14. | $\neg R(b,b)$ | $(\neg\exists,\ 11,\ [y/b])$ |
| 20. | $R(b,a)$ | $(\forall,\ 3,\ [x/b])$ |
| | $\times$ | (12, 20) |

| | | |
|---|---|---|
| 30. | $\neg\exists y R(c,y)$ | $(\neg\forall,\ 1,\ [x/c]*)$ |
| 31. | $\neg R(c,a)$ | $(\neg\exists,\ 30,\ [y/a])$ |
| 32. | $R(a,a)$ | $(\forall,\ 3,\ [x/a])$ |
| 33. | $\neg R(c,c)$ | $(\neg\exists,\ 30,\ [y/c])$ |
| 34. | $R(c,a)$ | $(\forall,\ 3,\ [x/c])$ |
| | $\times$ | (31, 34) |

| | | |
|---|---|---|
| 59. | $\neg\exists y R(d,y)$ | $(\neg\forall,\ 1,\ [x/d]*)$ |
| 60. | $\neg R(d,a)$ | $(\neg\exists,\ 59,\ [y/a])$ |
| 61. | $R(a,a)$ | $(\forall,\ 3,\ [x/a])$ |
| 62. | $\neg R(d,d)$ | $(\neg\exists,\ 59,\ [y/d])$ |
| 63. | $R(d,a)$ | $(\forall,\ 3,\ [x/d])$ |
| | $\times$ | (60, 63) |

| | |
|---|---|
| 8. | $\neg\exists y R(b,y)$ |
| 9. | $\neg R(b,b)$ |
| 10. | $R(b,b)$ |
| | $\times$ |

The tableau is potentially infinite:
The inference may or may not be refutable.
This computation took 0.364 seconds.