Enumeration of self-avoiding walks via the lace expansion

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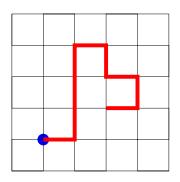
The self-avoiding walk model

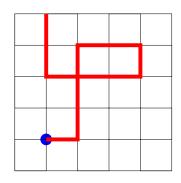
- A self-avoiding walk (SAW) is a path on a lattice, which starts at the origin and hops successively to neighbouring lattice sites without intersecting itself.
- We count the number of SAWs of length n, cn, and in particular study the critical behaviour of the generating function

$$C(x) = \sum_{n=0}^{\infty} c_n x^n$$

■ For the simple cubic lattice $c_0 = 1$, $c_1 = 6$, $c_2 = 30$, $c_3 = 150$, $c_4 = 726$, $c_5 = 3534$, . . .

SAW example





■ SAW!

Not a SAW, due to self intersection.

Known results

■ For the square lattice c_n has been enumerated to very high order via the finite lattice method by Iwan Jensen:

$$c_{71} = 4190893020903935054619120005916$$

- Best results for d > 2 are via direct enumeration.
- MacDonald et al., simple cubic lattice

$$c_{26} = 549493796867100942$$

■ For d = 4, c_n is known to c_{19} , for d = 5 to c_{15} , and for d = 6 to c_{14} .

Known results

It is universally believed (but not proven) that for dimensions d < 4 asymptotically</p>

$$c_n \sim A n^{\gamma-1} \mu^n [1 + \text{corrections}]$$

and the mean-square end-to-end distance behaves asymptotically as

$$\bar{\rho}_n \sim D \, n^{2\nu} \mu^n [1 + \text{corrections}]$$

Improved enumerations allow better estimation of A, μ , γ , D and ν .

Known results for $d \ge 3$

- For the cubic lattice, no exact results available; best estimates $\mu = 4.68404(9)$ (from exact series), $\gamma = 1.1575(6)$, $\nu = 0.5874(2)$ (from Monte-Carlo).
- $ightharpoonup \gamma = 1$ for d = 4, but with logarithmic corrections.
- For $d \ge 5$ it has been proven that $\gamma = 1$ using the lace expansion.

The two-step method

For brute force enumeration of SAWs of length *n* expect

$$\tau(n) \sim c_n \approx \mu^n$$

- Reduce the time taken by reducing the complexity.
- Take two steps at once, and represent walks which have the same endpoint by a single configuration.
- Sequence of endpoints defines a 2-step walk, Ω.

LThe two-step method



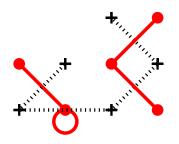
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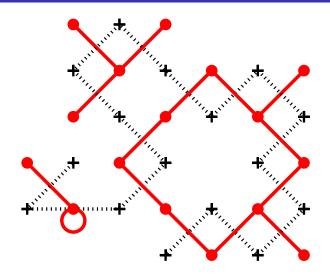


LThe two-step method





L_The two-step method



The two-step method

- Weight of a 2-step walk can be calculated in linear time in the size of the allocation graph.
- By counting many configurations at once, the two-step method reduces the complexity.
- Straightforward to prove this for $d \ge 3$ (also true for d = 2, but no proof yet).

The lace expansion

- The lace expansion, originally due to Brydges and Spencer, is a method that has been used to study the critical behavior of SAWs and related models above their critical dimension.
- Has not been applied to an enumeration problem before.

The lace expansion

■ The number of SAWs of length *n* may be obtained from the following recursion relation,

$$c_n = 2dc_{n-1} + \sum_{m=2}^n \pi_m c_{n-m}.$$

where π_m can be expressed in terms of the number of lace graphs of length m with N loops:

$$\pi_m = \sum_{N} (-1)^N \pi_m^{(N)}$$

Lace graphs are less numerous, and therefore easier to count!

The lace expansion

- First of these graphs are paths that avoid themselves until they return to the origin, i.e. graphs which form a single loop.
- Then graphs with 2, 3, 4, ... loops, represented by the following diagrams.

 $\pi^{(1)}$: start at the origin, must return to the origin.



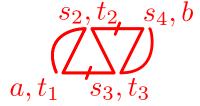
 $\pi^{(2)}$: start at the origin, return to the origin, continue until the first loop is intersected.

$$a, t_1 \bigcirc s_2, t$$

 $\pi^{(3)}$:

$$a, t_1 \stackrel{s_2, t_2}{\sum}_{s_3, b}$$

π⁽⁴⁾:



Why are there fewer lace graphs?

- Lace graphs can be thought of as generalised polygons; there are less of them because they are less spatially extended than SAWs.
- But only by a polynomial factor! Algorithmic complexity is unchanged, and is still given by the connective constant μ. i.e. for any N

$$\pi_m^{(N)} \sim \mu^m$$

■ For d = 3, n = 30,

$$c_{30}\approx 525\pi_{30}$$

Enumeration results

Simple cubic lattice

$$c_{30} = 270569905525454674614$$
 $c_{26} = 549493796867100942$
 $c_{30}/c_{26} = 492.3 \cdots$

■ Hypercubic lattice, d = 4

$$c_{24} = 124852857467211187784$$

 $c_{19} = 8639846411760440$
 $c_{24}/c_{19} = 14450.8\cdots$

Enumeration results

■ Hypercubic lattice, d = 5

$$c_{24} = 63742525570299581210090$$
 $c_{15} = 192003889675210$
 $c_{24}/c_{15} = 3.3 \times 10^{8}$

■ Hypercubic lattice, d = 6

$$c_{24} = 8689265092167904101731532$$

 $c_{14} = 373292253262692$
 $c_{24}/c_{14} = 2.3 \times 10^{10}$

Our enumerations for n = 24 allow us to calculate c_{24} for any dimension.

1/d expansion for μ

■ 1/d asymptotic expansion for the connective constant, with error estimate,

$$\mu = 2d - 1 - \frac{1}{2d} - \frac{3}{(2d)^2} - \frac{16}{(2d)^3} - \frac{102}{(2d)^4}$$

$$- \frac{729}{(2d)^5} - \frac{5533}{(2d)^6} - \frac{42229}{(2d)^7} - \frac{288761}{(2d)^8}$$

$$- \frac{1026328}{(2d)^9} + \frac{21070667}{(2d)^{10}} + \frac{780280468}{(2d)^{11}} + O\left(\frac{1}{(2d)^{12}}\right)$$

Similar formulae for A and D.

Analysing the series

- Complicated, because series are far from the asymptotic regime, and the corrections to scaling are large.
- Used differential approximants and direct fitting of the asymptotic form.
- For d=3, we obtain the most accurate value available for $\mu=4.684043(12)$ (c.f. $\mu=4.68404(9)$). Estimates of critical exponents of comparable accuracy to the best available Monte-Carlo estimates.
- Estimates of μ for d ≥ 4 are competitive with those obtained via the PERM algorithm by Owczarek and Prellberg.

- Significantly extended SAW series for $d \ge 3$ by combining two algorithmic improvements, the lace expansion and the two-step method.
- Obtained improved estimates of critical parameters through detailed series analysis.
- Both the lace expansion and the two-step method can be reformulated and applied to enumeration problems on arbitrary graphs.

- Preprint available at http://www.math.ubc.ca/~slade/se.pdf (submitted to J. Phys. A).
- All enumeration data at http://www.math.ubc.ca/~slade/lacecounts