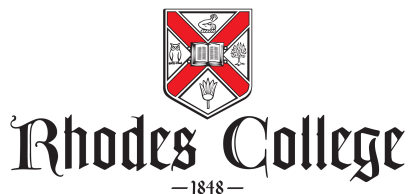


COMP 231-01

Introduction to Computer Organization

Final Exam Review Solutions



For this exam, the main addition to the exam is **functions** in assembly language and the register conventions we talked about in class. When you're designing assembly language programs on the test, you're responsible for using registers correctly according to the conventions. I recommend studying and reviewing the **register conventions** section of the assembly textbook and then working on coding the following exercises. We're also covering the FDXW cycle and memory, but those questions will appear as short answer problems.

To study for this exam, I recommend reviewing your work for the course, most especially the exam reviews and tests, as well as this final exam review. This exam is cumulative and will include questions from across the semester.

- The function detailed below is a recursive C function for calculating the result of an exponentiation operation. Note that this function is designed for integers greater than or equal to 0.

```
int exponentiate(int base, int power)
{
    if(power <= 0)
    {
        return 1;
    }
    else
    {
        return base * exponentiate(base, power-1);
    }
}
```

Translate this function into assembly language, using the calling conventions.

```

#Note that this function assumes the stack pointer has already been
#initiated appropriately in the main function.
#base is passed in through r4 and power through r5, by the register conventions
exponentiate:
    bgt r5, r0, ONERETURN
    subi sp, sp, 8 #Reserve stack space for two values

    subi r5, r5, 1 #We can pass this without saving its original value since we don't use it aft

    stw r4, 0(sp)
    stw ra, 4(sp)

    call exponentiate

    ldw ra, 4(sp)
    ldw r4, 0(sp)

    mul r2, r2, r4

    addi sp, sp, 8 #Deallocating stack space.

    ret

ONERETURN:
    movi r2, 1
    ret

```

- The code below is an assembly language program that has been improperly translated and does not follow the calling conventions. Find and correct the errors in this program. (There are intended to be 8 errors. Some will be related to the calling conventions, but others will be related to other aspects of the program.)

```

\#This function takes the values in memory location 1000 and applies the MATHFUNC to it.
\#Then, this function adds that same value (from MEM 1000) with the results of the MATHFUNC operation
\_start
ldw r9, 1000

call MATHFUNC

add r0, r9, r2

MATHFUNC: \#This function triples the argument value and adds 2 to it, then returns that value.
call FUNC
addi r9, 2
ret

FUNC: \#This function takes in an argument, triples the given value, and returns that value.
multi r9, r9, 3

```

Error list:

- `_start` doesn't have a colon (`_start:`).
- `ldw` takes in an offset/register value; this could be written correctly as `ldw r9, 1000(r0)`.
- `r9` is a caller-save register, and we don't save it before making a call!
- `r0` can not be written to; it always contains 0.
- `MATHFUNC` - `r9` is a caller-save register, and we don't save it before making a call!
- `MATHFUNC` - the argument (`r9`) is not passed in! `r9` should be stored in `r4`.
- `FUNC` - the argument (`r9`) is not passed in! `r9` should be stored in `r4`.
- `MATHFUNC` - the return value is not passed out in `r2`.
- `FUNC` - the return value is not passed out in `r2`.
- `MATHFUNC` does not save the `ra`! This causes an infinite loop.
- `addi` requires three arguments; this could be correctly written `addi r9, r9, 2`.
- `FUNC` - there is no return statement!
- There is no branch statement after `main` and before `MATHFUNC`...and so execution will continue into `MATHFUNC`, which is not intended.

This program ended up with a lot more errors than just 8 (13, in total). Give yourself a brief round of applause if you caught 10 or more of these errors!

- The function detailed below is a recursive C function for calculating the result of a factorial operation. Note that this function is designed for integers greater than or equal to 0.

```
int factorial(int input)
{
    if(input <= 1)
    {
        return 1;
    }
    else
    {
        return input * factorial(input - 1);
    }
}
```

Write an assembly program to calculate the factorial value of 10. To do this, write the main function and also translate this into assembly language, using the calling conventions.

#Note that this function assumes the stack pointer has already been

#initiated appropriately in the main function.

#base is passed in through r4 and power through r5, by the register conventions

_start:

movia sp, stack #Initialize the stack. Remember to set up the corresponding part at the end

movi r4, 10

call factorial

#We now have the result of 10 factorial stored in r2!

br END #Prevent accidental execution of factorial!

factorial:

blt r5, r0, ONERETURN

subi sp, sp, 4 #Reserve stack space for one value

subi r4, r4, 1 #We can pass this without saving its original value since we don't use it after

stw ra, 0(sp)

call factorial

ldw ra, 0(sp)

mul r2, r2, r4

addi sp, sp, 4 #Deallocating stack space.

ret

ONERETURN:

movi r2, 1

ret

```
END:
.skip 500 #These two lines set up a 500-byte storage space for the stack.
stack:
.end
```

- What is meant by the term spatial locality?
Variables located near each other in memory (ex., values stored in the same row of an array)
- What is meant by the term temporal locality?
A variable that has been used recently is likely to be used again.
- What is meant by the term pipelining?
Pipelining refers to the process by which commands are staged; instead of completing a whole command before moving to the next one, the command is split into individual parts so that resources can be maximized. In this approach, the fetch component of one command might be in progress while the ALU addition component of another command is executing.
- Why is out of order processing used in assembly language?
OOP (out-of-order processing) is used to speed up execution of code; if the CPU can be working on more than one instruction at once, the average time for an instruction to be completed can be reduced.
- What is the inherent conflict between branches and pipelining?
Pipelining seeks to improve run time by working on several instructions at once. This process is hampered if the CPU doesn't know which branch will be executed, though there are several strategies to counter this.

NIOS II Instruction Reference

This is a list of the syntax for the instructions that you may use from the NIOS instruction set:

instruction	usage
ldw	ldw rB, off(rA)
ldb	ldb rB, off(rA)
ldbu	ldbu rB, off(rA)
ldbu	ldbu rB, off(rA)
ldh	ldh rB, off(rA)
ldhu	ldhu rB, off(rA)
stw	stw rB, off(rA)
stb	stw rB, off(rA)
sth	stw rB, off(rA)
add	add rC, rA, rB
sub	sub rC, rA, rB
mul	mul rC, rA, rB
div	div rC, rA, rB
addi	addi rC, rA, IMMED16
subi	subi rC, rA, IMMED16
muli	muli rC, rA, IMMED16
divu	divu rC, rA, rB
and	and rC, rA, rB
or	or rC, rA, rB
xor	xor rC, rA, rB
andi	andi rC, rA, IMMED16
ori	ori rC, rA, IMMED16
xori	xori rC, rA, IMMED16
andhi	andhi rC, rA, IMMED16
orhi	orhi rC, rA, IMMED16
xorhi	xorhi rC, rA, IMMED16
mov	mov rC, rA
movi	movi rB, IMMED16
movui	movui rB, IMMED16
srl	srl rC, rA, rB
srli	srli rC, rA, IMMED5
sra	sra rC, rA, rB
srai	srai rC, rA, IMMED5
sll	sll rC, rA, rB
slli	slli rC, rA, IMMED5
jmp	jmp rA
br	br LABEL
beq	beq rA, rB, LABEL
bne	bne rA, rB, LABEL
blt	blt rA, rB, LABEL
ble	ble rA, rB, LABEL
bgt	bgt rA, rB, LABEL
bge	bge rA, rB, LABEL
bltu	bltu rA, rB, LABEL
bleu	bleu rA, rB, LABEL
bgtu	bgtu rA, rB, LABEL
bgeu	bgeu rA, rB, LABEL
call	call LABEL
ret	ret