

Work out this homework on paper and then turn it in either in-person or through Canvas. If turn it in on Canvas, please make sure that all pages of your work are in one file.

1. Big-oh ranking

Order the following big-oh complexities in order from slowest-growing to fastest-growing (this is not the same as slowest to fastest running time!) It is possible some of them are actually in the same big-oh category. If that is the case, make it clear which ones have the same complexity.

n^2 , 3^n , \sqrt{n} , 1 , $n \cdot \log(n)$, 2^n , $n!$, $2^{\log(n)}$, n^3 , n , $n^2 \log(n)$, $\log(n)$, 2^{n+1}

2. Big-oh complexity

Assume each formula below represents the running time $T(n)$ of some algorithm. For each formula, give the lowest big-oh complexity possible (the tightest bound).

Here, \log represents the base-2 logarithm.

(a) $5 + n^2 + 25n$

(b) $50n + 10n^{1.5} + 5n \cdot \log(n)$

(c) $3n + 5n^{1.5} + 2n^{1.75}$

(d) $n^2 \log(n) + n \cdot \log(n) + n \cdot (\log(n))^2$

(e) $2^n + n^{10}$

(f) $n \cdot \log(n) + 8n + n \cdot (\log(n))^2$

3. Big-oh complexity proof

Assume we have analyzed an algorithm, and its run time is determined to be

$$T(n) = n^3 + 2n + 3$$

(a) What is the big-oh running time of this algorithm? (This should be easy.) Call this function $f(n)$.

(b) Now, prove your answer.

In other words, find constants $c > 0$ and $n_0 > 0$ such that for all numbers $n \geq n_0$, $T(n) \leq c \cdot f(n)$. You may prove that your c and n_0 work by drawing a graph showing $T(n)$ and $c \cdot f(n)$. Label your axes and where n_0 is.

4. Big-oh complexity analysis of code

Determine the big-oh running time for the following algorithms in terms of n . (No justification needed.)

a. Matrix addition:

```
for (int i = 0; i < n; i++)
{
    for (int j = 0; j < n; j++)
    {
        c[i][j] = a[i][j] + b[i][j];
    }
}
```

b. Matrix multiplication:

```
for (int i = 0; i < n; i++)
{
    for (int j = 0; j < n; j++)
    {
        c[i][j] = 0;
        for (int k = 0; k < n; k++)
        {
            c[i][j] += a[i][k] * b[k][j];
        }
    }
}
```

c.

```
counter = 0;
while (n >= 1)
{
    n = n/2;
    counter++;
}
```

d.

```
counter = 0;
x = 1;
for (int i = 0; i < n; i++)
{
    for (int j = 0; j < x; j++)
    {
        counter++;
    }
    x = x * 2;
}
```

5. Recall that in class we wrote an `RArrayList` class that works similarly to the built-in Java `ArrayList` class. Suppose we want to add a method to our `RArrayList` class called `duplicate()` that will copy all the items in the list a given number of times, and append them to the end of the list. The function will take a parameter, `howmany`, that specifies the number of times the items should be copied and appended.

Example of use:

```
RArrayList mylist = new RArrayList();
mylist.append(1);
mylist.append(2);
mylist.append(3);
mylist.duplicate(3);
// mylist is now [1 2 3 1 2 3 1 2 3]
```

In this problem, you will fill in the starter code below. You may assume that there is enough room in the `RArrayList` to fit "`howmany`" total copies of the existing data. The code that you add should do the copying and appending part. Do not call any other functions; you should write all of your code inside this function. In other words, don't call `append` or anything like that. You may assume that `howmany` is an integer > 1 .

```
public void duplicate(int howmany) {
    // Remember, this code lives inside RArrayList, so you have access to
    // the data[] array, data.length, and the size variable.

    // Assume that the data[] array has enough room for the necessary copies.

    // YOUR CODE HERE: Copy the existing data in the array the correct number of
    // times to the end of the data[] array. Don't forget to update the size
    // variable when you're done.
}
```

You do not need to recopy the entire function on your answer sheet; just provide the code that would go in the `YOUR CODE HERE` section.

6. Suppose we have a singly-linked list class with just a head pointer (no tail pointer) like this:

```
public class Node {
    public int data;
    public Node next;
};

public class SList {
    private Node head;
    // Assume there are more methods defined to add items to the list, etc.
}
```

Write a member function for SList called `getSize()` that calculates and returns the size (number of elements) in the linked list.

Here is the skeleton method:

```
public int size()
{
    // Remember, this method lives inside SList, so you have access to the
    // head variable defined above. You should assume if head is null, the
    // list is empty.

    // YOUR CODE HERE: Traverse the linked list and count the number of elements.
    // Return your answer.
}
```

You do not need to recopy the entire function on your answer sheet; just provide the code that would go in the YOUR CODE HERE section.

7. This question uses the same SList singly-linked list class from the previous problem.

Assume we add the following method to the SList class:

```
public void strange() {
    Node curr = head;
    while (curr != null) {
        curr.next = curr.next.next;
        curr = curr.next
    }
}
```

(a) Suppose we make an SList `mylist` and add the numbers 1, 2, 3, 4, 5, and 6 to `mylist` (so `mylist` consists of those six numbers in that order). What does `mylist` look like after calling `mylist.strange()`? (You can just write down the items in the list from left to right.)

(b) Suppose we make an SList `mylist2` and add the numbers 1, 2, 3, 4, and 5 to `mylist2` (so `mylist2` consists of those five numbers in that order). What happens when calling `mylist2.strange()`?

(c) In general, describe the difference in behavior when running this function on a list with an even number of items versus a list with an odd number of items.

8. Many linked list functions can be written **recursively** as well as iteratively. For instance, here's a recursive function to print out the items in a linked list (using the same SList class from above).

```
public void print(Node curr) {  
    if (curr != null) {  
        System.out.println(curr.data);    // (1)  
        print(curr.next);                // (2)  
    }  
}
```

The code above would be called as `print(head)` where `head` would be a reference to the first node in the linked list.

(a) What would the `print()` function above do if you switched the order of the lines marked (1) and (2) in the code?

(b) Write a recursive function, similar to the one above, that returns the sum of all the numbers in the linked list. Hint: use this skeleton code:

```
public int sum(Node curr) {  
    if (curr == null) {  
        return _____;  
    }  
    else {  
        // call sum recursively, add something to that value, and return it  
    }  
}
```

Your code, like `print()`, should not use any loops.

(c) Compute the big-oh time of the `print()` function (or the `sum` function; their big-oh times are the same). Use the strategy we learned in class to compute big-oh times of recursive functions, and *show all your work*. See the class website for the handout if you forget how to do this. "*n*" here should represent the number of nodes/items in the list.