# **Table of Contents**

Introduction	3
Cheat Sheet	3
Basic syntactic concepts	4
Constants and expressions	4
Assigning values	4
Weak typing	
Variables and stem variables	
Tail substitutions	8
About that trailing period	9
Applying scalars to stems	
Default values for stems	
The environment and the lifecycle of variables	11
Visibility during function evaluation	
Control structures	12
The ifthenelse statement	12
The switch statement	13
Error handling	13
Related functions	14
raise_error	14
Looping	
Related Functions	
break	15
check after	16
continue	16
for_keys	
for_next	
Defining functions	
Examples	
Related functions	
return	
Help	
Modules	
Local variables and functions vs. modules ones	
Related Functions	
import	
load module	
Built-in function reference.	
String functions.	
contains	
index_of	
insert	
to_lower, to_upper	
replace	
substring	
tokenize	
trim	

vdecode	33
vencode	
Math functions.	
abs	
date_ms, date_iso	
decode b64	
encode b64	
from_hex.	
hash	
mod	
random	
random string	
to_hex	
Stem specific functions	
box	
exclude_keys	
— v	
from_json	
list_keys	
has_keys	
include_keys	
indices	
is_list	
keys	
list_append	
list_copy	
list_insert_at	
list_subset	
mask	
rename_keys	
set_default	
size	
to_json	
to_list	56
unbox	57
union	57
Input and output	58
Dir	58
rm	60
print, say	60
read_file	61
scan	63
vfs_mount	64
write_file	65
Scripts	
execute	
load_script	
run_script	
General functions	
is_defined	

is_function	68
remove	68

# **QDL** Specification

#### Introduction

This document is the specification for the QDL (pronounced "quiddle") **Q**uick and **D**irty **L**anguage, which is specific to the OA4MP system and is used for all manner of server side-scripting. The aim of it is to have a reasonable language that is as minimal as possible and allows for a wide range of operations on claims and server-side management of clients. It is very lean. The meaning of the name is simple. Quick means that it should be quick and easy to code most mundane tasks. Dirty means that there are minimal constructs available. There is a single type of loop you must manage, for instance and exactly two types of variables. Yes, in this day and age, this is a largely procedural language because it is a scripting tool.

A bit of motivation. A very common situation is having to configure many rules for doing fairly simple tasks, e.g., if a user has logged in through an institution, belongs to any of a few groups, has an affiliation with a third institution, then some subset of information should be returned. In large blocks

huge rats nest of conditions

==> do something conceptually easy but very repetitive to a large set of data

So this language needs to have a lot of horsepower for decision making and very simple ways to work on large data sets, On top of this, since these scripts frequently run on server, speed is essential and implicit looping (discussed later with stem variables) makes this quite a snap. Also, the few times some other construct is needed (such as a switch statement or a loop), the needs are quite basic. So this set the parameters for the language. There are 4 basic constructs (while, if/then, try/catch, while). That's it. You can learn this language in about 20 minutes and use it quite effectively. Since it is interpreted, you can bash along merrily trying things out until you decide what works.

### **Cheat Sheet**

Scalar (aka primitive, simple) types: null, boolean, number, string

Variable types: simple or compound (aka stem variables)

Control structures: if..then...else, while loop, switch statement, try...catch.

Encapsulation: module

# Basic syntactic concepts.

- 1. Weakly typed
- 2. Simple and stem variables
- 3. A full set of logical operations
- 4. A full but minimal set of control structures

# **Constants and expressions**

There are boolean, number (integers are (64 bit), decimal are unlimited), string and stems. Valid identifiers a-z, upper and lower case, \$\_ and digits. Variables may not start with a digit. There is not limit to the length of a variable name. These are all fine:

```
a___$
$holyCow
__internal_function
```

Note that in some cases, the dollar sign may be used for escaping disallowed characters. See vencode and vdecode for the particulars.

QDL supports the following standard algebraic operations:

```
+ - * / % ^
```

Note that the operator % is integer division. So 42/9 = 4.666... and 42%9 = 4. It also supports the following logical operators

```
< <= > >= && || !
and the following increment and decrement operators, which may be used before or after variables.
-- ++
```

The limited character set for *symbols* is intentional since this sidesteps running it on systems that may have character encoding issues or more usually, working on a system where the supported terminal types are very limited. Generally however, the contents of a string may be UTF-8 with no issues.

### **Assigning values**

There is a specific *assignment operator*, denotes := which is used when setting a variable. This is a fine example:

```
c := 4;
```

As we will see later, you may also use this to assign a value to a stem variable by including the final period:

```
a. := indices(3);
```

You may chain these

```
a := b := c := 1;
```

will assign each variable a, b, and c to the value of 1.

There are also shorthand assignment operators for each basic operation of  $+ - * / \% \land$ . So if *op* is one of these operators then

```
A op = B
```

is identical to issuing

```
A := A op B
```

Example:

```
a := 3;
a ^= 2
a
```

Another example.

```
A := 'a'
B := 'b'
q := A += B += 'c'
q
abc
A
abc
B
bc
```

# Weak typing

As we have said, there are 3 primitive or *scalar* types: boolean, number and string. Booleans are either **true** or **false**, e.g.

```
a := true;
```

Numbers are of two sorts which are used seamlessly as needed. Integers are 64 bits and require no special handling. Decimals are more or less arbitrary precision. We say more or less because if you give the decimal, it will be exact, but in operations (well, division only) where the decimal can't be exact, it is kept to 50 decimals precision.

The first result is exact because we specified the number of digits. In the second case, there is no exact decimal representation of 1/3, so it is truncated.

Strings are single quote delimited and you may embed single quotes by escaping with a \'. So here is a string:

```
my_string := 'abcd\'efg';
say(my_string);
abcd'efg
```

Since there are many time external programs use double quotes and one of the aims of QDL is to make it interoperate nicely with this saves a lot of time dealing with niggling issues about where an extra double quote crept in.

```
You may also concatenate strings easily using the + operator, so say('abc'+ '123'); abc123

Similarly, the "-" works on strings too and removes the right elements from the left:

say('abcdeababghabijab' - 'ab'); cdeghij

Strings may be compared using == and !=, so
```

would be true.

Unlike many languages, there is no explicit type set forth for most languages and indeed, you may even change the type on the fly without penalty. For instance, this causes no error:

'foo' == 'foo'

```
my_var := 'Avast ye scurvy dogs!';
my_var := 42;
```

Where in many languages this would raise an exception. This is the "dirty" part of the name: the onus is on the programmer to keep this straight. Variables may contain the letters (upper or lower case), digits, underscore and dollar sign. Variables are case sensitive, so do be careful.

### Variables and stem variables.

A *Simple* variable, also called a *scalar* consists of primitive types, which are boolean, number (both integer and decimal) and strings. These look just like any other variable from most programming languages (the ":=" is the assignment operator). So for instance

```
a := 'foo';
my_boolean := true;
my_integer := 123;
my_decimal := 432.3454;
```

are all valid simple variables.

A *Compound* variable is embodied in what is termed *stem variables*. These are of the form head.tail

(Geeky stuff, the period is actually called the *aggregation operator*.) Remember that the variable is **head.** (note the trailing period!!) and the tail – which may be complicated -- is just indexing. This effectively means that a stem can be something as simple as a list or quite a complex data structure indeed. As a matter of fact, any data structure an be modelled with a stem. The index is any string and we usually refer to them as keys. Some definitions:

- A *list* is a stem which contains indices 0,1,...n
- Two stems are *conformable* if they have the same keys
- **Subsetting** is in effect for most stem operations. This means that if the result is a stem, it contains only the keys common to its arguments. If two stems are conformable, there is no subsetting.
- *Tail resolution* means that if a stem has many indices, like a.b.c.d, then it is resolved from right to left, with the system checking each index to see if that variable has been defined, then substituting.

Note that this is one of the very few strict pattern enforced in QDL – a stem must end in a period, so issuing something like this (to make a list of integers)

```
a := indices(5);
```

fails with a message like "Error: You cannot set a scalar variable to a stem value". This is because the item on the right is an aggregate, aka a stem and on the one on the left is a simple scalar.

There are two common use patterns. The first pattern is to invoke a function on a stem variable which alters every element and returns a stem variable with the same keys, each element having been transformed. Here is an example. This uses integer indices and this makes it a list. To populate it you would just set the elements:

```
myList.0 := 'the';
myList.1 := 'quick';
myList.2 := 'brown';
```

This effectively is an array (which is just a map whose keys are integers). The second use case comes from having indices that are not indices, which allows you to make maps on the fly, with the index being the key:

```
myMap.idp := 'http://idp.bigstate.edu/saml';
myMap.port := 636;
myMap.eppn := claims.sub + '@bigstate.edu';
```

In this case, there is now a map with keys, idp, port and eppn whose values are as above. Stem variables have their own section later with more details. There are several operations supported on them.

It is certainly possible to have mixed data, for instance

```
my_stem.help := 'this is my stem'
list_append(my_stem., indices(5))
{0=0, help=this is my stem, 1=1, 2=2, 3=3, 4=4}
```

which shows that this stem includes a list and has another entry called *help*.

#### Tail substitutions.

If you have defined a variable, say

```
k := 3;
my_var. := indices(5);
and issue the following
my_var.k := 'foo';
```

Then this will result in my\_var.3 being equal to the string 'foo'. In short if the tail has a value at that point, this is used first. If not, then the tail itself as a string is used. This lets you do things like

```
i := 0;
while[
    i < 5
]do[
    say('the value = ' + my_var.i);
    i++;
];
which prints out

the value = 0
the value = 1
the value = 2
the value = foo
the value = 4</pre>
```

Moreover, substitutions happen from right to left (!! backwards from reading order), so if you set

```
x := 0;
y.0 := 1;
z.1 := 2;
w.2 := 3;
```

Then you could reference a value like

```
W.Z.Y.X
```

which would resolve to

```
w.2 == w.z.y.x == 3
```

This permits more readable values, e.g. **time.manner.place** is a perfectly fine reference. **HOWEVER** the stem is always the leftmost symbol. The rest are just a very compact way to index it.

So why do this? Because it allows for something very powerful: implicit looping. You can have stems do a tremendous amount of work without ever really having to access the elements.

Example.

```
a. := abs(mod(random(1000000),1000000));
a.42
769942
```

That's 1 million random numbers in the range of 0 to 1000000 and we showed the 42<sup>nd</sup> value just because. Here's a polynomial:

```
a. := a.^2 + 3*a. -4;
a.42
592812993186
and if you insist, here is a check
769942^2 + 3*769942 - 4
592812993186
```

Stems can indeed to a tremendous amount of implicit looping.

### About that trailing period.

To refer to a stem as an aggregate (everything) you must include the period. This is a perfectly fine example

```
a.0 := 'foo';
a.1 := 'bar';
a.2 := 'baz';

a := 2; // so this is a scalar - no period at the end
say(a.a);
```

This just points out that a and a. are considered to be wholly unrelated. Here the scalar a has value of 2 which is substituted as the tail of the stem, so the answer printed is the value of a.2.

You may nest stems as well, so

```
a. := 3 * indices(5); // count by 3's, so 0,3,6, ...,12
b. := 10 * indices(5);// count by 10's so 0,10, ...,40
a.b. := b.;
```

works just fine. You can access the value as

```
c. := a.b.; // note the trailing . on the variable c to show it is stem
```

but you cannot issue a.b.3 and expect to get anything back (b.3 == 30 which is not an index in a.) unless you have set it explicitly. This is because, again, it is included in the stem variable a. as an entry and you must refer to it by its proper index. TL;DR version is that tail resolution happens for scalars.

### Applying scalars to stems

A scalar is a simple value. All of the basic operations in QDL work on stems as aggregates. So for instance, if you needed to make a list counting by 3's, you could issue

```
a. := 3*indices(5);
    say(a.)
{0=0, 1=3, 2=6, 3=9, 4=12}

(The output is of the form {key=value, key=value,...}). So lets say we have the following:
    ring.find := 'One Ring to find them';
    ring.rule := 'One Ring to rule them all';
    ring.bring := 'One Ring to bring them all';
    ring.bind := 'and in the darkness bind them';

Let's find every element that contains the word 'One", respecting case:
    say(contains(ring., 'One'));
{bind=false, find=true, bring=true, rule=true}
Or for that matter, doesn't contain this word:
```

The output of these functions is a boolean-valued stem and there is a very useful function called *mask* which simply will return the elements that have corresponding true values.

```
say(mask(ring., contains(ring., 'One')));
{find=One Ring to find them,
bring=One Ring to bring them all,
rule=One Ring to rule them all}
```

{bind=true, find=false, bring=false, rule=false}

say(!contains(ring., 'One'));

And of course you could just find the ones that don't have the word "One" in them:

```
say(mask(ring., !contains(ring., 'One')));
{bind=and in the darkness bind them}
```

This points out that using stems can do a tremendous amount of work for you. Since QDL is interpreted (each line is read, then parsed and executed) having as much happen as possible with a command improves both performance and efficiency. Besides, working with aggregates is often much ore intuitive than slogging through each element.

#### Default values for stems

You may set a default value for stems – this is a very nice thing indeed so you don't have to initialize every element in one before using it. The way it works is you issue

```
set_default(stem., scalar);
```

where the *scalar* is any scalar. Then from that point forward, any time a value is accessed, if it has not been explicitly set, the default is returned. If the stem does not exist, it will be created. Note that subsequent operations on the stem do *not* alter the default value.

```
E.g.
    set_default(stem., 2);
    say(stem.woof);
2
    say(has_key(stem., 'woof'));
false
```

# The environment and the lifecycle of variables.

Every script has both global and local variables associated with it. Variables defined in a block (such as in the body of a loop or in the body of a conditional statement) are local to that block. Variables defined outside that block are global to the program. Modules are completely self-contained.

So what if you need to have variable defined and accessible outside a block but need to set the value inside one, like

```
if[
  // nasty conditions
]then[
  // lots of stuff
  a := 'foo';
]else[
  // lost of stuff
  a := 'bar';
];
// trying to use a will result in errors
```

What you can do is set the value to null outside the block. So do this.

```
a := null;
// same code as above
```

Now attempts to use the variable will work properly. You may also check if the value has been set by checking if it is null:

```
if[
   a != null
]then[
  // lost of stuff
];
```

# Visibility during function evaluation

Generally functions inherit the values of their parent state, but they do not inherit imported symbols. Following Leibnitz' lead, a function is to relate *cause* (the inputs) to *effect* (the output). Therefore, functions should have their values passed to them and not draw them in willy-nilly. You may, of course, just import modules in the body of the function or pass the values in as arguments.

### **Control structures**

There are 4 basic control structures. All of them are delimited with brackets [].

- 1. The basic conditional of if[ *condition* ]then[ ... ]else[ ... ]; Note that the else clause is optional.
- 2. Switch statements (which are lists of conditionals) of switch[...]
- 3. Try ... catch for error handling
- 4. Looping construct while [condition ]do[...];

#### The if..then..else statement

```
The basic format is
```

```
if[
   condition
]then[
   //statememts
]else[
   // more statements
];
```

Note that **if**[, **]then**[ and **]else**[ are treated as statements, so cannot be split. The else clause is optional. And of course, a simple example is in order:

```
)buffer on
Buffer mode on.
j :=5;
if[
    j <5 || 5 < j
]then[
    say(j + ' is not 5');
]else[
    say('j is ' + j);
];
.
j is 5
)buffer off
Buffer mode off.</pre>
```

(Rather than stick this in a script, we used the interpreters )buffer command to keep all of the lines and then execute it when there is a single period as the first character on a line. Oh and interpreter commands which start with a ')' still work in buffer mode. Consults the work space documentation for more.)

#### The switch statement

Branched decision-making is a basic construction for most languages. In the case of QDL, there is the switch[]; construct. The basic format is

```
switch[
  if[condition1]then[body1];
  if[condition2]then[body2];
  if[condition3]then[body3];
//... arbitrarily many
];
```

The execution is each condition is checked and as soon as one returns *true* that body is executed and the construct returns.

#### **Examples**

An example of a switch statement might be

```
i := 11;
    switch[
        if[i<5]then[var.foo := 'bar';];
        if[5=i]then[var.foo := 'fnord';];
        if[5<i]then[var.foo := 'blarf';];
    ];
    say(var.foo);
blarf</pre>
```

Note that the elements of the switch statement are if..then blocks (including final semicolon). No else clauses will be accepted. Whitespace aside of strings, of course, is ignored.

# **Error handling**

There is a try ... catch block construction. Its format is

```
try[
  // statements;
]catch[
  // statements;
];
```

You enclose statements in a try block and call raise\_error if there is an exceptional case. Note that unlike many languages, all exception are fail-fast, so there is no way to hop back at the point of the error and resume processing. An example

```
j := 42;
try[
    remainder := mod(j, 5);
    if[remainder == 0]then[say('A remainder of 0 is fine.');];
```

```
if[remainder == 4]then[say('A remainder of 4 is fine.');];
  if[remainder == 1]then[raise_error(j + ' not divisible by 5, R==1', 1);];
  if[remainder == 2]then[raise_error(j + ' not divisible by 5, R==2', 2);];
  if[remainder == 3]then[raise_error(j + ' not divisible by 5, R==3', 3);];
]catch[
  if[error_code == 1]then[say(error_message);];
  if[error_code == 2]then[say(error_message);];
  if[error_code == 3]then[say(error_message);];
]; // end catch block
42 not divisible by 5, R==2
```

So the way these work is if a certain condition is met, you call the raise\_error function with a message and optional numeric value. These are available in the catch block so you can figure out which error was raised and deal with it. Not much else to them.

#### **Related functions**

#### raise error

#### Description

Raises an error conditions and passes control to the catch block. Note that raising this outside of a try ... catch block will have no effect, since there is nothing to catch it.

#### Usage

```
raise_error(message, [code]);
```

### **Arguments**

message a human readable message that describes this error

code a user-defined integer that tells what this code is.

### Output

None directly. The values are set to error\_message and error\_code (if present) and are accessible in the catch block. Typically, you define the error code and look for it in the catch block.

# Examples

```
try[
    if[ 3 == 4]then[raise_error('oops', 2);];
]catch[
    say(error_message);
    switch[
        if[error_code == 2]then[...];
        // maybe a bunch of other errors
    ]; // end switch
]
oops
```

# Looping.

The basic structure of a loop is

```
while[
   logical condition
]do[
   statements
];
```

For example, to print out the numbers from 0 to 5:

```
i := 0;
while[
    i < 5
    ]do[
    say(i++);
];
0
1
2
3</pre>
```

This is known as 'yer basic loop'.

There are 3 functions that can help loop.

#### **Related Functions**

#### break

### Description

Interrupt loop processing by exiting the loop.

### Usage

break();

# Arguments

None.

### Output

None.

### Examples

In this example, the loop terminates if the variable equals 3.

#### check\_after

#### Description

Sometimes only a post-positional loop will do – this means that the loop executes at least once. This is not often the case, but is very hard to replicate. Invoking this function will do just that. Your condition will be checked post-loop.

#### Usage

```
check_after(condition);
```

### **Arguments**

The argument is a logically -valued expression.

### Output

None. This exists only in looping statements

```
a := 0;
while[
  check_after(a != 0)
  ]do[
  say(a);
];
0
```

#### continue

### Description

During loop execution, skip to the next iteration.

#### Usage

```
continue();
```

#### **Arguments**

None.

#### Output

None.

#### Examples

In this example, the loop skips to the next iteration is the variable is 3.

### for\_keys

### Description

This is a non-deterministic loop over the keys in a stem variable. All the keys will be visited, but there is no guarantee of the order.

### Usage

```
for_keys(var, stem.);
```

### **Arguments**

var is a simple variable and will contain the current key during the loop. If it has already been defined, it's values will be over-written.

stem is a stem variable. The keys of this stem will be assigned to the *var* and may be accessed.

### Output

Nothing. This is only used in looping constructions.

#### Example

```
my.foo := 'bar';
my.a := 32;
my.b := 'hi';
my.c := -0.432;
while[for_keys(j,my.)]do[say('key=' + j + ', value=' + my.j);];
key=a, value=32
key=b, value=hi
key=c, value=-0.432
key=foo, value=bar
```

#### for next

#### Description

This allows for a deterministic loop and will run through a set of integers.

#### Usage

```
for_next(var, stop_value, [start_value, increment]);
```

#### **Arguments**

Only the first two are required.

*Var* the variable to be used. As the loop is executed, this value will change.

*stop\_value* the final value for the loop. When the variable acquires this value, the loop is terminated (so the loop body does **not** execute with this value!

*start\_value* (optional, default is 0). The first value assigned to *var*.

*increment* (optional, default is 1). How much the loop variable should be incremented on each iteration.

### Output

None. This only is used in loops.

### Examples

```
A simple loop
while[
  for_next(j,5)
]do[
  say(j);
];
0
1
2
```

```
3
4
while[
  for_next(k, 5, 0, -1)
  ]do[
    say(k);
];
5
4
3
2
1
```

And here is an example of looping through the elements of a stem variable.

```
// Set the values initially
my_stem.0 := 'mairzy';
my_stem.1 := 'doats';
my_stem.2 := 'and';
my_stem.3 := 'dozey';
my_stem.4 := 'doats';
while[
  for_next(j,5)
 ]do[
  say(my_stem.j);
]
mairzy
doats
and
dozev
doats
```

# **Defining functions**

The syntax for defining a function is very simple:

```
define[
    signature
]body[
    >> useful comments
    statements
];
```

Note that within the body, any variables defined are local to that block unless they are save, *e.g.* in the environment. Previously defined values are still available. The variables in the signature are populated with the values (which are copied) . To return a value, invoke the method return — which is only allowed inside function definitions, by the way. No return statement implies there is no output from the function.

The *signature* of a function is

```
name(arg0, arg1, arg2,...)
```

this means that the function is defined by its name. Of course, if you define it in a module, you may have to qualify it.

An example

```
define[
    sum(a, b)
  ]body[
    >> add a pair of integers and return the sum.
    c:= a+b;
    say('the sum is ' + c);
    return(c); // this terminates execution and returns the value.
  ];
say('the sum of 3 and 4 is ' + sum(3,4));
This prints
the sum of 3 and 4 is 7
```

You may use functions any place in your code once they have been defined, so it is a good practice to put them at the beginning.

#### **Examples**

Here is a QDL program to find the *Armstrong numbers* in the range of 100 - 1000. A number, xyz is an Armstrong number iff  $xyz = x^3 + y^3 + z^3$ .

```
define[
  armstrong(m)
  ]\bod\[
     >> An Armstrong number is a 3 digit number that is equal to the sum of its cubed
digits.
     \rightarrow This computes them for 100 < n < 1000.
     >> So for example 407 is an Armstrong number since 407 = 4^3 + 0^3 + 7^3
     if[ m < 100]then[say('sorry, m must be 100 or larger'); return();];</pre>
     if[1000 < m]then[say('sorry, m must be less than 1000'); return();];</pre>
     sum := 0;
     while[
          for_next(j, m)
        ]do[
           n := j;
           while[
                0 < n
           ]do[
                b := mod(n, 10);
                sum := sum + b^3;
                n := n%10; // integer division means n goes to zero
           ]; //end inner while
           if[sum == j]then[say(sum);];
           sum := 0;
        ]; // end while
  ]; // end define
```

#### **Related functions**

#### return

#### Description

Return a value or none.

#### Usage

```
return([value]);
```

#### **Arguments**

One. The value to be returned.

#### Output

The value to be returned.

#### Examples

Here is a cheerful little program that ignore everything and just returns "hello world".

```
define[
   hello_world(x)
]body[
  return('hello world!');
];
  say(hello_world(42));
hello world!
```

# Help

Functions have a very specific documentation feature. At the very top of the body, you may enter documentation, each line is of the form

```
>> text
```

And these will be taken when the function definition is read and kept available for consultation. You may get a full listing of every user-defined (meaning, not built in) function the workspace knows about. Each is printed as

```
)help list fib2(1): This will compute the n-th element of a Fibonacci sequence. f(1): This is a comment
```

where there is the name, (number of arguments) and the first line of the comment (which should be meaningful and explain to the user what the function does.) The first is a good comment. The second is not.

Note that this is only for functions. A common convention is workspaces is to have a stem variable named help. each of whose keys contains a description of something useful.

```
)help fib2 1
fib2(1):
This will compute the n-th element of a Fibonacci sequence.
A Fibonacci sequence, a_, a_1, a_2, ... is defined as
    a_n = a_n-1 + a_n-2
Acceptable inputs are any positive integer.
This function returns a stem list whose elements are the sequence.
```

### **Modules**

One of the most basic ideas in programming is *encapsulation* that is to say, a group of statements with their own state (so the variables and functions know about each other and their workings are independent of the rest of the environment). There is, of course, the concept of a *module* in QDL that does this. Very simply, a module is defined using a unique identifier that **must** be a URN and an alias. The alias is just a regular name like a variable. One of the great evils of many scripting languages is that there is no encapsulation – every variable is global and the net effect is extremely hard to find bugs.

Let us say that we had the following script in the file my-module.qdl:

```
module[
    'a:a','a'
]body[
    foo := 'abar';
    define[f(n)]body[return(n+1);];
]; // end

The syntax is clear here:

module[
    uri, alias
    ]body[
    // any statements except another module definition
];
```

The URN puts this in a unique namespace. There can be no conflicts. The alias is used as a shorthand to access it. URNs many be quite complex. Generally though you should not put in query parameters and such. To access anything in a module explicitly you must *qualify* the name *i.e.*,

alias#name

It is easiest to see this in action, so in the interpreter (clear state)

```
)funcs
)vars
```

Note that there is nothing in this session so these show nothing. Let's import our module above in a file called my-module.qdl. We are going to import it with another alias just to show we canonical

```
load('my-module.qdl');
import('a:a', 'b');
     )vars
b#foo
     )funcs
b#f(1);
```

Loading the module makes the session aware of it, but you must also import it. If we did not specify the second argument, then the default alias in the module definition, 'a', would be used. Now at this point, there is nothing else in the session and there are no name clashes possible, so we can use the unqualified name:

```
say(foo);
abar

And we can invoke the function too
    say(f(5));
6
```

### Local variables and functions vs. modules ones

If we had something like

```
foo := 5;
    load('my-module.qdl');
    import('a:a', 'b');
    )vars
b#foo, foo
```

This shows us that there is a session variable with the same name. If we try to access the unqualified name we get

```
say(foo);
Error: The variable "foo" exists in multiple modules. You must qualify which one
you wish to use.
```

So this means that there is no way for the interpreter to resolve this. The solution is that you have to qualify the session variable with a simple #:

```
say(#foo);
6
```

There is no namespace in effect for a basic session. This is the most straightforward way to handle name clashes, *viz.*, if there is any conflict, require the user resolve it rather than having the interpreter doing it and risking getting it wrong. Part of the philosophy with QDL is that is avoids a lot of the sorts of things that cause hard to find errors and tells you up front how to fix it.

#### **Related Functions**

#### import

#### Description

Import a module.

#### Usage

```
import(urn[, alias]);
```

#### **Arguments**

urn = a Uniform Resource Name (as per RFC 2141 and 1737, if you track such things).

Alias = a string that conforms to the requirements of a QDL variable name.

Note that both of these are passed in as strings and the URN is checked for form. Also, when a module is defined, it has an alias given. This is the default, so you may import the module with only the URN. If you specify another alias, that is used.

### Output

Nothing.

# Examples

```
import('my:/thingie', 'ahab');
```

Would locate the module with URN *my:/thingie* and let you use *ahab* as the local alias. Remember that the module must be loaded first or you will get an error.

### load\_module

### Description

Load a module.

### Usage

```
load_module(file);
```

# **Arguments**

file – the full path to the file that contains the module. This is loaded then run in its own environment which is then added to your session.

# Output

Nothing

# Examples

load\_module('/home/bob/qdl/modules/mega\_module.qdl');

Would load the given module.

### **Built-in function reference**

There are several built in functions in various categories. All of these can take simple or compound variables.

# **String functions**

#### contains

#### Description

To find if a string contains another string

#### Usage

```
contains(source, snippets [, case_sensitive])
```

#### **Arguments**

source – a string or stem of strings that is the target of the search.

snippets – a string or stem of strings that are what are being search for

case\_sensitive — (optional) if this is **true** (default) then the check is done respecting case. If it is **false** then the matching is done after converting the arguments to lower case. (Note that the original values are never altered.)

### Output

A scalar or stem (if the arguments were stems) if the snippet(s) was (were) found. Note that

- source a scalar, snippet a scalar → result is a simple boolean
- source a scalar, snippet a stem → result is a stem with identical keys to the snippet and with boolean entries
- source a stem, snippet a stem, → result is a stem with identical keys to the source and boolean entries.
- Both stems, the result is conformable to the left argument and the right. In other words, to be in the result, only entries with matching keys are tested.

### Examples

#### Example 1.

```
a := 'What light through you window breaks?';
contains(a , 'Juliette');
```

returns false, since there is no string 'Juliette' in the first string.

#### Example 2.

```
source := 'the rain in Spain';
    snippet.article := 'the';
    snippet.1 := 'in';
    snippet.2 := 'Portugal';
    output. := contains(source, snippet.);
    output.article == true;
    output.1 == true;
    output.2 == false;
Example 3.
    source.foo := 'bar';
    source.fnord := 'baz';
    source.woof := 'arf';
    snippet.foo := 'ar';
    snippet.fnord := 'y';
    output. := contains(source., snippet.);
    output.foo == true;
    output fnord == false;
```

In this case, only the corresponding keys are checked if **both** arguments are stem variables.

#### index\_of

### Description

This finds the position of the target in the source. If the target is not in the source, then the result is -1.

### Usage

```
index_of(source, snippet [,caseSensitive])
```

### **Arguments**

source – a scalar or stem variable.

snippet – a scalar or stem variable.

case\_sensitive – (Optional) a boolean that if **true** will check for case and if false will check against the arguments as all lower case.

### Output

if both are strings, the result is the first index of where the snippet starts in the source. If one is a stem and the other a scalar, the result is conformable to the stem and the operation is applied to each element of the stem. If both are stems, then only corresponding keys are checked.

#### **Examples**

```
sourceStem.rule := 'One Ring to rule them all';
sourceStem.find := 'One Ring to find them';
sourceStem.bring := 'One Ring to bring them all;
sourceStem.bind := 'and in the darkness bind them';

targetStem.all := 'all';
targetStem.One := 'One';
targetStem.bind := 'darkness';
targetStem.7 := 'seven';
result. = index_of(sourceStem., targetStem.);
```

returns the following stem variable, which has one entry since two stem variable are only checked against corresponding keys:

```
result.bind == 11;
```

#### insert

#### Description

Insert a given string at a given position of another string

#### Usage

```
insert(source, snippet, index)
```

### Arguments

```
{\tt source-the\ string\ to\ be\ updated}
```

snippet - the string to insert

index – the position in the source string to insert the snippet

### Output

The updated string.

### Examples

```
say(insert('abcd', 'foo', 2));
abfoocd
```

This also works for stem variables

#### to\_lower, to\_upper

#### Description

This will convert the case of a string to all upper or lower case respectively.

#### Usage

```
to_lower(arg), to_upper(arg)
```

#### **Arguments**

arg – either a string or a stem variable of strings. Non-strings are ignored.

#### Output

A conformable argument of strings.

#### Examples

```
a := 'mairzy doats';
b := to_upper(a);
say(b);
MAIRZY DOATS
```

### replace

### Description

Replace every occurrence of a string by another

### Usage

```
replace(source, old, new)
```

### **Arguments**

```
source – the original string or stem of strings
```

```
old – the current string
```

```
new – the new string.
```

There is an statement about conformability. In this case if 2 or three of the arguments are stems, then only matching keys get replaced – the same key must be in all arguments or this is skipped. If exactly one of the arguments is a stem, then the replacement is made one each element with same arguments – in effect they are turned in to stem variables with constant entries. If all three are scalars it is just a standard replacement.

#### Output

The updated string

#### Examples

And example with two stem variables and a simple string.

```
sourceStem.rule := 'One Ring to rule them all';
sourceStem.find := 'One Ring to find them';
sourceStem.bring := 'One Ring to bring them all;
sourceStem.bind := 'and in the darkness bind them';

old.all := 'all';
old.one := 'one';
old.bind := 'darkness';
old.7 := 'seven';
newValue := 'two';
output. := replace(sourceStem., old., newValue);
```

The resulting output is a stem (because an input is) and in this case it is

```
output.bind == 'and in the two bind them';
```

Why is this? Because the only key that the two stems have in common is 'bind' and that is applied to replace 'darkness' with the new value of 'two'.

#### substring

### Description

Return the substring of an argument beginning at the *n*th position.

### Usage

```
substring(arg, n [,length] [,padding])
```

### **Arguments**

```
arg - the string or stem of strings to be acted up n - the start position in each string
```

length (optional) – the number of characters to return. Note that if this is omitted, the rest of the string is returned. If it is longer than the length of the string, only the rest of the string is returned *unless* the *pad* argument is given.

padding – a string that is used cyclically as the source for padding.

### Output

The substring. Notice that this behaves somewhat differently than in some other languages in that it may be used to make results longer than the original argument.

#### **Examples**

A basic example. Remember that the first index of a string is 0, so n = 2 means the substring starts on the *third* character.

```
a := 'abcd';
  say(substring(a,2));
cd
```

To use the padding feature

```
say(substring(a,3,10,'.'));
d.....
```

And do note that the *padding* need not just be a character, but will be repeated as needed:

```
say(substring(a,1,20,'<>'));
bcd<><><><><
```

Finally, a stem example. Note that the padding option makes all results the same length:

```
b.0 := 'once upon';
b.1 := 'a midnight';
b.2 := 'dreary';
d. := substring(b., 0, 15,'.');
while[for_next(j,3)]do[say(d.j);];
once upon.....
a midnight.....
dreary.......
```

#### tokenize

### Description

This will take a string and and delimiter then split the string using the delimiter.

### Usage

```
tokenize(arg, delimiter)
```

### **Arguments**

arg – either a string or stem of strings

### Output

If *arg* is a string, then a list of tokens. If *arg* is a stem, then a stem of stems. Remember that the keys are preserved if the argument is a stem and a simple list (keys are 0, 1,...) if a string.

### Examples

A simple example

```
say(tokenize('ab,de,ef,',','));
{0=ab, 1=de, 2=ef}
```

An example tokenizing a stem variable.

```
q.foo := 'asd fgh';
q.bar := 'qwe rty';
say(tokenize(q., ' '));
{bar={0=qwe, 1=rty}, foo={0=asd, 1=fgh}}
```

Tokenizer only works on strings. Here is the result of attempting to tokenize an integer.

```
say(tokenize(12345, '1'));
12345
```

In this case, the argument is returned unchanged.

#### trim

#### Description

Trim trailing space from both ends of a string. One point to note is that since stem variables can contain stem variable, this only operates at the top-level and if you wish to trim included stems you must do so directly. This prevents "predictable but unwanted behavior."

#### Usage

```
trim(arg)
```

### Arguments

arg is either a string or a stem of strings. This function has no effect on non-strings.

### Output

A result conformable to its argument.

### Examples

An example

```
a := ' blanks ';
'blanks' == trim(a);
```

Another example, using a mixed stem variable.

```
my_stem.0 := ' 'foo';
my_stem.1 := -42;
my_stem. := trim(my_stem);
my_stem.0 == 'foo';
```

```
my_stem.1 == -42; // unchanged.
```

#### vdecode

#### Description

Decode and encoded variable name. All escaped characters are unescaped. **Note** this does support UTF-8 encodings of all characters implicitly.

#### Usage

vdecode(string)

#### **Arguments**

String – the String to be decoded. Note that not every string can be decoded! Since escaped characters are of the form \$xy where x and y are hexadecimal numbers, it is possible to get illegal

### Output

The decoded string.

#### Examples

```
vdecode('$26$2A$28$26$25$23')
&*(&%#
```

An example of a name that cannot be decoded. This is, of course, because escaped values are of the form \$xy where x and y are hexadecimal numbers:

```
vdecode('$$foo')
Error: Could not decode string:Invalid escape sequence
```

Here is an example showing the UTF-8 implicit encodings of non-ASCII characters

```
vencode('你浣')
$E4$BD$A0$E6$B5$A3
vdecode('$E4$BD$A0$E6$B5$A3')
你浣
'你浣' == vdecode('$E4$BD$A0$E6$B5$A3')
true
```

(This example was chosen because these two characters can get swapped if the encodings are not carefully handled, so this shows that a well-known edge case is handled correctly. And no, I have no idea what these characters represent.)

#### vencode

#### Description

Encode a string, escaping all characters that are not legal for a variable name. This is done by using the standard URL escape specification (where every escaped character is replaced with a % and s 2 digit hexadecimal code), but using a \$ in place of the %. We have to do this in variable names since % represents integer division in QDL. So for instance, the blank ' ' escapes to %20 which in turn we would represent as \$20.

The reason for this is interoperability, mostly with JSON. If we are going to export variables (such as stems) the stem names are, in point of fact, legal variable names and will be resolved against the symbol table. In the case of JSON, these correspond to the property names, which are simple strings and may contain illegal characters. So if we have a JSON object

```
{"p:/*/q#","woof"}
```

This cannot be turned in to a stem without some change to the property name of p:/\*/q#. This function will escape everything and return p\$3A\$2F\$2A\$2Fq\$23 which can be turned back in to the original using the vdecode function.

#### Usage

vencode(arg)

### **Arguments**

arg is the string for which all illegal variable characters will be escaped.

### Output

A that is a legal variable name in QDL. Note that every string input can be encoded this way.

### Examples

vencode('&\*(&%#') \$26\$2A\$28\$26\$25\$23

### **Math functions**

#### abs

### Description

Find the absolute value of a number

#### Usage

```
abs(arg)
```

#### **Arguments**

arg – a number or a stem filled with numbers.

#### Output

If a single number, the absolute value of that number. If a stem of numbers, the absolute value of all of them.

#### Examples

```
say(abs(-123));
123
```

# date\_ms, date\_iso

### Description

Compute and convert dates between milliseconds and the ISO 8601 standard format.

#### Usage

```
date_ms([arg])
date_iso([arg])
```

### **Arguments**

either none, an argument or stem of arguments.

None:

date\_ms() returns the current time in milliseconds

date\_iso() - returns the current time in ISO 8601 format.

A single argument

```
date_ms(arg) -- if arg is in ms, return it, otherwise convert it to ISO format
```

date\_iso(arg) - if *arg* is ISO format, convert it to ms. Otherwise return it.

### Output

The date in the appropriate format.

### Examples

```
say(date_iso());
```

```
2020-01-18T22:10:38.250Z

say(date_ms('2020-01-18T22:10:38.250Z'));

1579385438250

say(date_ms(1579385438250));

1579385438250
```

#### decode\_b64

#### Description

Decode an encoded string. The result will be a simple string, so if the original is binary, you will see gibberish.

#### Usage

decode\_b64(arg)

#### **Arguments**

arg – may be a string or a stem of strings. Each string will be decoded.

#### Output

The decoded string

### Examples

 $say(decode\_b64('VGhlIHF1aWNrIGJyb3duIGZveCBqdW1wcyBvdmVyIHRoZSBsYXp5IGRvZw'); The quick brown fox jumps over the lazy dog$ 

### encode b64

### Description

Encode a string in base 64. The returned string is URL safe.

### Usage

encode\_b64(arg)

### **Arguments**

arg − a string or it may be a stem of strings.

### Output

If a single argument, it will be base 64 encoded. If a stem, each element will be. Non-strings are not changed.

#### **Examples**

say(encode\_b64('The quick brown fox jumps over the lazy dog'); VGhlIHF1aWNrIGJyb3duIGZveCBqdW1wcyBvdmVyIHRoZSBsYXp5IGRvZw

#### from hex

### Description

Take a hexadecimal representation of a string and convert it back

#### Usage

from\_hex(arg)

#### **Arguments**

arg -- a hexadecimal string or a stem of them.

#### Output

The string that corresponds to this argument or a stem of them (if the argument was a stem).

#### **Examples**

say(from\_hex('54686520717569636b2062726f776e20666f78206a756d7073206f76657220746865206c617a7920646f67')); The quick brown fox jumps over the lazy dog

#### hash

### Description

Calculates the SHA-1 digest of the arguments and returns the value as a hex string. This means that the result is a fixed 20 byte result (and the resulting output is a hexadecimal string exactly 40 characters long, regardless of the size of the input string). This is used in various cryptographic applications.

## Usage

hash(arg)

## Arguments

arg – either a single string or a stem of strings.

### Output

A hex string that is the hash. Note that while this is a hex string, it is most emphatically not the same as the output from the to\_hex function.

#### **Examples**

```
say(hash('the quick brown fox jumps over the lazy dog'));
2fd4e1c67a2d28fced849ee1bb76e7391b93eb12
say(hash('the quick brown fox jumps over the lazy do');
6186ce3119913cfabfe4b7952ba765b132948dd2
```

A point to make about SHA-1 hashes is that even a tiny change in the input completely changes the output in very hard to guess ways. This is what makes them so useful in *e.g.*, security.

#### mod

#### Description

Compute the modulus, *i.e.*, the remainder after long division, of two integer. Since there are currently only integers as allowed numbers, this is needed in cases where the remainder is required.

### Usage

```
mod(a,b)
```

#### **Arguments**

a and b may be either scalars or stems.

### Output

## Examples

To compute the simple remainder

```
say(mod(27,4));
3
```

And sure enough 27/4 = 6 with a remainder of **3**.

A stem example. Computer the remainder of a bunch of values

```
a.0 := 11;
a.1 := 20;
say(mod(a.,4));
{0=3, 1=0}
```

In this case, since 20 is evenly divisible by 4, the modulus (aka remainder) is 0.

#### random

### Description

Generate either a single signed 64 bit random number (no argument) or a list of them.

#### Usage

```
random([n])
```

#### **Arguments**

number (optional) -- the number of random values you want to generate.

### Output

If there is no argument, a single random 64 bit number. If there is a number, n, supplied. Then a stem variable with indices 0,1,... n-1 containing 64 bit integers.

### Examples

```
say(random());
8781275837297675785

random(5);
{0=992008599211413363,
1=-923325524698236965,
2=-538113564003975731,
3=4402717987856179625,
4=-7602891879941326663}
```

Here is an example of generating a list of 10 random numbers between 1 and 10:

```
x. := 1+ abs(mod(random(10), 11));

say(x.);

\{0=4, 1=2, 2=5, 3=2, 4=5, 5=10, 6=9, 7=5, 8=9, 9=1\}
```

## random\_string

### Description

Generate a random hexadecimal string.

## Usage

```
random_string(n)
```

### **Arguments**

n (optional) – gives the size in bytes. The default is 16 bytes = 128 bits.

### Output

A string in hex format, all upper case.

### Examples

```
say(random_string(32));
```

#### to hex

#### Description

Convert a string to its hexadecimal representation

### Usage

to\_hex(arg)

#### **Arguments**

arg – a string or a stem of strings.

### Output

A string that is the hexadecimal representation of the underlying bytes for the string.

#### **Examples**

```
say(to_hex('The quick brown fox jumps over the lazy dog'));
54686520717569636b2062726f776e20666f78206a756d7073206f76657220746865206c617a7920646
f67
```

# **Stem specific functions**

#### box

### Description

Take any set of variables and turn them in to a stem, their names becoming the keys. This removes them from the symbol table so the only access afterwards is as part of the stem. See the function *unbox* for the inverse of this.

### Usage

```
box(var0, var1, ...);
```

### Arguments

There must be at least on argument. The arguments are variables that have been defined. These will be put in to a stem and removed from the symbol table.

### **Ouput**

A *true* if this succeeded.

#### **Examples**

```
a. := -5 + indices(5);
b. := 5 + indices(5);
)vars
a., b.
c. := box(a., b.);
    )vars
c.
```

So a. and b. no longer are in the symbol table, but are in the stem:

```
c.
{
b.={0=5, 1=6, 2=7, 3=8, 4=9},
a.={0=-5, 1=-4, 2=-3, 3=-2, 4=-1}
}
```

### exclude\_keys

#### Description

Remove a set of keys from a stem.

### Usage

```
exclude_keys(stem1,, stem2.)
```

### **Arguments**

stem1. = the target of this operation.

stem2. = a list of keys to be removed. Note that the *values* of this stem are what are to be removed from the target. There is no assumption that the keys of stem2 are integers, for instance.

## Output

A new stem that contains none of the keys in stem 2.

## Examples

```
a.foo := 'q';
a.bar := 'w';
b.w := 42;
b.a := 17;
say(exclude_keys(b., a.));
{a=17}

a.rule := 'One Ring to rule them all';
a.find := 'One Ring to find them';
a.bring := 'One Ring to bring them all';
a.bind := 'and in the darkness bind them';
```

```
list.0 := 'rule';
list.1 := 'bring';

say(exclude_keys(a., list.));
{bind=and in the darkness bind them,
find=One Ring to find them}
```

### from\_json

#### Description

Take a string that represents an object in JSON (JavaScript Object Notation) and return a stem representation. JSON is quite popular, (as in, it is inescapable anymore) but do remember it is a notation for objects that live in Javascript. To be blunt, it was designed to send information in REST-ful applications but because it is easy to use, is now being used by many as a general data description format. It was intended to tightly couple in-memory data structures on a server to be processed by a browser, so using it generally is arguably a bad idea. And yet, here we are. If you stick to the intended original purpose, it works great.

Note that there are some incompatibilities. First off, there is no actual standard way for JSON to represent all data, so you have to know what the structure of a JSON object is before you get it. (The simple example of something common that is impossible to represent a set in JSON without structuring it somehow.) Also, how to represent certain common things like dates varies by library. If you get something very odd where you expect a date (normally some large object with a ton of entries in it), then the suggestion is to, if at all possible, replace any dates with ISO 8601 format dates, which are the standard or perhaps integers that represent the time in milliseconds. Also, since stems end in a period, if an entry in a JSON object is another JSON object, that key will get a required period. If you use the to\_json command, this will always return valid JSON with stem markers. The names of properties will be escaped.

Stems are more general data structures than JSON, and if you really need to serialize an object to JSON (as it is properly termed) then you might want to serialize your stem in sub-units rather than send over a single massive object. Keep it simple.

**Note:** It should always work turning a JSON object in to a stem. Turning it back might not but the former is all we can guarantee. Now, if the stem follows a few guidelines of

- do not replicate integer indices as lists (it's fine to have stem.0. be an entry and the stem entry will be deserialized) but not both stem.0. and stem.0 and as long as this is observed
- No odd characters in the names of JSON properties

then there should be no problems mapping stems to JSON and back and forth.

### Usage

```
from_json(var)
```

#### **Arguments**

var = the string-valued variable to be converted

#### Output

A stem that contains the information in the original.

from\_json('{"woof":"arf","0":0,"1":1,"2":2}')

Generally a stem list corresponds to a JSON array, do if you need to convert back and forth, it is best not to overload the stem with non-list values. Just keep everything as simple as possible.

Note that JSON properties will be turned in to valid QDL variable names by escaping them as needed. This permits good interoperability.

#### Examples

#### A simple example.

```
{0=0, 1=1, 2=2, woof=arf}
Reading a file
// here is a large, messy example
  b := read_file('/tmp/my_json.json');
// (some indenting was done manually to keep this vaguely readable)
{"a":"b","s":"n","d":"m"
 "foo":{"tyu":"ftfgh","rty":"456","ghjjh":"456456",
 "woof":{"a3tyu":"ftf222gh","a3rty":"456222","a3ghjjh":"422256456"},
"0":"qwe","1":"eee","2":"rrr"},
"0":"foo", "1":"bar"}
  b. := from_json(b)
  say(b., true)
{
 a=b,
 s=n,
 d=m,
 foo.= {
  tyu=ftfgh,
  rty=456,
  ghjjh=456456,
  woof.= {
   a3tyu=ftf222gh,
   a3rty=456222,
   a3ghjjh=422256456
  },
  0=qwe,
  1=eee,
  2=rrr
 },
 0=foo,
 1=bar
```

```
}
// And just to check it really is a stem
  b.foo.woof.
{a3tyu=ftf222gh, a3rty=456222, a3ghjjh=422256456}
```

#### A simple example of a stem them cannot be represented in JSON

In this case, there is an index clash and an index error is generated

```
stem.0. := indices(3)
stem.0 := 5
// So let's show out stem. Note the there is both a scalar and list entry:
stem.
{0=5, 0.={0=0, 1=1, 2=2}}
to_json(stem.)
Error: The stem contains a list element "0" and a stem entry "0.". This is not convertible to a JSON Object
```

#### Escaping of JSON keys.

In this example, a simple JSON list is given and the key for this list is not even remotely a valid variable in QDL. In this case, the name is escaped.

```
a. := from_json('{"#$rt":[0,1,2]}');
a.
{$23$24rt.={0=0, 1=1, 2=2}}
a.$23$24rt.2 := 5
a.
{$23$24rt.={0=0, 1=1, 2=5}}
to_json(a.)
{"#$rt":[0,1,5]}
```

#### A less contrived example.

```
a. := from_json('{"Jäger-Groß":"enrolled"}')
a.
{J$C3$A4ger$2DGro$C3$9F=enrolled}
```

This may seem a bit hard to deal with, but do remember that there are many functions such as for\_keys that allow you to simply run over all stem values without having to know how the escaping turned out and if there is a question, you can use *vencode* or *vdecode* plus the fact that variables are resolved.

```
while[
   for_keys(key, a.)
]do[
   if[
        key == vencode('Jäger-Groß')
      ]then[
        // ... do stuff with this
   ];
];
// or just grab a values directly
x := vencode('Jäger-Groß');
```

attributes. := a.x; // grabs this by key and the values are now accessible.

#### list\_keys

#### Description

Return a list of the keys for a given stem variable.

#### Usage

```
list_keys(arg.)
```

#### **Arguments**

A stem variable. Remember that the entire stem is referenced by just the head + "."

#### Output

A list of keys where the new keys are cardinals and the values are the keys in the original stem. See also the *keys()* function.

#### **Examples**

Example. Let's say you have the following stem variable"

```
sourceStem.rule :='One Ring to rule them all;
sourceStem.find := 'One Ring to find them';
sourceStem.bring := 'One Ring to bring them all';
sourceStem.bind := 'and in the darkness bind them';
and you issue
var. := list_keys(sourceStem.);
The result would be
  var.0 == bind
  var.1 == find
```

**Note** that there is no canonical order to keys, so the keys to sourceStem. Can appear in any order in the result.

Another example.

var.2 == bring
var.3 == rule

Let us say that you got the above key set. How might you use it? In a loop:

```
while[
   for_keys(j, var.)
]do[
   sourceStem.var.j // ... do stuff with this
];
```

which loops through all the values in source stem.

### has\_keys

### Description

Check is a list of keys is in a target stem.

#### Usage

```
has_keys(target., list.)
```

### **Arguments**

target. — the target of this operation list. — a list of keys.

### Output

A boolean list with true as the value if the target contains the key and false if it does not.

### Examples

Just because it is easy to do, I am going to make a stem filled with 5 random integers, then a list of 10 indices. Obviously only the first 5 indices in *w*. will be in *var*.:

```
var. := random(5);
w. := indices(10);
say(has_keys(var., w.));
{0=true, 1=true, 2=true, 3=true, 4=true, 5=false, 6=false, 7=false, 8=false, 9=false}
```

## include\_keys

## Description

Take a stem, *a*. and a list of indices, *list*. and return the values of *a*. that have the same indices as *list*.

## Usage

```
include_keys(var., list.)
```

## **Arguments**

```
var. – a stem
```

list. - a list of keys. These will be the keys of the result, *var.* will supply the values.

#### Output

A stem with the keys from the *list*. and the corresponding values from the *var*.

### Examples

```
a.rule := 'One Ring to rule them all';
a.find := 'One Ring to find them';
a.bring := 'One Ring to bring them all';
a.bind := 'and in the darkness bind them';

list.0 := 'rule';
list.1 := 'bring';

say(include_keys(a., list.));
{bring=One Ring to bring them all,
rule=One Ring to rule them all}
```

#### indices

### Description

Make a list of indices. This is very useful in conjunction with looping.

### Usage

```
indices(arg);
```

### **Arguments**

arg is a number. This will be the size of the resulting index set.

### Output

A stem variable whose keys are the integers and whose entries are the same.

### Examples

```
x. := indices(3);
results in

x.0 == 0
x.1 == 1
x.2 == 2

So a common pattern is
while[
    for_keys(j, x.)
]do[
    myStem.j := // other stuff
    // myStem.x.j is an equivalent reference.
```

#### is list

### Description

Determine if the argument is precisely a list. That means, that it is a stem with only integer indices.

### Usage

```
is_list(stem.)
```

### **Arguments**

stem. The stem to check

### Output

A boolean that tells if this is a list.

### Examples

```
my_stem.help := 'this is my stem'
list_append(my_stem., indices(5))
{0=0, help=this is my stem, 1=1, 2=2, 3=3, 4=4}
is_list(my_stem.)
false
```

Why is this false? Because it has a non-integer index. The function tells you if the object is a list and only a list. Compare with

```
is_list(indices(10))
true
```

### keys

### Description

Return a stem of the keys for a given stem variable.

## Usage

```
keys(arg.)
```

### **Arguments**

A stem variable. Remember that the entire stem is referenced by just the head + "."

## Output

A stem of keys where every key has itself as the value.

#### **Examples**

Example. Let's say you have the following stem variable"

```
sourceStem.rule :='One Ring to rule them all';
sourceStem.find := 'One Ring to find them';
sourceStem.bring := 'One Ring to bring them all';
sourceStem.bind := 'and in the darkness bind them';
and you issue
  keys(sourceStem.)
{bind=bind,
  find=find,
  bring=bring,
  rule=rule}
```

**Note** that there is no canonical order to keys, so the keys to sourceStem. Can appear in any order in the result.

This is extremely useful with e.g. the rename function. So you can get all the keys, change their values and rename them.

An example to rename keys

Let us say we had the following stem with these keys (which were generated someplace else and we imported, e.g., from JSON):

```
b.0A2_foo := 'a';

b.0A2_woof := 'b';

b.0A2_arf := 'c';

b.fnord := 'd';

b.

{0A2_arf=c,

0A2_foo=a,

0A2_woof=b,

fnord=d}
```

The keys() command gives the following

```
keys(b.)
{0A2_arf=0A2_arf,
fnord=fnord,
0A2_foo=0A2_foo,
0A2_woof=0A2_woof}
```

To rename all the keys so that any with the OA2\_ prefix are changed, issue

```
rename_keys(b., keys(b.)-'0A2_')
{arf=c,
foo=a,
fnord=d,
woof=b}
```

### list\_append

#### Description

Append a list stem another list stem.

#### Usage

```
list_append(stem1., stem2. | value)
```

### **Arguments**

```
stem1. = the stem to which the values will be appended
stem2. = the stem from which values will be taken or
value = the scalar that will be appended.
```

### Output

A stem, the same as the first argument but with the appended values.

#### Examples

```
stem1. := indices(5);
stem1.
{0=0, 1=1, 2=2, 3=3, 4=4}
stem2. := 5*indices(6);
stem2.
{0=0, 1=5, 2=10, 3=15, 4=20, 5=25}
list_append(stem1., stem2.)
{0=0, 1=1, 2=2, 3=3, 4=4, 5=0, 6=5, 7=10, 8=15, 9=20, 10=25}
list_append(stem2., 'woof')
{0=0, 1=5, 2=10, 3=15, 4=20, 5=25, 6=woof}
```

### list\_copy

### Description

Copy from one list stem to another.

## Usage

```
list_copy(source., start_index, length, target., target_index)
```

### **Arguments**

```
source. = the stem that is the source of the copy.
start_index = the index in the source where the copy starts
```

length = how many elements to copy

```
target. = the target stem of the copy
```

target\_index = the index in the target that will receive the copy. Note that any elements already in these locations will be replaced. If you need to insert elements, consider using the list\_insert command.

#### Output

The updated target stem. Note that the target *is* modified in this operation.

#### **Examples**

```
source. := indices(5)+10
target. := indices(6) - 50
list_copy(source., 2, 3, target., 4)
{0=-50, 1=-49, 2=-48, 3=-47, 4=12, 5=13, 6=14}
```

So this took the 3 elements from source. starting at index 2 and copied them to target. starting at index 4 there.

#### list insert at

### Description

insert a sublist into another list, starting at a given point. All the indices in the target list are shuffled to accommodate this.

#### Usage

```
list_insert_at(source., start_index, length, target., target_index);
Arguments
```

### Examples

```
source. := indices(5) + 20
target. := indices(6) - 100
list_insert_at(source., 2, 3, target., 4)
{0=-100, 1=-99, 2=-98, 3=-97, 4=22, 5=23, 6=24, 7=-96, 8=-95}
```

So this inserted 3 elements from the source starting at index 2 in the source and placed them at index 4 in the target, moving everything else.

### list\_subset

## Description

Grab a subset of a given list.

### Usage

```
list_subset(source., start_index [, length]);
```

#### **Arguments**

```
source. = the stem list to take a subset of
start_index = where to start in the source stem
```

length (option) = how many elements to take. Omitting this means take the rest of the stem

### Output

The altered stem. Note that this does nothing to the original stem and the indices are adjusted accordingly, so that the first index of the output is 0.

### Examples

```
stem. := indices(5) + 20;
// Just grab the tail of this list
list_subset(stem., 2)
{0=22, 1=23, 2=24}
// grab some stuff in the middle.
list_subset(stem., 1, 3)
{0=21, 1=22, 2=23}
```

#### mask

### Description

Take a boolean mask of a stem.

### Usage

```
mask(target. bit_stem.)
```

### Arguments

target – the stem variable to acted upon.

bit\_mask – a stem variable with the same keys as target and boolean values. If the value is **true** then the entry is kept in the result and if **false** it is not. Note that if there are missing keys then these will not be returned either (so subsetting is still in effect), essentially making them equivalent to **false** entries.

## Output

A subset of the target.

#### **Examples**

```
header.transport := 'ssl';
header.iss := 'OA4MP_agent';
header.idp := 'http://oa4mp.org/idp/secure';
header.login_allowed := 'true';
// case insensitive match
header. := mask(header., !contains(header., 'oa4mp', false));
// This removes every entry containing 'oa4mp'
say(size(header.);
```

#### rename\_keys

### Description

Rename the keys in a stem. See also the *keys()* command.

#### Usage

```
rename_keys(target., new_keys.);
```

#### **Arguments**

target. - the stem to be altered

new\_keys. - a stem of keys. The keys in it are the ones in target. To be altered to their values.

### Output

A new stem whose keys have been altered as per the second list.

## Examples

```
a.foo := 42;
a.bar := 43;
a.baz := 44;
key_list.foo := 'a';
key_list.bar := 'b';
key_list.baz := 'woof';
rename_keys(a., key_list.)
say(a.);
{a=42, b=43, woof=44}
```

Note that since this changes the keys in the target., the key\_list. Unrecognized keys are skipped. A subset is fine too:

```
a.foo := 42;
a.bar := 43;
a.baz := 44;
key_list.foo := 'a';
key_list.bar := 'b';
rename_keys(a., key_list.)
```

```
say(a.);
{a=42, b=43, baz=44}
```

#### set\_default

### Description

Set the default value for a stem. If a kys is requested but has not been set, the default value is returned. This allows you initialize a stem without having to explicitly fill in every value. Note especially that the default value is not figured in to other calculations.

### Usage

```
set_default(target., scalar);
Arguments
target. – the stem
scalar – the default value
```

### Output

This returns the default value set.

### Examples

```
set_default(x., 1);
say(x.);
{}
```

So no values have been defined. Let's set one and check it:

```
x.0 := 10;
say(x.);
{0=10}
```

And if we needed to access a value of x. that has not been set

```
say(x.1);
1
```

Just to emphasize, default values are not used in most stem operations.

```
say(get_keys(x.));
{0=0}
```

#### size

#### Description

Return the size of the argument

### Usage

```
size(var)
```

### **Arguments**

var – any variable or argument

### Output

This varies.

- stem the number of keys (this does not check if there are stems as values)
- string the length of the string
- boolean, integer, decimal zero, since these are scalars.

### Examples

```
say(size(42));

say(size('abcd'));
4
```

## to\_json

### Description

Convert a stem to a JSON string. Note that JSON = JavaScript Object Notation is a common way to represent objects and is treated as a notation, not a data structure. See the extended note in the **from\_json()** section.

## Usage

```
to_json(stem. [,convert? | indent, indent])
```

### **Arguments**

stem. = the stem to represent in JSON notation

indent (optional) = whether or not to indent the resulting string to make it more readable. This controls how much whitespace is added. The higher the number, the more space in the result. Usually a value of 1 or 2 is sufficient for most cases.

So these are valid calls

```
to_json(stem., false) – do not convert the names
```

to\_json(stem. 2) – indent the output with a spacing of 2

to\_json(stem., false, 2) – do not convert the stem variable names and indent with a spacing of 2.

#### Output

A string in JSON which represents the argument.

#### **Examples**

```
a. := indices(3)
a.woof := 'arf'
to_json(a.)
{"woof":"arf","0":0,"1":1,"2":2}
// and just to show how to indent the result
to_json(a.,1)
{
  "woof": "arf",
  "0": 0,
  "1": 1,
  "2": 2
}
```

Large JSON objects are often best handled through files or other means rather than directly.

```
claims. := from_json(read_file('/tmp/claims.json'));
  size(claims.)
137
```

Again, JSON is a notation for an object and you must know what the structure of the object is and all the particulars about it to do anything useful with it.

### to\_list

### Description

Take an arbitrary list of things and put them all in to a stem list.

## Usage

```
to_list(arg0, arg1, arg2, ...)
```

#### **Arguments**

Any number of arguments of any type are allowed, including stems.

### Output

A stem list

### Examples

In this example, a few values are given, including a stem that is made with the *indices* command. Note that each element of the new list is one of the arguments.

```
say(to_list(3,45.34,'abc',indices(3)));
{0=3, 1=45.34, 2=abc, 3={0=0, 1=1, 2=2}}
```

#### unbox

### Description

Takes a stem variable and splits it up, turning each key in to a variable. Note that this worls only as long as the keys are valid variable names.

### Usage

```
unbox(stem.);
```

### **Arguments**

stem. - the stem to unbox

## Ouput

A true if the result worked.

### Examples

#### union

# Description

Take a set of stems and put them all together in to a single stem

#### Usage

```
union(stem1., stem2., ,,,);
```

#### Arguments

The arguments are either stems or variables that point to stems.

### Output

The output is a new stem that contains all of the keys. Note that if there are multiple keys then the *last* argument with that key is what is set. The result is guaranteed to have every key in all the arguments in it.

#### Examples

```
a. := -5 + indices(10);
b. := 5 + indices(5);
a.arf := 'woof';
b.woof := 'bow wow';
c. := -20 + indices(3);
union(a., b., c.)
{0=-20, 1=-19, 2=-18, 3=8, 4=9, arf=woof, 5=0, 6=1, 7=2, 8=3, woof=bow wow, 9=4}
```

Note that in this example these stems have some keys contained in the previous one.

## Input and output

#### Dir

# Description

List a directory content

### Usage

dir(arg)

### **Arguments**

arg is the path to a directory. It may also be in a virtual file system

## Output

A stem list of the elements of the directory

### Examples

```
dir('qdl-vfs#/zip/root')
{0=scripts/, 1=other/, 2=readme.txt}
```

This lists the given directory in the mounted VFS. Note that in this case it so happens to be a zip archive of a file, mounted at qdl-vfs#/zip/.

#### mkdir

### Description

Make a set of directories in a file system

#### Usage

mkdir(arg)

### **Arguments**

arg is a path. All of the intermediate paths will be created as needed.

### Output

A boole, true if the operation succeeded and false otherwise.

### Examples

An example of trying to make a directory in a read-only VFS will fail:

```
mkdir('qdl-vfs#/zip/foo')
Error: You do not have permissions make directories in the virtual file system
```

Making a directory in a system that is writeable works fine:

```
mkdir('qdl-vfs#/pt/woof-123')
true
```

#### rmdir

### Description

Remove an empty directory from a file system.

## Usage

rmdir(arg)

#### **Arguments**

arg is a path in a file system to an empty directory. You must remove all files and sub-directories for this to work. Also, unlike mkdir, this will only remove the last component.

### Output

A true if this succeeded and a false otherwise.

### Examples

#### rm

### Description

Remove a single file from a directory.

### Usage

rm(arg)

### **Arguments**

arg is the full path to the file

## Output

A true if this succeeded, false otherwise

## Examples

## print, say

## Description

Print out the argument to the console. The two functions *say* and *print* are synonyms. Use whichever you prefer for readability.

## Usage

```
say(arg [,prettyPrintForStems])
print(arg [,prettyPrintForStems])
```

#### **Arguments**

```
arg – anything.
```

prettyPrintForStems (optional) -*IF* arg is a stem, try to print a pretty version of it, defined as being more vertical.

#### Output

The printed representation of the argument will be put o the console **and** the value returned is whatever was printed (so you can embed it in other statements – a very useful debugging trick.)

### Examples

The momentous entire "Hello World" program in QDL:

```
say('Hello World');
Hello World
```

And since 42 is the answer to all Life's questions (as per the Hitchhiker's Guide To The Galaxy)

```
print(42);
42
    say(42);
42
```

Here is an example of how to use this to intercept and print out an intermediate result,

```
a := say(432 + 15);
447
```

Pretty print only applies to stems. It attempts to make a somewhat more human readable version

```
f. := indices(6);
  say(f., true);
{
  0=0,
  1=1,
  2=2,
  3=3,
  4=4,
  5=5
}
```

Compare this with just plain printing it:

```
say(f.);
{0=0, 1=1, 2=2, 3=3, 4=4, 5=5}
```

### read file

# Description

Read a file. The result is always a string

#### Usage

```
read_file(file_name [, to_list || is_binary])
```

#### **Arguments**

file\_name — the full path to the file. This may be in a virtual file system too.

The next argument is optional and is an integer

```
to_list == 0 - return the result as a stem list each line separate
```

is\_binary == 1 - return the result as a base 64 encoded string of bytes.

If no second argument is given, the result is simply a string of the entire contents of the file.

#### Output

Either a simple string (only file name is given), a stem if it is flagged as a list or a base64 string if it si flagged as binary.

#### Examples

```
cfg. := read_file('/var/lib/tomcat/conf/server.xml', 0);
```

Would read in the file /var/lib/tomcat/conf/server.xml and return a stem. Each line in the file is in order in *cfg.0*, *cfg.1*, ... Compare this with

```
big_string := read_file('/var/lib/tomcat/conf/server.xml');
```

Which reads the same file and puts the entire thing in a single string.

```
my_b64 := read_file('/var/lib/crypto/keystore.jks' , 1);
```

this reads the keystore.jks file (which is binary) and base64 encodes it, storing it in the *my\_b64* variable. QDL does not have the capacity to do low-level operations on binary data, but it can move them where they need to go faithfully.

A couple of more examples:

```
// read a file as a stem
    say(read_file('/home/ncsa/dev/ncsa-git/security-lib/ncsa-qdl/src/test/
resources/hello_world.qdl',0));
{0=/*, 1= The expected Hello World program. , 2= Jeff Gaynor, 3= 1/26/2020,
4=*/, 5=say('Hello world!');}

// read the exact same file and turn the bytes into a base 64 string.
    say(read_file('/home/ncsa/dev/ncsa-git/security-lib/ncsa-qdl/src/test/
resources/hello_world.qdl',1));
LyoKICBUaGUgZXhwZWNOZWQgSGVsbG8gV29ybGQgcHJvZ3JhbS4gIAogIEplZmYgR2F5bm9yCiAgMS8yNi8
yMDIwCiovCnNheSgnSGVsbG8gd29ybGQhJyk7Cg
```

say(decode\_b64('LyoKICBUaGUgZXhwZWN0ZWQgSGVsbG8gV29ybGQgcHJvZ3JhbS4gIAogIEp1ZmYgR2F5bm9yCiAgMS8yNi8yMDIwCiovCnNheSgnSGVsbG8gd29ybGQhJyk7Cg'));

```
/*
  The expected Hello World program.
  Jeff Gaynor
  1/26/2020
*/
say('Hello world!');
```

(This is the sample hello world program for qdl).

#### scan

#### Description

Prompt a user for input

### Usage

```
scan([prompt])
```

#### **Arguments**

prompt (optional) – something to print out, probably cuing the user.

#### Output

Whatever the user types in as a string. There is no end of line marker returned.

# Examples

```
response := scan('do you want to continue?(y/n):');
do you want to continue?(y/n):y
    say(response);
y
```

So the user sees the prompt (in this case "do you want to continue?(y/n):") and types in the response of "y", which is stored in the variable *response*. In this next example we input a loop in buffer mode, then execute it.

```
)buffer on
stop_looping := 'n';
while[
        stop_looping != 'y'
        ]do[
        stop_looping := scan('stop looping? (y/n):');
        ]; //end do
        .
stop looping? (y/n):foo
stop looping? (y/n):bar
stop looping? (y/n):y
```

Only when we enter the expected response of "y" does it stop.

#### vfs mount

#### Description

Mount a virtual file system

#### Usage

```
vfs_mount(root_dir, scheme, mount_point [, permissions])
```

### **Arguments**

```
root_dir = the local directory that is the new root
```

scheme = the scheme (label) for this system

mount\_point = the internal path (starts with a /) for programs to refer to.

permissions (optional) = the permissions the VFS has. These are 'r' for read and 'w' for write. If omitted, the VFS is mounted in read-only mode.

#### Output

A 0 if there was no problem.

### Examples

In this example, we will mount a local file system and read a file. We mount the VFS for both reads and writes. You refer to a file in the vfs seamlessly using the scheme to prefix it.

```
vfs_mount('/home/ncsa/dev/qdl/scripting','qdl-vfs','/', 'rw')
0
    read_file('qdl-vfs#/client.xml')
<?xml version="1.0" encoding="UTF-8" standalone="no"?>
<!DOCTYPE properties SYSTEM "http://java.sun.com/dtd/properties.dtd">
comment>OA4MP stream store</comment>
... lots more
```

So the given file is loaded and read. All file operations behave normally. The reason for virtual file systems is two-fold. First off, if QDL is running in server mode, directories may be mounted in read only fashion to provide access to libraries, modules and such in a completely installation independent way. Secondly, when QDL is running in server mode, all standard file operations are prohibited but you may still have virtual ones. This allows a server to, for instance, mount a jar file with libraries in it.

(The reason for this on a server is security: Often servers run with enhanced privileges which may be inherited by applications. QDL always seeks to be a good citizen and only allows what is specifically granted to it.)

#### write file

#### Description

write contents to a file

### Usage

```
write_file(file_name, contents [,is_base64])
```

### **Arguments**

file\_name - the name of the file.

contents = A string or a stem list. If the stem is not a list (so indices 0, 1, ...) then this will fail/

is\_base64 (optional) – if **true** then try to decode this into binary and write it as a binary object.

### Output

Returns **true** if this succeeded.

#### Examples

```
hello_world :=
'LyoKICBUaGUgZXhwZWN0ZWQgSGVsbG8gV29ybGQgcHJvZ3JhbS4gIAogIEplZmYgR2F5bm9yCiAgMS8yNi
8yMDIwCiovCnNheSgnSGVsbG8gd29ybGQhJyk7Cg';
   write_file('/tmp/test.qdl', hello_world, true);
```

This is just the base64 encoded hello world.qdl script from the read file example.

# **Scripts**

A script is simply a sequence of QDL commands in a file. You may run scripts in a variety of ways and these commands let you do it from in the workspace.

#### execute

## Description

Send a string to the interpreter and evaluate it.

### Usage

execute(string)

### **Arguments**

string – the string to be executed. This must be a valid QDL statement or it will fail

### Output

Whatever the executes statement outputs.

### **Examples**

```
execute('say(2 + 2);'));
4

execute('say(\'abc\' + \'def\');');
abcdef
```

### load\_script

### Description

Read a script from a file and execute it in the current environment. This means that any variables it sets of functions it defines are now part of the active workspace. Caution that this will overwrite whatever you have if there is a name clash.

#### Usage

```
load_script(file_name)
```

### **Arguments**

file\_name — the fully qualified path to the file

## Output

None

### Examples

```
load_script('/home/bob/qdl/math_util.qdl');
```

# run\_script

### Description

Read a script from a file and execute it in a completely new environment. The output of the file is piped to the current console.

## Usage

```
run_script(file_name)
```

#### **Arguments**

file\_name - the fully qualified path to the file

### Output

none

### Examples

```
run_script('/home/bob/qdl/format_reports.qdl');
```

# **General functions**

These are functions that are generally applicable and do not fall in to the other categories.

### is\_defined

### Description

A scalar-only function that will return if a given variable is defined, i.e., has been assigned a value.

### Usage

```
is_defined(var)
```

### **Arguments**

var is the variable. Remember that stem variables end with a period if you are addressing the entire thing.

# Output

A boolean.

## Examples

```
a := 'foo';
    say(is_defined(a));
true

    say(is_defined(b));
false

    b. := make_index(4);
    say(is_defined(b.));
true

    say(is_defined(b.1));
true
```

```
say(is_defined(b.woof));
false
```

This last example shows that if a stem is defined, then you can use this to check the elements as well.

#### is function

### Description

Checks if a symbol is a function.

#### Usage

```
is_function(var [, argCount])
```

#### **Arguments**

var is a string that contains the name of the function.

argCount (optional) — This is the number of arguments that the function may accept. If you omit it than the result will reflect if there is *any* such named function, regardless of argument count.

### Output

A boolean which is **true** if the function is defined in the current scope.

### Examples

In this case, a function, *f* is defined in a module called *mytest:functions* 

```
say(is_defined('f'));
false
   import('mytest:functions');
   say(is_defined('f'));
true
```

This also works with stem elements, so

```
is_defined(t.x)
```

would return true if the stem t. contained the element x.

#### remove

### Description

Remove a stem and its values from the symbol table

#### Usage

```
remove(var)
```

#### **Arguments**

var – a scalar or stem to be removed.

#### Output

none

false

stem.

 $\{0.=\{0=0, 1=1, 2=2\}\}$ 

#### Examples

Here we define a stem and check is defined, then remove it.

```
t. := indices(5);
    say(is_defined(t.));
true
    remove(t.);
    say(is_defined(t.));
false
Here we set a variable then remove it.
    p := 'abc';
    say(is_defined(p));
true
    remove(p);
    say(is_defined(p));
false
Also, this will remove entries to stems, so
remove(t.b)
will remove the entry with index b from the stem t. Similarly
remove(t.x.)
will remove the entire sub-stem x. Use with care!
  stem.0. := indices(3)
  stem.0 := 5
  stem.
\{0=5, 0.=\{0=0, 1=1, 2=2\}\}
  is_defined(stem.0)
true
  remove(stem.0)
  is_defined(stem.0)
```

is\_defined(stem.0.)
true