

Multiparameter Persistent Homology

Topological Data Analysis and Neuroscience

<https://github.com/ndag/TDA-and-Neuro>

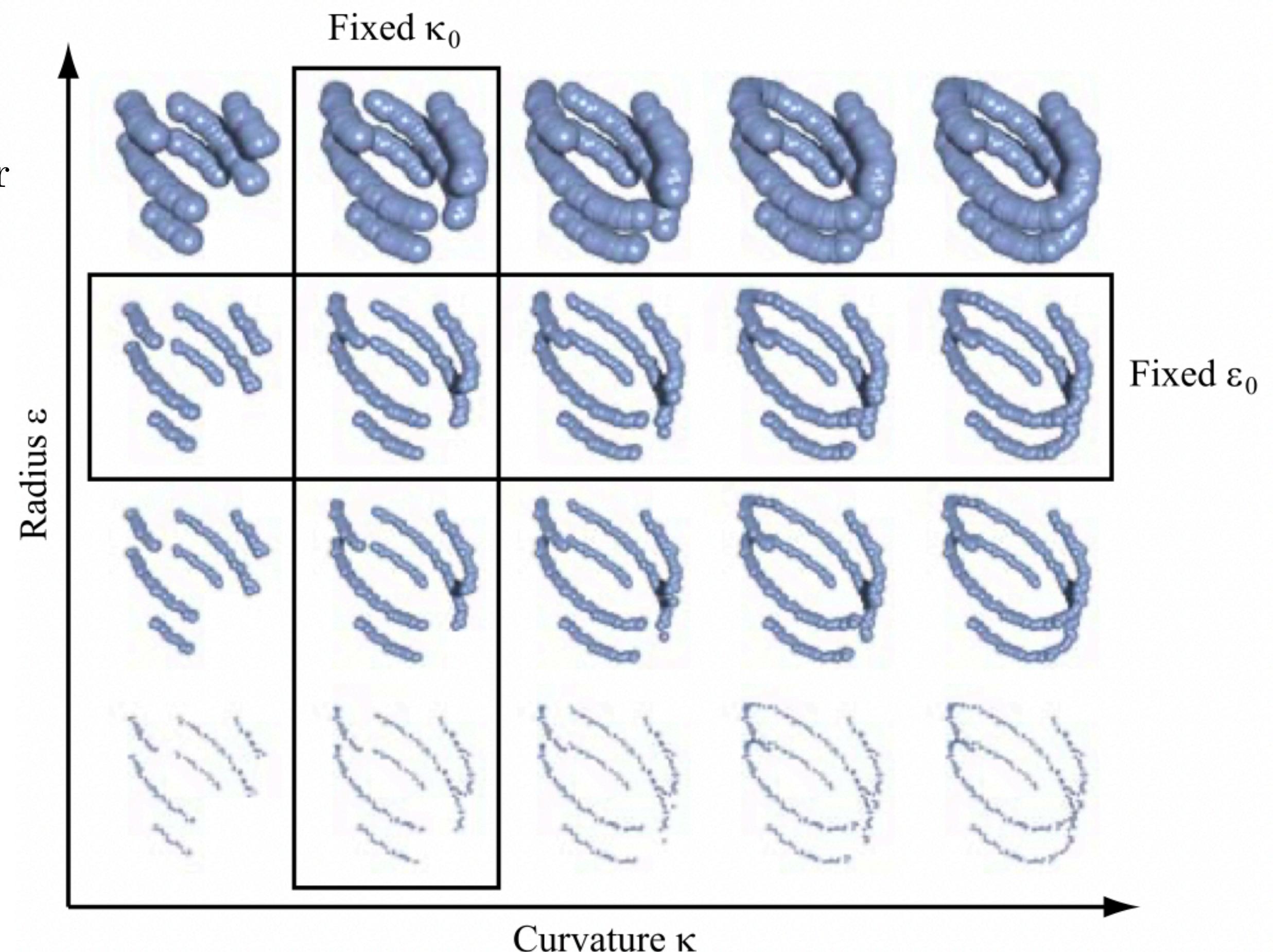
Alex McCleary

The Main Idea

- The goal of 1-parameter persistence is to break down filtrations into their fundamental parts, called indecomposables.
- Unfortunately, this goal has not been realized in a practical way. Either the indecomposables used are so complicated they are impossible to work with or they lose information.
- While this situation is not ideal, we can still learn a lot from many of the invariants proposed.

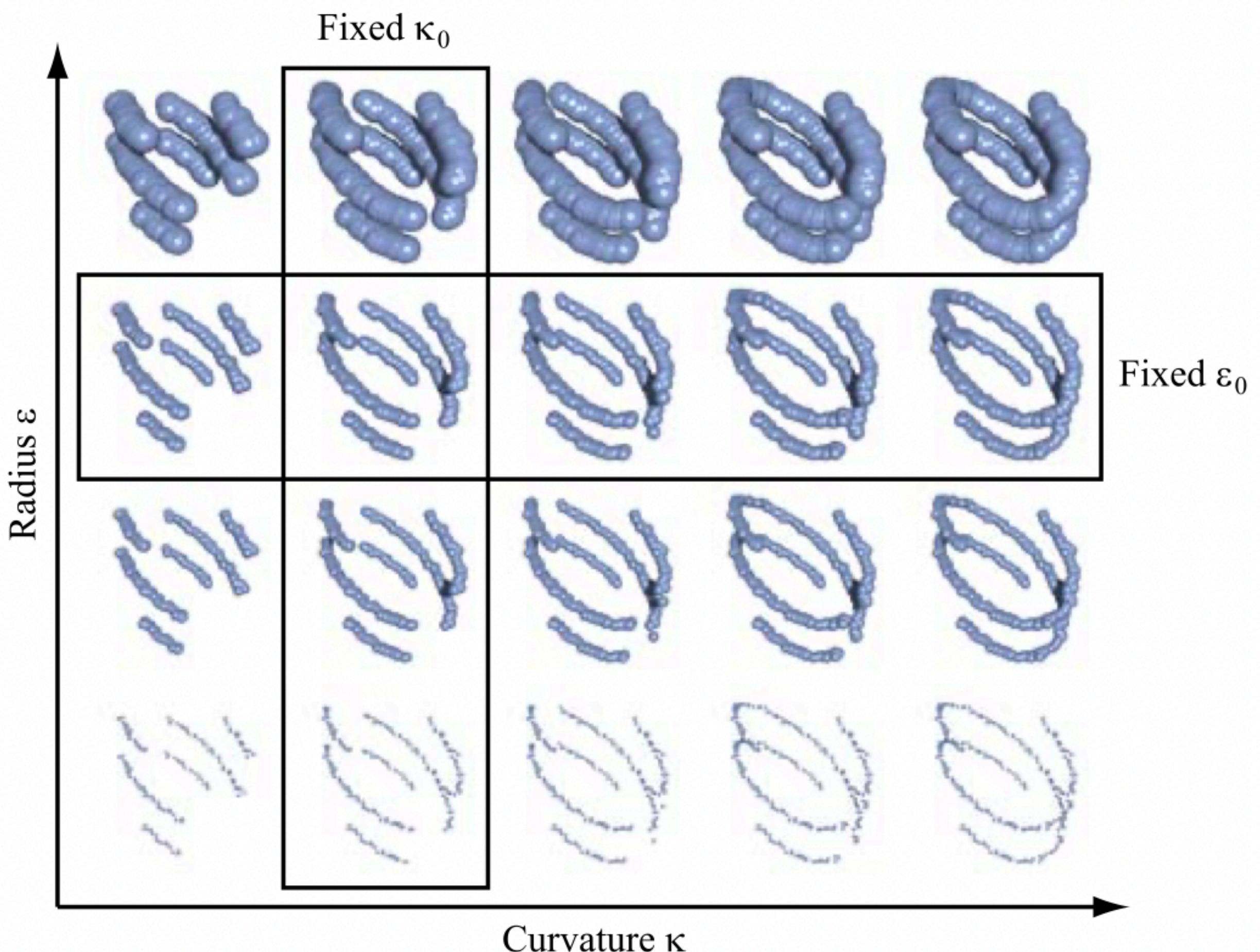
Multiparameter Filtrations

- A multiparameter filtration \mathcal{F} consists of a topological space $\mathcal{F}(\vec{a})$ for each $\vec{a} \in \mathbb{R}^n$ and an inclusion $\mathcal{F}(\vec{a}) \hookrightarrow \mathcal{F}(\vec{b})$ whenever $\vec{a} \leq \vec{b}$.



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- The main sources of examples for multiparameter filtrations are:
 - Augmented metric spaces. That is, metric spaces with real valued functions on them.
 - Sublevel set filtrations.
 - Level set filtrations.

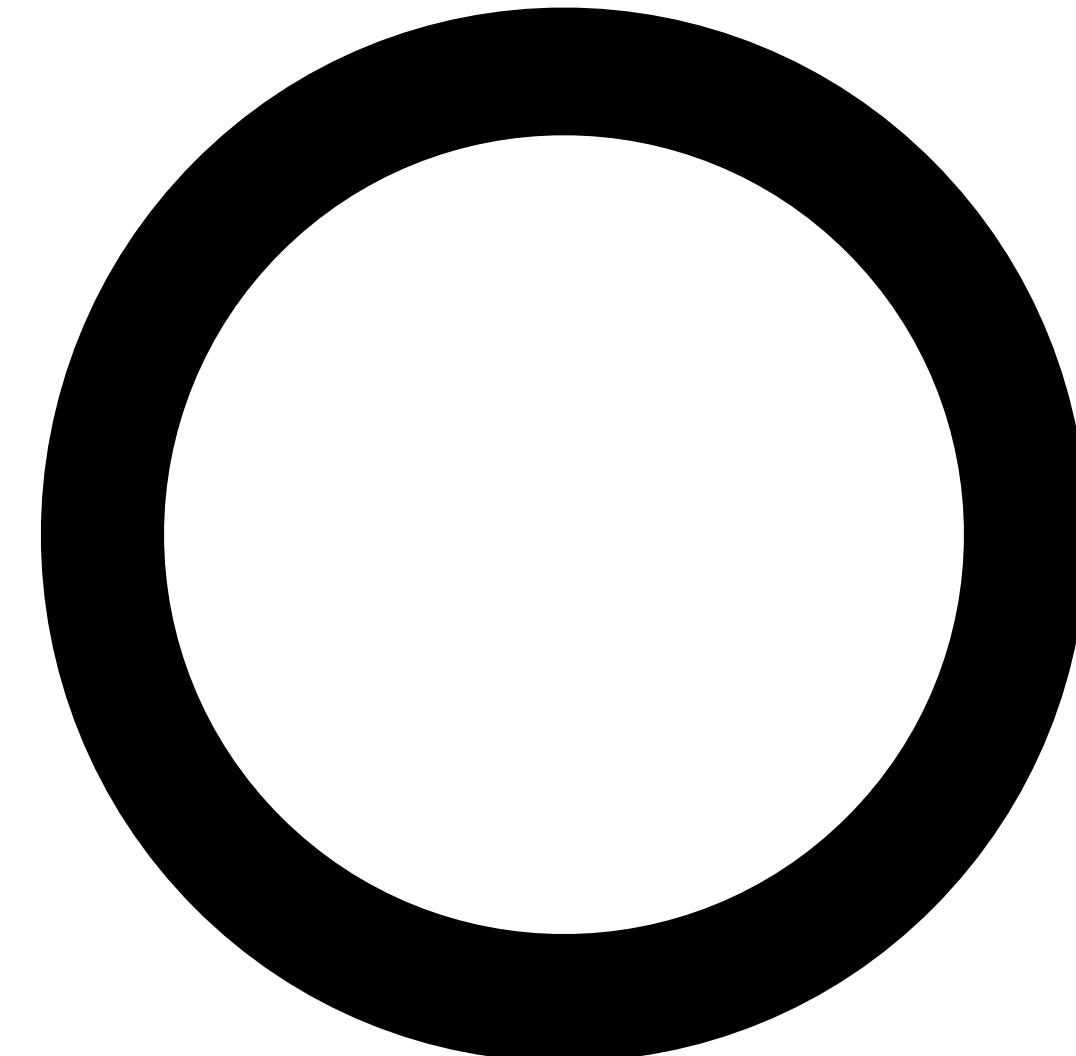


Example: Sublevel Set Filtrations

Let X be a topological space (or simplicial complex). Every function $f : X \rightarrow \mathbb{R}^n$ induces an n -parameter filtration M by $\mathsf{M}(\vec{a}) = f^{-1}((-\infty, \vec{a}])$

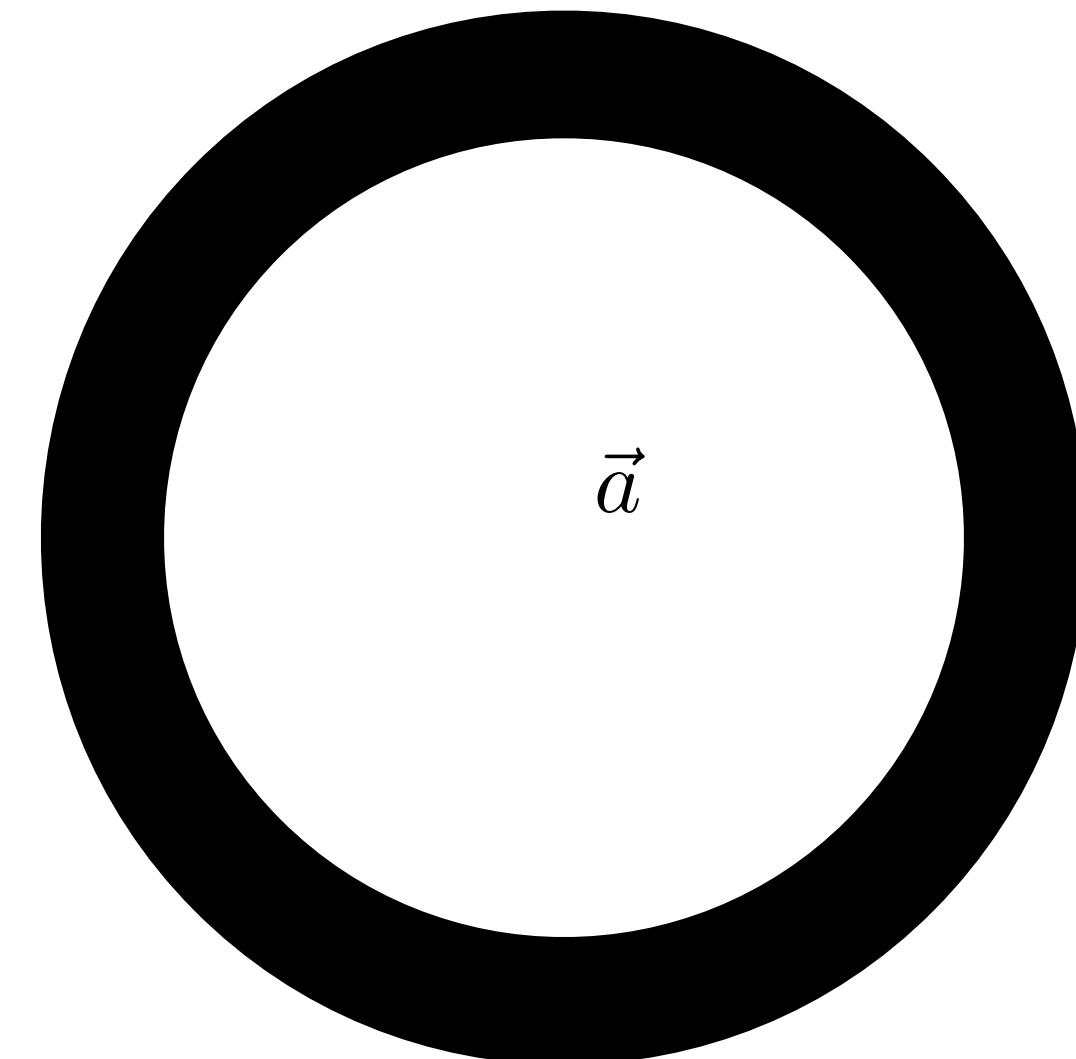
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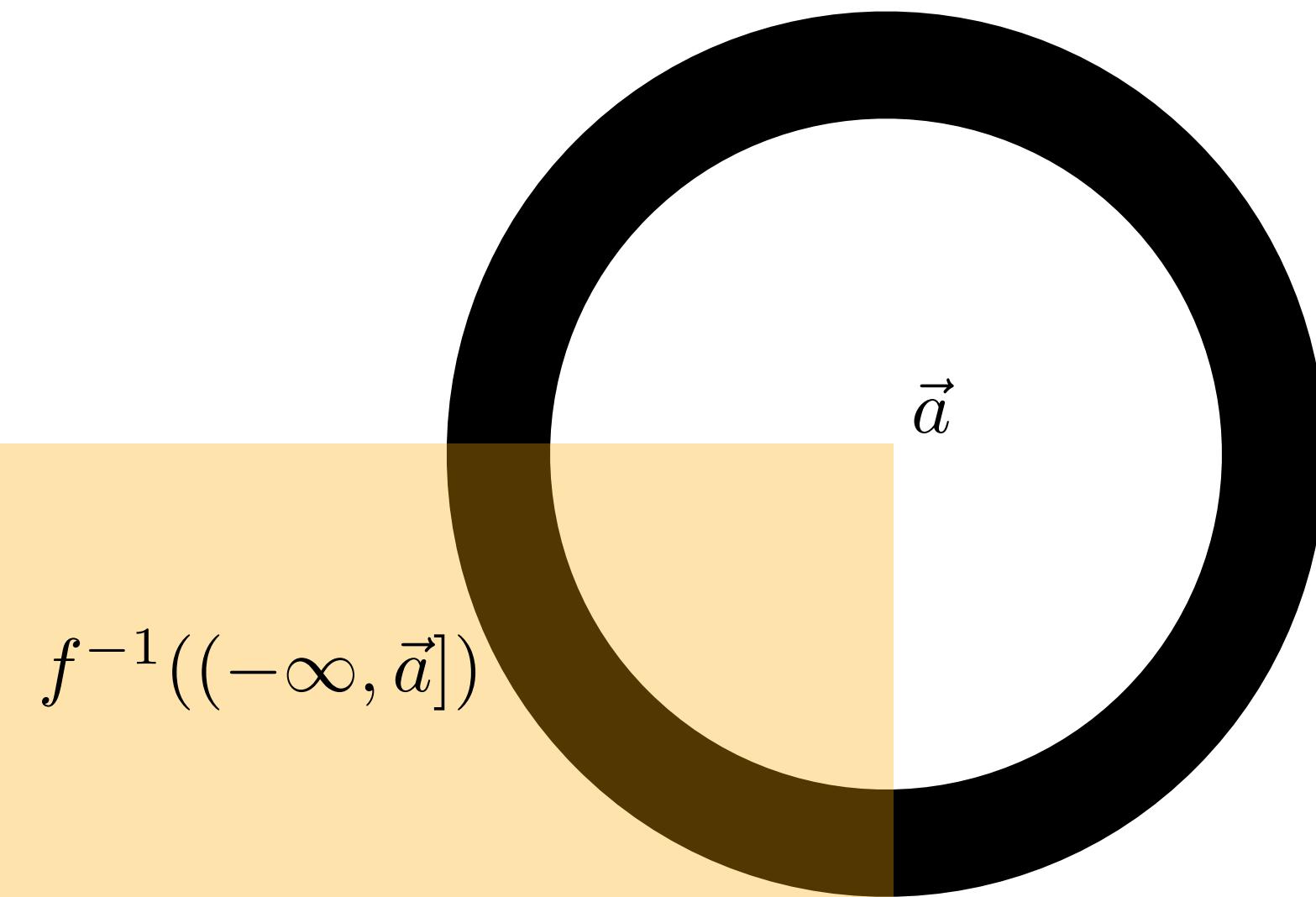
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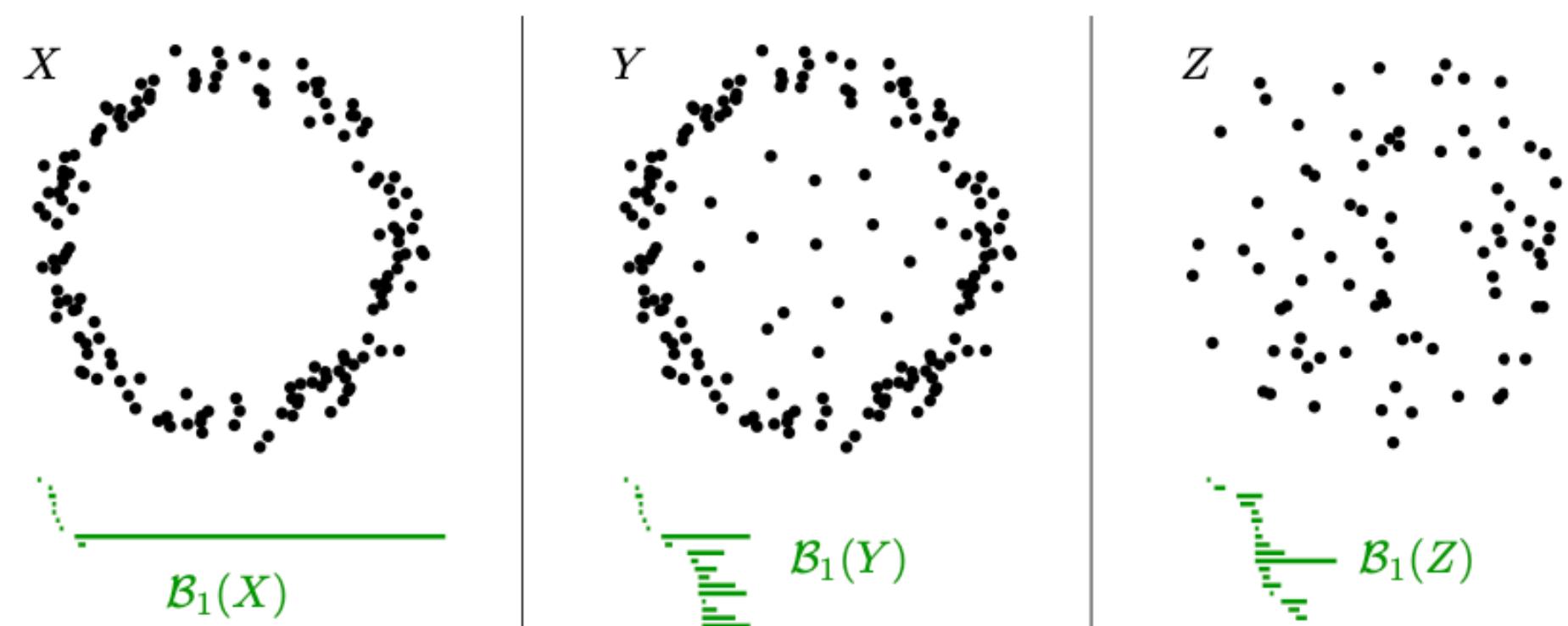
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Persistence Modules

Multiparameter Persistence Modules

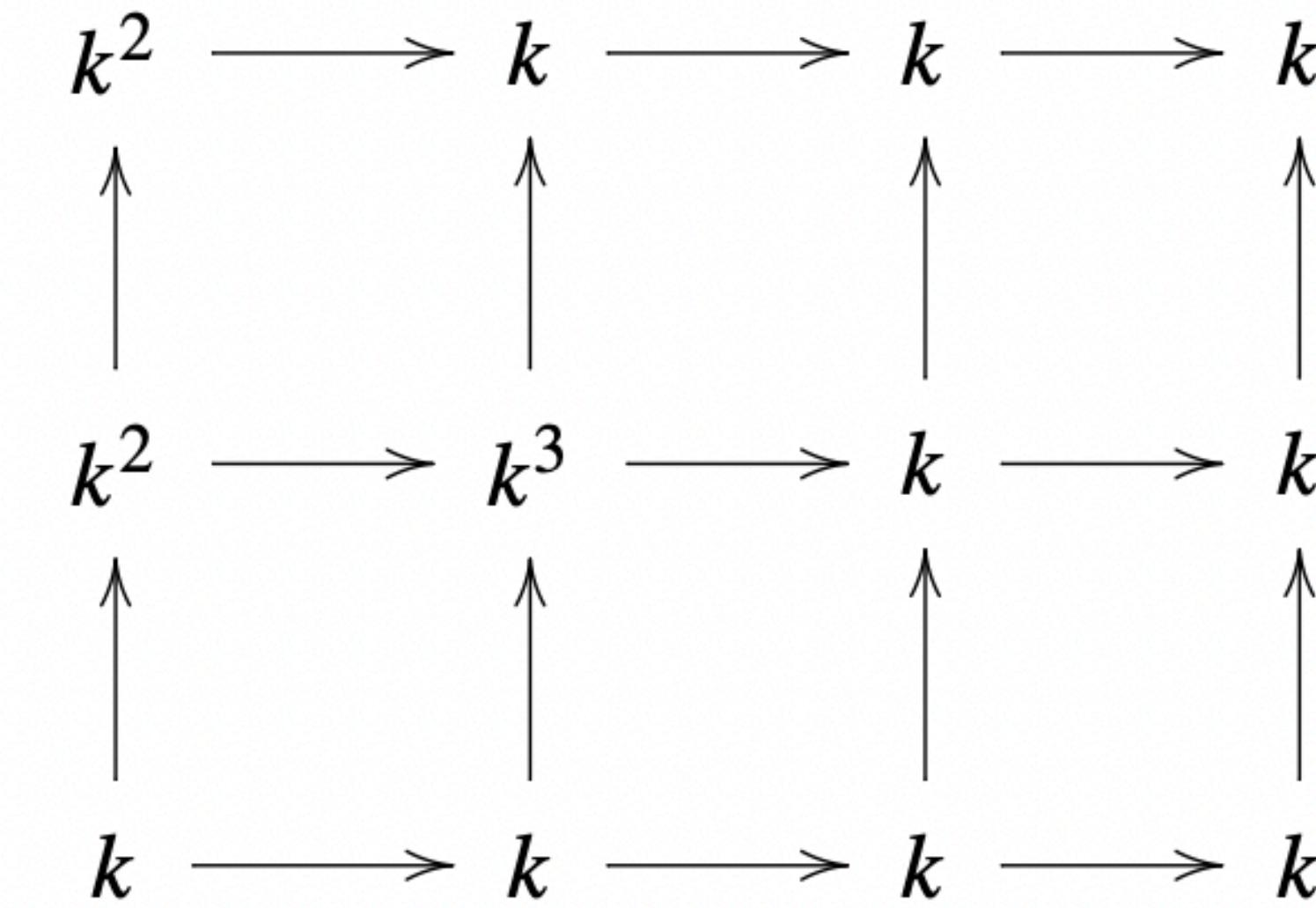
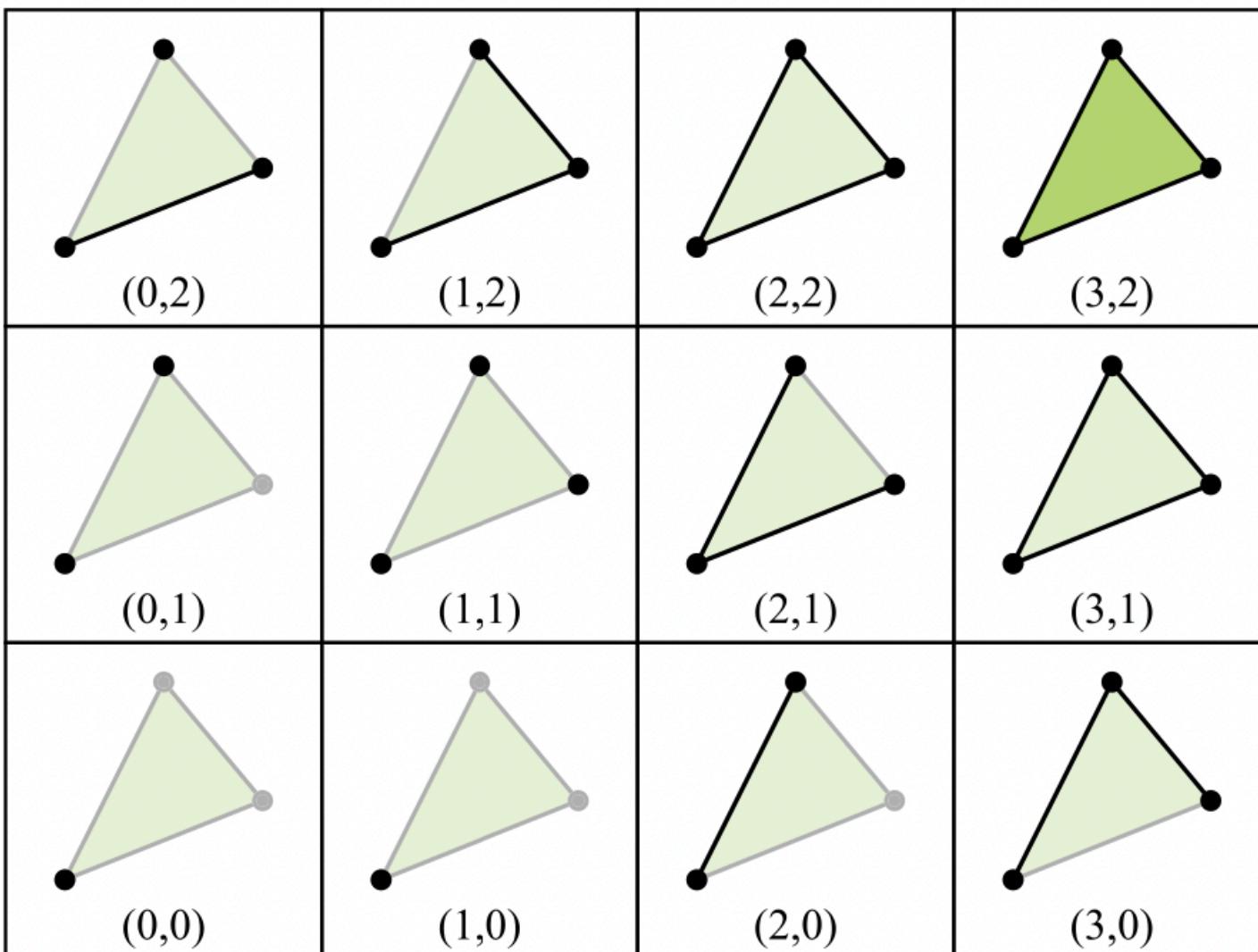
- A multiparameter **persistence module** consists of a vector space $\mathbf{M}(a)$ for each $a \in \mathbb{R}^n$ and a linear transformation $\mathbf{M}(a \leq b) : \mathbf{M}(a) \rightarrow \mathbf{M}(b)$ for each $a \leq b$.

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- An ε -interleaving between M and N consists of linear maps $\varphi_{\vec{a}} : M(\vec{a}) \rightarrow N(\vec{a} + \vec{\varepsilon})$ and $\psi_{\vec{a}} : N(\vec{a}) \rightarrow M(\vec{a} + \vec{\varepsilon})$ for each $a \in \mathbb{R}^n$ such that the following diagrams commute

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$$\begin{array}{ccc} M(\vec{b}) & \xrightarrow{\varphi_{\vec{b}}} & N(\vec{b} + \vec{\varepsilon}) \\ \uparrow & & \uparrow \\ M(\vec{a}) & \xrightarrow{\varphi_{\vec{a}}} & N(\vec{a} + \vec{\varepsilon}) \end{array}$$

$$\begin{array}{ccc} M(\vec{a} + 2\vec{\varepsilon}) & & N(\vec{a} + 2\vec{\varepsilon}) \\ \uparrow & \swarrow \psi_{\vec{a} + \vec{\varepsilon}} & \uparrow \\ M(\vec{a}) & \xrightarrow{\varphi_{\vec{a}}} & N(\vec{a} + \vec{\varepsilon}) \\ & \uparrow & \uparrow \\ & M(\vec{a} + \vec{\varepsilon}) & N(\vec{a}) \end{array}$$

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Interval Decomposable Persistence Modules

- A subset $I \subseteq \mathbb{R}^n$ is an interval if:
 - For any $a, c \in I$ and any $b \in \mathbb{R}^n$ with $a \leq b \leq c$, $b \in I$
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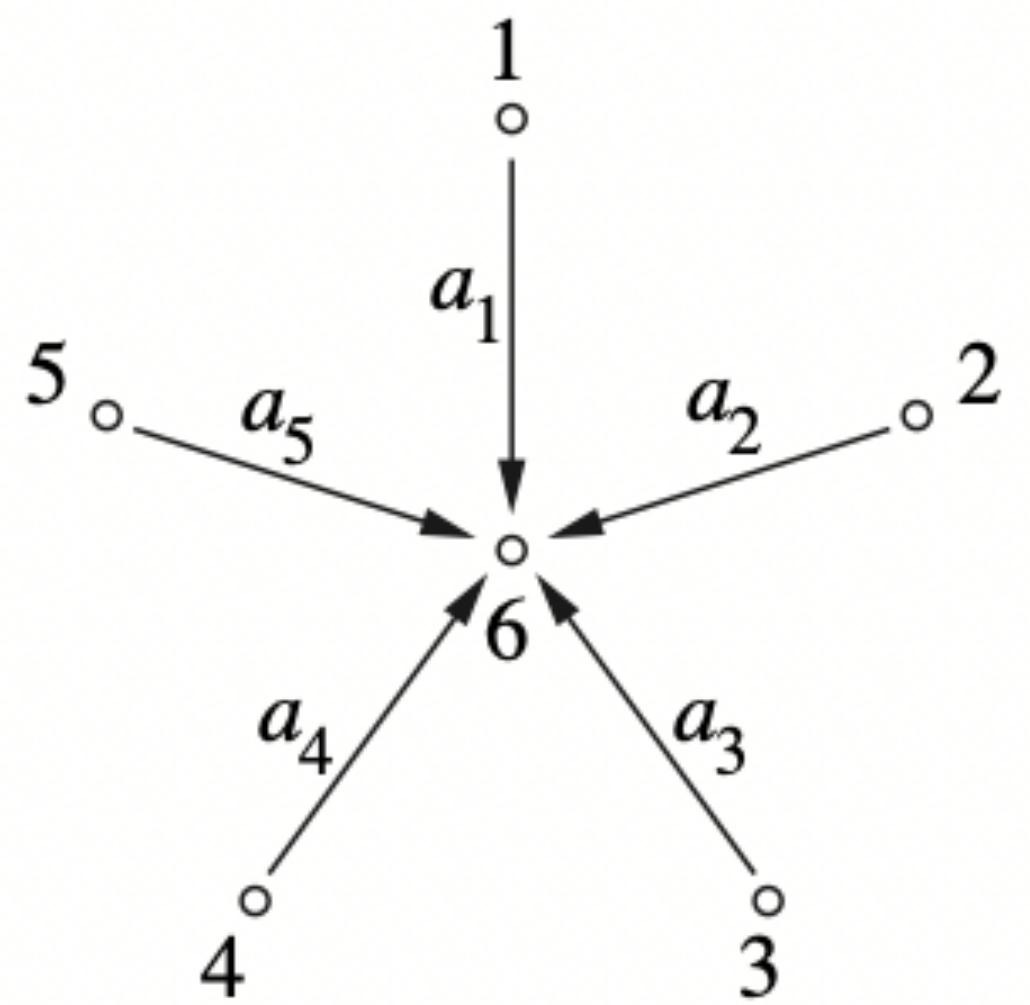
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$$\begin{array}{ccc} & & k \\ & \nearrow (x, y) \mapsto x + y & \\ k & \xrightarrow{x \mapsto (x, 0)} & k^2 \\ & \downarrow x \mapsto (0, x) & \\ & & k \end{array}$$

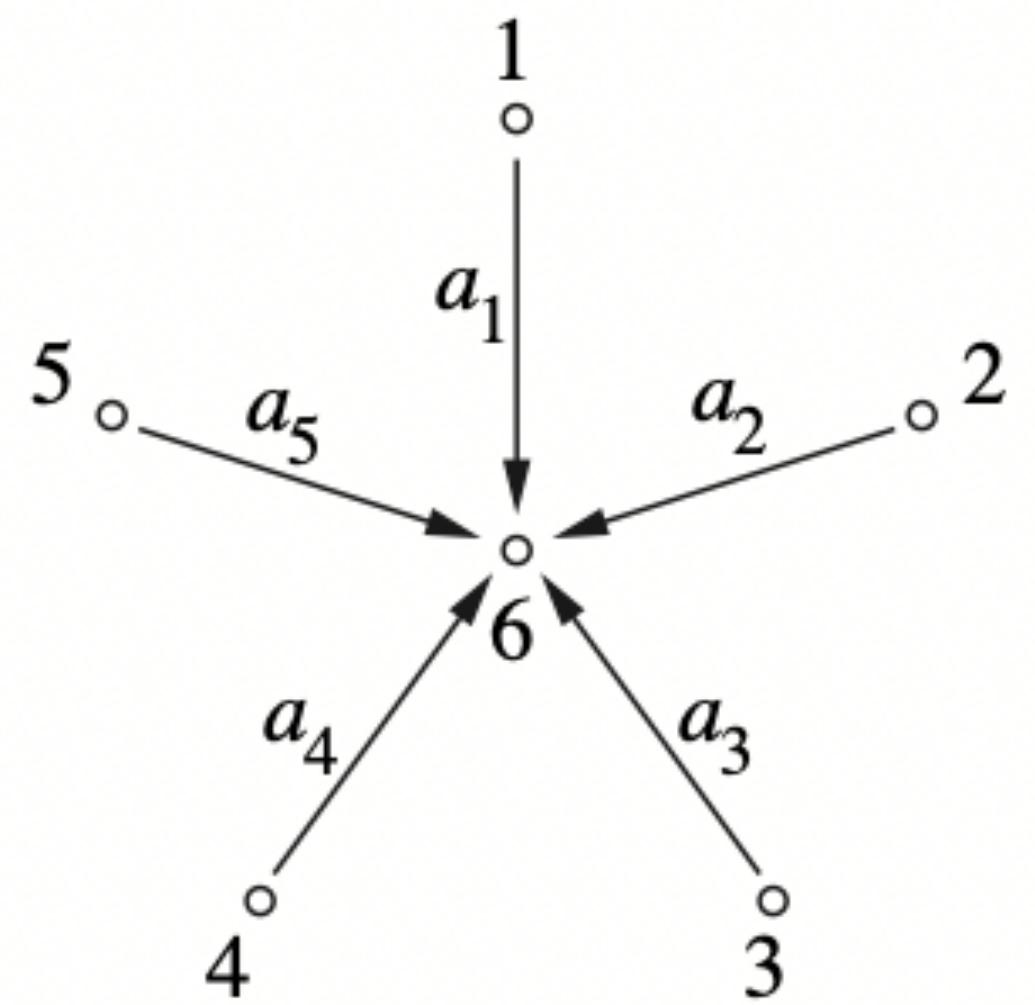
Quivers

- A quiver Q is a directed multigraph. That is, a quiver is a (finite) set of vertices Q_0 and a (finite) set of edges Q_1 with two functions $h, t : Q_1 \rightarrow Q_0$.



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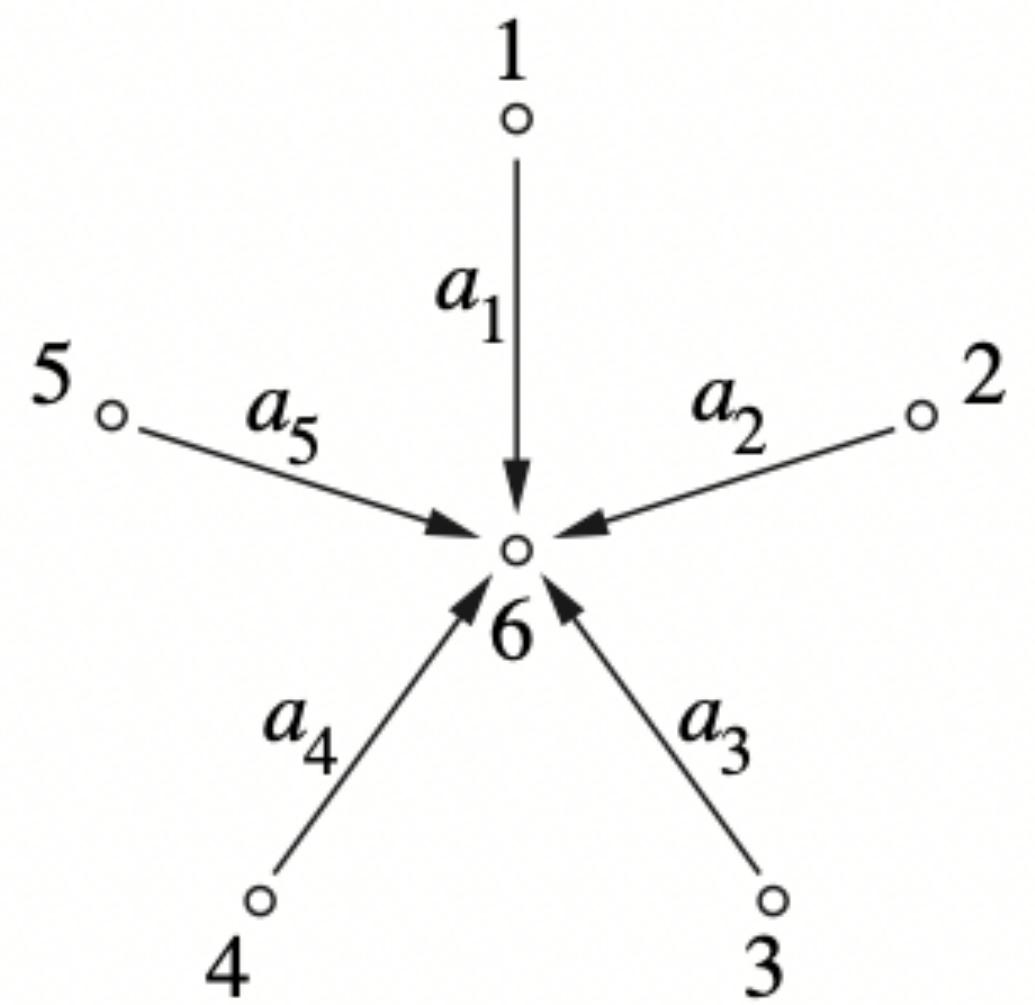
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- Let V and W be representations of Q . A morphism from V to W consists of linear maps $\varphi_x : V_x \rightarrow W_x$ for each $x \in Q_0$ such that the following diagram commutes for any $a \in Q_1$

$$\begin{array}{ccc} V_{t(a)} & \xrightarrow{V_a} & V_{h(a)} \\ \varphi_{t(a)} \downarrow & & \downarrow \varphi_{h(a)} \\ W_{t(a)} & \xrightarrow{W_a} & W_{h(a)} \end{array}$$



Quiver Representations

- A morphism $\varphi : V \rightarrow W$ between two quiver representations is an isomorphism if φ_x is an isomorphism for each $x \in Q_0$.

$$\begin{array}{ccc} k & \xrightarrow{x \mapsto (x, 0)} & k^2 \\ & \uparrow x \mapsto (0, x) & \\ & k & \end{array}$$

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A commutative diagram showing a representation of a quiver. The bottom row consists of two nodes labeled k . An arrow points from the left k to the right k , labeled $x \mapsto (x, 0)$. Above this, the top row also consists of two nodes labeled k^2 . An arrow points from the left k^2 to the right k^2 , labeled $(x, y) \mapsto x + y$. A vertical arrow connects the two k nodes, labeled $x \mapsto (0, x)$. The label k is placed below the rightmost k^2 node.

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- Any multiparameter persistence module can be viewed as a quiver representation so a classification of quiver representations would provide a way to decompose a persistence module.

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A commutative diagram showing a mapping from k to k^2 . The horizontal arrow is labeled $x \mapsto (x, 0)$. A vertical arrow from k^2 back to k is labeled $x \mapsto (0, x)$. A diagonal arrow from k to k^2 is labeled $(x, y) \mapsto x + y$. The label k is placed at the bottom right corner of the diagram area.

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- For some quivers, this is easy, for others, this is nearly impossible.

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Indecomposable Representations

- Let V and W be representations of a quiver Q . The direct sum of V and W is the quiver representation $V \oplus W$ defined by $(V \oplus W)_x = V_x \oplus W_x$ with linear maps defined componentwise.

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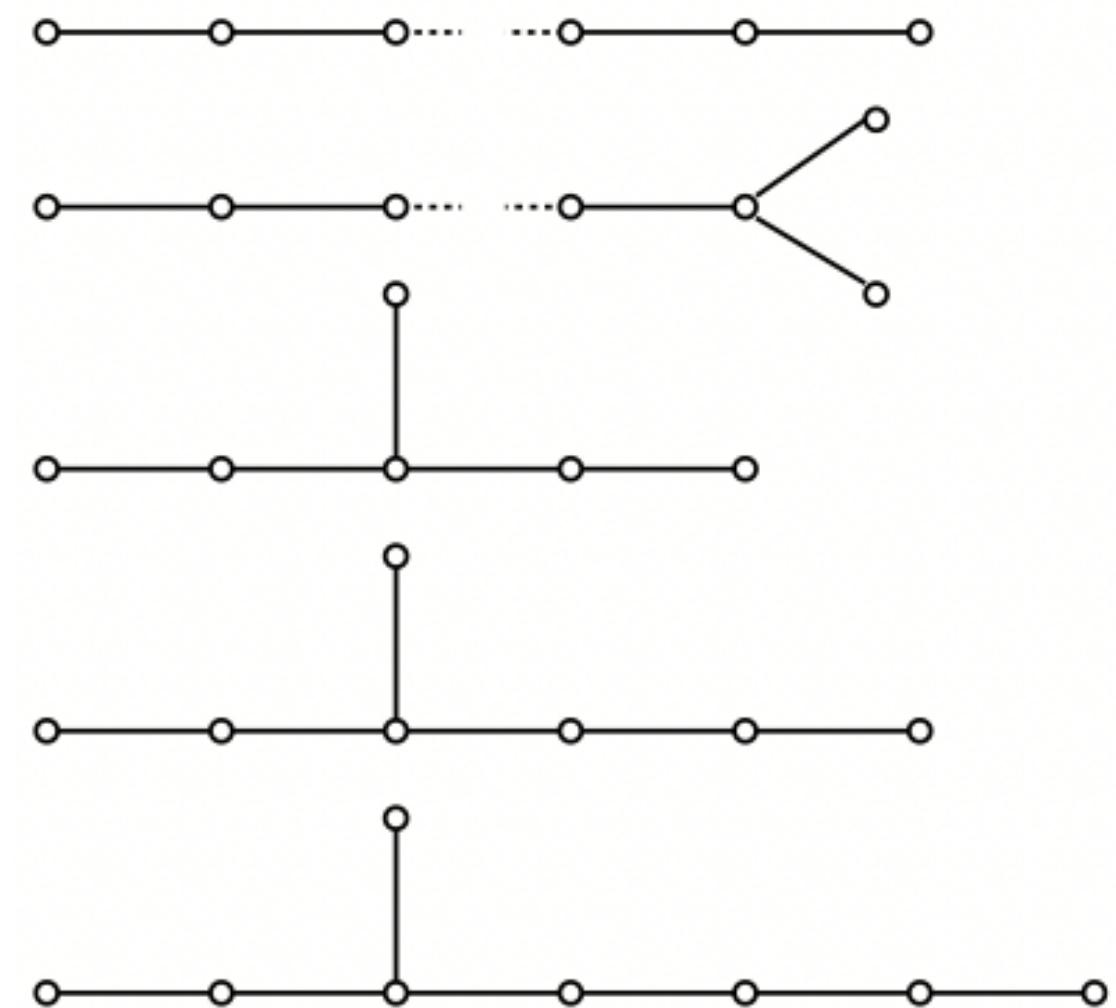
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- A nice classification of indecomposable quiver representations would provide an answer to multiparameter persistence.

Gabriel's Theorem

A quiver Q is of finite type if it has only finitely many indecomposable representations.

Theorem[Gabriel, 1972] A quiver Q is of finite type if and only if the underlying graph is a union of the following graphs:

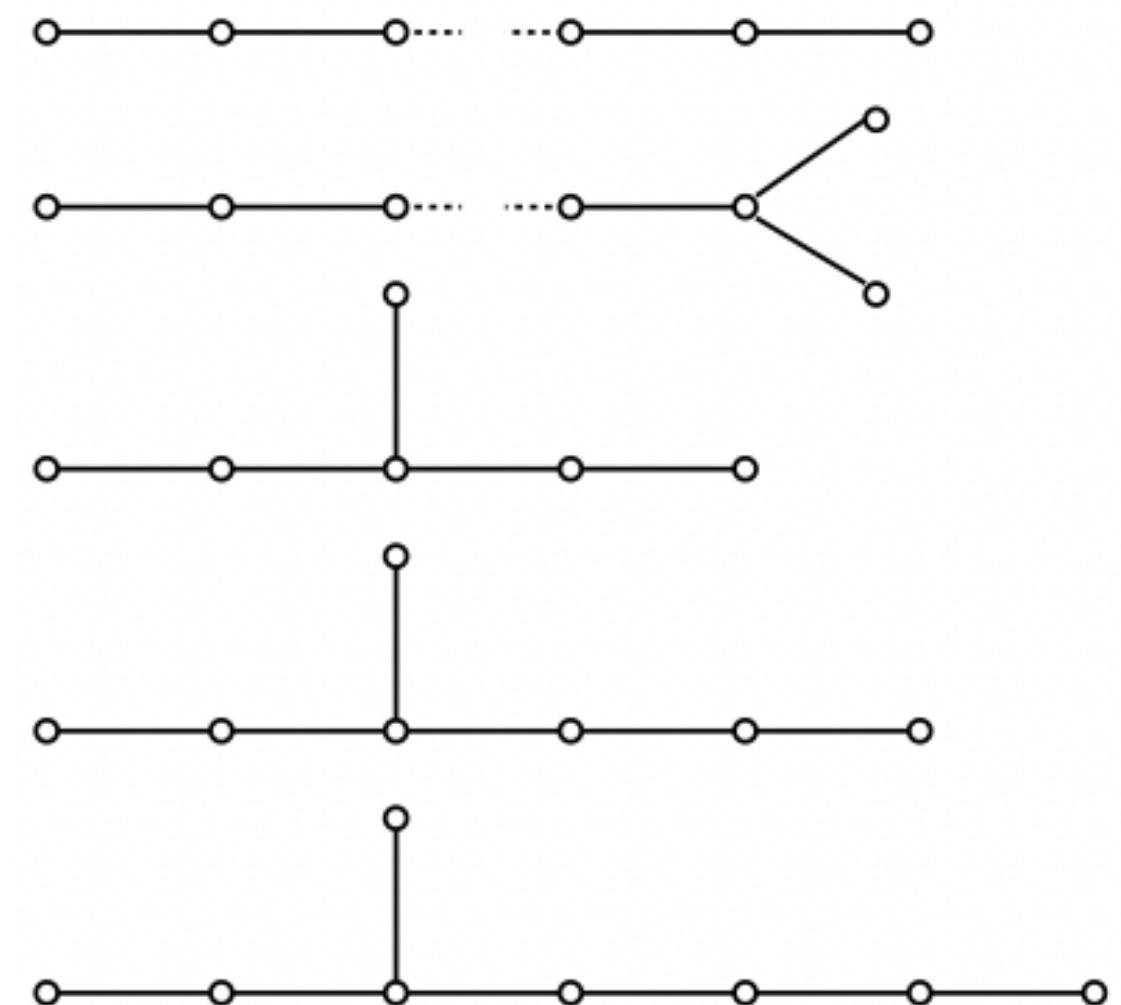


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Gabriel's Theorem implies that there are infinitely many indecomposables for most multiparameter persistence modules.



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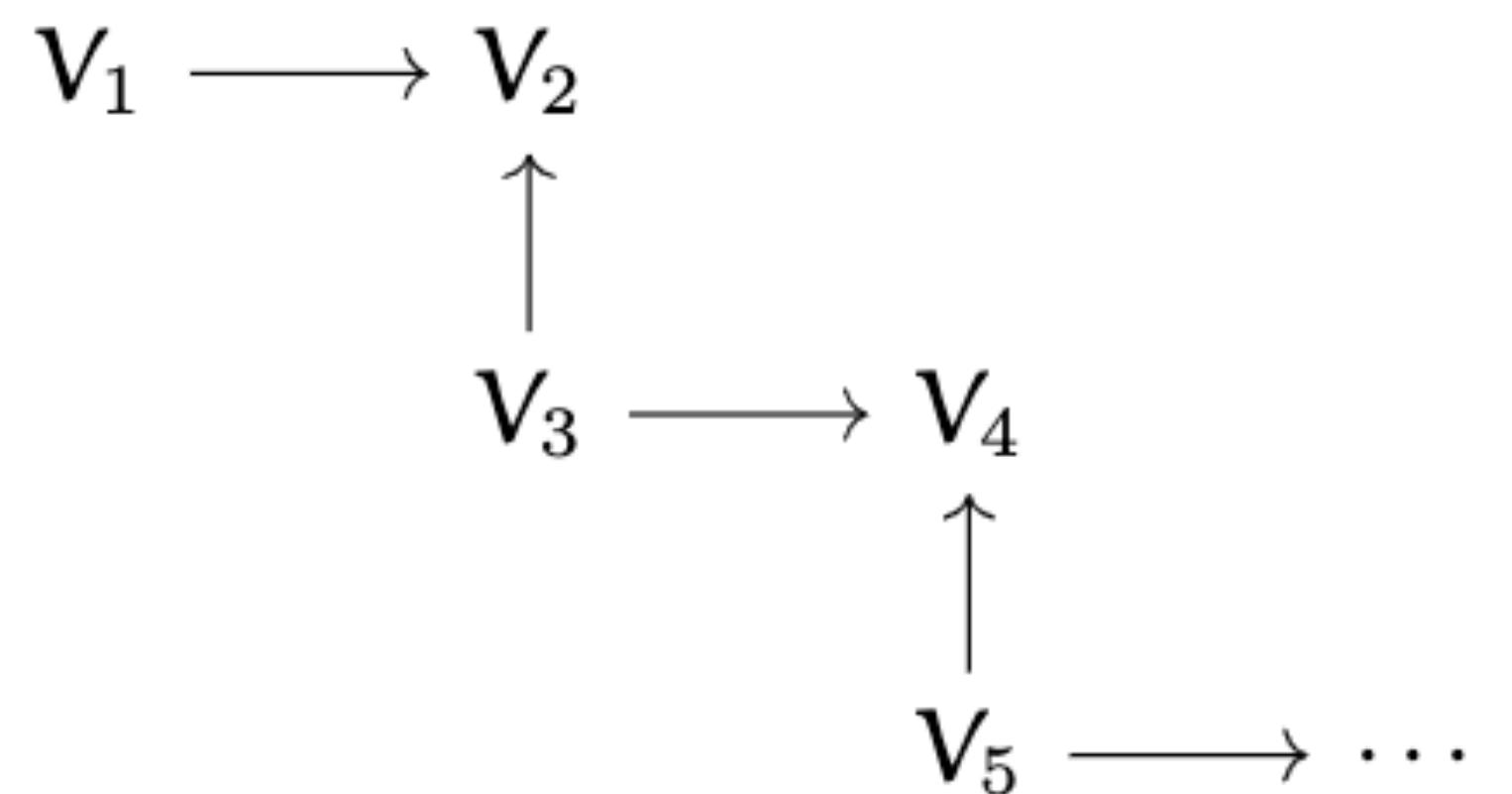
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- The indecomposables of a zigzag persistence module are of the form $k^{[b,d]}$ where $k_a^{[b,d]} \cong k$ if $a \in [b, d]$ and $k^{[b,d]} \cong 0$ otherwise. The maps between adjacent elements in $[b, d]$ are identity maps and all other maps are 0.

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- This means we can decompose any zigzag persistence module M as

$$\bigoplus_{I \in \mathcal{I}} k^I$$

where \mathcal{I} is some multiset of intervals.



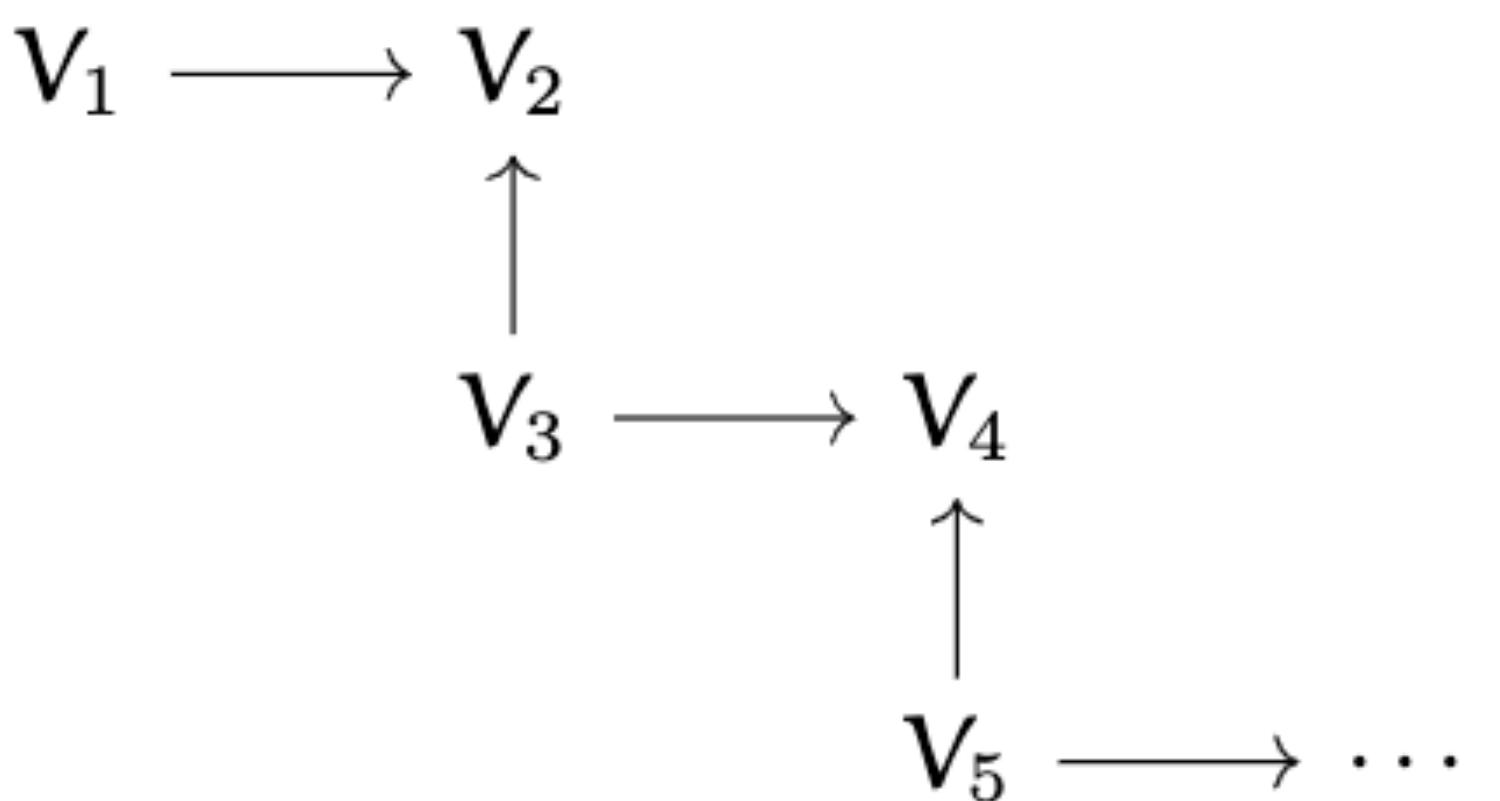
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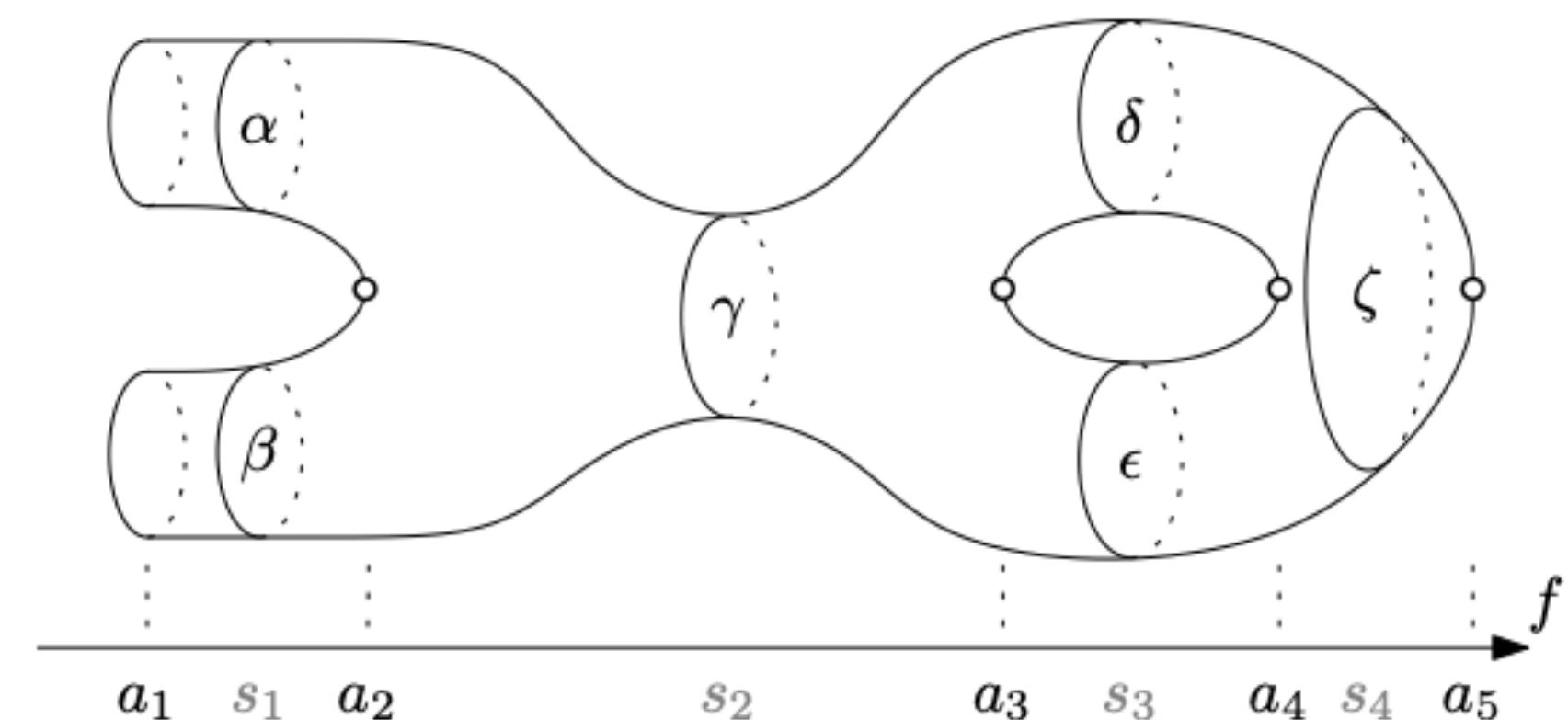
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- We call such a decomposition a barcode of M



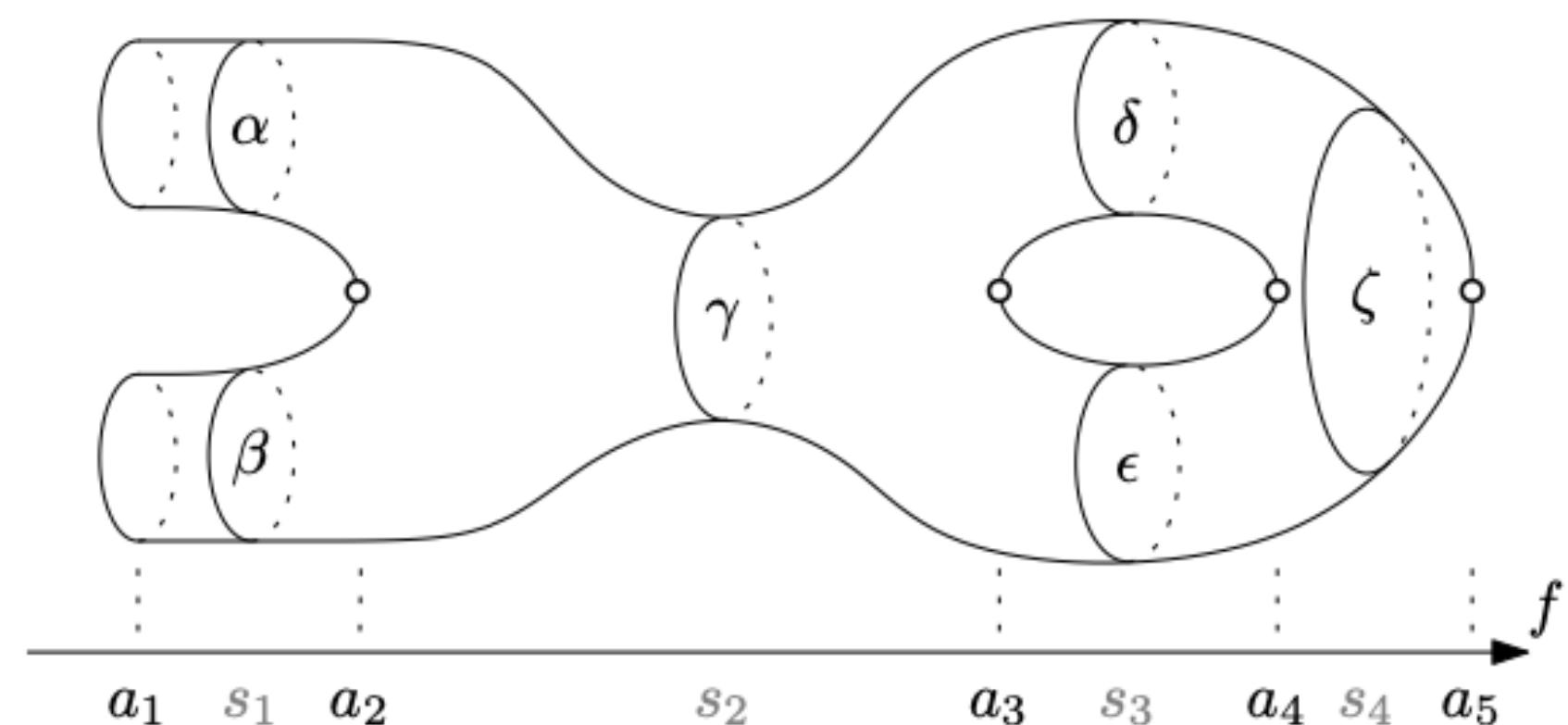
Level Set Persistence

- Let $f : X \rightarrow \mathbb{R}$ be a “nice” continuous map from a topological space X to \mathbb{R} . The level set of f at $t \in \mathbb{R}$ is $f^{-1}(t) \subseteq X$.



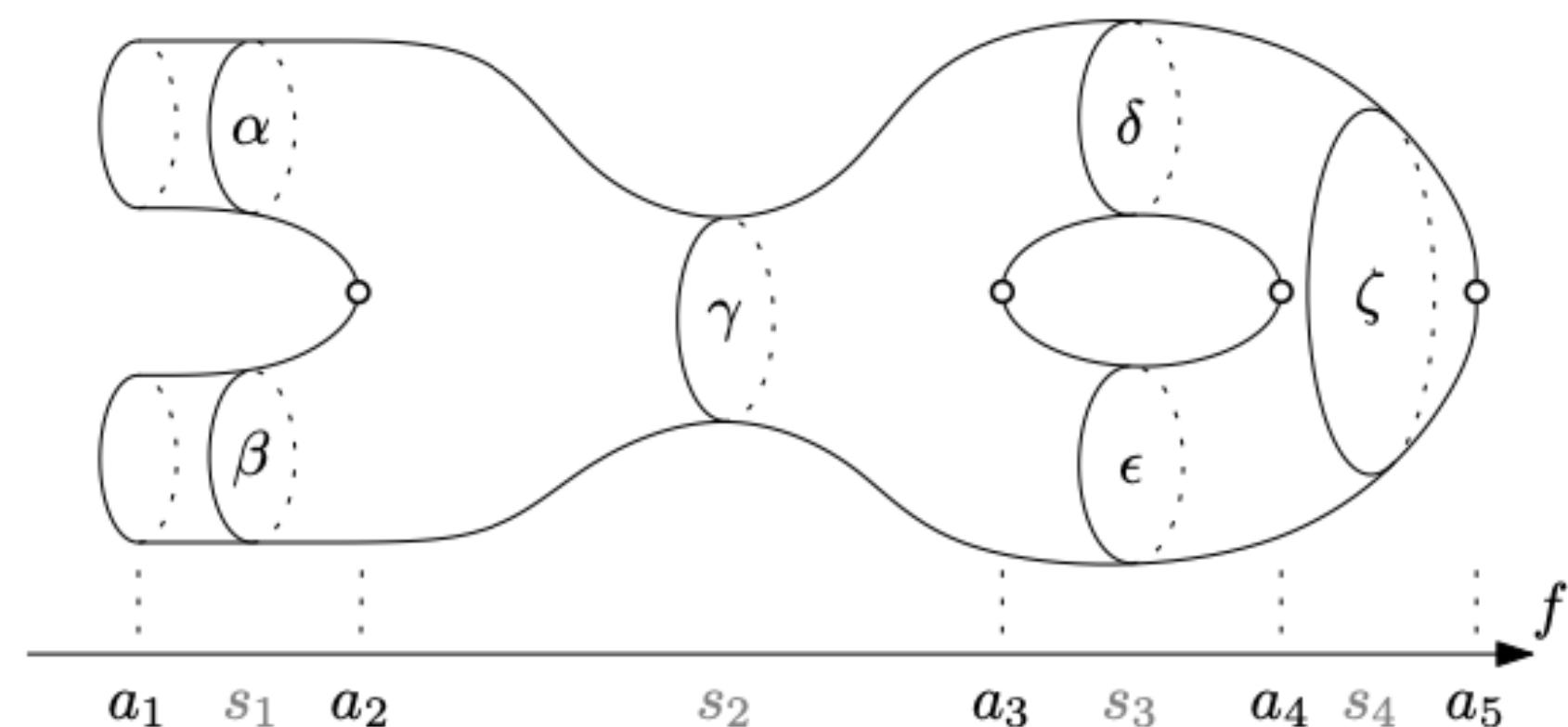
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- As t varies, there are only finitely many places where $f^{-1}(t)$ changes homotopy type. We call these critical points.



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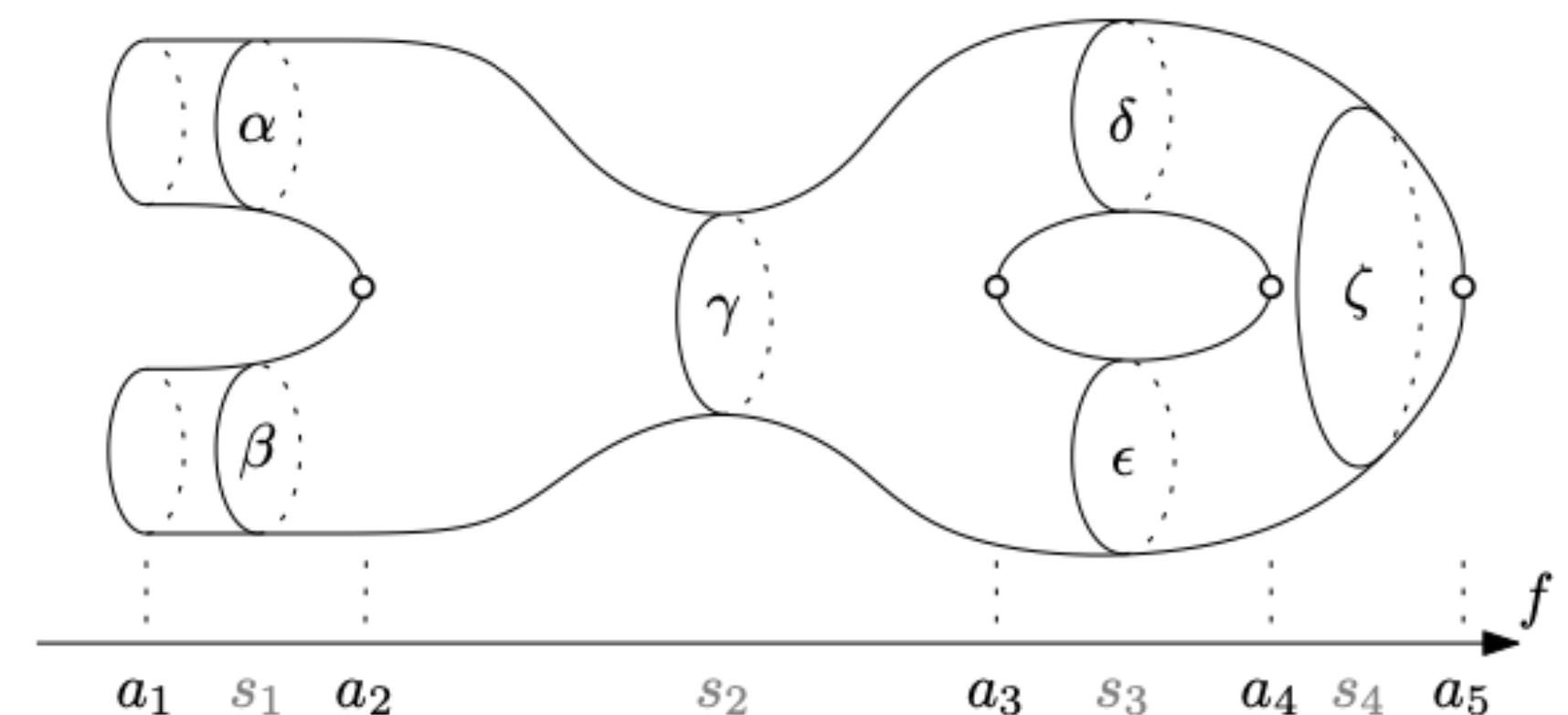
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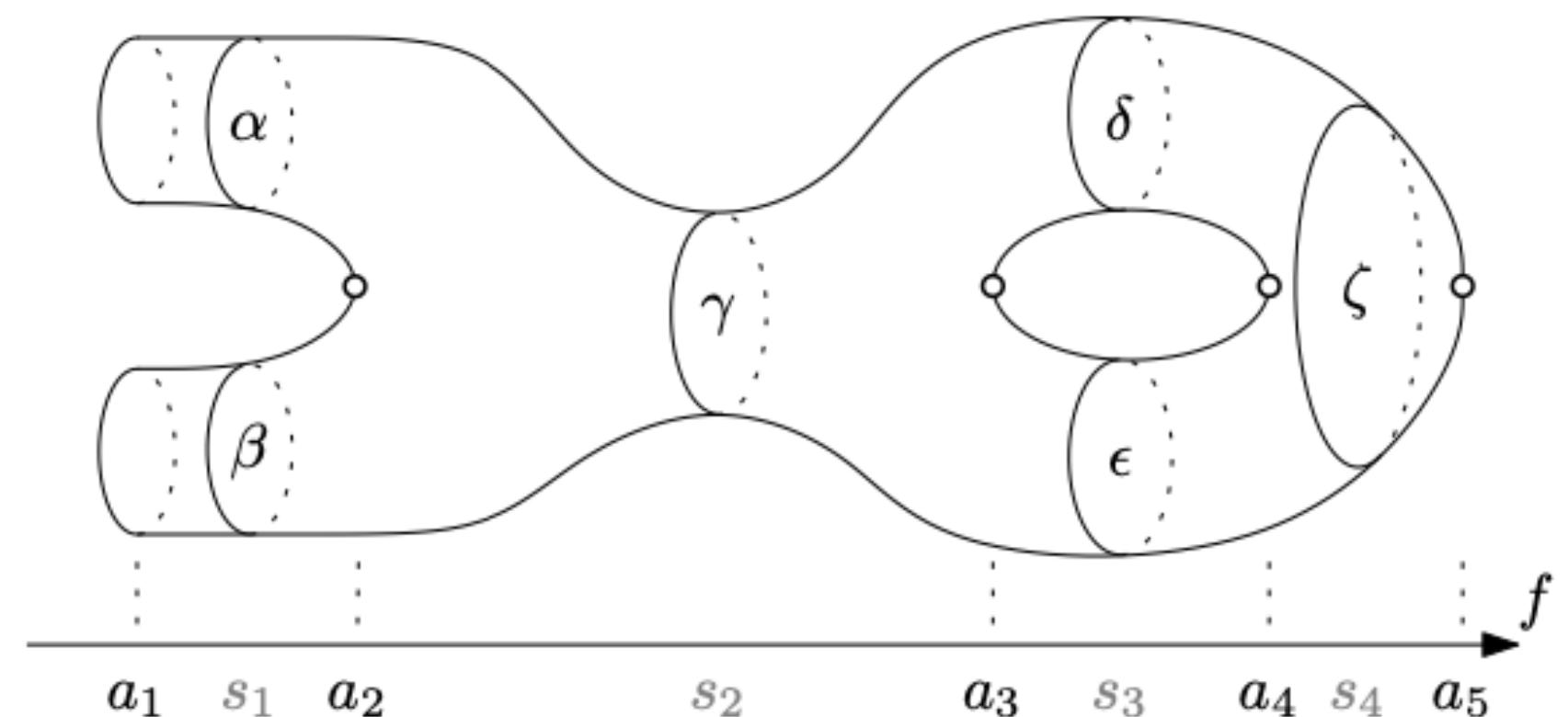


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- Applying homology to this sequence of topological spaces induces a zigzag persistence module

$$H_i(f^{-1}(s_0)) \longrightarrow H_i(f^{-1}([s_0, s_1])) \longleftarrow H_i(f^{-1}(s_1)) \longrightarrow H_i(f^{-1}([s_1, s_2])) \longleftarrow \dots$$



Invariants

The Rank Function

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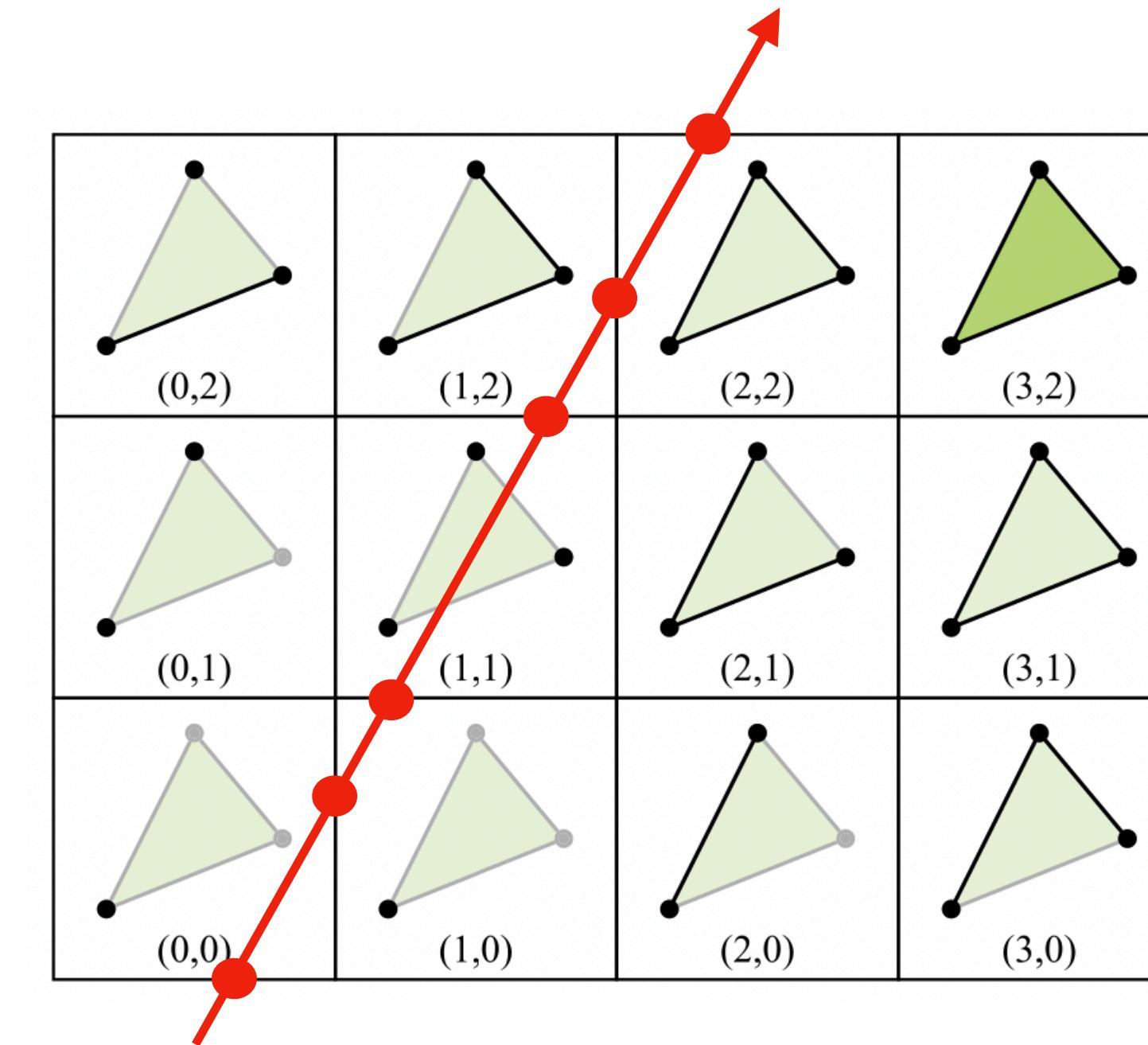
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- In the 1-d setting, the rank function is equivalent to the persistence diagram.

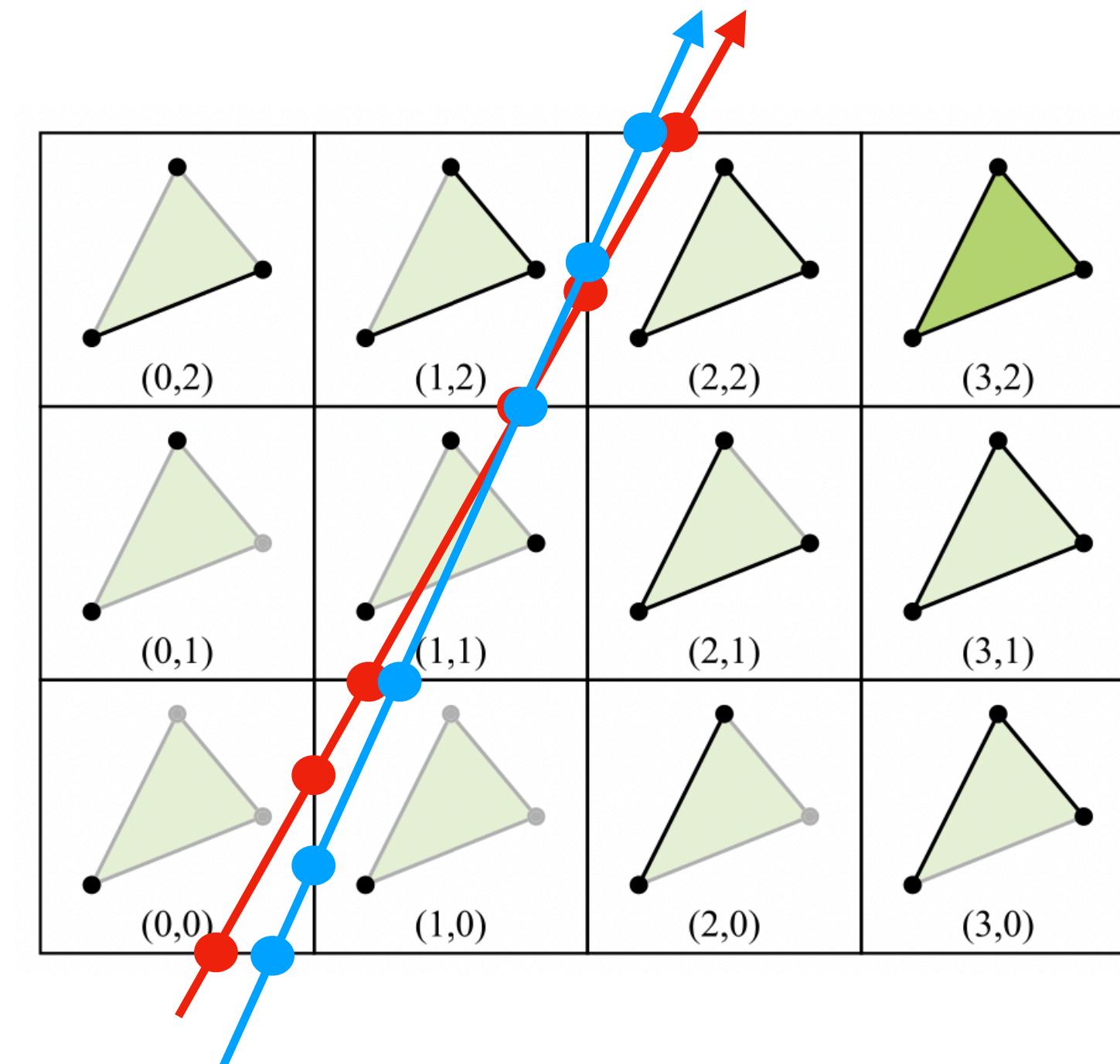
Slicing Multiparameter Persistence Modules

- Given a multiparameter persistence module, we can slice it to obtain 1-parameter persistence modules.
- This allows us to use the ordinary persistent homology algorithm to obtain a barcode for each slice.
- The collection of all these barcodes is known as the fibered barcode.
- This idea was developed in [4]. They prove that the fibered barcode is equivalent to the rank invariant.



Slicing Continued

- The persistence diagrams for each slice do not need to be computed independently.
- Nearby lines will have similar barcodes.
- The vineyard algorithm can be used to compute the barcode of a nearby line from the barcode of a fixed line in linear time.
- The fibered barcode is a stable invariant [4].

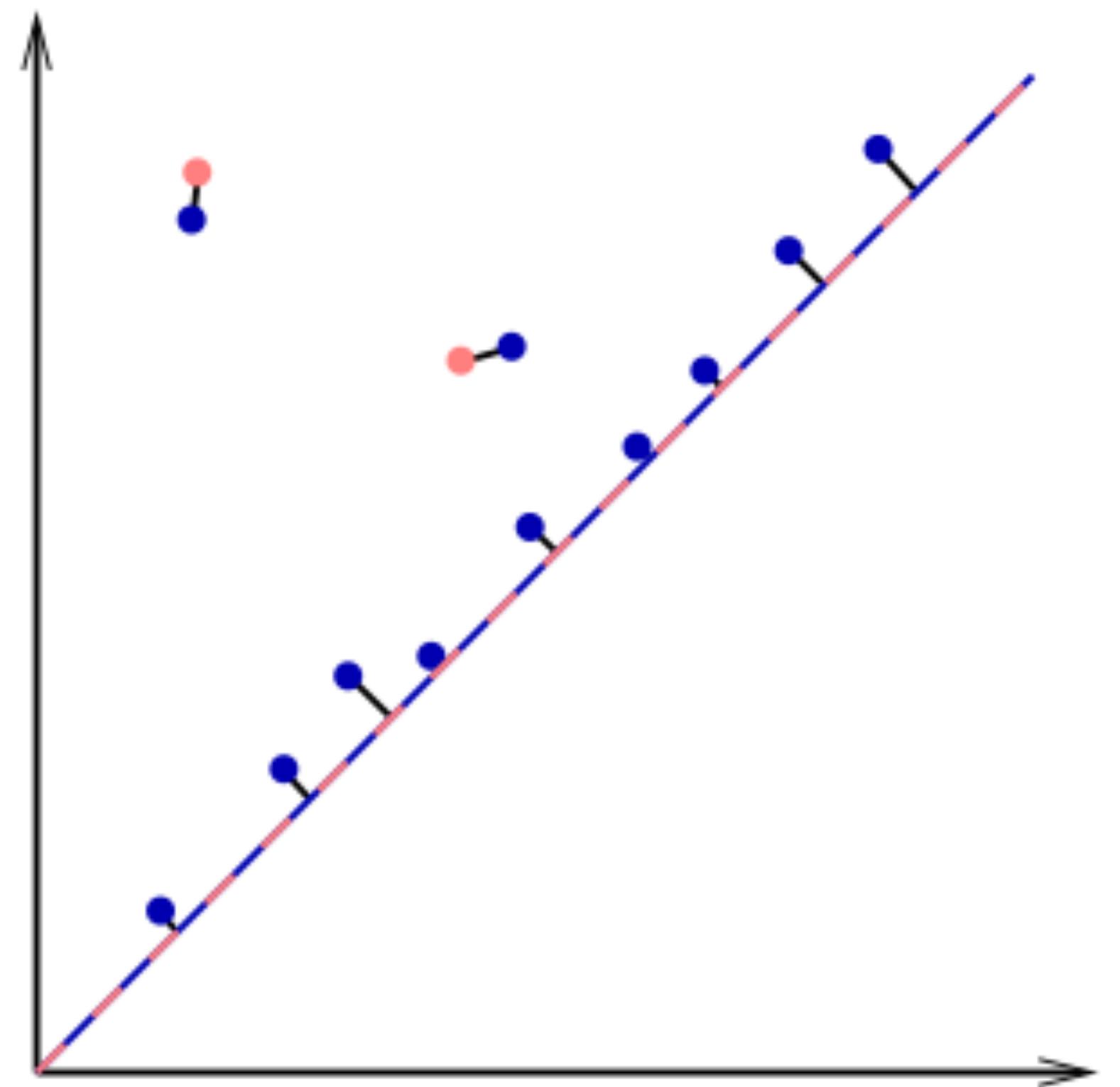


The Bottleneck Distance

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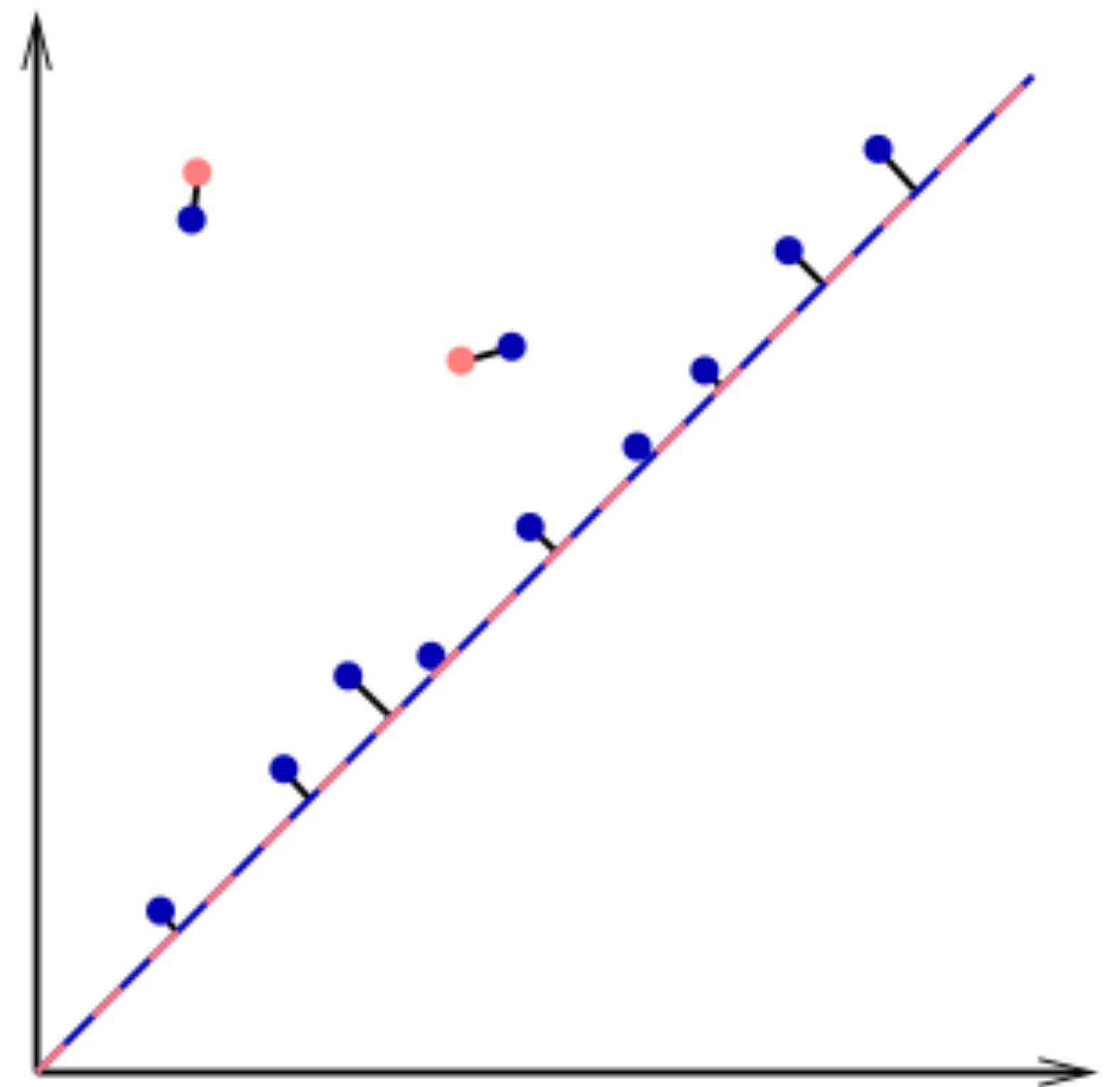
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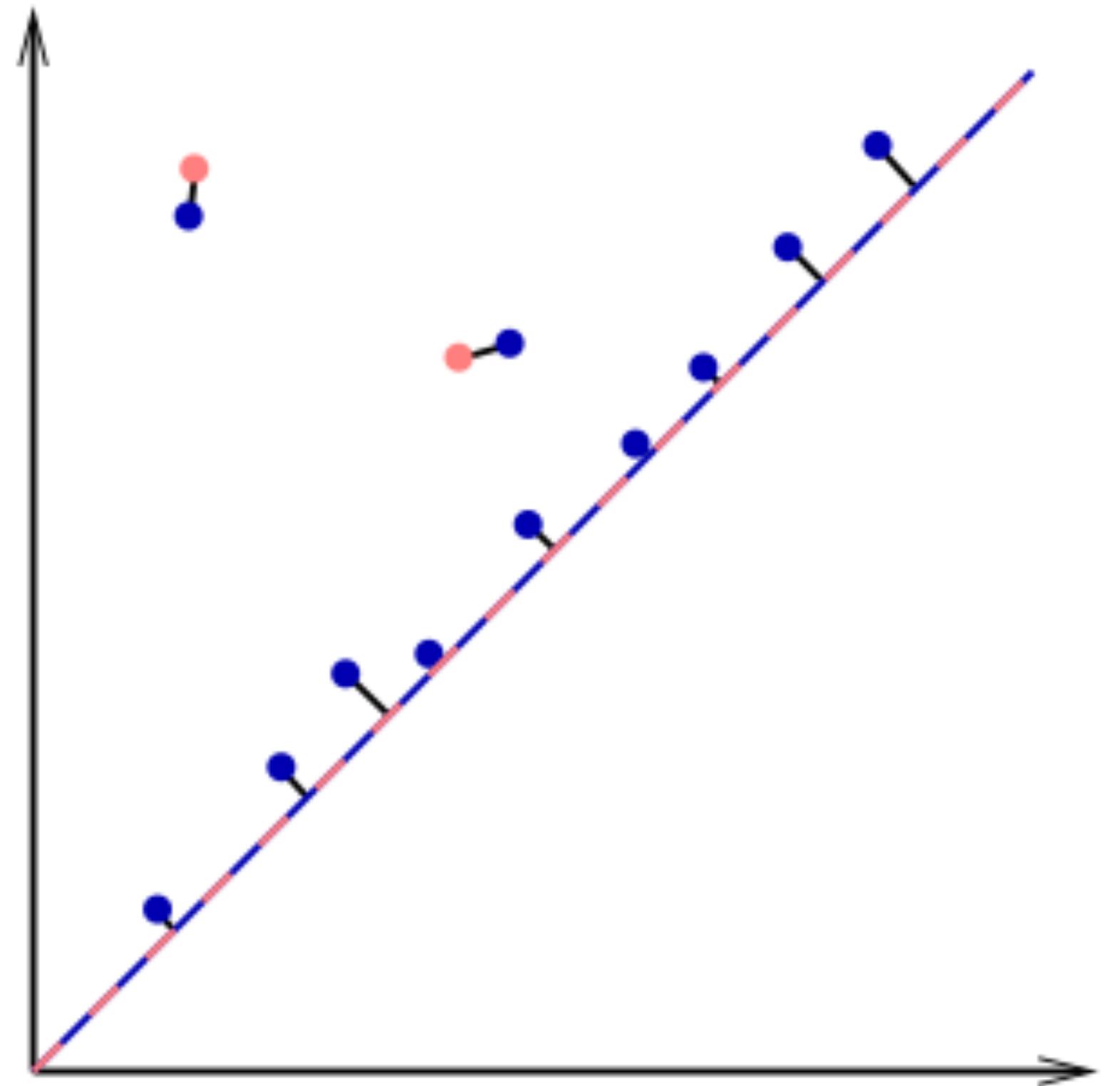


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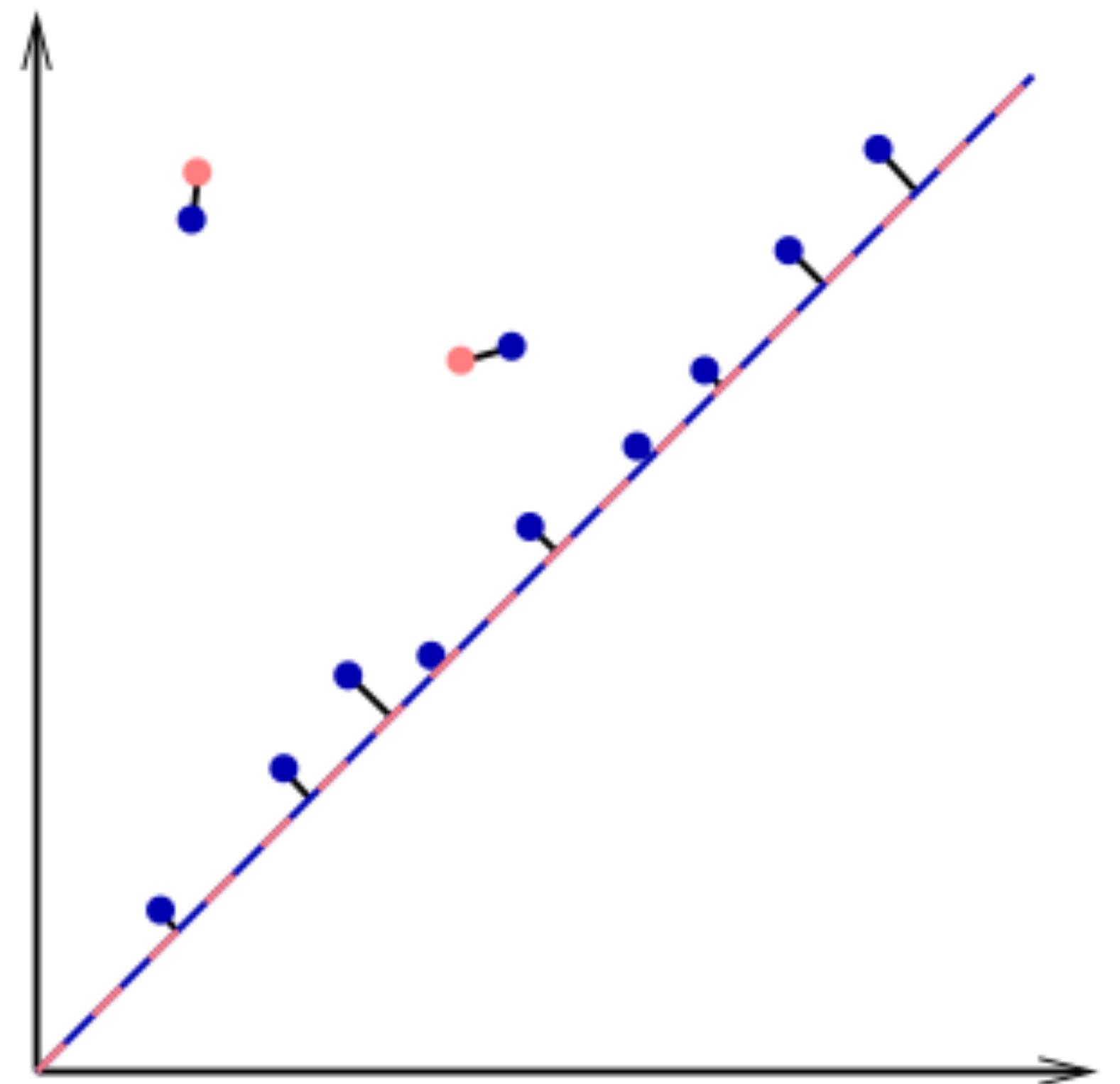
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- The bottleneck distance extends to a distance between fibered barcodes by maximizing the bottleneck distance between each sliced barcode.



The Matching Distance

- Let $M, N : \mathbb{R}^2 \rightarrow \text{Vec}_k$ be persistence modules and \mathcal{L} be the set of all lines in \mathbb{R}^2 with positive slope.

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- The matching distance between the fibered barcodes of M and N is

$$\sup_{\ell} \omega_\ell d_B(BM_\ell, BN_\ell)$$

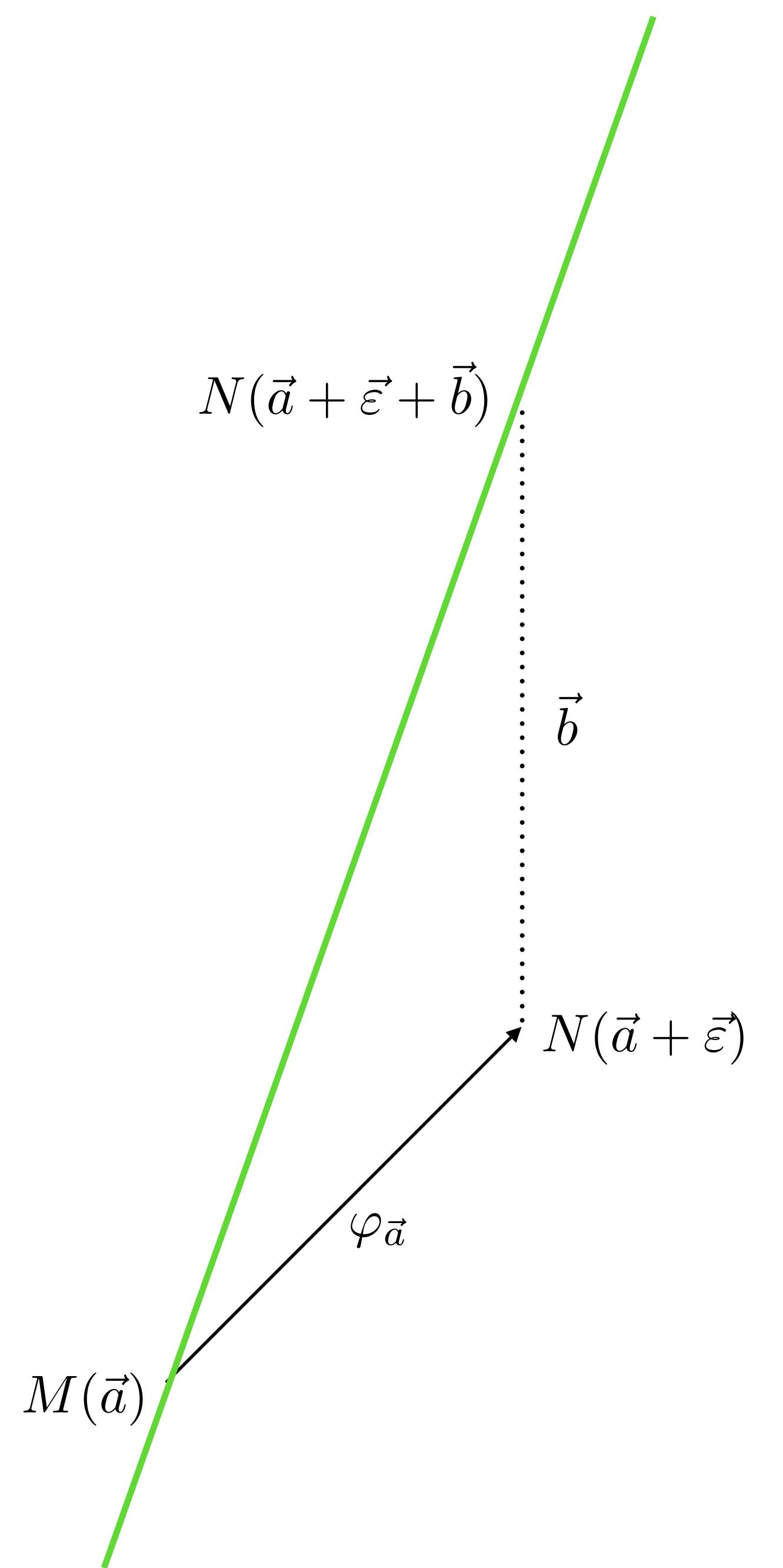
where $\omega_\ell = \frac{1}{\sqrt{1+m^2}}$ when $\ell = mx + b$

Stability of the Fibered Barcode

- Suppose $M, N : \mathbb{R}^2 \rightarrow \text{Vec}_k$ are ε -interleaved.

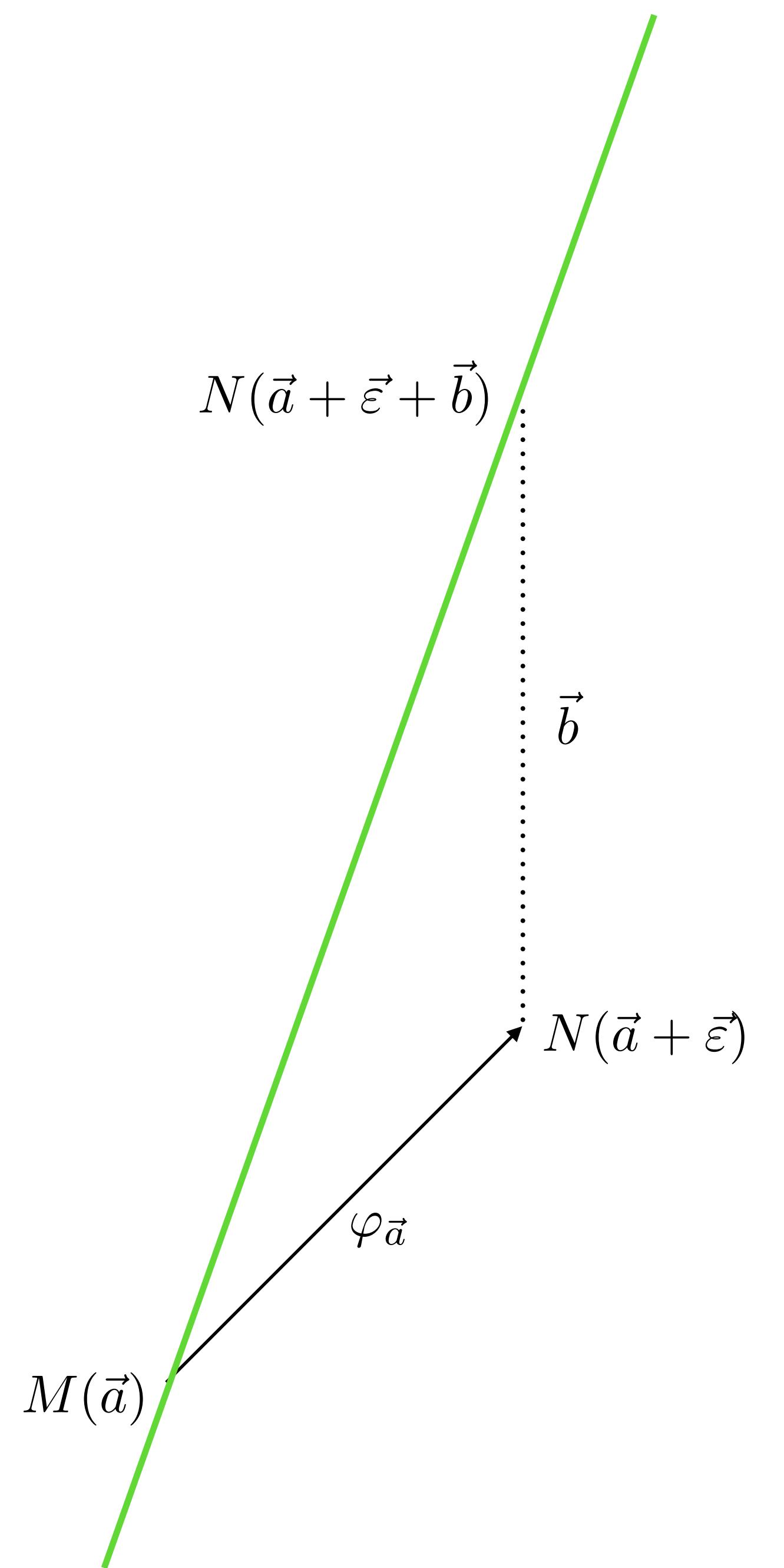
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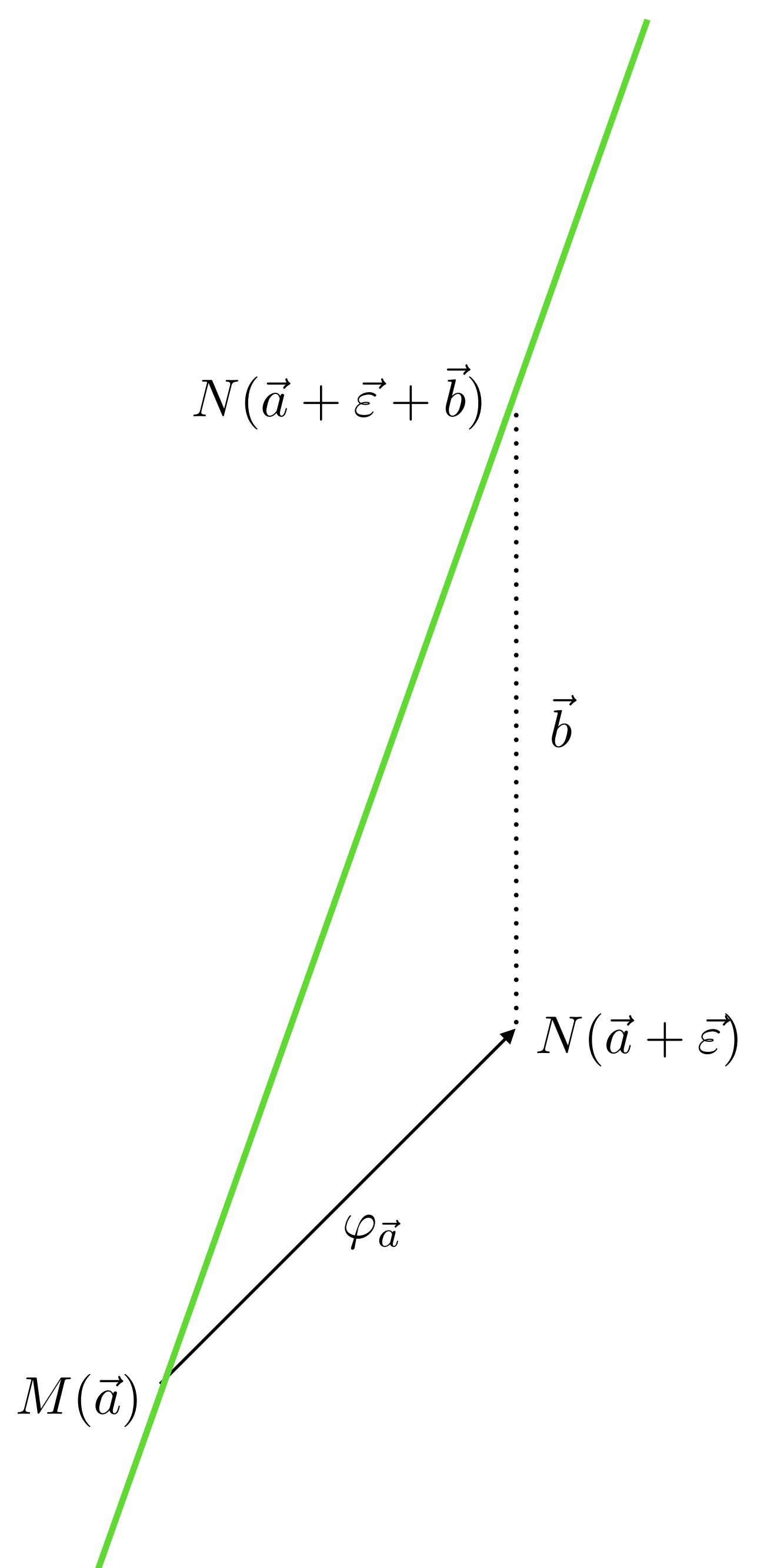
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- It follows that the matching distance between the fibered barcodes of M and N is at most ε .



Algorithms

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- Essentially, RIVET allows you to input a multiparameter persistence module, select a line, and it will compute the persistence diagram of the module restricted to that line.
- Unfortunately, RIVET is too slow for use with large data sets, running in time $O(n^8)$ [6].

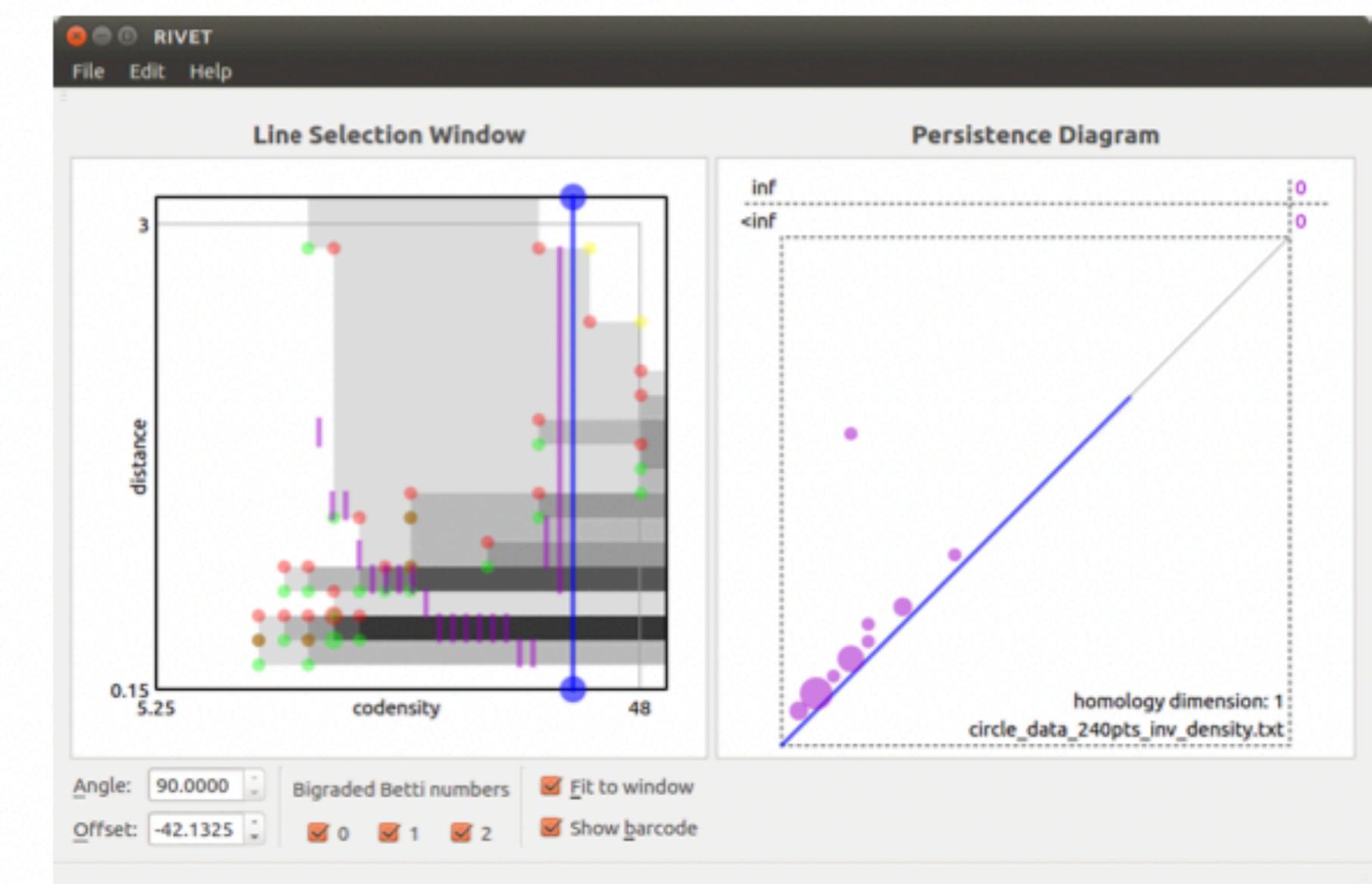
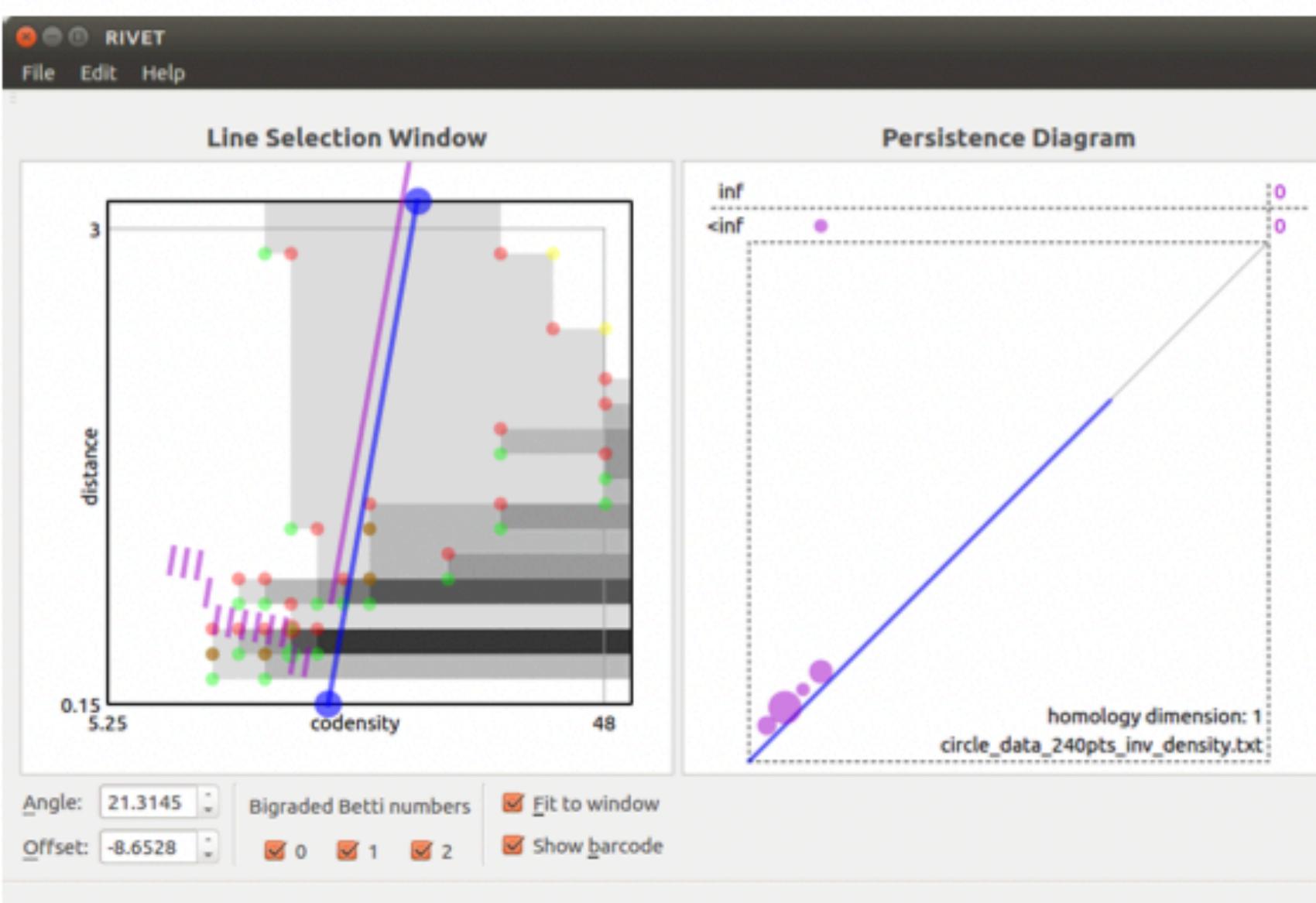
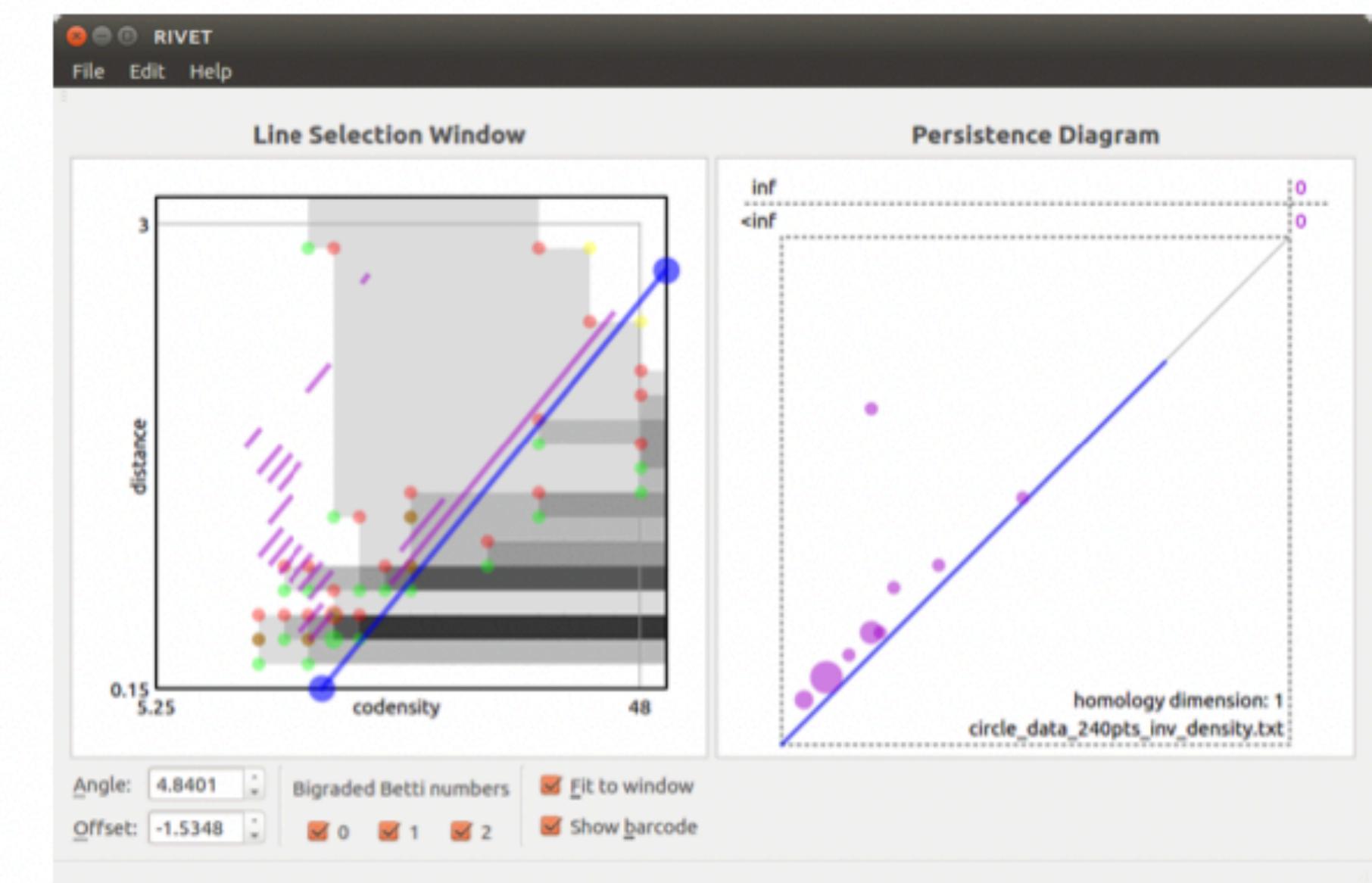
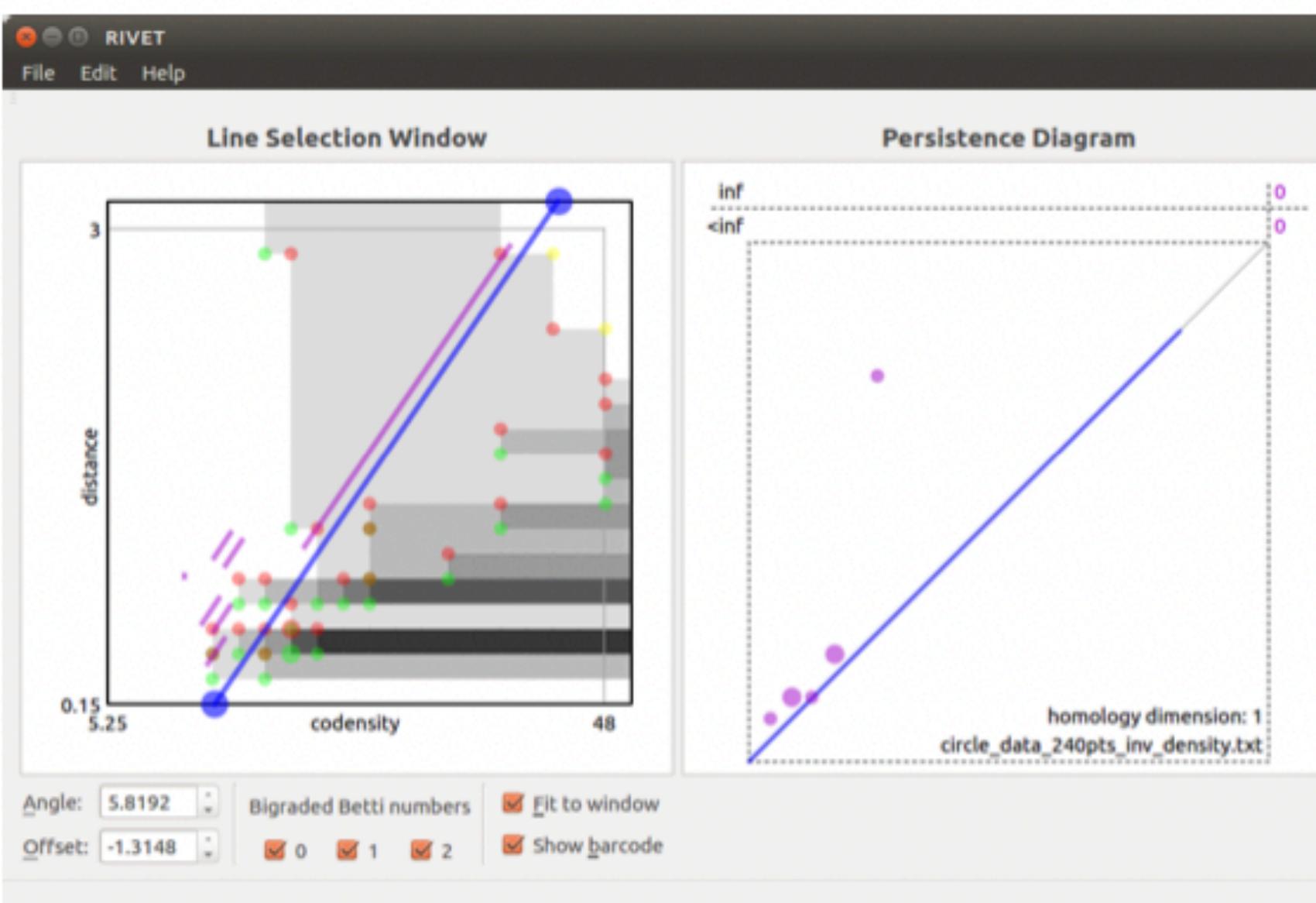


Figure 3: Screenshots of RIVET for a single choice of 2-D persistence module M and four different lines L . RIVET provides visualizations of the dimension of each vector space in M (greyscale shading); the 0th, 1st, and 2nd bi-graded Betti numbers of M (green, red, and yellow dots); and the barcodes of the 1-D slices M^L , for each L (in purple).

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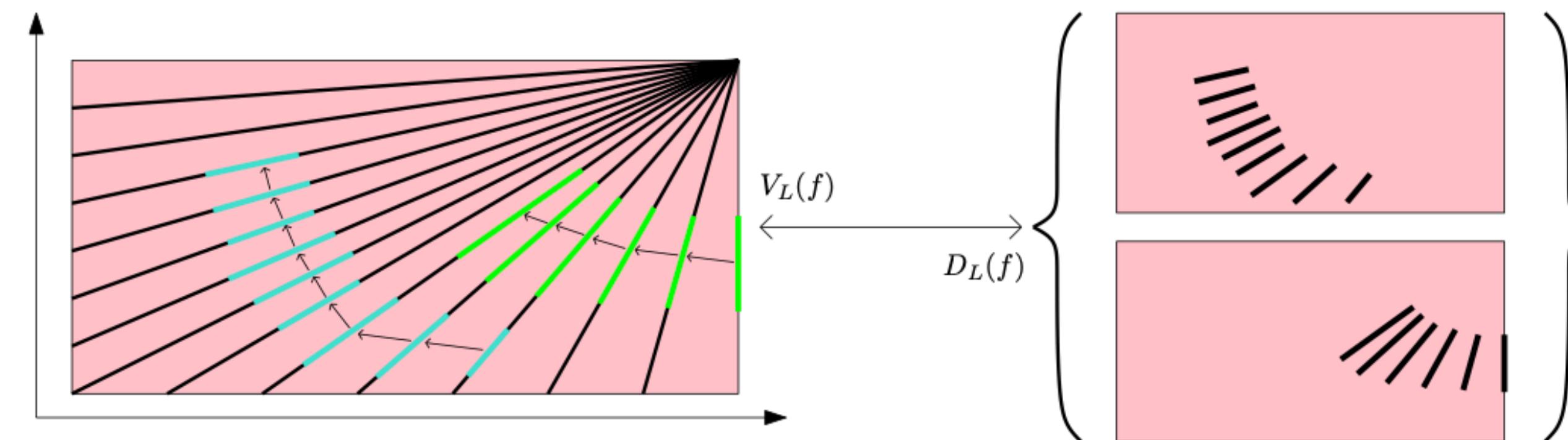
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- Elder rule staircodes are able to determine whether or not certain persistence modules are interval decomposable.

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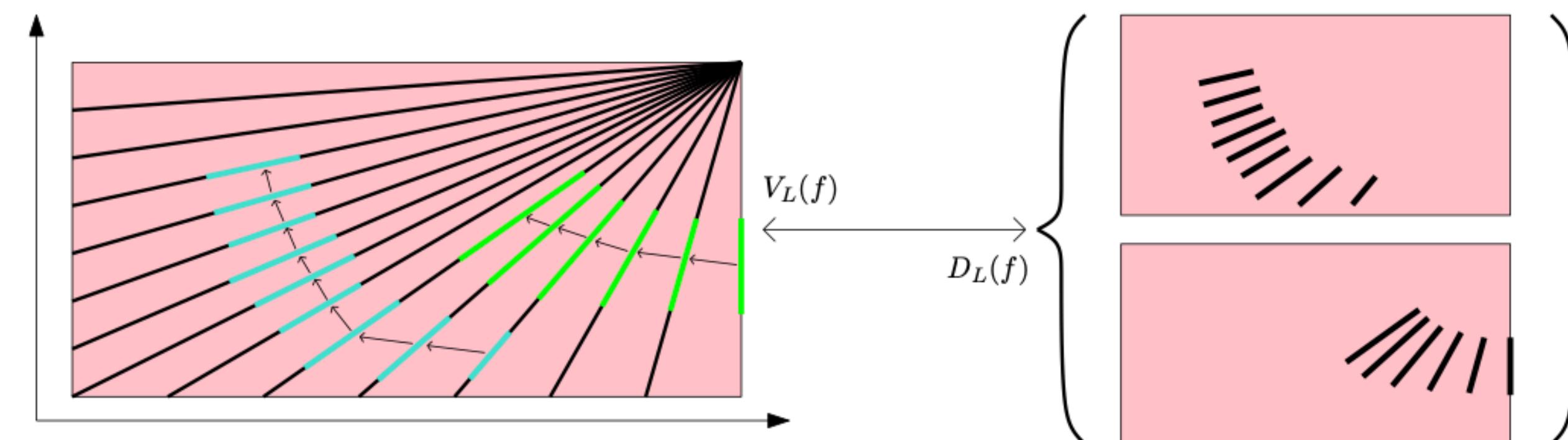
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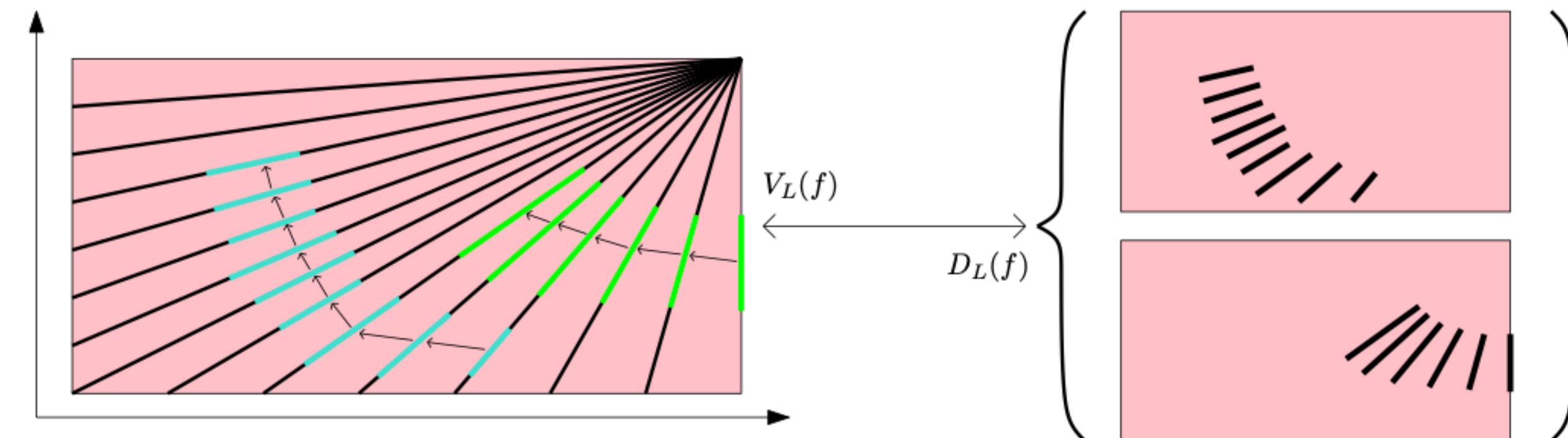
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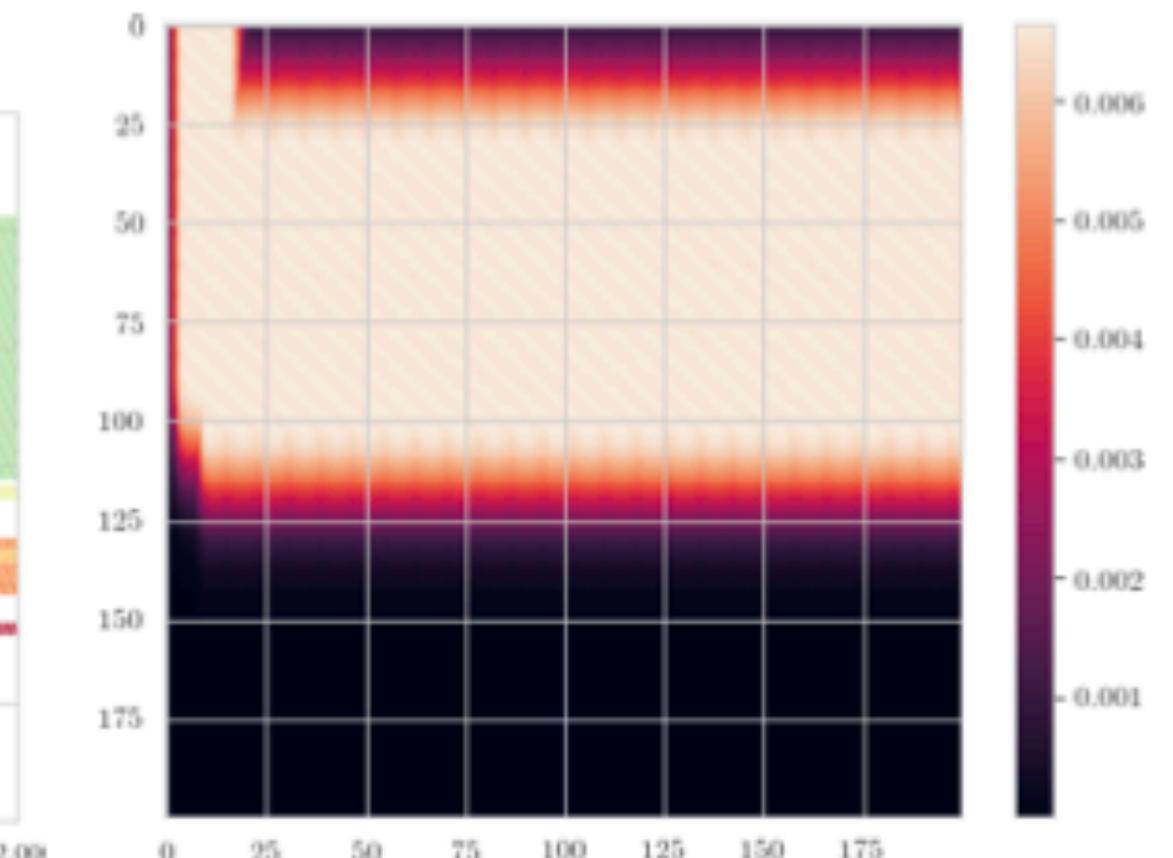
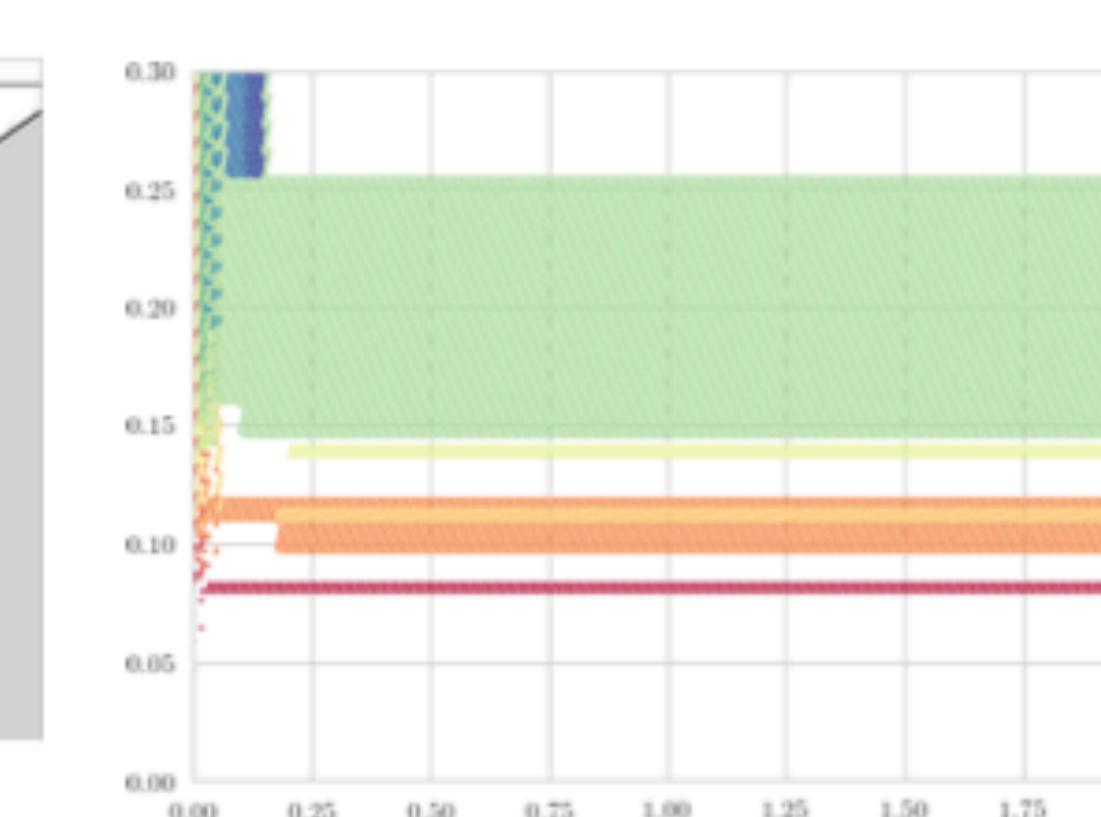
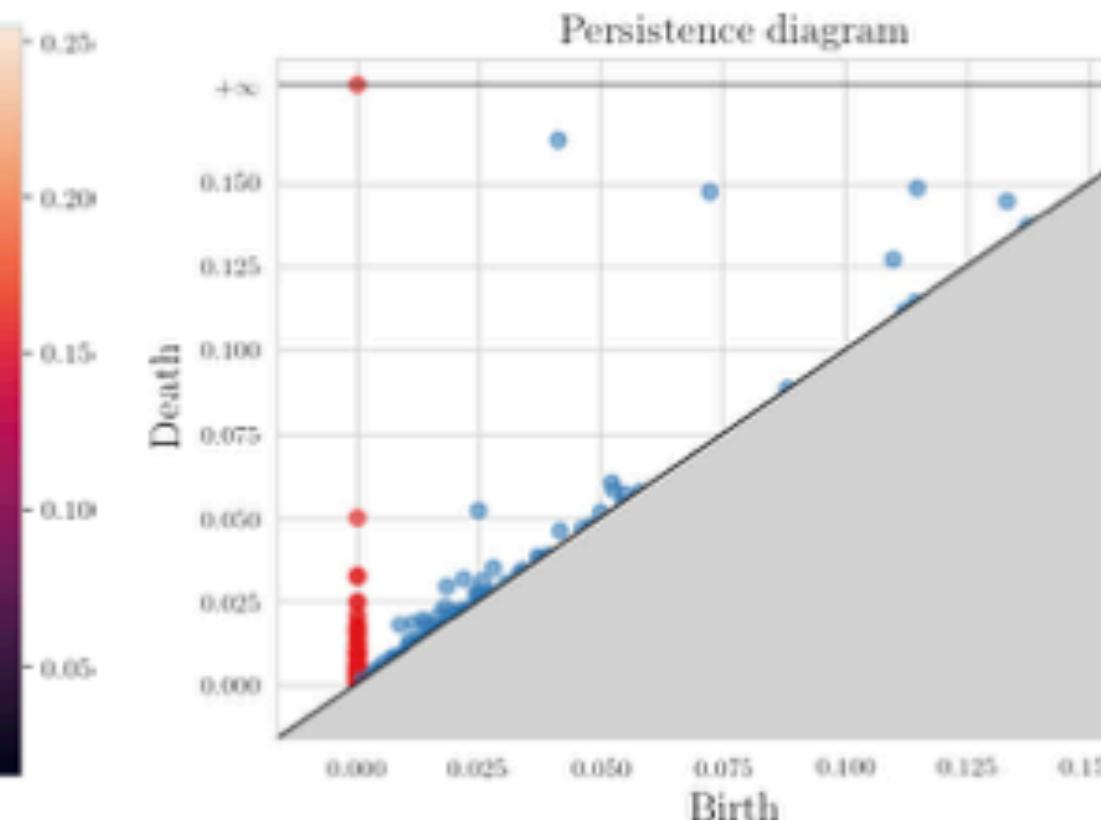
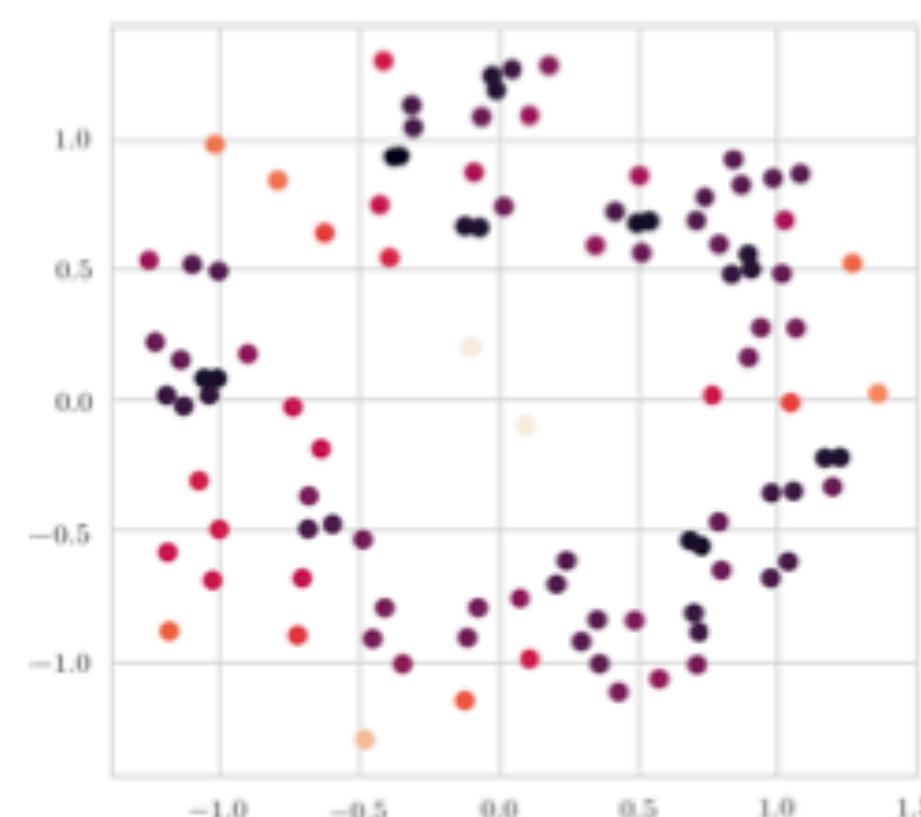
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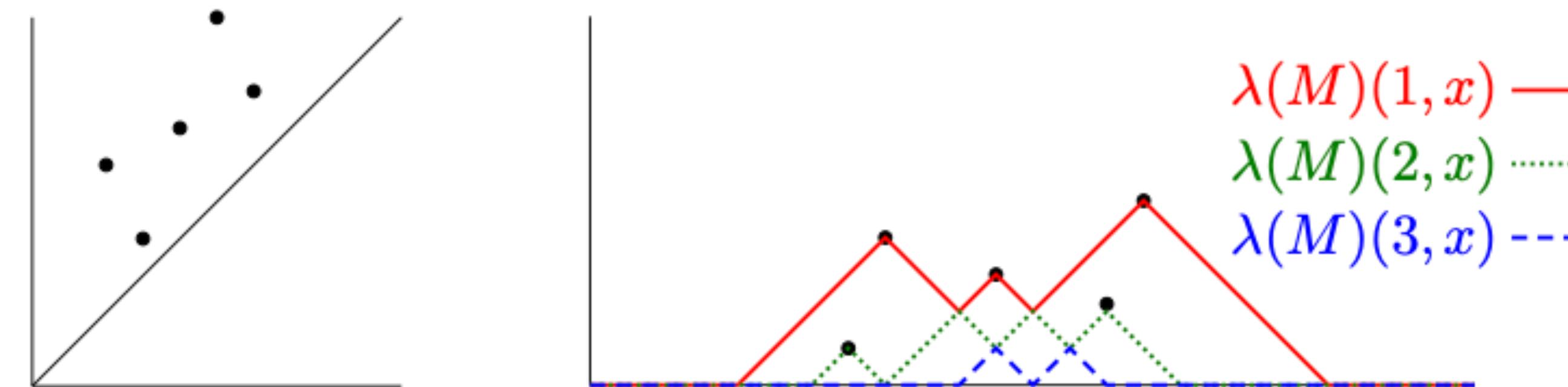
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Statistics

- Even in the 1-parameter setting, barcodes do not necessarily have well defined means.
- Persistence landscapes and persistence images fix this issue, producing vectors in euclidean space where an arsenal of statistics techniques can be applied.
- Multiparameter persistence landscapes in particular obey a strong law of large numbers and a central limit theorem [7].

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- Multiparameter persistence landscapes are built off of the fibered barcode and thus require a program such as RIVET first.
- Multiparameter persistence images are generally faster to compute than fibered barcodes, with a complexity depending on the resolution one wishes to calculate with.

Summary

- Multiparameter persistence is hard. There's no clear answer unlike the 1-parameter case.
- Even though we don't have a complete invariant, we can still compute invariants that lose information.
- The most widely used invariant for multiparameter persistence is the fibered barcode (AKA the rank invariant).
- The fibered barcode is stable to noise.
- Persistence images and landscapes are useful for machine learning applications and statistics.
- The elder rule staircode is the best bet for 0-dimensional homology with augmented metric spaces.

Coming Weeks

- Next week Nate Clause will go more in depth into the algorithms for multiparameter persistence including zigzag persistence and a demo of RIVET.
- The following week Brantley Vose will cover “The Importance of Forgetting” and show how zigzag persistence appears in neuroscience.

References

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