

View By:  Groups:

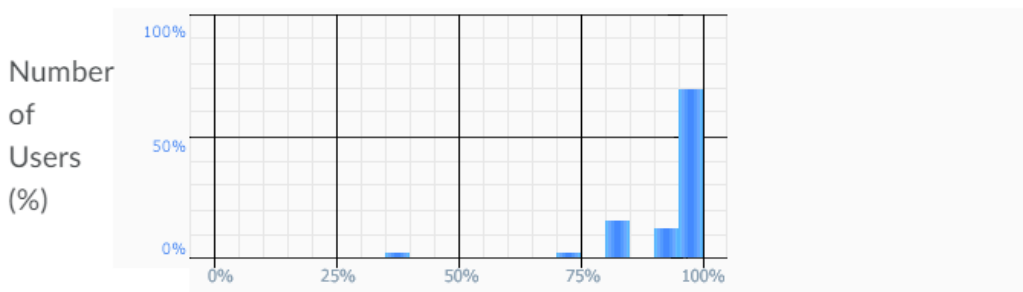
## MQ2 (Wed Sep 15) Class Statistics

Number of submitted grades: 59 / 61

Minimum:  35 %  
Maximum:  100 %  
Average:  94.15 %  
Mode: 100 %  
Median: 100 %

Standard Deviation: 11.21 %

### Grade Distribution



View By:  Groups:

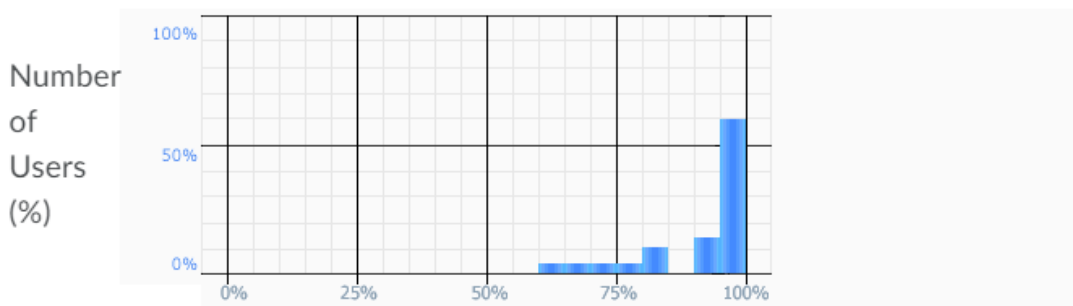
## MQ2 (Wed Sep 15) Class Statistics

Number of submitted grades: 50 / 53

Minimum:  60 %  
Maximum:  100 %  
Average:  91.4 %  
Mode: 100 %  
Median: 100 %

Standard Deviation: 12.41 %

### Grade Distribution

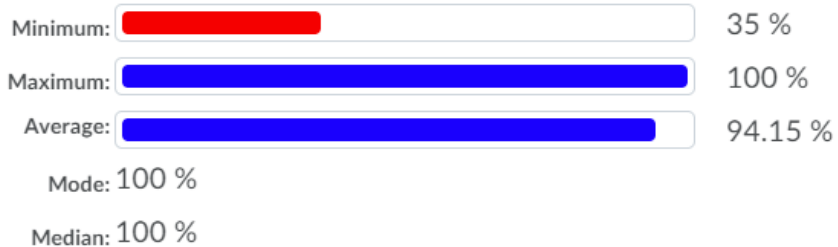



- Friday 17 September:
  - Mobius Quiz 3
- Sunday 19 September:
  - Complete reading up to the end of Section 0.2 (Polynomials)
  - There's a question on next MQ based on the polynomials reading!
- Monday 20 September:
  - Mobius Quiz 4
- Wednesday 22 September:
  - **Complete Written Assignment 2: WA2**
- Wednesday 22 September:
  - Mobius Quiz 5
- Thursday 23 September:
  - WA02 solutions will be posted, hopefully before 12pm: **Check the solutions in detail!**
- Friday 24 September before class:
  - Complete reading Chapter 2 of the course notes
- Friday 24 September before class:
  - Mobius Quiz 6

View By:  Groups:

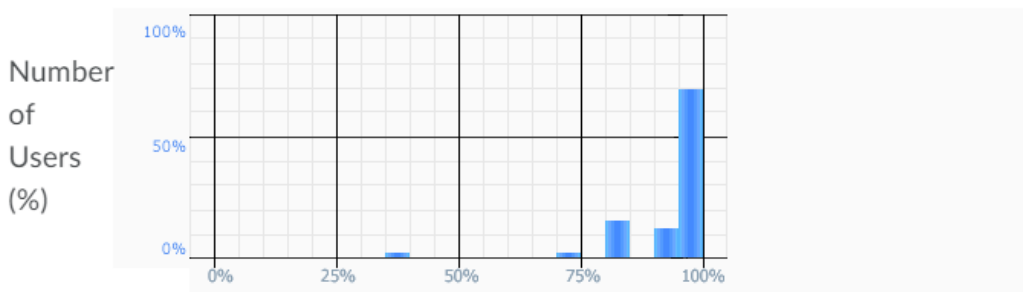
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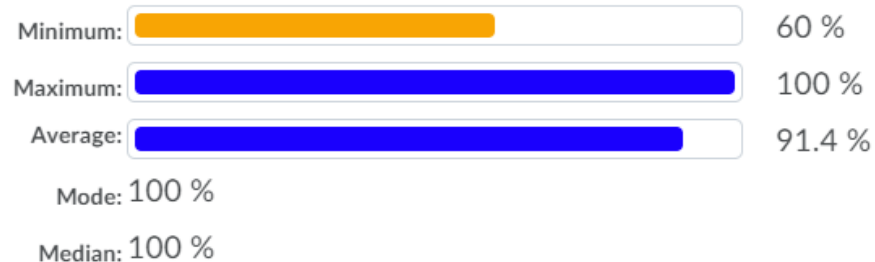
## Grade Distribution




View By:  Groups:

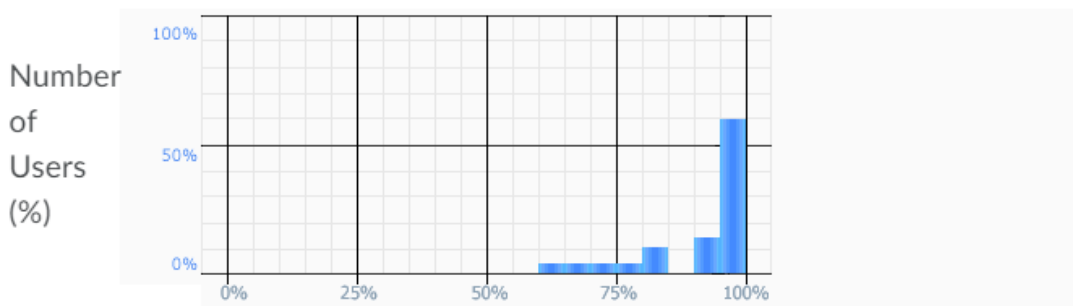
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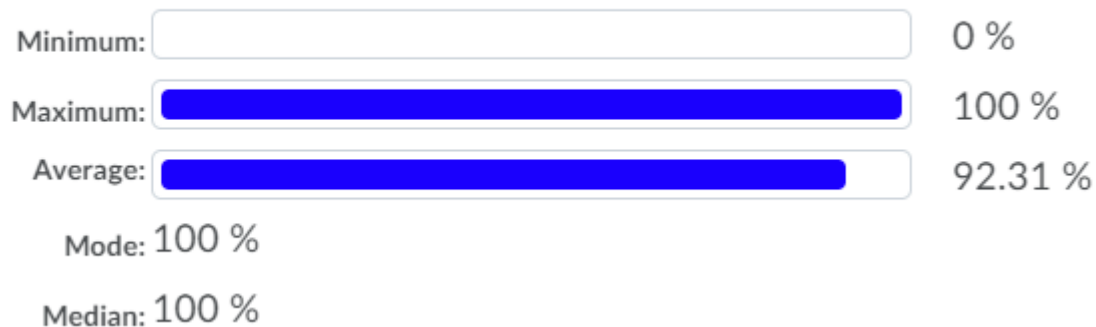


Standard Deviation: 12.41 % 

## Grade Distribution



Number of submitted grades: 1,206 / 1,269



# MATH 135: Lecture 5

Dr. Nike Dattani

17 September 2021

$A$	$B$	$A \wedge B$	$A \vee B$	$\neg A$	$\neg B$	$(\neg A) \wedge (\neg B)$	$(\neg A) \vee (\neg B)$
T	T	T	T	F	F	F	F
T	F	F	T	F	T	F	T
F	T	F	T	T	F	F	T
F	F	F	F	T	T	T	T



$$\neg(A \wedge B) \equiv (\neg A) \vee (\neg B)$$

$$\neg(A \vee B) \equiv (\neg A) \wedge (\neg B)$$

De Morgan's laws!

# Assignment 2!

- Do not skip steps!
  - How many columns will you have in the truth table for proving this:

$$(P \Rightarrow Q) \wedge (\neg Q \Rightarrow R) \equiv Q \vee \neg(R \Rightarrow P)$$

- $\neg(a < b < c)$ ?
- $\neg((a < b) \wedge (b < c))$

How is this simplified?

De Morgan's Laws!

Hint:



“I come to campus only if I have to teach MATH 135”

Do I come to campus if there's no MATH 135?

If there's MATH 135 then do I come to campus?

If I have to teach MATH 135 then I come to campus:  $A \Rightarrow B$

I come to campus only if I have to teach MATH 135:  $B \Rightarrow A$

I come to campus iff I have to teach MATH 135:  $B \Rightarrow A$  (if is from only if)  
 $B \Leftarrow A$  (only if is from if/then)

I have to teach MATH 135 iff I come to campus:  $A \Leftarrow B$  (if is from only if)  
 $A \Rightarrow B$  (only if is from if/then)

# Assignment 2!

“The only positive integers are those in  $\mathbb{Q}$ .”

- Is this iff?
- Positive integer  $\Rightarrow \mathbb{Q}$  (it's the only way something can be a positive integer)
- $\mathbb{Q} \Rightarrow$  Positive integer? (Some  $\mathbb{Q}$  are not integers)
- Q2c on WA02 had a typo.
  - It was written as only  $A \Rightarrow B$ , but should be  $A \Leftrightarrow B$ .
- Remember that:  $A \Leftrightarrow B$  can be turned into:  $(A \Rightarrow B) \wedge (B \Rightarrow A)$ .



# Assignment 2!

Prove *without using a truth table*:

$$\neg (A \wedge (\neg B)) \equiv B \vee (\neg A)$$

For statement variables  $A$  and  $B$ , prove that

$$\neg(A \wedge (\neg B)) \equiv (B \vee (\neg A))$$

without using a truth table.

**Solution:** Starting with the logical expression to the left of the  $\equiv$  sign, we have

$$\begin{aligned}\neg(A \wedge (\neg B)) &\equiv (\neg A) \vee (\neg(\neg B)), \\ &\equiv (\neg A) \vee B, \\ &\equiv B \vee (\neg A),\end{aligned}$$

using De Morgan's Laws,  
using double negation,  
using Commutative Laws.

# Prove *without using a truth table*

What are some of the properties you can use?

- De Morgan's laws!
- Double negation
- Commutativity
- Associativity
- Distributivity
- Negation of an implication:  $\neg (A \Rightarrow B)$
- Negation of an *iff*:  $\neg (A \Leftrightarrow B) = \neg (A \Rightarrow B \wedge B \Rightarrow A)$

Do not skip any steps !!!

If  $x$  and  $y$  are in  $Z$ , then  $x$  is in  $Z$  and  $y$  is not in  $Z$

Hypothesis?

Conclusion?

Converse?

Contrapositive?

Negation?

$$\neg(A \wedge B) = \neg A \vee \neg B$$



What's the last one called?

# Q5: Justification is necessary!

Given  $\mathbf{x}, \mathbf{y}$ , there exists some  $\mathbf{r}$  that has a certain relation to  $\mathbf{x}, \mathbf{y}$

- How would you justify that it's true?
- An easy way is often to just find the  $\mathbf{r}$ !
- How would you justify that it's **not** true?
- An easy way is to find an  $(\mathbf{x}, \mathbf{y})$  pair for which it's impossible to have such an  $\mathbf{r}$

# Review of last lecture

- If the hypothesis  $A$  is false, is  $B$  true?
- No! Only  $(A \Rightarrow B)$  is true!

Hypothesis FALSE, but Conclusion also FALSE

Hypothesis (A)

Conclusion (B)

$$\forall x \in \mathbb{R}, ((x > 2) \Rightarrow (x^2 \geq 4)).$$

Implication always TRUE if Hypothesis is FALSE

Implication (A  $\Rightarrow$  B)

$x$	$x > 2$	$x^2 \geq 4$	$(x > 2) \Rightarrow (x^2 \geq 4)$
$x > 2$	T	T	T
$x = 2$	F	T	T
$-2 < x < 2$	F	F	T
$x \leq -2$	F	T	T

- We will use the convention that  $(A \Rightarrow B)$  is **true** if A is **false**. In this case we say it's "***vacuously true***".
- This way we don't have to spend time checking cases that ***do not impact*** the open sentence.
- This convention *might* not be followed in some types of [non-classical logic \(click for link to Wikipedia page!\)](#).