

5374 Fall '22

Numerical Linear Algebra

Lecture 21

Today:

- * Eigenvectors: power iteration / QR iteration
- * The Singular Value Decomposition

Last time we talked about eigenvectors and the eigenvector decomposition of a square matrix A :

$$A = V D V^{-1}$$

where V is an invertible $n \times n$ matrix and D is a diagonal $n \times n$ matrix.

The columns of V are eigenvectors of A and the elements along the diagonal of D are the eigenvalues of A .

We also learned that if A is symmetric then A has an eigenvector decomposition with V an orthogonal matrix.

Today: How can we compute either an eigenvector of A or a full decomposition?

Literature: Solomon's chapter on eigenvalues has a good practical review of up-to-date algorithms.

Bau-Trefethen: A more in-depth discussion of the underlying theory, and in particular of an important topic called Krylov subspace methods.

Now let's review two basic methods.

1. Power Iteration

Assume A admits a basis of eigenvectors:

$$v_1, \dots, v_n$$

with corresponding eigenvalues $\lambda_1, \dots, \lambda_n$ s.t.

$$|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|$$

$$\text{def } \mu := \frac{|\lambda_2|}{|\lambda_1|} \quad (\text{so } \mu \in (0,1))$$

The power iteration method produces a fast approximation to v_1 , and it works as follows:

Given $x \in \mathbb{R}^n$, we know that we can write

$$x = c_1 v_1 + \dots + c_n v_n \quad \left(\begin{array}{c} \text{for some} \\ c_1, \dots, c_n \end{array} \right)$$

We know this is so, even if we don't know anything about v_1, \dots, v_n (other than they exist). Let's assume we have an x such that $c_i \neq 0$. Then, observe that

$$Ax = \lambda_1 c_1 v_1 + \lambda_2 c_2 v_2 + \dots + \lambda_n c_n v_n$$

\vdots

$$A^k x = \lambda_1^k c_1 v_1 + \lambda_2^k c_2 v_2 + \dots + \lambda_n^k c_n v_n$$

We can write the RHS as:

$$\lambda_1^k \left(c_1 v_1 + \left(\frac{\lambda_2}{\lambda_1} \right)^k c_2 v_2 + \dots + \left(\frac{\lambda_n}{\lambda_1} \right)^k c_n v_n \right)$$

The idea is that $\left| \frac{\lambda_2}{\lambda_1} \right|^k \rightarrow 0$ as $k \rightarrow \infty$ and the same for $\left| \frac{\lambda_j}{\lambda_1} \right|^k$ for $j=3, 4, \dots, n$.

So for large k , the above vector has almost the same direction as v_1 .

Observe:

$$\begin{aligned} \|A^k x - \lambda_1^k c_1 v_1\|_2 &= \lambda_1^k \left\| \left(\frac{\lambda_2}{\lambda_1}\right)^k c_2 v_2 + \dots + \left(\frac{\lambda_n}{\lambda_1}\right)^k c_n v_n \right\| \\ &\leq \lambda_1^k \sum_{j=2}^n \left| \frac{\lambda_j}{\lambda_1} \right|^k |c_j| \\ &\leq \lambda_1^k \left| \frac{\lambda_2}{\lambda_1} \right|^k \sum_{j=2}^n |c_j| \end{aligned}$$

Remark: While we don't know $\sum_{j=2}^n |c_j|$ directly, it is independent of k . Moreover, if A is symmetric so that the v_i 's are orthonormal, we have

$$\begin{aligned} \sum_{j=2}^n |c_j| &\stackrel{\text{(Cauchy-Schwarz)}}{\leq} \left(\sum_{j=2}^n 1^2 \right)^{1/2} \left(\sum_{j=2}^n |c_j|^2 \right)^{1/2} \\ &\leq \sqrt{n-1} \|x\|_2 \end{aligned}$$

On the other hand, we have the following basic inequality: if $x_1, x_2 \neq 0$ then

$$\left\| \frac{x_1}{\|x_1\|} - \frac{x_2}{\|x_2\|} \right\| \leq \frac{2 \|x_1 - x_2\|}{\max\{\|x_1\|, \|x_2\|\}}$$

$$\left(\frac{x_1}{\|x_1\|} - \frac{x_2}{\|x_2\|} \right) = \frac{x_1}{\|x_1\|} - \frac{x_2}{\|x_1\|} + \frac{x_2}{\|x_1\|} - \frac{x_2}{\|x_2\|}$$

$$\left\| \frac{x_1}{\|x_1\|} - \frac{x_2}{\|x_2\|} \right\| \leq \frac{\|x_1 - x_2\|}{\|x_1\|} + \|x_2\| \left| \frac{1}{\|x_1\|} - \frac{1}{\|x_2\|} \right|$$

$$\leq \left| \frac{\|x_1 - x_2\|}{\|x_1\|} + \frac{\|x_2\|}{\|x_1\| \cdot \|x_2\|} \right| \frac{\|x_2\| - \|x_1\|}{\|x_1\| \cdot \|x_2\|}$$

Since $|\|x_2\| - \|x_1\|| \leq \|x_2 - x_1\|$ (by the triangle inequality)

$$\leq \left(\frac{\|x_1 - x_2\|}{\|x_1\|} + \frac{\cancel{\|x_2\|} \|x_1 - x_2\|}{\|x_1\| \cancel{\|x_2\|}} = \frac{2\|x_1 - x_2\|}{\|x_1\|} \right)$$

Let's combine,

$$\|A^k x - \lambda_1^k c_1 v_1\| \leq |\lambda_1|^k \left| \frac{\lambda_2}{\lambda_1} \right|^k \sum_{j=1}^n |c_j|$$

and

$$\left\| \frac{x_1}{\|x_1\|} - \frac{x_2}{\|x_2\|} \right\| \leq \frac{\|x_1 - x_2\|}{\max\{\|x_1\|, \|x_2\|\}}$$

We get

$$\left\| \frac{A^k x}{\|A^k x\|} - v_1 \right\| \leq \frac{|\lambda_1|^k \left| \frac{\lambda_2}{\lambda_1} \right|^k \sum_{j=1}^n |c_j|}{\| \lambda_1^k c_1 v_1 \| = |\lambda_1|^k |c_1|}$$

(since v_1 is a unit vector)

Then,

$$\left\| \frac{A^k x}{\|A^k x\|} - v_1 \right\| \leq \frac{\cancel{|\lambda_1|^k} \left| \frac{\lambda_2}{\lambda_1} \right|^k \sum_{j=1}^n |c_j|}{\cancel{|\lambda_1|^k} |c_1|}$$

In conclusion,

$$\left\| \frac{A^k x}{\|A^k x\|} - v_1 \right\| \leq \mu^k \left(\frac{1}{\|c\|} \sum_{j=1}^n |c_j| \right)$$

This shows how $\frac{A^k x}{\|A^k x\|}$ can be a good approximation to the "top" eigenvector (v_1) of A .

How good an approximation? , choosing x at random we can be confident - if the v_k 's are orthonormal - that

$$\frac{1}{\|c\|} \sum_{j=1}^n |c_j|$$

won't be too large. If μ is say, ≤ 0.5 , then

$$\mu^k \leq 10^{-15} \quad \text{if } k \geq 50.$$

Even if $\mu = 0.9$, the convergence is still pretty fast! ($\mu^{300} \leq 10^{-13}$)

2. QR Iteration (see Solomon for more details)

Observation: If A and B are such that

$$A = Q B Q^t \quad (Q \text{ orthogonal})$$

Then A and B have the same eigenvalues.

Take A , and compute its QR decomposition

$$A = Q_1 R_1$$

Observe
$$Q_1^{-1} A Q_1 = \cancel{Q_1^{-1}} (\cancel{Q_1} R_1) Q_1$$
$$= R_1 Q_1 =: A_2$$

So $A_2 = Q_1^{-1} A Q_1$, and A_1 and A_2 have the same eigenvalues.

Take the QR decomposition of A_2 ,

$$A_2 = Q_2 R_2$$

$$2r \quad A_3 := R_2 Q_2 \quad (= Q_2^{-1} A Q_2)$$

:

and so on.

We end up with sequen A_k, Q_k, R_k where

$$A_k = Q_k R_k \quad (QR \text{ decomp. of } A_k)$$

$$A_{k+1} := R_k Q_k \quad (= Q_k^{-1} A_k Q_k)$$

The matrices A_k all have the same eigenvalues.

Suppose A_k, Q_k, R_k all converge to matrices $A_\infty, Q_\infty, R_\infty$. Then we have

$$A_\infty = Q_\infty R_\infty = R_\infty Q_\infty$$

Observation: In this situation if v is an eigenvector of R_∞ , then $Q_\infty v$ is also an eigenvector of R_∞ , and with the same eigenvalue! i.e.

$$R_\infty v = \lambda v_\infty \Rightarrow Q_\infty R_\infty v = \lambda Q_\infty v$$

$$\Rightarrow (\text{since } Q_\infty R_\infty = R_\infty Q_\infty) \quad R_\infty(Q_\infty V) = \lambda(Q_\infty V)$$

so $Q_\infty V$ is also an eigenvector of R_∞ with the same eigenvalue.

If the eigenvalues of R_∞ are all different, we would conclude that

$$Q_\infty V = \pm V$$

for any eigenvector of $R_\infty \Rightarrow$ the eigenvalues of R_∞ are the same as those of $R_\infty Q_\infty = A_\infty$

$$\text{Now if } A_\infty V = \lambda V \Rightarrow \lambda V = (R_\infty Q_\infty)V \\ = \pm R_\infty V$$

This shows that up to signs the eigenvalues of A_∞ (so, of A) are given by the diagonal elements of R_∞ .

Krylov Space Methods ...

It has been found that a powerful approach to study a matrix A is to pick a generic vector x_0 and study the matrix:

$$\begin{pmatrix} Ax_0 & A^2 x_0 & \dots & A^n x_0 \end{pmatrix}$$

(see Bar and Trefethen's latter chapters)

The Singular Value Decomposition

Let A be a $m \times n$ matrix. Then ^(with non-negative entries) there exists a diagonal $m \times n$ matrix Σ , an orthogonal matrices U and V such that

$$A = U \Sigma V^t$$

This generalizes the eigenvector decomposition and it holds for all matrices, the price we pay is having two different orthogonal matrices U and V . Their columns are called the left and right singular vectors, and the diagonal elements of Σ are called the singular values.

Observe (previous): If $A = U \Sigma V^t$, then

$$\begin{aligned} AA^t &= U \Sigma V^t V \Sigma^t U^t & A^t A &= V \Sigma^t U^t U \Sigma V^t \\ &= U \Sigma \Sigma^t U^t & &= V \Sigma^t \Sigma V^t \end{aligned}$$