Lincheck: Testing Concurrent Data Structures in Java

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- Teaching concurrent programming course @ ITMO University
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- PhD student @ IST Austria

Writing concurrent code is pain

Writing concurrent code is pain

... testing it is not much easier!

var i = 0 i.inc() i.inc()

```
var i = 0
i.inc() // 0 i.inc() // 0
```

We do not expect this!

Sequential model

sequential specification on operations

Concurrent model



Linearizability

(usually)

Execution *is linearizable* ⇔ ∃ equivalent *sequential* execution wrt *happens-before* order (a bit harder)

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```
val q = MSQueue<Int>()

q.add(1)
q.poll(): 1
q.poll(): 2
q.add(2)
```

Execution *is linearizable* ⇔ ∃ equivalent *sequential* execution wrt *happens-before* order (a bit harder)

This counter is not linearizable











```
class MSQueueTest {
  val q = MSQueue<Int>()
Initial state
```

```
class MSQueueTest {
  val q = MSQueue<Int>()
   @Operation fun add(element: Int) =
         q.add(element)
  @Operation fun poll() = q.poll()
```

Operations on the data structure

```
class MSQueueTest {
  val q = MSQueue<Int>()

  @Operation fun add(element: Int) =
     q.add(element)

@Operation fun poll() = q.poll()
```

Operation parameters can be non-fixed!

```
class MSQueueTest {
  val q = MSQueue<Int>()
   @Operation fun add(element: Int) =
        q.add(element)
                                            The Magic
  @Operation fun poll() = q.poll()
                                              Button
  @Test fun runTest() =
         LinChecker.check(QueueTest::class)
```

```
class MSQueueTest {
                                      Do we have such
  val q = MSQueue<Int>()
                                        instrument?
  @Operation fun add(element: Int) =
         q.add(element)
  @Operation fun poll() = q.poll()
  @Test fun runTest() =
         LinChecker.check(QueueTest::class)
```

```
class MSQueueTest {
                                     Do we have such
  val q = MSQueue<Int>()
                                       instrument?
  @Operation fun add(element: Int) =
        q.add(element)
                                              YEEES!
  @Operation fun poll() = q.poll()
  @Test fun runTest() =
        LinChecker.check(QueueTest::class)
```

Lin-Check Overview

Lincheck = Linearizability Checker (supports not only linearizability)
https://github.com/Kotlin/kotlinx-lincheck

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- 1. Generates a random scenario
- Executes it a lot of times
- 3. Verifies the results

Lin-Check Overview

Lincheck = Linearizability Checker (supports not only linearizability)
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1. Generates a random scenario

2. Executes it a lot of times

3. Verifies the results

ScenarioGenerator

Runner

Verifier

Invalid Execution Example

How to check results for correctness?

Simplest solution:

- 1. Generate all possible sequential histories
- 2. Check whether one of them produces the same results

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2 threads x 15 operations ⇒ OutOfMemoryError

How to check results for correctness?

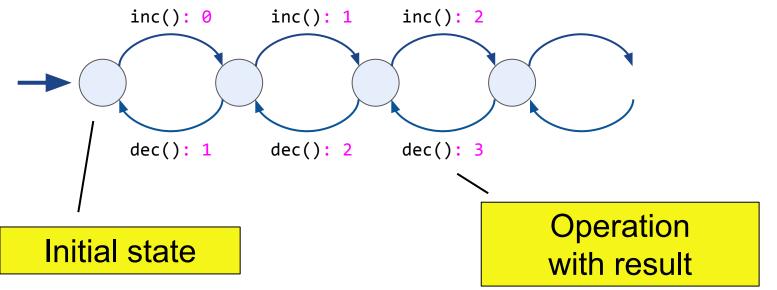
Simplest solution:

- 1. Generate all possible sequential histories
- 2. Check whether one of them produces the same results

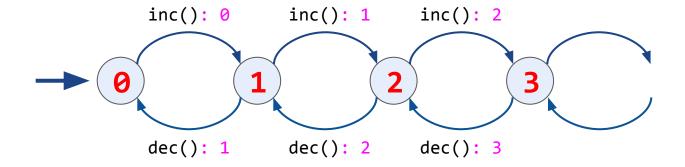
Smarter solution: Labeled Transition System (LTS)

LTS (Labeled Transition System)

LTS is infinite



LTS (Labeled Transition System)

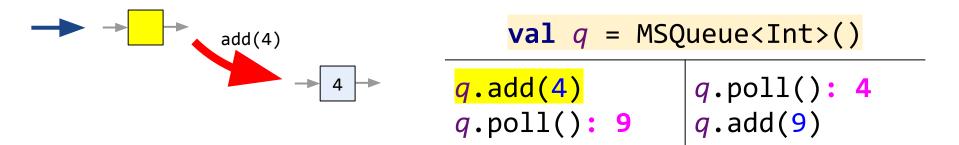


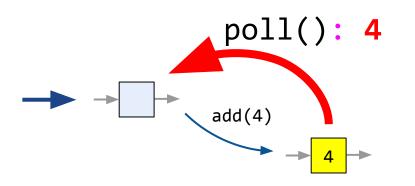
LTS-based verification



```
val q = MSQueue<Int>()

q.add(4)
q.poll(): 9
    q.add(9)
```

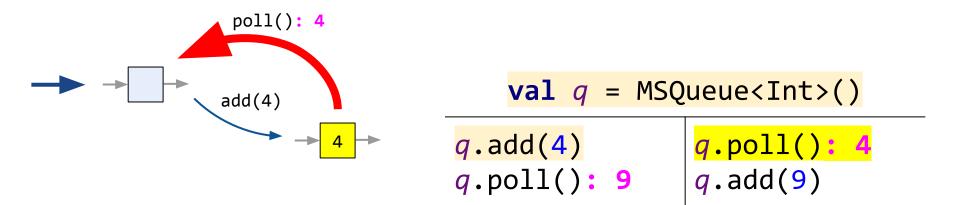


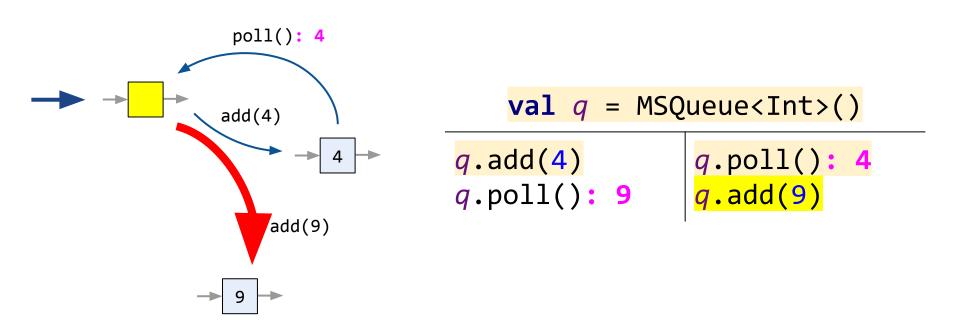


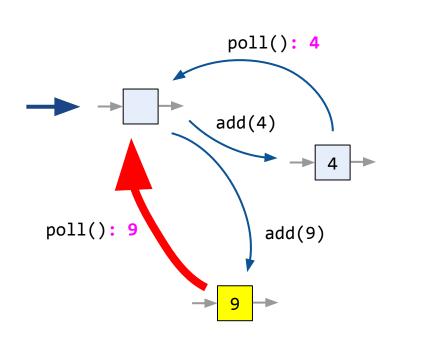
```
val q = MSQueue<Int>()

q.add(4)
q.poll(): 9
q.add(9)
```

Result is different

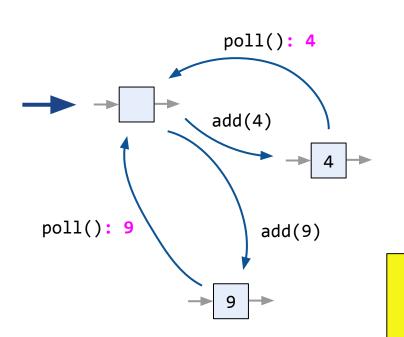






```
val q = MSQueue<Int>()

q.add(4)
q.poll(): 4
q.poll(): 9
```



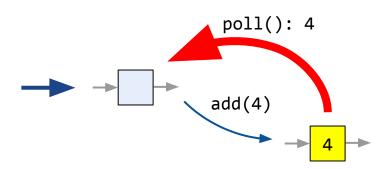
```
val q = MSQueue<Int>()

q.add(4)
q.poll(): 4
q.add(9)
```

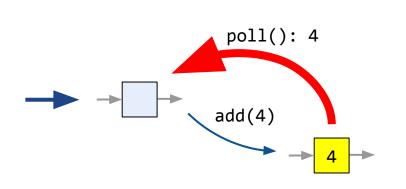
A path is found ⇒ correct

- We build LTS lazilly, like on the previous slides
- We use sequential implementation

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- We use sequential implementation
- Equivalence via equals/hashcode implementations

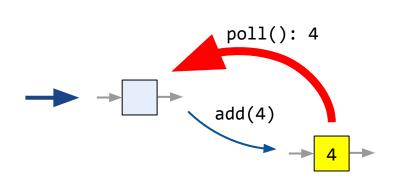


```
class MSQueueTest {
   val q = MSQueue<Int>()

   // Operations here

   override fun equals(other: Any?) = ...
   override fun hashCode() = ...
}
```

- We build LTS lazilly, like on the previous slides
- We use sequential implementation
- Equivalence via equals/hashcode implementations



```
class MSQueueTest: VerifierState() {
   val q = MSQueue<Int>()

   // Operations here

   override fun generateState() = q
}
```

What if our data structure is blocking by design?

send waits for receive and vice versa

```
Producer 1
   val elem = ...
                                Consumer
    c.send(elem)
                                    while(true) {
                                       val elem = c.receive()
                                       process(elem)
Producer 2
   val elem = ...
    c.send(elem)
                        val c = Channel()
```

```
Producer 1
                                                Has to wait for send
   val elem = ...
                                 Consumer
    c.send(elem)
                                     while(true) {
                                       val elem = c.receive()
                                        process(elem)
Producer 2
   val elem = ...
    c.send(elem)
                         val c = Channel()
```

```
Producer 1
   val elem = ...
    c.send(elem)
Producer 2
   val elem = ...
    c.send(elem)
```

```
Consumer

while(true) {

val elem = c.receive()

process(elem)

}
```

```
val c = Channel()
```

```
Producer 1
   val elem = ...
    c.send(elem)
Producer 2
   val elem = ...
    c.send(elem)
```

```
Consumer

while(true) {

val elem = c.receive()

process(elem)

}
```

```
val c = Channel()
```

```
val elem = ...
```

2 c.send(elem)

Producer 2

```
val elem = ...
c.send(elem)
```

Rendezvous!

```
Consumer
   while(true) {
        val elem = c.receive()
        process(elem)
```

```
val c = Channel()
```

```
val elem = ...

2 c.send(elem)
```

Producer 2 val elem = ... c.send(elem)

Consumer

```
while(true) {
1  val elem = c.receive()
3  process(elem)
}
```

```
val c = Channel()
```

```
Producer 1
```

```
val elem = ...

2 c.send(elem)
```

```
val elem = ...
c.send(elem)
```

Consumer

```
while(true) {
1  val elem = c.receive()
3  process(elem)
}
```

```
val c = Channel()
```

```
Producer 1

val elem = ...
```

```
2 c.send(elem)
```

```
val elem = ...
```

4 c.send(elem)

Consumer

```
while(true) {
```

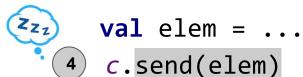
- 1 val elem = c.receive()
- process(elem)
 }

Has to wait for receive

```
val c = Channel()
```

```
Producer 1
```

```
val elem = ...
2 c.send(elem)
```



Consumer

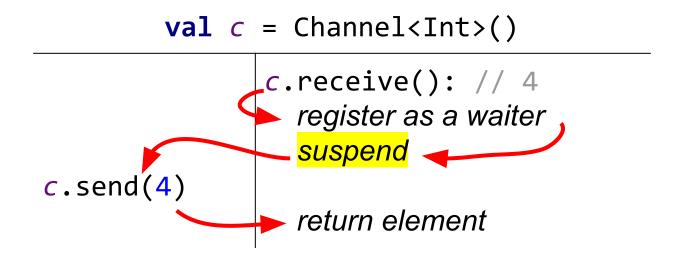
```
while(true) {
    val elem = c.receive()
    process(elem)
}
```

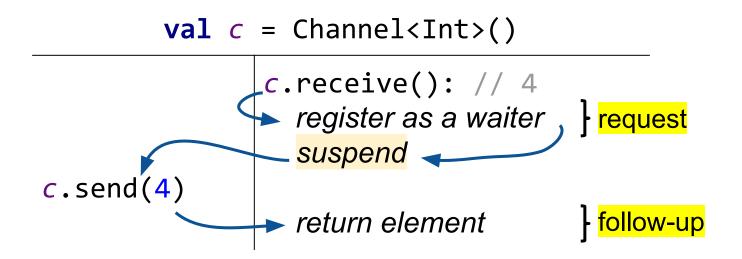
```
val c = Channel()
```

```
Producer 1
   val elem = ...
                                Consumer
 (2) c.send(elem)
                                    while(true) {
                                      val elem = c.receive()
                                     process(elem)
Producer 2
   val elem = ...
    c.send(elem)
                           Has to wait for receive
```

val c = Channel()

Non-linearizable because of suspension





```
val c = Channel<Int>()
c.receive<sup>REQ</sup>(): tik
c.send(4)
c.receive<sup>FUP</sup>(tik): 4
Unique ticket, ∈ℕ
```

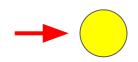
```
val c = Channel<Int>()
c.receive<sup>REQ</sup>(): tik
c.send(4)
c.receive<sup>FUP</sup>(tik): 4
```

Follow-ups should be invoked after the corresponding requests

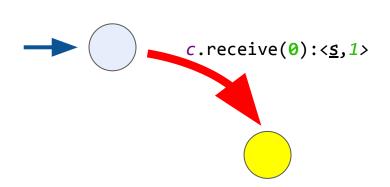
```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4,_>
```

Let's always pass tickets, for simplicity

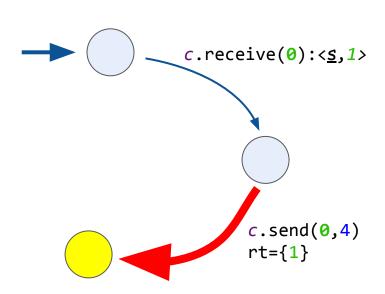
```
val c = Channel<Int>()
c.receive(0): <s,1>
c.send(0, 4)
c.receive(1): <4,_>
suspended
```



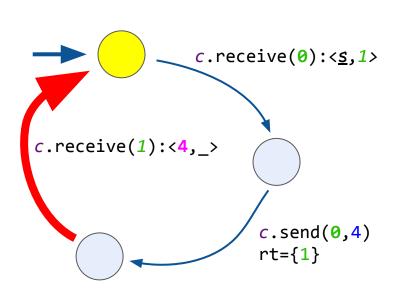
```
val c = Channel<Int>()
c.receive(0): <5,1>
c.send(0, 4)
c.receive(1): <4,_>
```



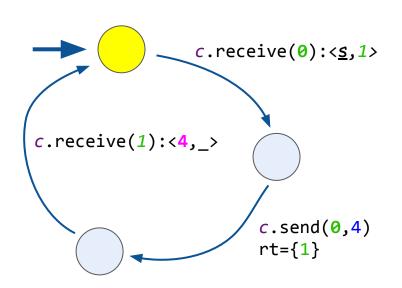
```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4,_>
```



```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4,_>
```

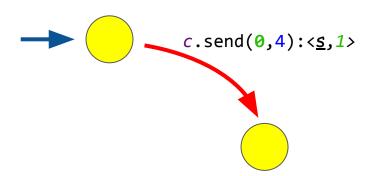


```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4,_>
```

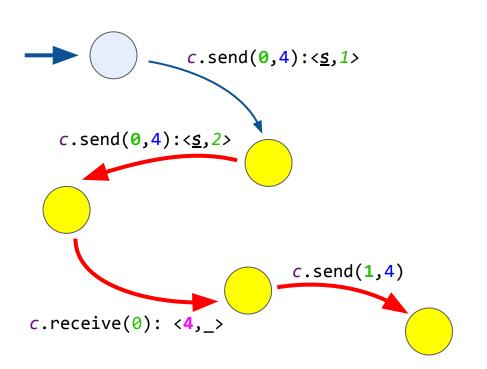


```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4,_>
```

Looks similar

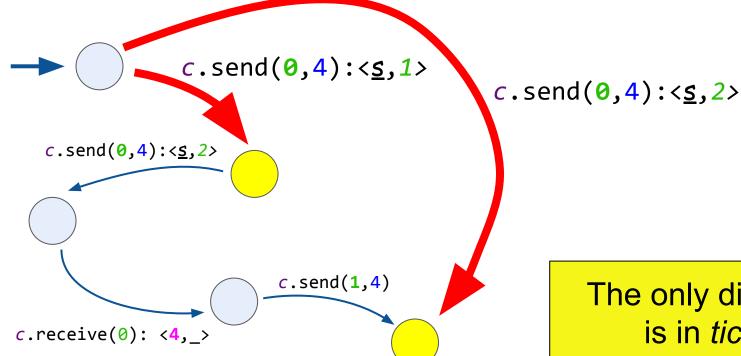


```
val c = Channel<Int>()
c.send(0, 4): <<u>s</u>,1>
```



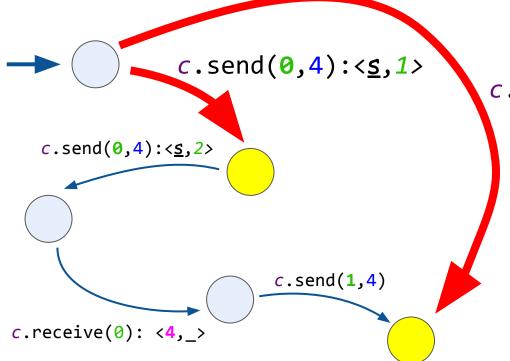
```
val c = Channel<Int>()
c.send(0, 4): <<u>s</u>,1>
c.send(0, 4): <<u>s</u>,2>
c.receive(0): <4,_>
c.send(1, 4)
```

LTS for Dual Data Structures



The only difference is in *tickets*

LTS for Dual Data Structures

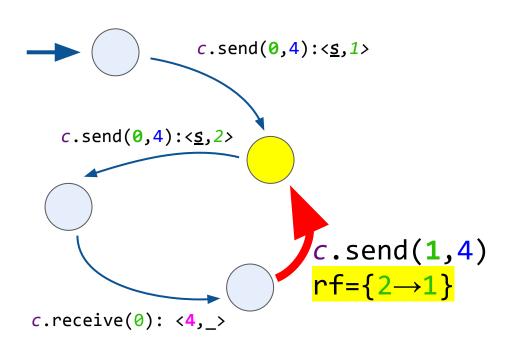


c.send(0,4):<<u>s</u>,2>

Let's forbid such duplicate transitions

The only difference is in *tickets*

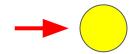
LTS for Dual Data Structures



```
val c = Channel<Int>()
c.send(0, 4): <<u>s</u>,1>
c.send(0, 4): <<u>s</u>,2>
c.receive(0): <4,_>
c.send(1, 4)
```

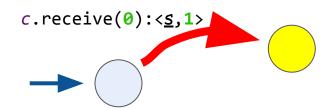
```
val c = Channel<Int>()

c.receive(): 4
c.receive(): 5
c.send(4): 5+Unit
```



```
val c = Channel<Int>()
```

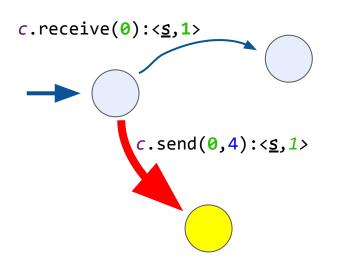
```
c.receive(): 4
c.receive(): 5
```



```
val c = Channel<Int>()
```

```
c.receive(): 4
c.receive(): 5
```

Results are different

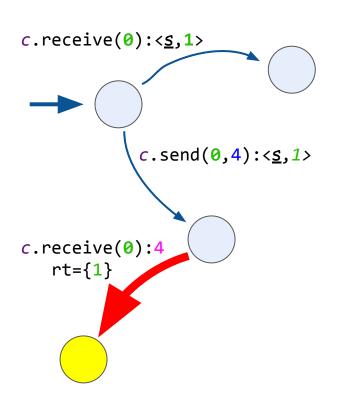


```
val c = Channel<Int>()
```

```
c.receive(): 4
c.receive(): 5
```

```
c.send(4): ≤+Unit
```

suspended, ticket 1

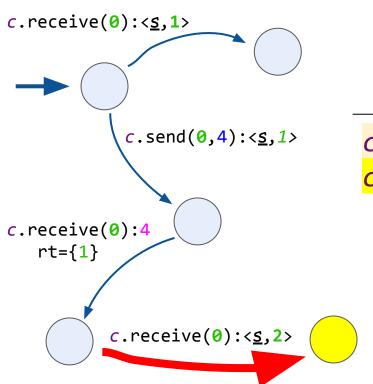


```
val c = Channel<Int>()
```

```
c.receive(): 4
c.receive(): 5
```

```
c.send(4): ≤+Unit
```

suspended, ticket 1 resumed

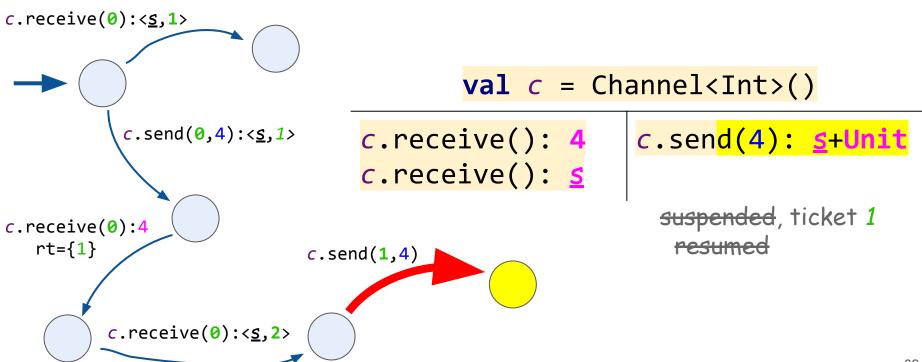


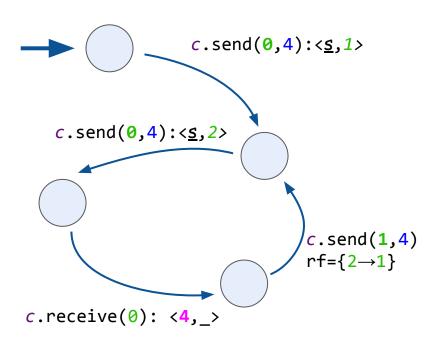
```
val c = Channel<Int>()
```

```
c.receive(): 4
c.receive(): 5
```

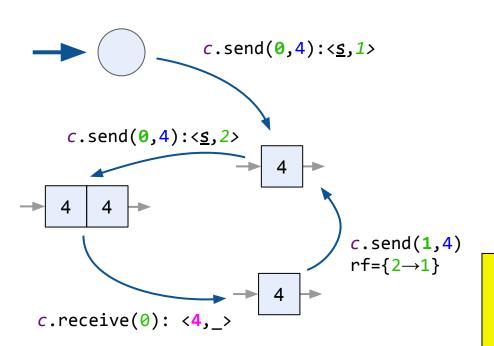
```
c.send(4): ≤+Unit
```

suspended, ticket 1 resumed



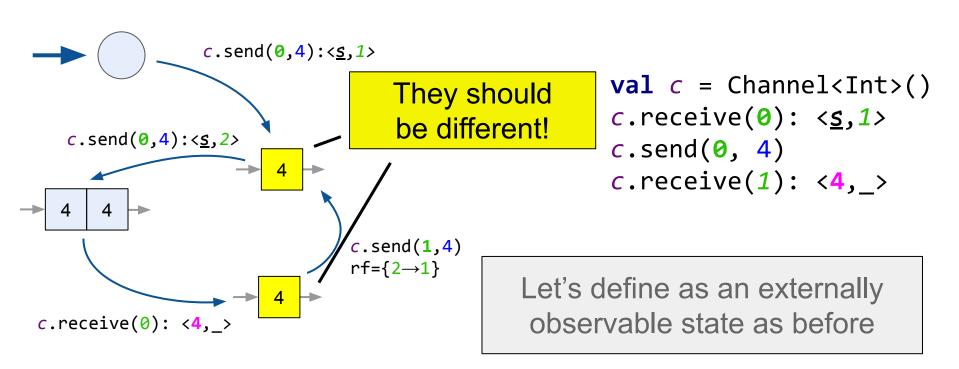


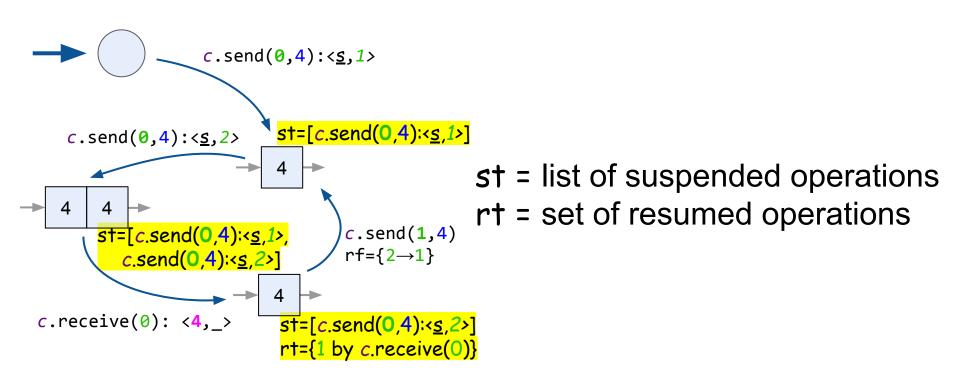
```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4,_>
```

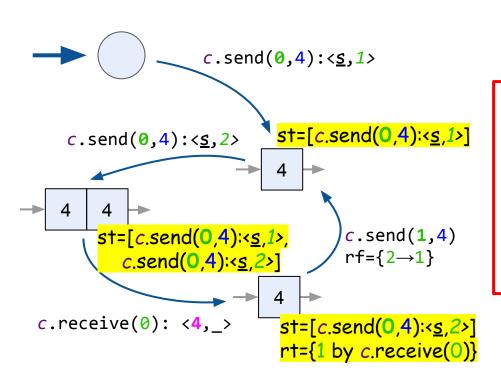


```
val c = Channel<Int>()
c.receive(0): <<u>s</u>,1>
c.send(0, 4)
c.receive(1): <4, >
```

Let's define as an externally observable state as before

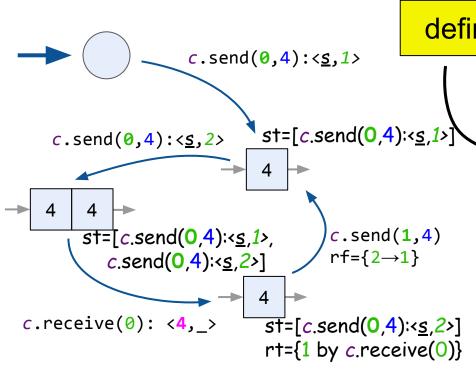






States are equal iff $\exists f: \mathbb{N} \rightarrow \mathbb{N}$ that

- 1. externally observable states
- 2. st-s wrt rf on tickets (as lists)
- 3. rt-s wrt rf on tickets (as sets) are equal



defined via equals/hashcode

States are equal iff $\exists f: \mathbb{N} \rightarrow \mathbb{N}$ that

- 1. externally observable states
- 2. st-s wrt rf on tickets (as lists)
- 3. rt-s wrt rf on tickets (as sets) are equal

maintained by Lin-Check

Channel Test Example

```
class RendezvousChannelTest: LinCheckState() {
   val c = Channel()

    @Operation    suspend    fun    send(x: Int) = c.send(x)
    @Operation    suspend    fun    receive(): Int = c.receive()

   override fun    generateState() = Unit
}
```

Channel Test Example

```
class BufferedChannelTest: LinCheckState() {
  val c = Channel()
  @Operation suspend fun send(x: Int) = c.send(x)
  @Operation suspend fun receive(): Int = c.receive()
   override fun generateState(): Any {
       val state = ArrayList<Int>()
       var x: Int?
      while(true) {
           x = c.poll()
           if (x == null) break
           state += x
       return state
```

Uncovered topics

- Verifiers for several relaxed contracts
- How to run scenarios in the most "dangerous" way
- API

Future plans

- Smart running strategies
- Supporting randomized relaxed contracts

Questions?

https://github.com/Kotlin/kotlinx-lincheck