

# How to find deadlock not hitting it

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# Real-life example

```
interface Con { // Connection
    void onRegister(ConnectionManager cm);
    void close();
}
```

```
class ConnectionManager {
    Set<Con> cons = new HashSet<>();

    synchronized void register(Con c) {
        cons.add(c);
        con.onRegister(this);
    }

    synchronized void unregister(Con c) {
        cons.remove(c);
    }

    ...
}
```

```
interface Con { // Connection
    void onRegister(ConnectionManager cm);
    void close();
}

class ConnectionManager {
    Set<Con> cons = new HashSet<>();

    synchronized void register(Con c) {
        cons.add(c);
        con.onRegister(this);
    }

    synchronized void unregister(Con c) {
        cons.remove(c);
    }

    ...
}

// ... ConnectionManager

synchronized int getConCount() {
    return cons.size();
}

synchronized void closeAll() {
    for (Con c: cons)
        c.close();
}
}
```

```
class ConnectionImpl implements Con {
    ConnectionManager cm;

    @Override
    public synchronized void onRegister(ConnectionManager cm) {
        this.cm = cm;
    }

    @Override
    public synchronized void close() {
        if (cm != null) {
            cm.unregister(this);
            cm = null;
        }
    }
}
```

Run functional test...

Works well!

Run stress test...

```
public void test() throws Exception {
    ConnectionManager cm = new ConnectionManager();
    new Thread(() -> {
        while (true) {
            Connection con = new ConnectionImpl();
            cm.register(con);
            con.close(); // lock con -> lock cm
        }
    }).start();
    for (int i = 1; i < 100_000; i++) {
        System.out.println("Iteration #" + i);
        cm.closeAll(); // lock cm -> lock con
    }
}
```

Java stack information for the threads listed above:

---

"Thread-0":

```
at a.ConnectionManager.unregister(Sample.java:20)
- waiting to lock <0x000000795b8ee60> (a a.ConnectionManager)
at a.ConnectionImpl.close(Sample.java:45)
- locked <0x000000795d72fd8> (a a.ConnectionImpl)
at a.StressTest.lambda$test$0(StressTest.java:16)
at a.StressTest$$Lambda$1/1297685781.run(Unknown Source)
at java.lang.Thread.run(Thread.java:745)
```

"main":

```
at a.ConnectionImpl.close(Sample.java:44)
- waiting to lock <0x000000795d72fd8> (a a.ConnectionImpl)
at a.ConnectionManager.closeAll(Sample.java:29)
- locked <0x000000795b8ee60> (a a.ConnectionManager)
at a.StressTest.test(StressTest.java:21) <27 internal calls>
```

**Found 1 deadlock.**



## ConnectionManager cm Connection con

---

```
1 MONITORENTER con
2 con.close()
5 MONITORENTER cm
6 cm.unregister(con)
```

```
3 MONITORENTER cm
4 cm.closeAll()
7 MONITORENTER con
8 con.close()
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## ConnectionManager cm Connection con

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1 MONITORENTER con  
2 con.close()  
  
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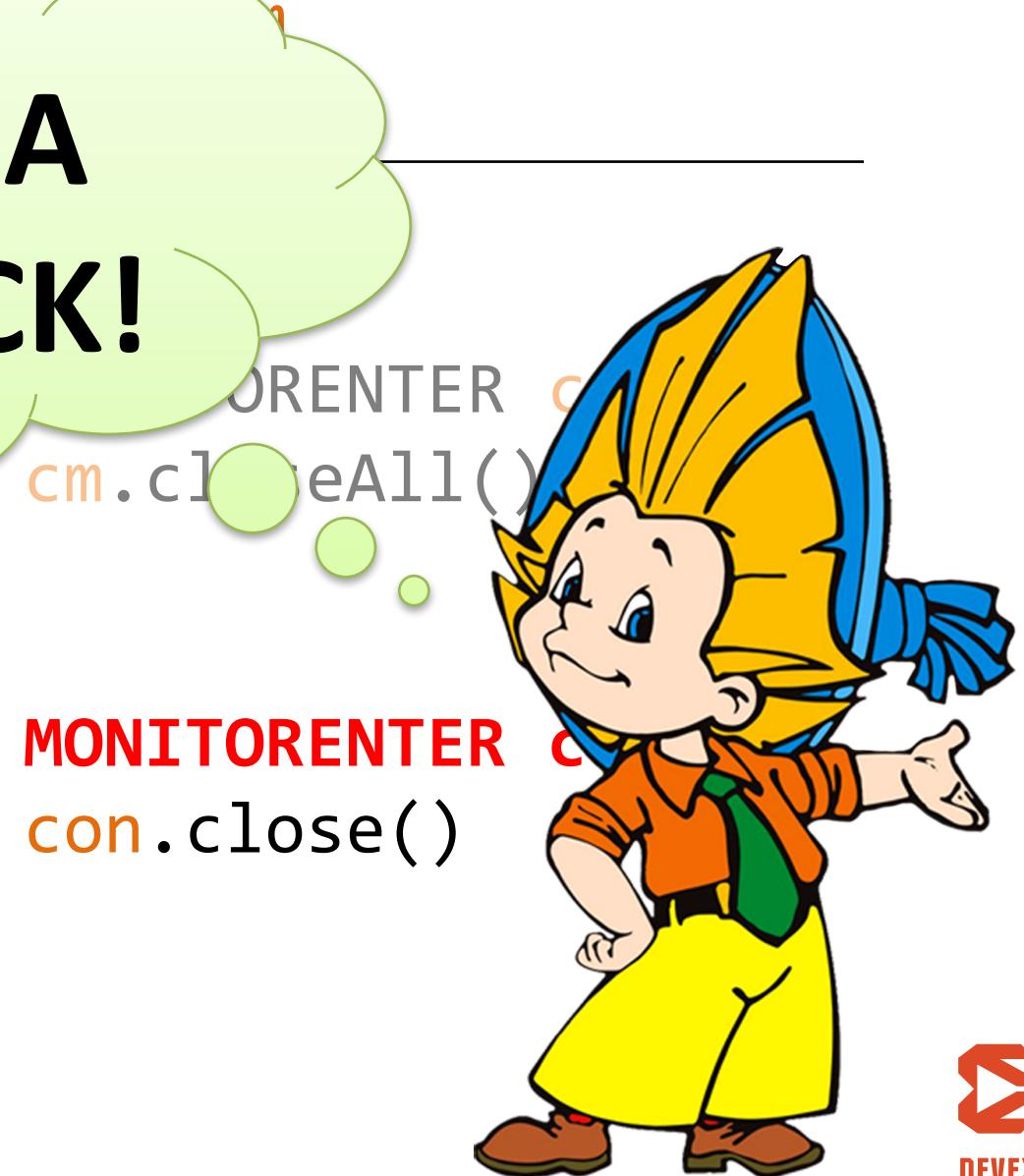
3 MONITORENTER cm  
4 cm.closeAll()  
  
7 MONITORENTER con  
8 con.close()

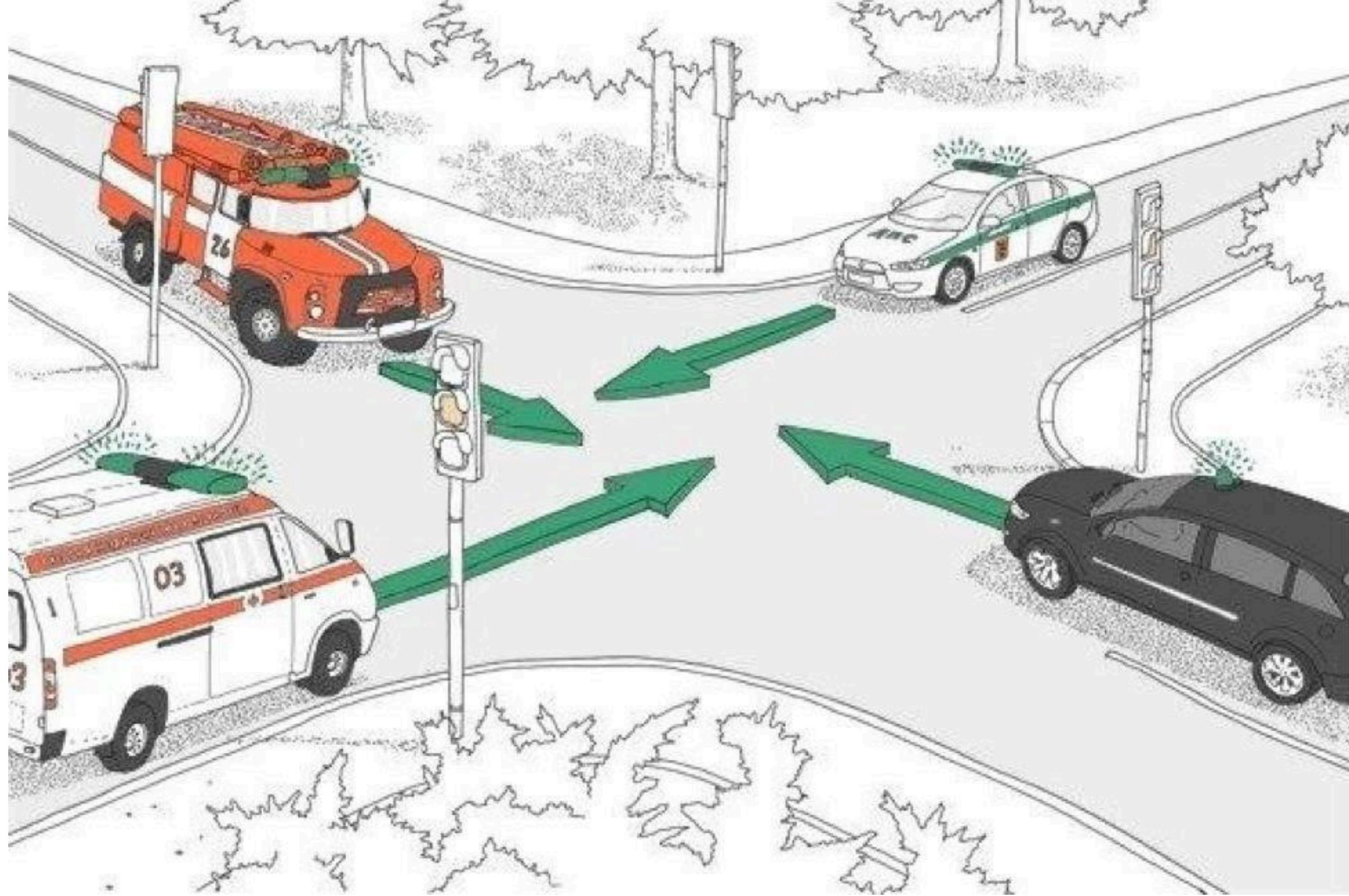
HERE IS A  
**DEADLOCK!**

1 MONITOR  
2 con.close()

5 MONITORENTER cm  
6 cm.unregister(con)

7 MONITORENTER c  
8 con.close()





# Program analysis approaches

- Static analysis
- Model checking
- Dynamic analysis

# Static analysis

- Analyzes source code without its execution
- Can guarantee deadlock-freedom
- A lot of false positives

# Model checking

- Analyzes program **model**
- Can guarantee deadlock-freedom
- No false positives
- Requires a lot of computational resources

# Dynamic analysis

- Depends on the execution
- Few false positives
- Applicable for large programs

# Dynamic analysis: techniques

- Collecting execution traces (off-line)
  - JCarder, MulticoreSDK
- **Detection immediately (on-line)**
  - All context information is available
    - Stacktrace, internal state
  - VisualThreads

# ALGORITHM

# Lock hierarchy

Any two locks should be acquired in the same order in the same thread

```
void transfer(Account from, Account to, int amount) {  
    // Order lock acquisitions by ids  
    if (from.id < to.id)  
        lock(from, to)  
    else  
        lock(to, from)  
    // Do transfer  
    unlock(from, to)  
}
```



# Potential deadlock

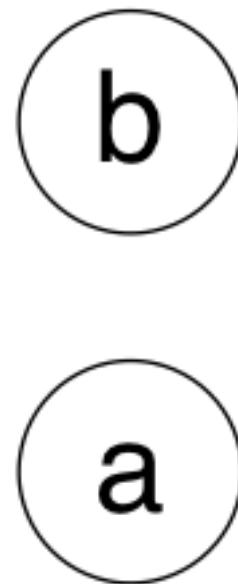
- Lock hierarchy compliance – primary method to avoid deadlocks
- *Potential deadlock* – this hierarchy violation

# Lock graph

- Directed graph
- Vertex  $\leftrightarrow$  lock
- Edge  $(a, b)$  means that  $a$  is acquired before  $b$
- Cycle = potential deadlock

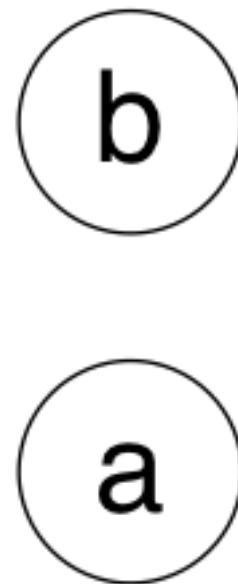
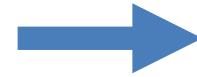
# Lock graph: example

Thread 1	Thread 2
1 a.lock()	
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3 unlock(a, b)	
	4 b.lock()
	5 a.lock()
	6 unlock(a, b)



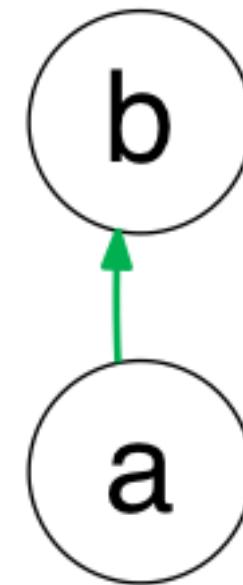
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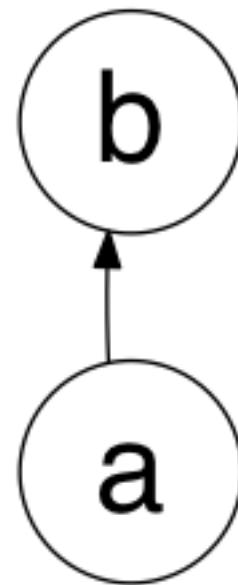
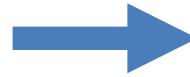
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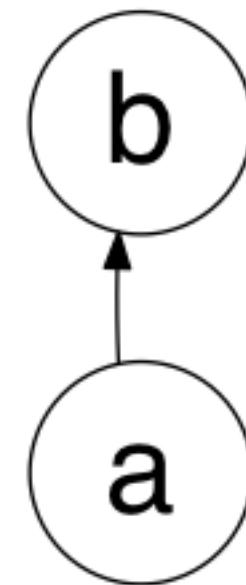
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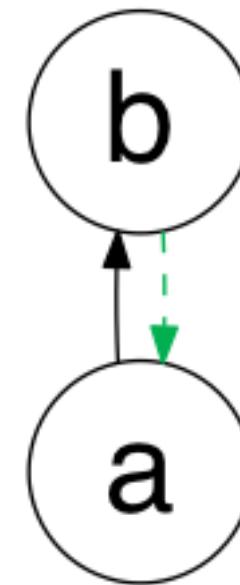
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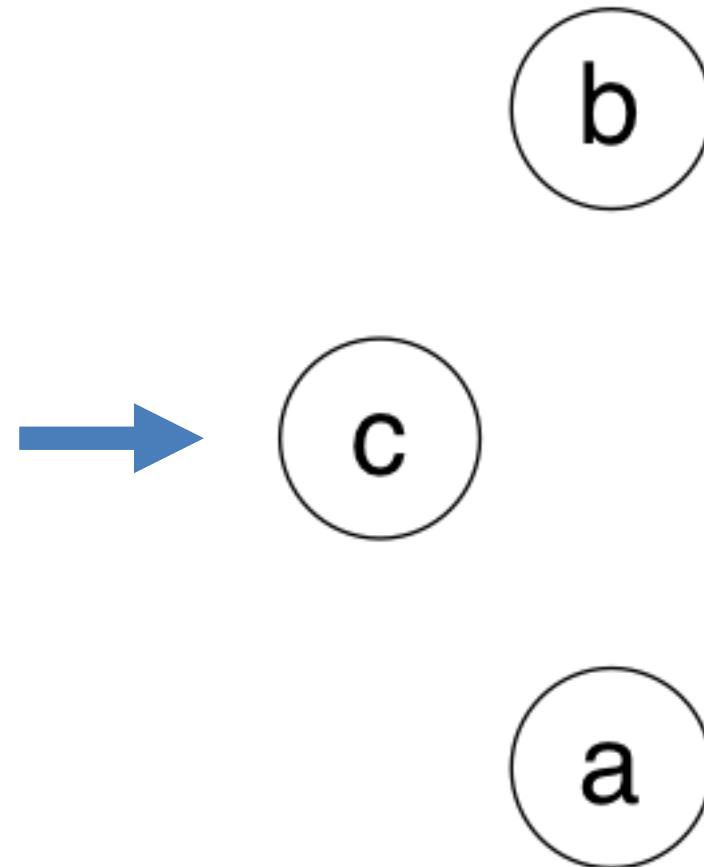
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# Minimization principle: problem

- Two cycles may refer to the same error

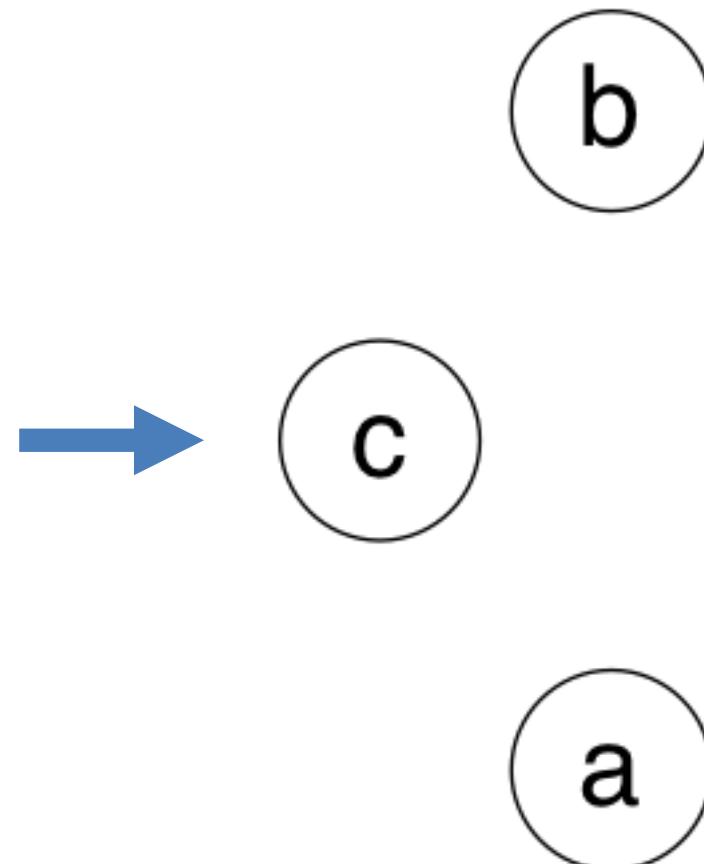
Thread 1	Thread 2
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# Minimization principle: problem (1)

- Two cycles may refer to the same error

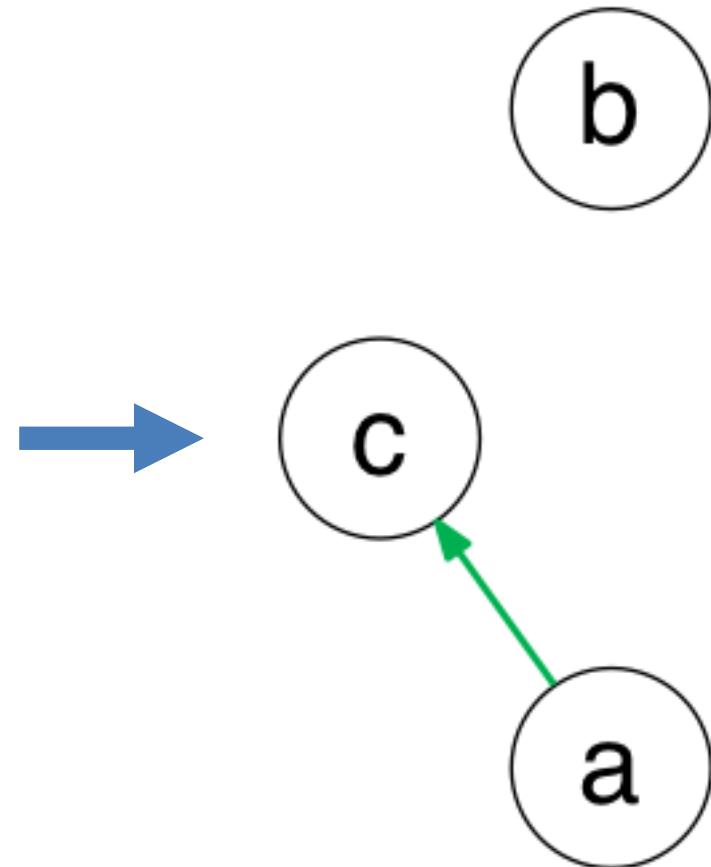
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# Minimization principle: problem (2)

- Two cycles may refer to the same error

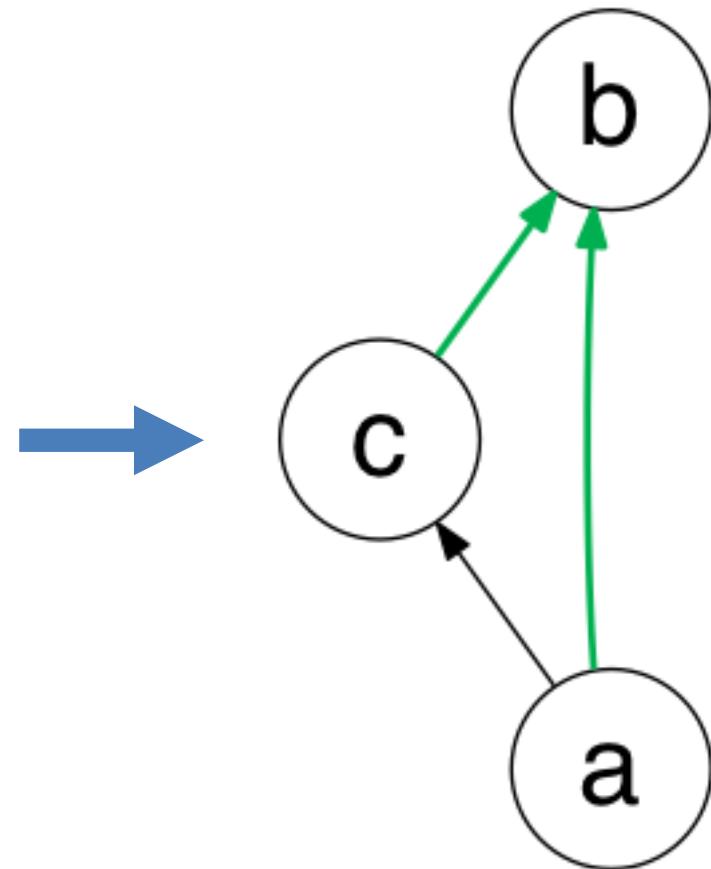
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# Minimization principle: problem (3)

- Two cycles may refer to the same error

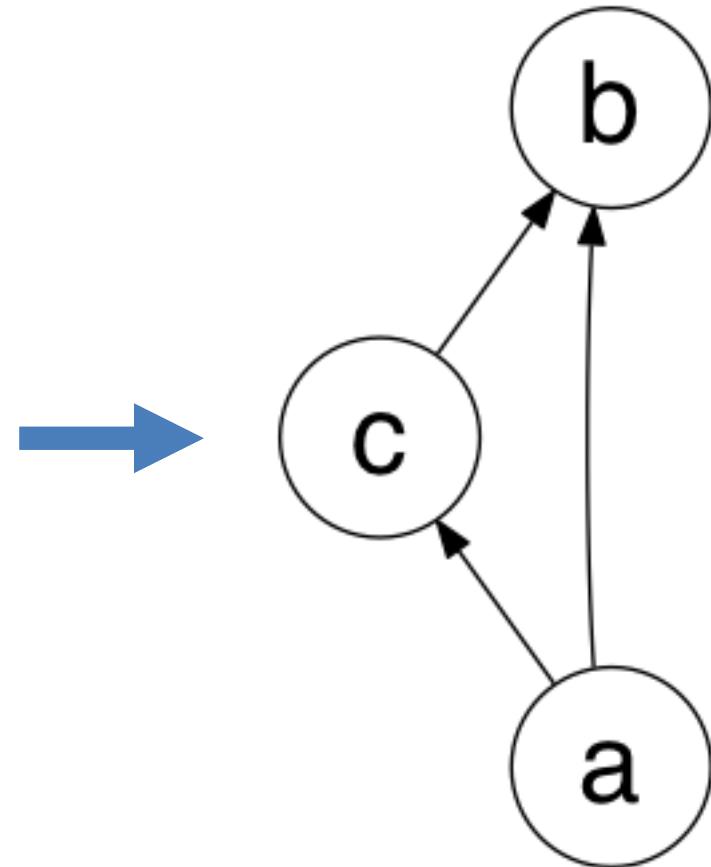
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3 b.lock()	
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# Minimization principle: problem (4)

- Two cycles may refer to the same error

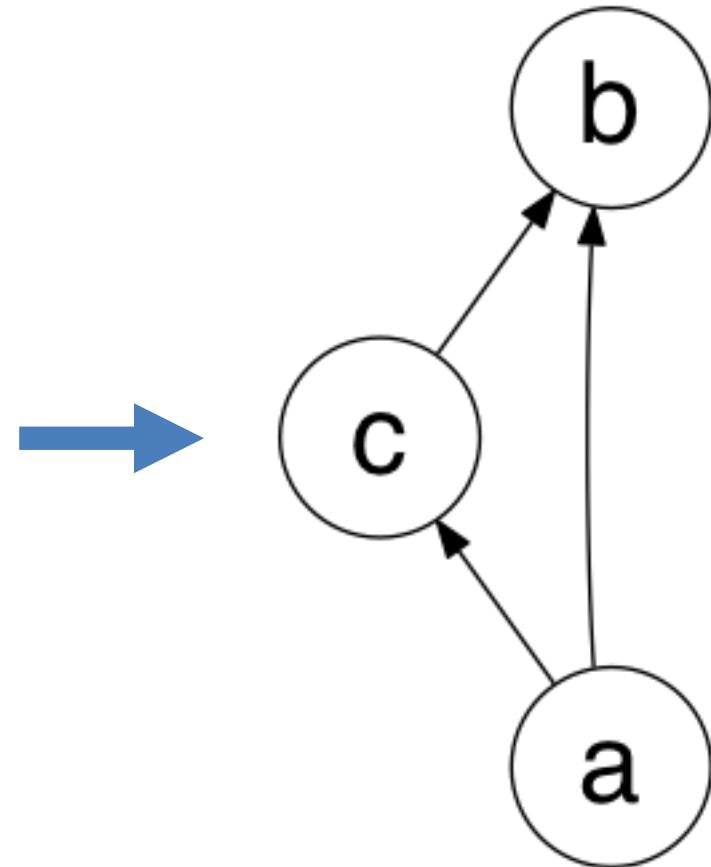
Thread 1	Thread 2
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2 c.lock()	
3 b.lock()	
4 <b>unlock(a,b,c)</b>	
	5 b.lock()
	6 a.lock()
	7 unlock(a,b)



# Minimization principle: problem (5)

- Two cycles may refer to the same error

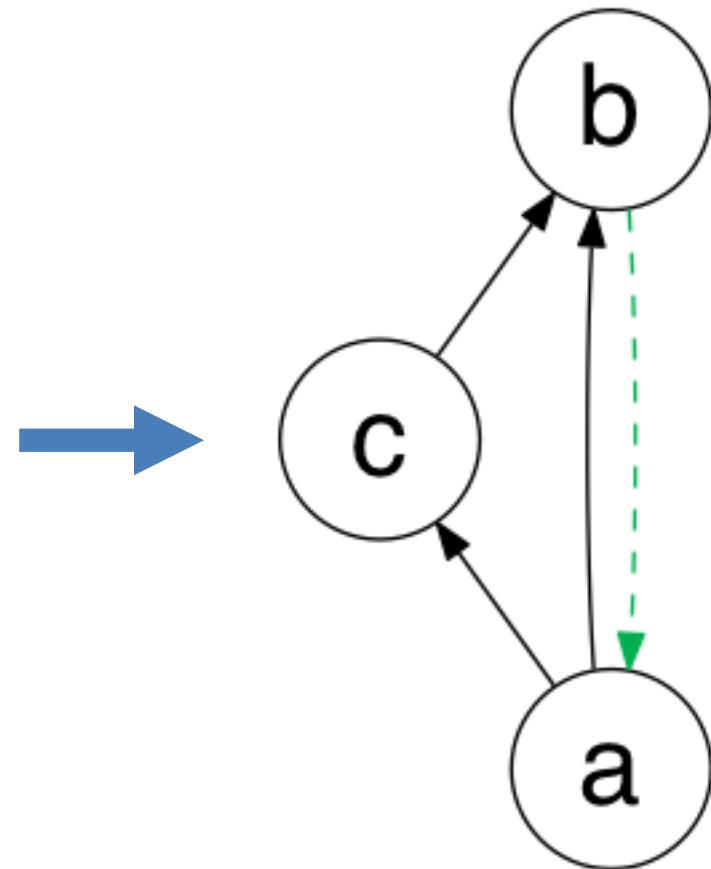
Thread 1	Thread 2
1 a.lock()	
2 c.lock()	
3 b.lock()	
4 unlock(a,b,c)	
	5 b.lock()
	6 a.lock()
	7 unlock(a,b)



# Minimization principle: problem (6)

- Two cycles may refer to the same error

Thread 1	Thread 2
1 a.lock()	
2 c.lock()	
3 b.lock()	
4 unlock(a,b,c)	
	5 b.lock()
	6 a.lock()
	7 unlock(a,b)



# Minimization principle

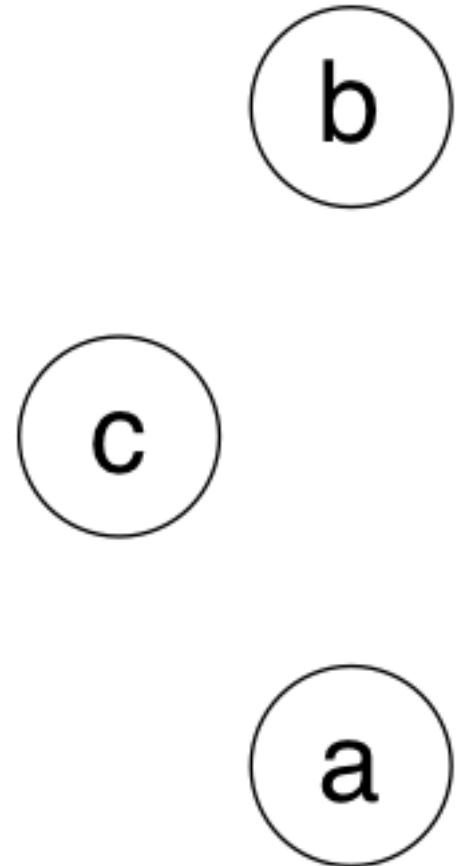
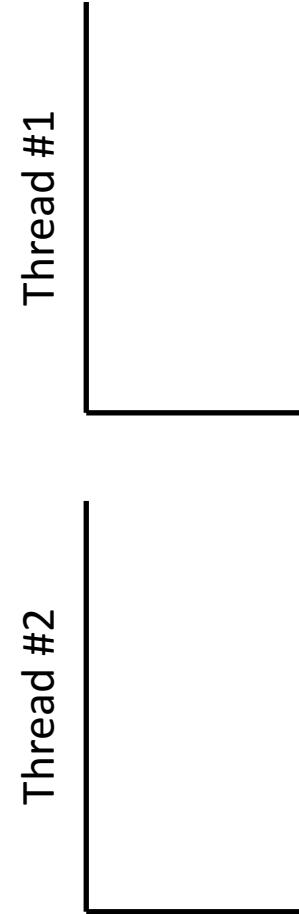
- Shorter cycle is more useful
- Notify about the shortest one only

# Algorithm layout

- Capture lock acquire and release operations
- On lock acquisition:
  - Add lock to the set of locks held by the current thread
  - Add new edges to the lock graph
  - Report new cycles
- On lock release:
  - Remove lock from the set of held locks

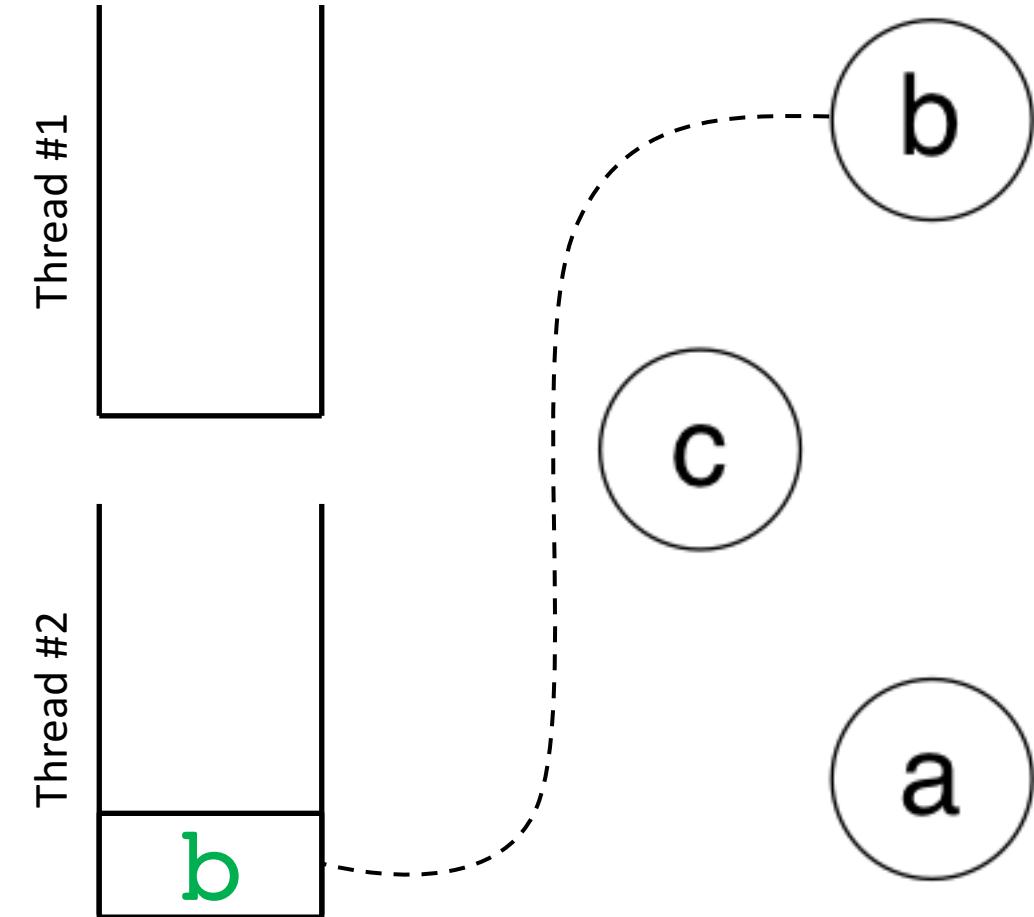
# Lock acquisition and release analysis

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
	6 b.unlock()
7 b.lock()	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	



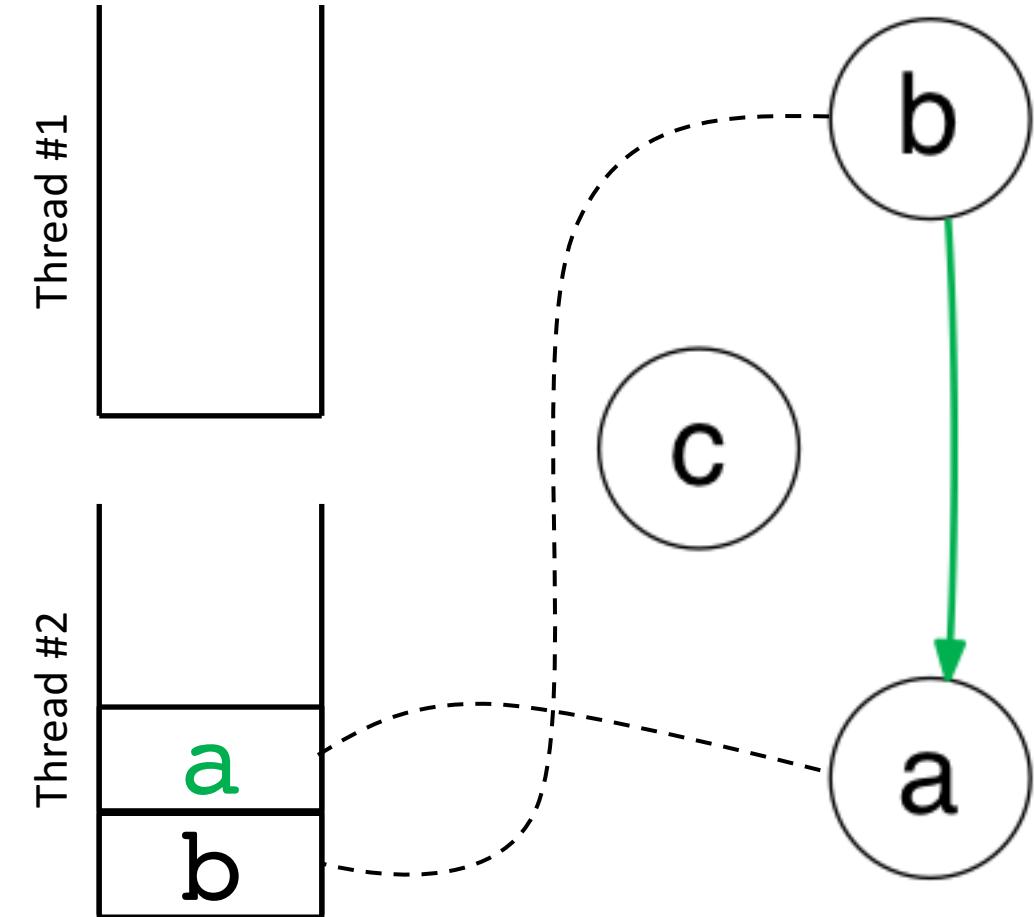
# Lock acquisition and release analysis (1)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
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7 b.lock()	
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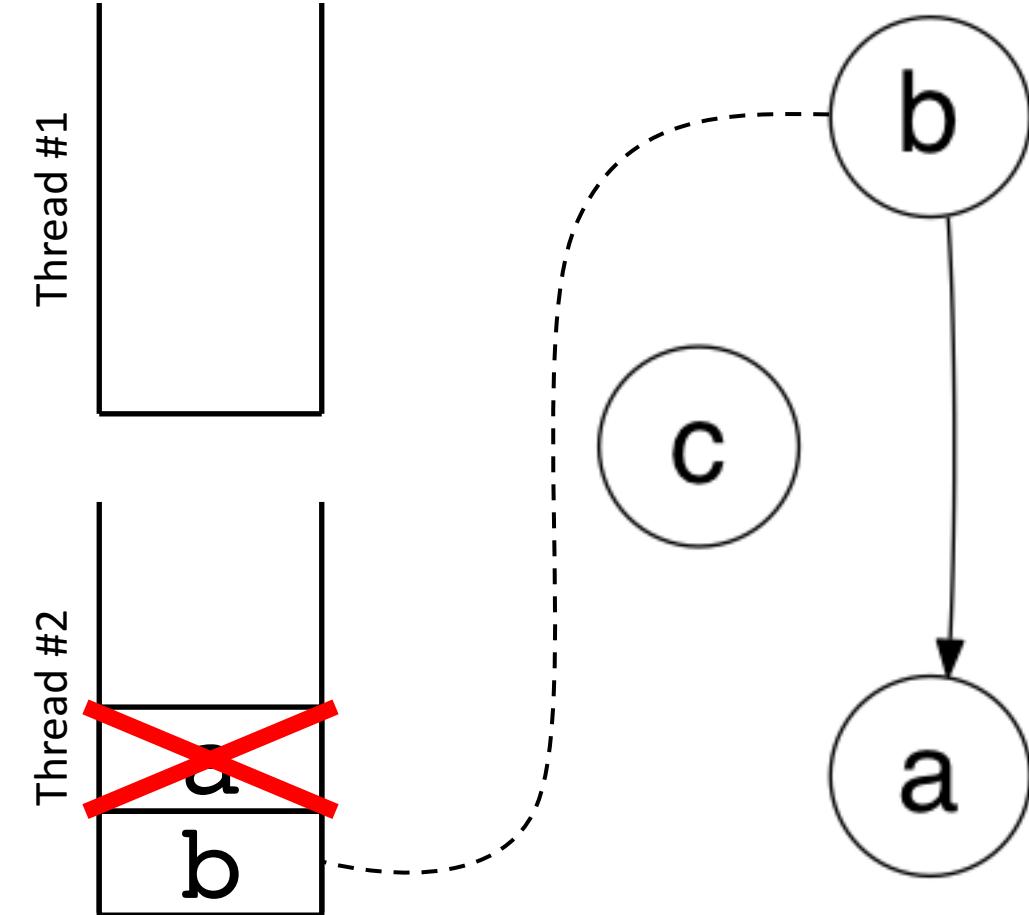
# Lock acquisition and release analysis (2)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
7 b.lock()	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	



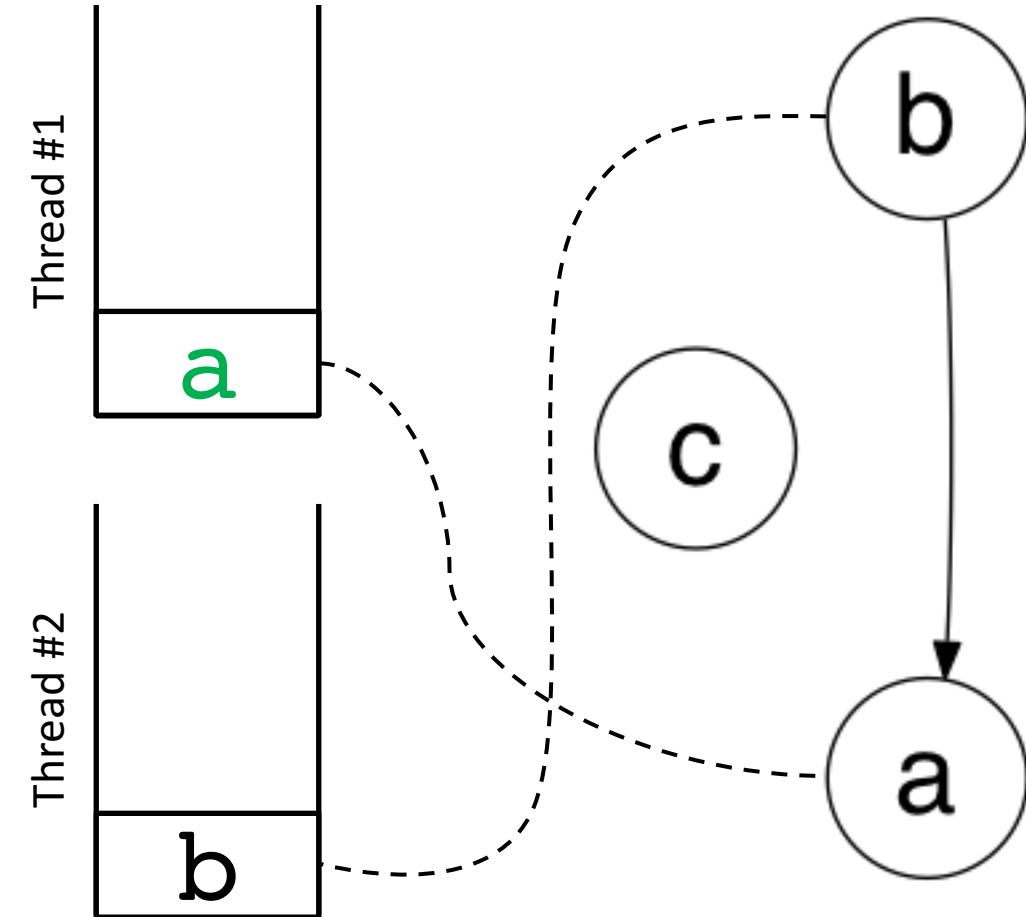
# Lock acquisition and release analysis (3)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
7 b.lock()	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	



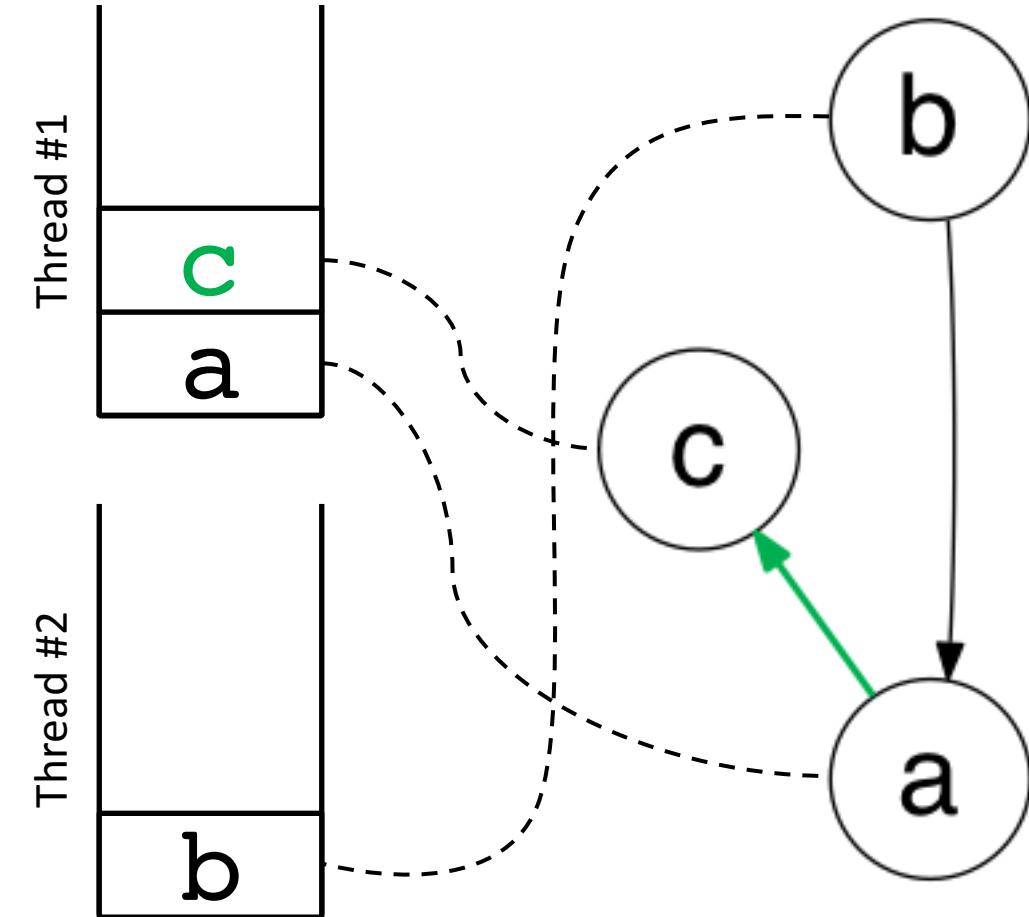
# Lock acquisition and release analysis (4)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
7 b.lock()	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	6 b.unlock()



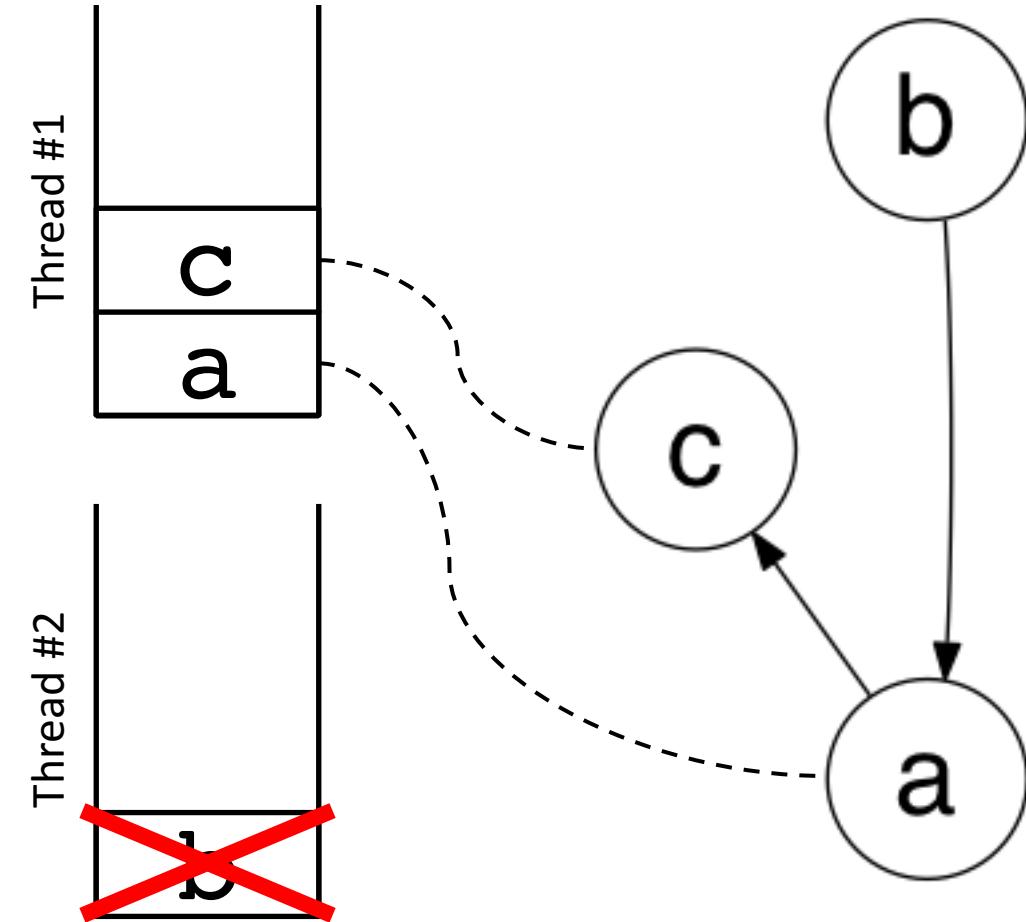
# Lock acquisition and release analysis (5)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	6 b.unlock()
7 b.lock()	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	



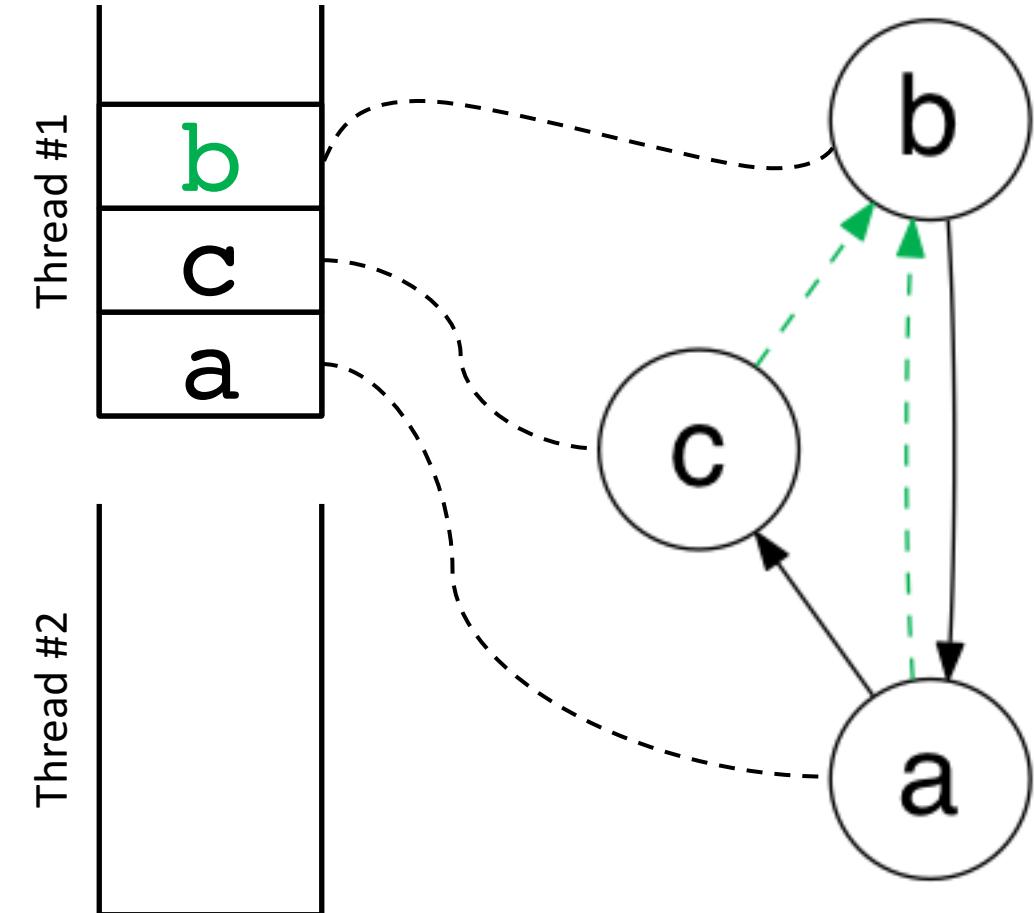
# Lock acquisition and release analysis (6)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
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5 c.lock()	
7 b.lock()	
8 a.unlock()	
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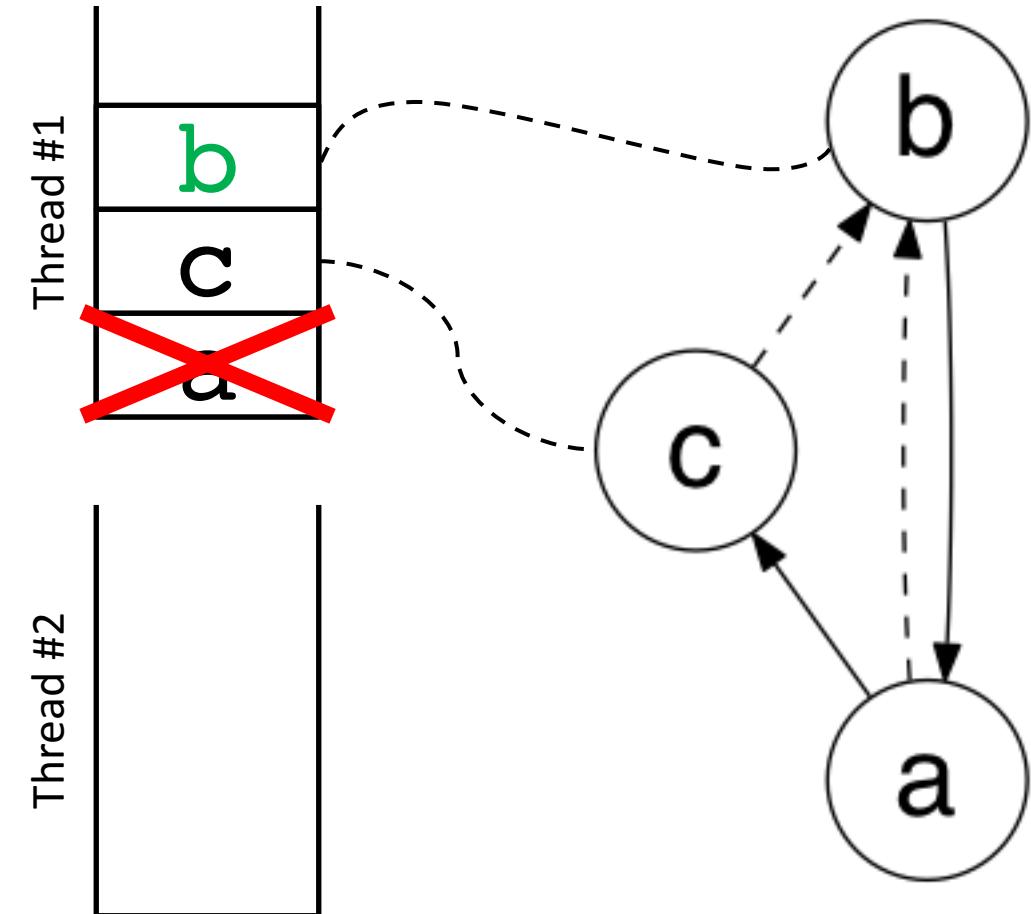
# Lock acquisition and release analysis (7)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
7 <b>b.lock()</b>	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	6 b.unlock()



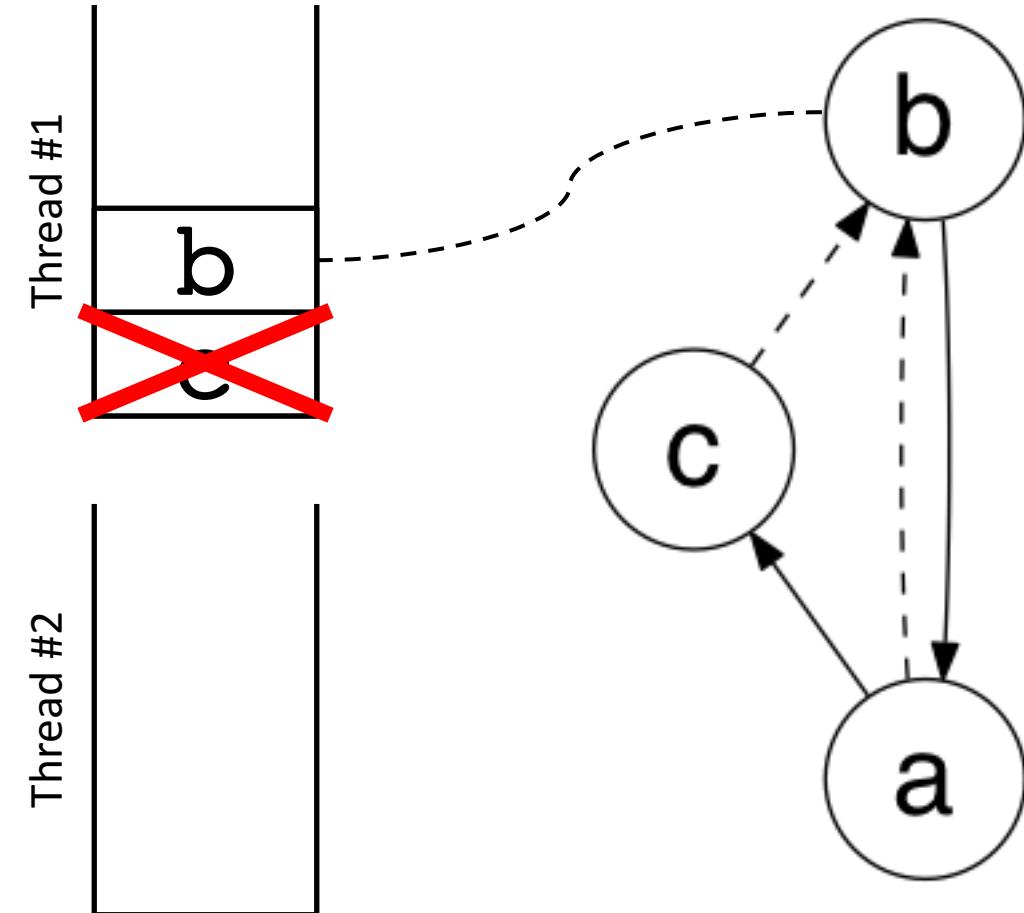
# Lock acquisition and release analysis (8)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
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5 c.lock()	
7 b.lock()	
8 <b>a.unlock()</b>	
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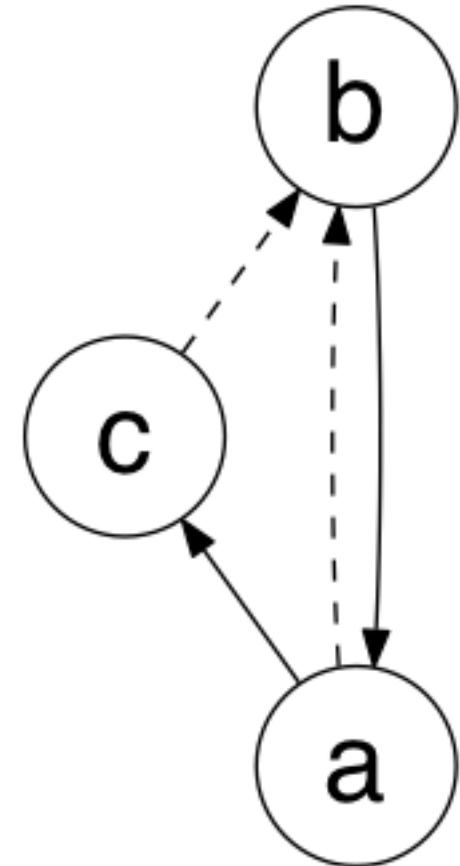
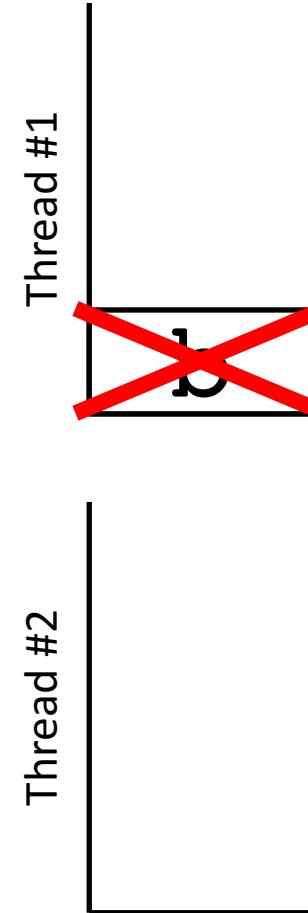
# Lock acquisition and release analysis (9)

Thread 1	Thread 2
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	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
	6 b.unlock()
7 b.lock()	
8 a.unlock()	
9 <b>c.unlock()</b>	
10 b.unlock()	



# Lock acquisition and release analysis (10)

Thread 1	Thread 2
	1 b.lock()
	2 a.lock()
	3 a.unlock()
4 a.lock()	
5 c.lock()	
	6 b.unlock()
7 b.lock()	
8 a.unlock()	
9 c.unlock()	
10 b.unlock()	



# Set of held locks

- Common lock usage patterns

```
synchronized(o) {  
    // Do something  
}  
  
synchronized void f() {  
    // Do something  
}
```

```
l.lock()  
try {  
    // Do something  
} finally {  
    l.unlock()  
}
```

- Locks are acquired and released in LIFO (stack) order

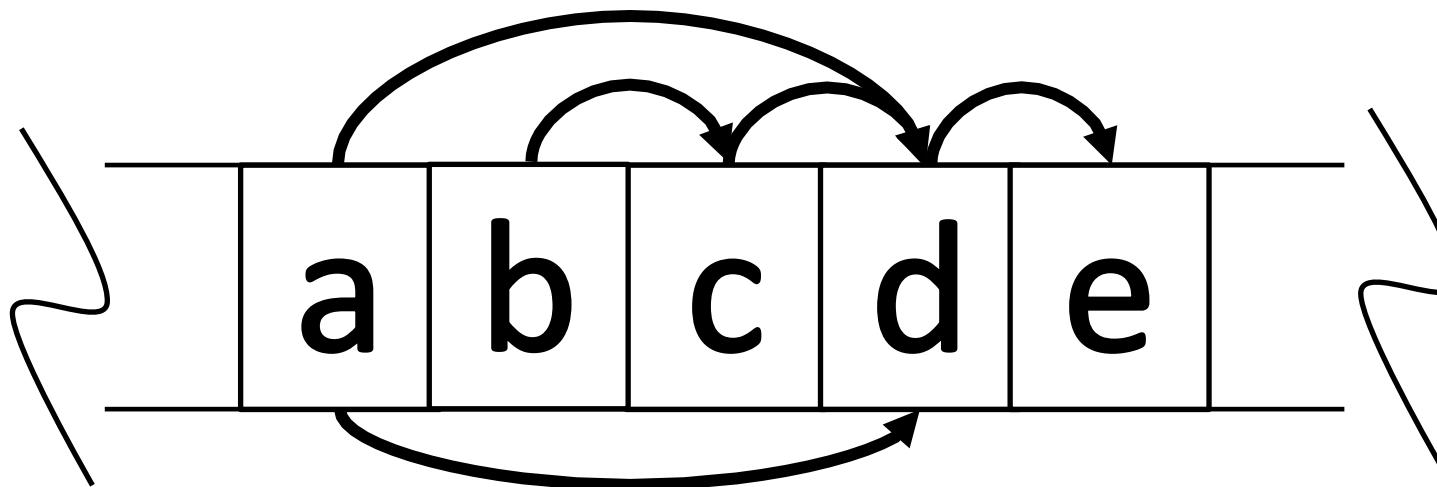


# Set of held locks

- Stack with possibility of removing from the middle
- Commonly works in  $O(1)$
- Works in  $O(B)$  at worst
  - $B$  – number of locks acquired by the current thread

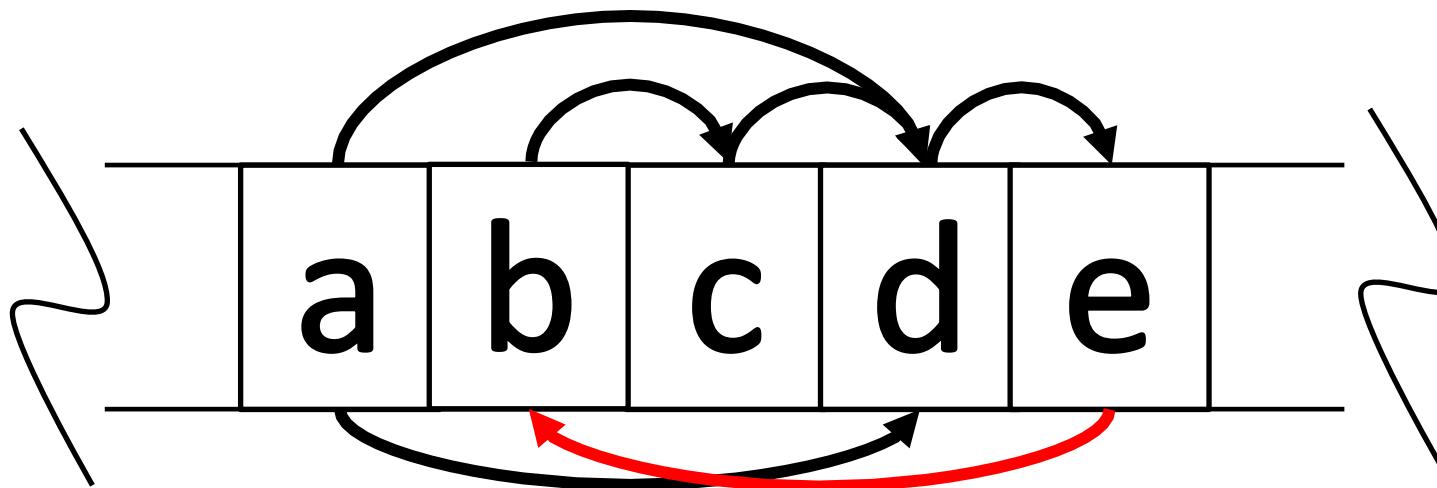
# Topological order

- Constructed by the edges of the directed graph
- No topological order  $\Leftrightarrow$  graph has a cycle



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- Constructed by the edges of the directed graph
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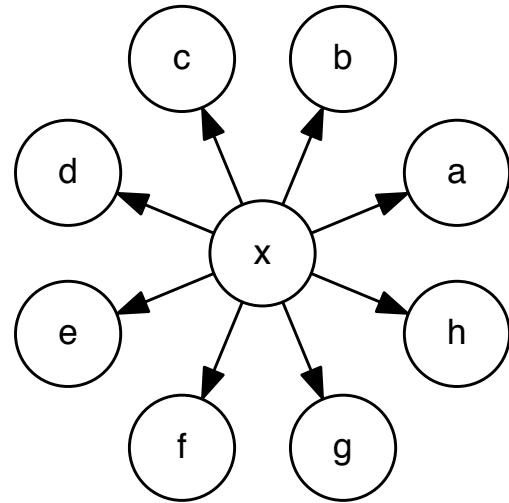


# Topological order

- The acyclic part of the lock graph is to be kept
- Edge  $(a, b)$  leads to a topological order violation  $\Rightarrow$  the shortest path  $b \rightsquigarrow a$  refers to the shortest cycle
- Incremental maintenance
  - MNR algorithm suggested by Marchetti-Spaccamela et al.

# New nodes processing

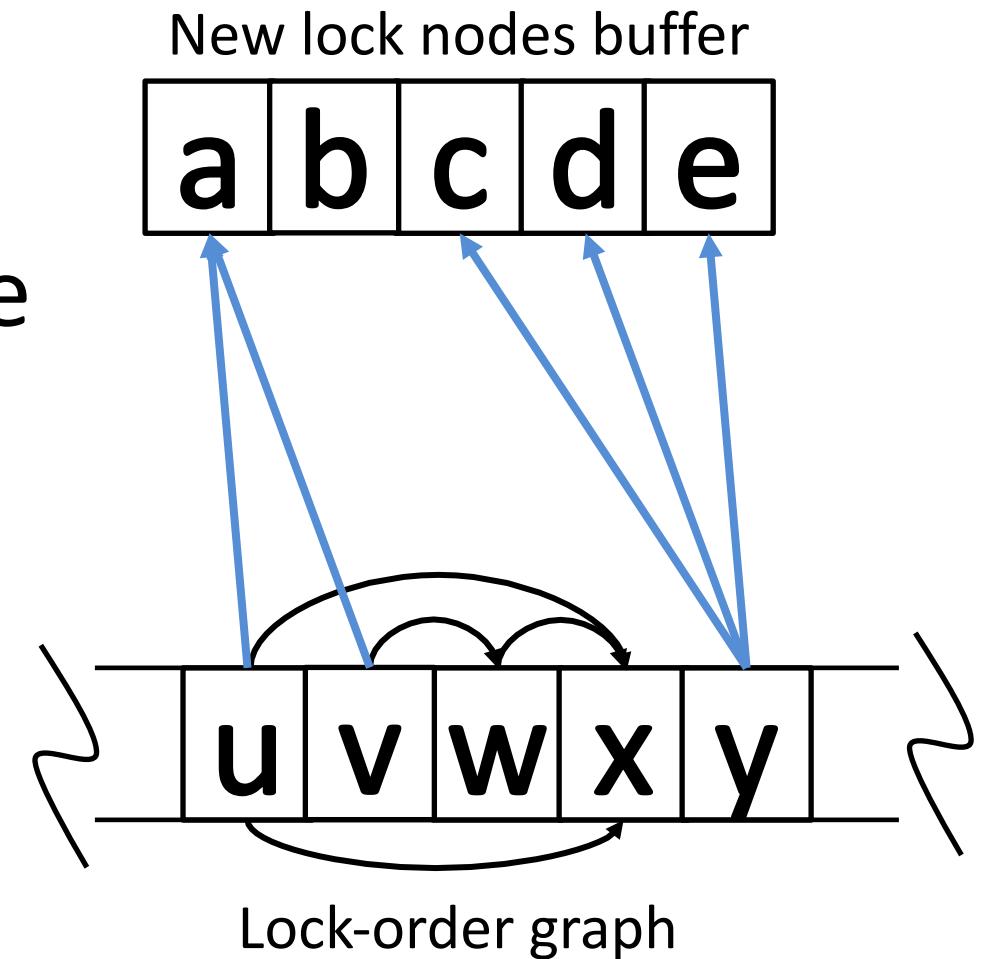
- Typical pattern in lock graphs



- One long-lived lock at the center and many hundreds of short-lived locks around it

# New nodes processing

- Lock-free buffer for new nodes
- Initialize topological order value
  - if buffer is full
  - when outgoing edge should be added



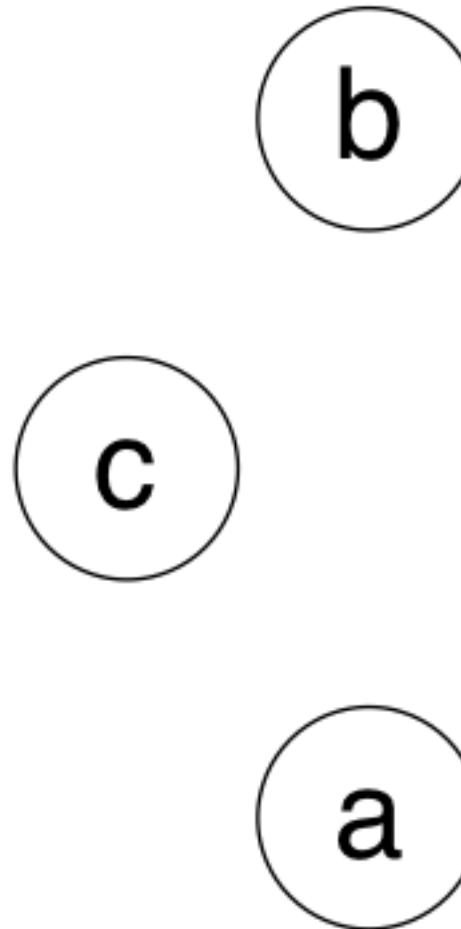
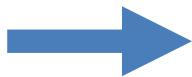
# Complexity

- Lock acquisition analysis
  - If acquisition does not produce new cycles
    - Typical case
    - $O(B + |V| + |E|)$  at worst
  - New cycle is detected
    - $O(B \cdot (|V| + |E|))$
- Lock release analysis
  - $O(B)$  at worst

# Limitation: cycle minimization

```
void transfer(Account from,  
             Account to,  
             int amount) {  
  
    from.lock()  
    to.lock()  
    // Do transfer  
    unlock(from, to)  
}
```

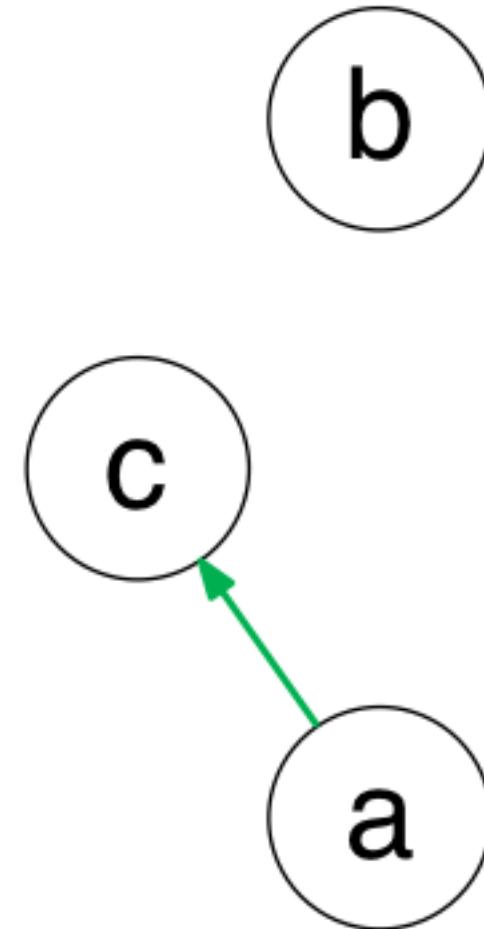
- 
- 1 transfer(a, c, ...)
  - 2 transfer(c, b, ...)
  - 3 transfer(b, a, ...)
  - 4 transfer(a, b, ...)



# Limitation: cycle minimization (1)

```
void transfer(Account from,  
             Account to,  
             int amount) {  
  
    from.lock()  
    to.lock()  
    // Do transfer  
    unlock(from, to)  
}
```

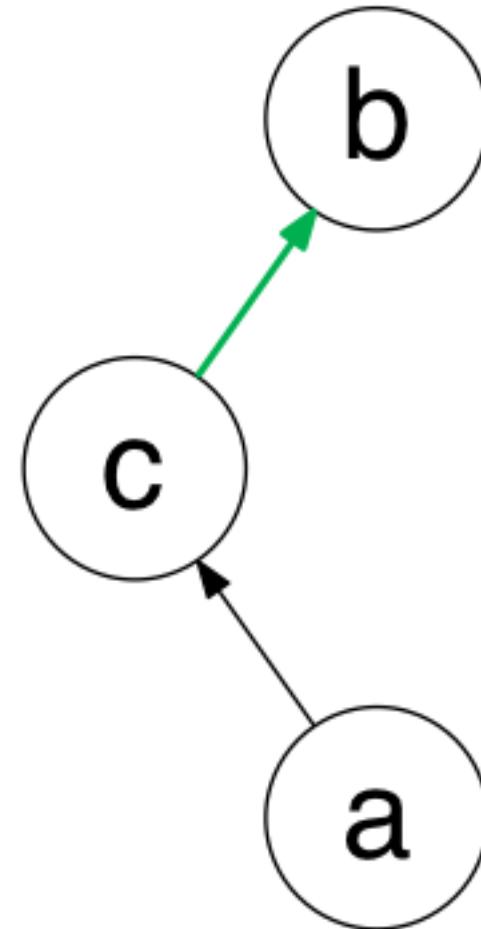
- 
- 1 transfer(a, c, ...)
  - 2 transfer(c, b, ...)
  - 3 transfer(b, a, ...)
  - 4 transfer(a, b, ...)



## Limitation: cycle minimization (2)

```
void transfer(Account from,  
             Account to,  
             int amount) {  
    from.lock()  
    to.lock()  
    // Do transfer  
    unlock(from, to)  
}
```

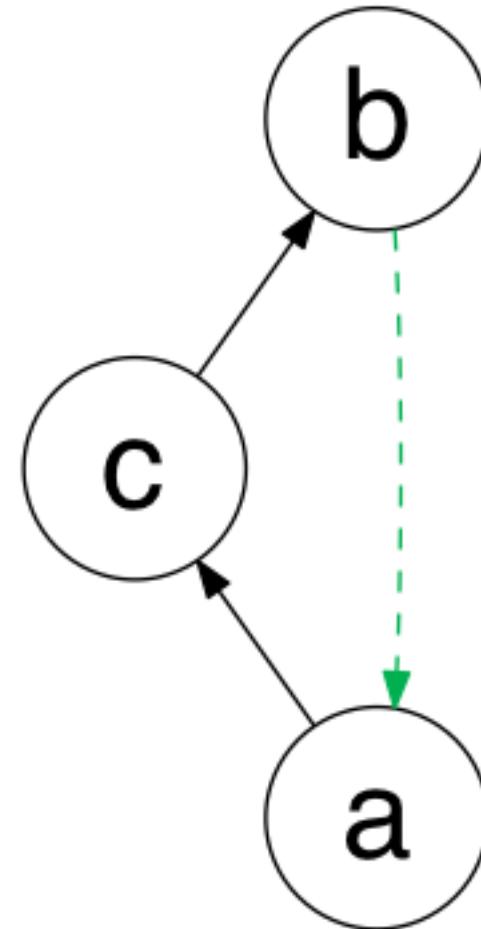
- 
- 1 transfer(a, c, ...)
  - 2 **transfer(c, b, ...)**
  - 3 transfer(b, a, ...)
  - 4 transfer(a, b, ...)



# Limitation: cycle minimization (3)

```
void transfer(Account from,  
             Account to,  
             int amount) {  
    from.lock()  
    to.lock()  
    // Do transfer  
    unlock(from, to)  
}
```

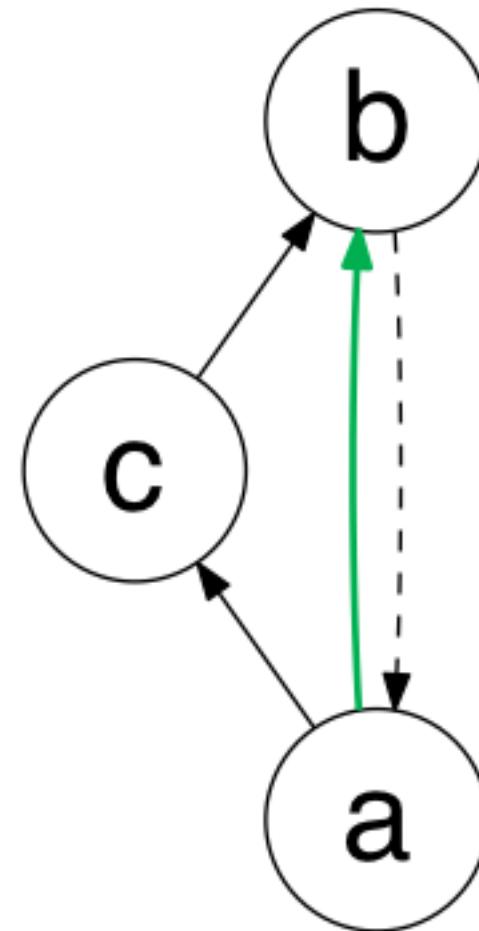
- 
- 1 transfer(a, c, ...)
  - 2 transfer(c, b, ...)
  - 3 **transfer(b, a, ...)**
  - 4 transfer(a, b, ...)



# Limitation: cycle minimization (4)

```
void transfer(Account from,  
             Account to,  
             int amount) {  
    from.lock()  
    to.lock()  
    // Do transfer  
    unlock(from, to)  
}
```

- 
- 1 transfer(a, c, ...)
  - 2 transfer(c, b, ...)
  - 3 transfer(b, a, ...)
  - 4 **transfer(a, b, ...)**



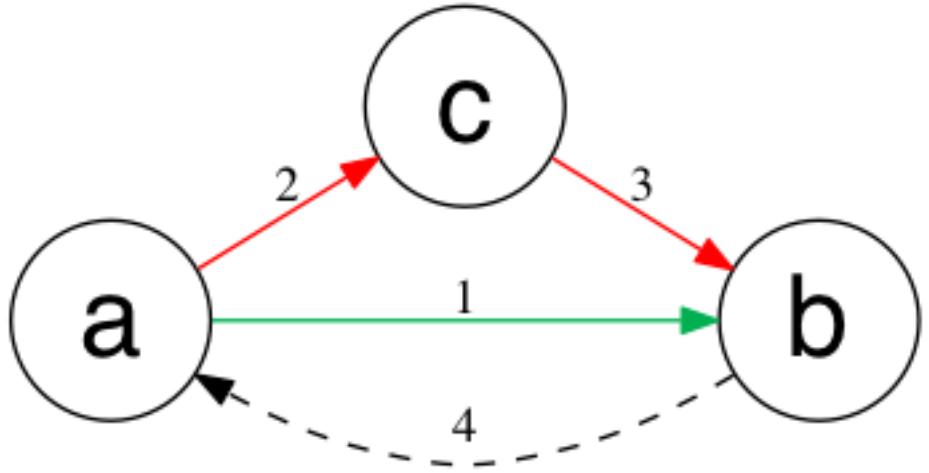
## Limitation: cycle minimization

- Adding edge  $(a, b)$  creates new cycle
- Minimization of  $\langle b, \dots, a \rangle \Leftrightarrow$  minimization of  $b \rightsquigarrow a$ 
  - Can be implemented with help of BFS
- Usually 2 or 3 vertexes in cycle
- Current implementation doesn't support minimization

# Limitation: independent cycles

- Logically independent cycles:

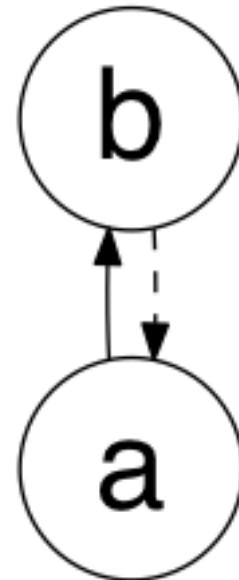
- $\langle a, c, b \rangle$
- $\langle a, b \rangle$



- Independent cycle is ignored due to minimization principle
- One useful cycle is better 😊

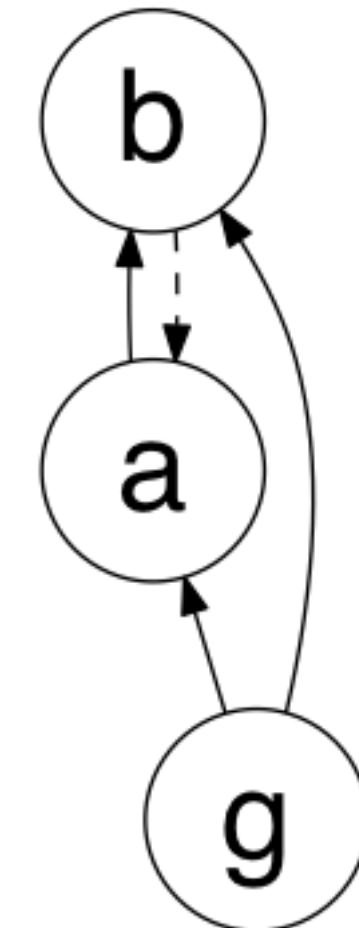
# Limitation: single threaded cycle

Lock a, b	
<hr/>	
1	a.lock()
2	b.lock()
3	unlock(a, b)
4	b.lock()
5	a.lock()
6	unlock(a, b)



# Limitation: guarded cycle

Thread 1	Thread 2
1 <b>g.lock()</b>	
2    a.lock()	
3    b.lock()	
4    unlock(a, b)	
5 <b>g.unlock()</b>	
	6 <b>g.lock()</b>
	7    b.lock()
	8    a.lock()
	9    unlock(a, b)
	10 <b>g.unlock()</b>



# IMPLEMENTATION

# DI-Check

- <https://github.com/Devexperts/dlcheck>
- Handy for testing
  - Throws an exception if `-Ddlcheck.fail=true`
  - Or write your own behavior
- Implemented as a Java agent
  - `java -javaagent:dlcheck.jar -jar your_app.jar`
  - Modifies byte-code within ASM framework



# Byte-code transformation

Lock acquisition and release analysis:

- First and last action in `synchronized` method
- `MONITORENTER` and `MONITOREXIT` instructions
- `java.util.concurrent.locks.Lock`
  - Method `lock`
  - Successful invocation of method `tryLock`
  - Method `unlock`

# Synchronized method: in simple terms

- Lock acquisition analysis
- Lock release analysis

```
synchronized void f() {  
    // Do something  
}
```

```
synchronized void f() {  
    onLockAcquire(this,  
        2 /*location id*/);  
    try {  
        // Do something  
    } finally {  
        onLockRelease(this);  
    }  
}
```

# Synchronized meth

What about...  
exceptions?

- Lock acquisition analysis
- Lock release analysis

`synchronized`

`// Do something`



```
    synchronized (...) {
        onLockAcquire(this,
                      2 /*location id*/);

        try {
            // Do something
        } finally {
            onLockRelease(this);
        }
    }
```

# Synchronized method: exception

WRONG

```
synchronized void f() {  
    onLockAcquire(this, 2);  
    try {  
        // Do something  
    } finally {  
        onLockRelease(this);  
    }  
}
```

CORRECT

```
synchronized void f() {  
    try {  
        onLockAcquire(this, 2);  
        // Do something  
    } finally {  
        onLockRelease(this);  
    }  
}
```

# Synchronized method: byte-code

- Lock acquisition analysis
- Lock release analysis

```
synchronized void g() {  
    try {  
        onLockAcquire(this, 2);  
        // Do something  
    } finally {  
        onLockRelease(this);  
    }  
}
```

```
synchronized void g();  
Code:  
0:  aload_0 // this  
1:  iconst_2 // location id  
2:  invokestatic #22 // onLockAcquire  
    // Do something  
5:  aload_0 // this  
6:  invokestatic #26 // onLockRelease  
9:  return  
10:  aload_0 // this  
11:  invokestatic #26 // onLockRelease  
14:  athrow  
Exception table:  
from  to   target  type  
  0    9      10  any
```

# Synchronized block

```
void f() {  
    synchronized (o) {  
        // Do something  
    }  
}  
  
void f() {  
    MONITORENTER o  
    try {  
        // Do something  
    } finally {  
        MONITOREXIT o  
    }  
}
```

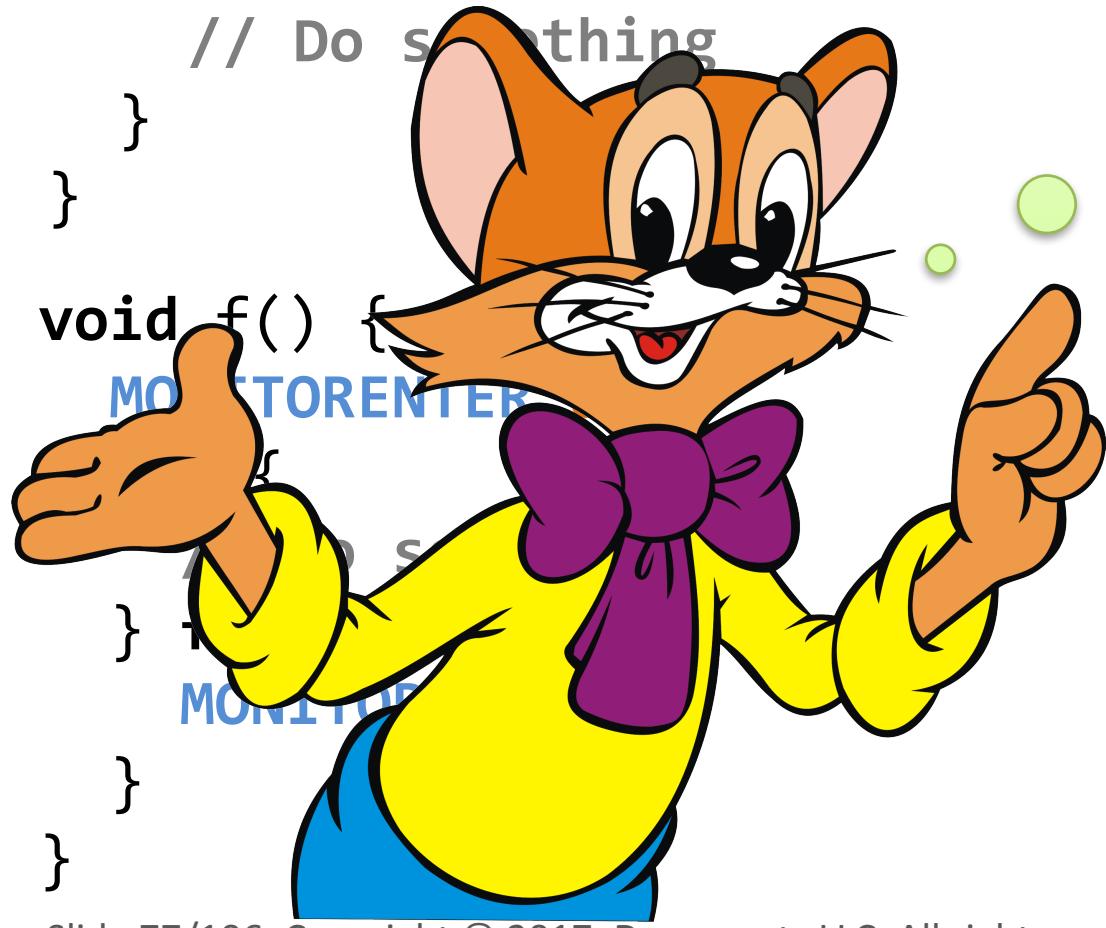
■ Lock acquisition analysis

■ Lock release analysis

```
void f() {  
    MONITORENTER o  
    try {  
        onLockAcquire(o, 42);  
        // Do something  
    } finally {  
        MONITOREXIT o  
        onLockRelease(o);  
    }  
}
```

# Synchronized block

```
void f() {  
    synchronized (o) {  
        // Do something  
    }  
}  
  
void f() {  
    MONITORENTER o  
    try {  
        onLockAcquire(o, 42);  
        // Do something  
    } finally {  
        MONITOREXIT o  
        onLockRelease(o);  
    }  
}
```



Not so simple!

```
void f() {  
    MONITORENTER o  
    try {  
        onLockAcquire(o, 42);  
        // Do something  
    } finally {  
        MONITOREXIT o  
        onLockRelease(o);  
    }  
}
```

# Synchronized block: backstage

```
void f() {  
    synchronized (o) {  
        // Do something  
    }  
}
```

Exception table:

from	to	target	type
7	9	12	any
12	15	12	any

```
void f();  
Code:  
try {  
    0: aload_0 // this  
    1: getfield #2 // monitor  
    4: dup  
    5: astore_1 // store monitor  
    6: monitorenter  
        // Do something  
    7: aload_1 // monitor  
    8: monitorexit  
    9: goto 17 // skip finally  
   12: astore_2 // exception  
   13: aload_1 // monitor  
   14: monitorexit  
   15: aload_2 // exception  
   16: athrow  
  17: return
```

# Synchronized block: exception

WRONG

```
void f() {  
    MONITORENTER o  
    try {  
        onLockAcquire(o, 42);  
        // Do something  
    } finally {  
        MONITOREXIT o  
        onLockRelease(o);  
    }  
}
```

CORRECT

```
void f() {  
    MONITORENTER o  
    try {  
        onLockAcquire(o, 42);  
        // Do something  
    } finally {  
        MONITOREXIT o  
        try {  
            onLockRelease(o);  
        catch (Throwable t) {}  
    }  
}
```



# Synchronized block: not a block

```
void f() {  
    MONITORENTER o  
    try {  
        // Do something  
    } finally {  
        MONITOREXIT o  
    }  
}
```

```
void f() {  
    MONITORENTER o  
    try {  
        onLockAcquire(o, 42);  
    } catch (Throwable t) {  
        MONITOREXIT o  
        onLockRelease(o);  
        throw t;  
    }  
    try {  
        // Do something  
        ...  
    }
```

Add try-catch

# j.u.c.l.Lock: common usages

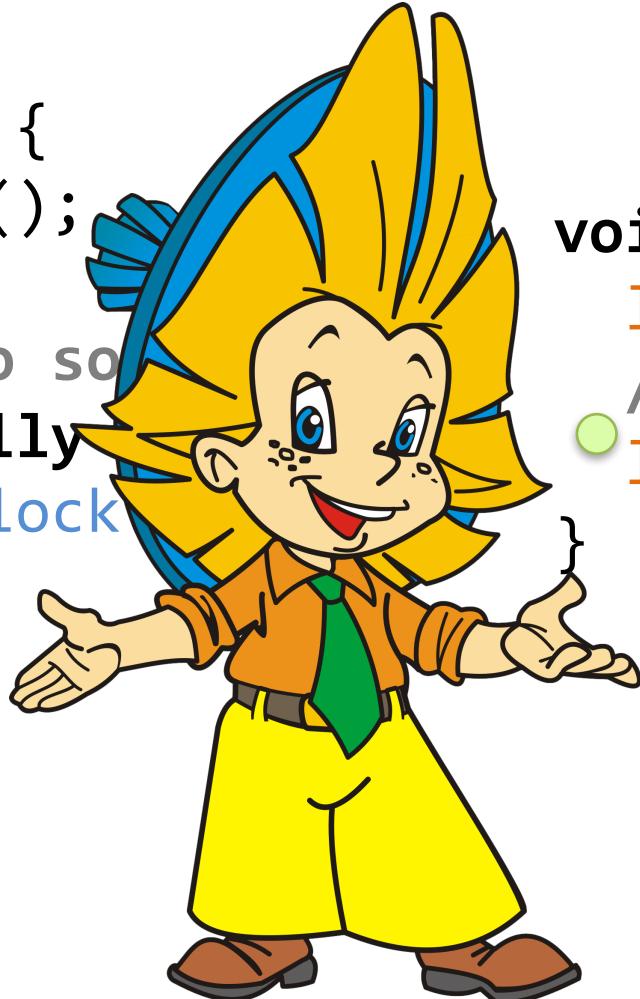
```
void f() {  
    l.lock();  
    try {  
        // Do something  
    } finally {  
        l.unlock();  
    }  
}
```

```
void g() {  
    l.lock();  
    // Do something  
    l.unlock();  
}
```

```
void h() {  
    b = l.tryLock();  
    // Do something  
    l.unlock();  
}
```

# j.u.c.l.Lock: general usages

```
void f() {  
    l.lock();  
    try {  
        // Do something  
    } finally {  
        l.unlock()  
    }  
}
```



```
void g() {  
    l.lock();  
    // Do something  
    l.unlock();  
}
```

It looks like unpaired  
**MONITORENTER**  
and **MONITOREXIT!**

```
tryLock();  
,, Do something  
l.unlock();  
}  
}
```

# j.u.c.l.Lock: byte-code

```
void g() {  
    l.lock();  
    // Do something  
    l.unlock();  
}
```

Code:

```
void g();  
Code:  
0: aload_0 // this  
1: getfield #3 // Lock 1  
4: invokeinterface #4, 1 // lock()  
   // Do something  
9: aload_0 // this  
10: getfield #3 // Lock 1  
13: invokeinterface #5, 1 // unlock()  
18: return
```

# j.u.c.l.Lock: tryLock

```
void h() {  
    b = l.tryLock();  
    // Do something  
    l.unlock();  
}
```



```
void h() {  
    b = l.tryLock();  
    if (b) {  
        try {  
            onLockAcquire(l, 2);  
        } catch (Throwable t) {  
            l.unlock();  
            onLockRelease(l);  
            throw t;  
        }  
        // Do something  
        l.unlock();  
        onLockRelease(l);  
    }  
}
```

# ClassTransformer

Code template:

```
public final byte[] transform(..., classFileBuffer) {  
    ClassReader cr = new ClassReader(classFileBuffer);  
    ClassWriter cw = new ClassWriter(...);  
    ClassVisitor classTransformer = new MyTransformer(cw, ...);  
    cr.accept(classTransformer, EXPAND_FRAMES);  
    return cw.toByteArray();  
}
```

Let's try to run...



# ClassTransformer: LinkageError

**java.lang.LinkageError**: loader (instance of sun/misc/Launcher\$AppClassLoader): attempted **duplicate class definition** for name: "com/intellij/util/lang/UrlClassLoader"  
at java.lang.ClassLoader.defineClass1(Native Method)  
at java.lang.ClassLoader.defineClass(ClassLoader.java:763)  
at java.security.SecureClassLoader.defineClass(SecureClassLoader.java:142)  
at java.net.URLClassLoader.defineClass(URLClassLoader.java:467)  
at java.net.URLClassLoader.access\$100(URLClassLoader.java:73)  
at java.net.URLClassLoader\$1.run(URLClassLoader.java:368)  
at java.net.URLClassLoader\$1.run(URLClassLoader.java:362)  
at java.security.AccessController.doPrivileged(Native Method)  
at java.net.URLClassLoader.findClass(URLClassLoader.java:361)  
at java.lang.ClassLoader.loadClass(ClassLoader.java:424)  
at sun.misc.Launcher\$AppClassLoader.loadClass(Launcher.java:331)  
at java.lang.ClassLoader.loadClass(ClassLoader.java:357)  
...



# ClassTransformer: LinkageError

```
java.lang.LinkageError: loader found multiple incompatible class definitions while loading / instance of sun/misc/Launcher$AppClassLoader for name: "com/intel/asm/ClassWriter" at java.lang.ClassLoader.defineClass1(Native Method) at java.lang.ClassLoader.defineClass(ClassLoader.java:763) at java.security.SecureClassLoader.defineClass(SecureClassLoader.java:142) at java.net.URLClassLoader.findClass(URLClassLoader.java:467) at java.net.URLClassLoader.access$100(URLClassLoader.java:72) at java.net.URLClassLoader$1.run(URLClassLoader.java:365) at java.net.URLClassLoader$1.run(URLClassLoader.java:363) at java.security.AccessController.doPrivileged(Native Method) at java.net.URLClassLoader.findClass(URLClassLoader.java:361) at java.lang.ClassLoader.loadClass(ClassLoader.java:357) at sun.misc.Launcher$AppClassLoader.loadClass(Launcher.java:357) at java.lang.ClassLoader.loadClass(ClassLoader.java:357) ...
```

F\*\*\*ing  
ASM!



# Problem ClassWriter

```
protected String getCommonSuperClass(String type1,  
                                    String type2) {  
    Class<?> c1, c2;  
    ClassLoader loader = getClass().getClassLoader();  
    try {  
        c1 = Class.forName(type1.replace('/', '.'), false, loader);  
        c2 = Class.forName(type2.replace('/', '.'), false, loader);  
    } catch (Exception e) {  
        throw new RuntimeException(e);  
    }  
    if (c1.isAssignableFrom(c2))  
        return type1;  
    ...  
}
```

# Write it yourself!

```
@Override  
protected String getCommonSuperClass(String type1,  
                                     String type2) {  
    ClassInfo c = ciCache.getOrBuildClassInfo(type1, loader);  
    ClassInfo d = ciCache.getOrBuildClassInfo(type2, loader);  
    if (c.isAssignableFrom(d, ciCache, loader))  
        return type1;  
    if (d.isAssignableFrom(c, ciCache, loader))  
        return type2;  
    ...
```

ClassInfo is available BEFORE transformation





**Hooray! It works!**

# JAgent framework

- This problem is solved in JAgent
  - <https://github.com/Devexperts/jagent>
- Makes java agents writing more simple

# Dependency conflict

- Several agents are used:

```
java -javaagent:agent1.jar -javaagent:agent2.jar  
-jar my_app.jar
```



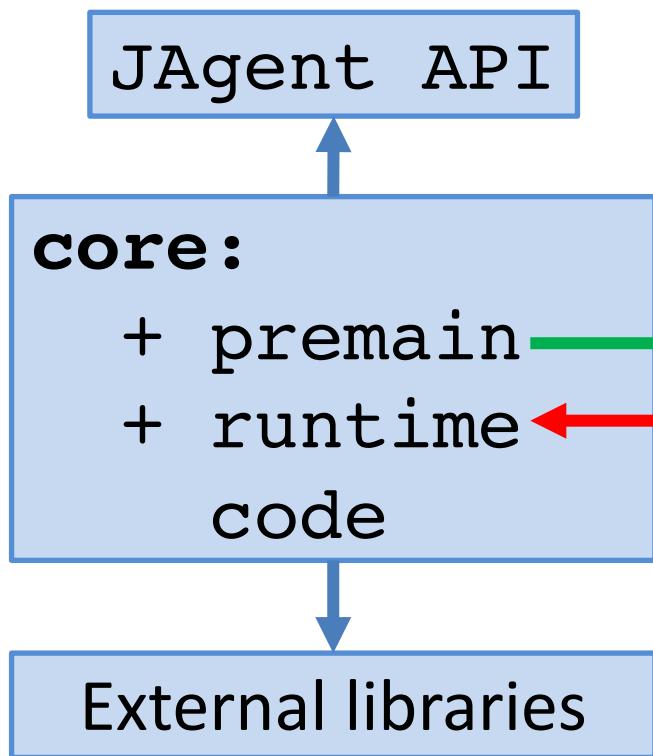
- Dependencies should be isolated

# Isolation techniques

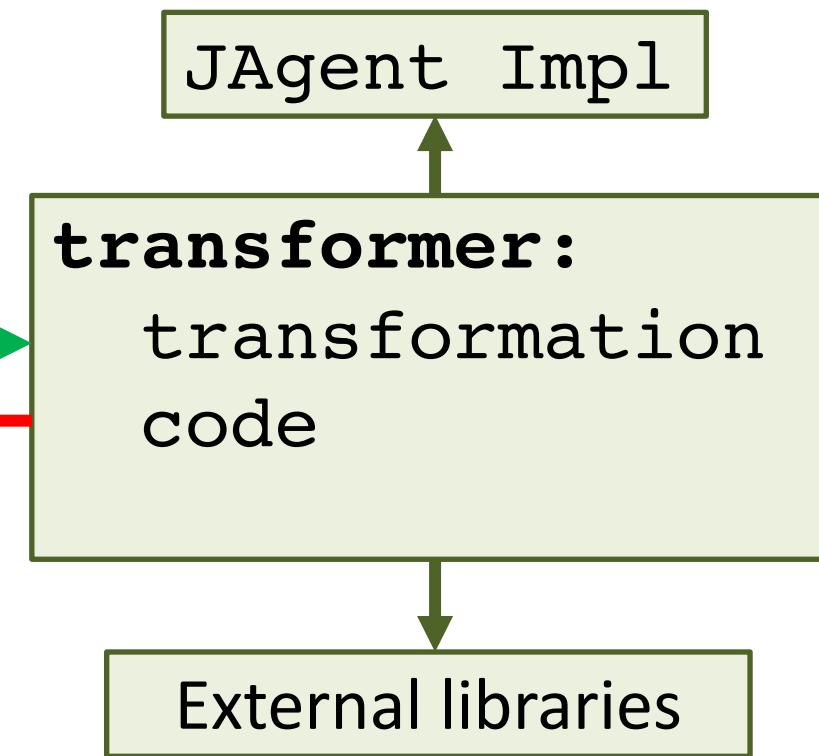
- Repackaging
  - maven-shade-plugin (or analogue)
  - Needs to know transitive dependencies and their structure
  - Sometimes doesn't work as well (reflection)
- Separate ClassLoader
  - Cannot be used for injected code

# JAgent will help us

## System ClassLoader



## JAgent ClassLoader

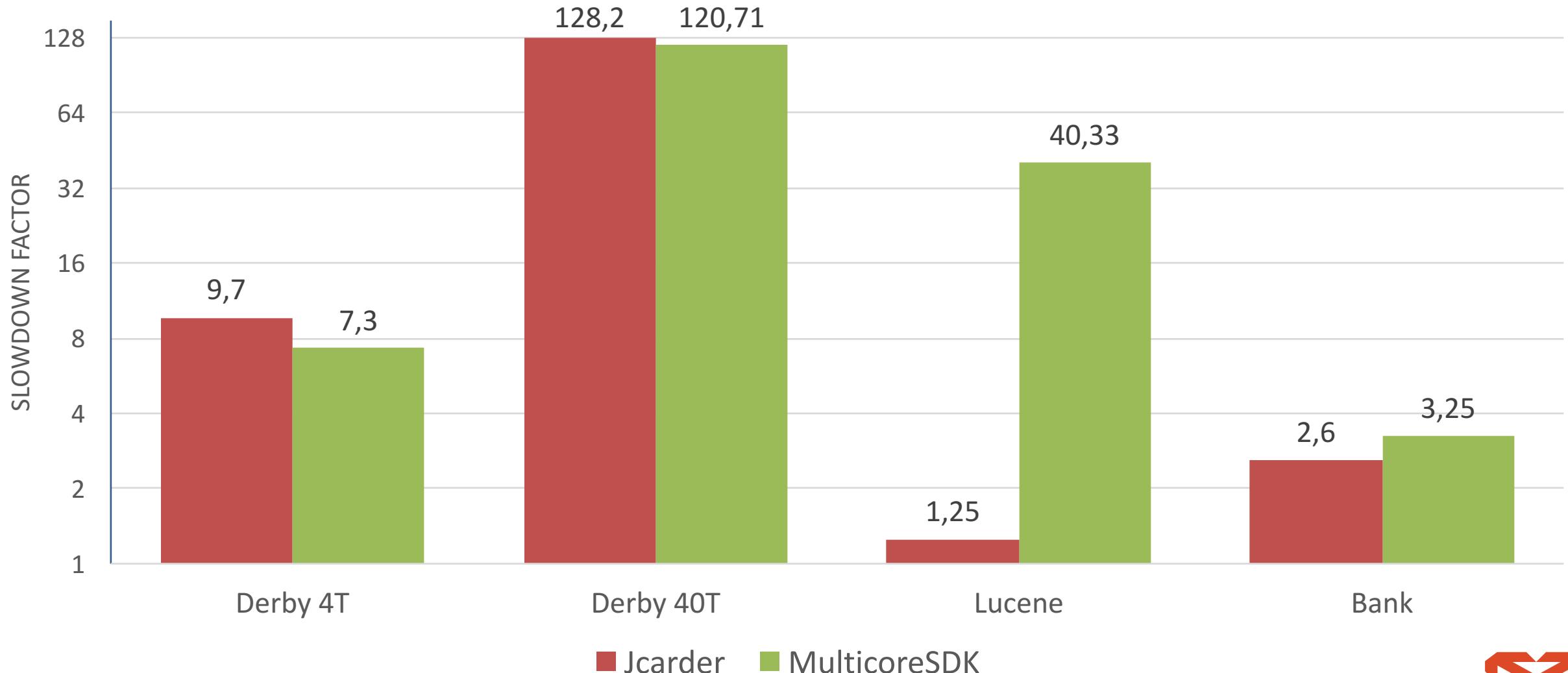


# EVALUATION

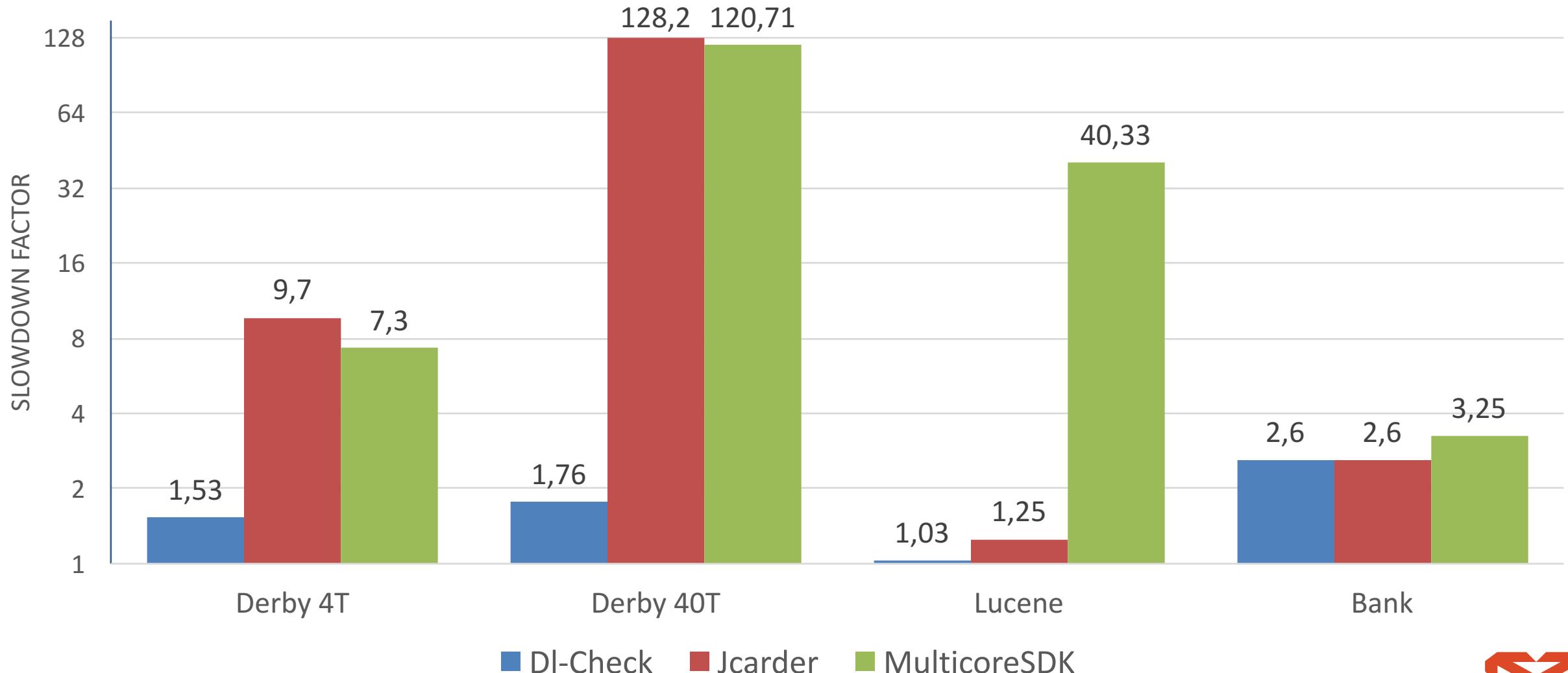
# Benchmarks

- SpecJVM2008: Apache Derby (4 and 40 threads)
- DaCapo: Apache Lucene
- DaCapo: Banking application
- ...

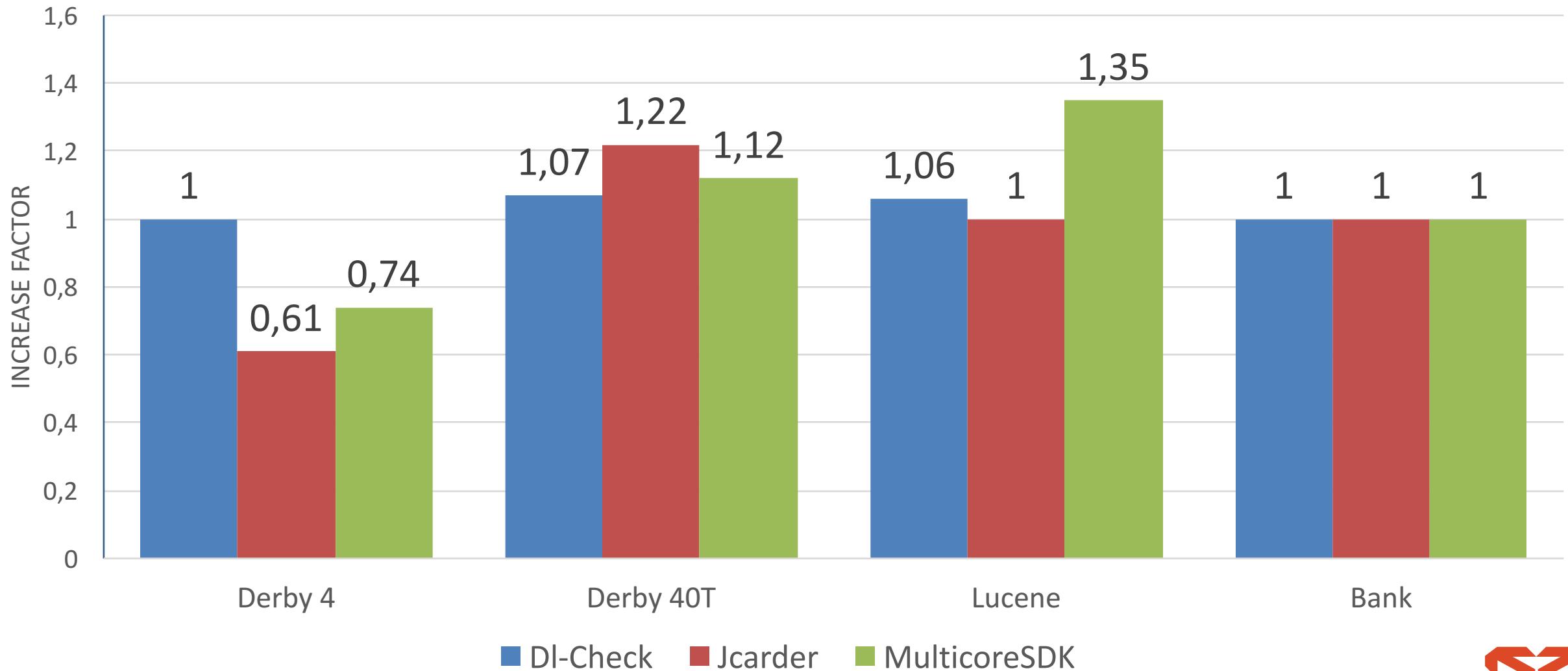
# Performance impact



# Performance impact



# Memory usage impact



# Real-life example

```
public void testClose() {  
    cm.register(con);  
    // lock cm  
    // lock con for con.onRegister()  
    assertEquals(1, cm.getConCount());  
    con.close();  
    // lock con  
    // lock cm for cm.unregister()  
    assertEquals(0, cm.getConCount());  
}
```



=====

!!! Potential deadlock !!!

=====

### Cycle in lock graph: ###

Lock **ConnectionImpl@d70c109** was acquired at:

- a.ConnectionImpl.onRegister([Sample.java:39](#))
- a.ConnectionImpl.close([Sample.java:44](#))

Lock **ConnectionManager@2be94b0f** was acquired at:

- a.ConnectionManager.getConCount([Sample.java:24](#))
- a.ConnectionManager.register([Sample.java:15](#))
- a.ConnectionManager.unregister([Sample.java:20](#))

### Current lock stack: ###

Lock **ConnectionImpl@d70c109** was acquired at:

- a.ConnectionImpl.onRegister([Sample.java:39](#))
- a.ConnectionImpl.close([Sample.java:44](#))

### Current stacktrace: ###

- a.ConnectionManager.unregister([Sample.java:20](#))
- a.ConnectionImpl.close([Sample.java:45](#))
- a.FunctionalityTest.testClose([FunctionalityTest.java:22](#))

...

# Evaluation in real projects

- Internal projects
- IntelliJ IDEA
- Kotlin
- ...

# Summary

- On-line algorithm for detecting potential deadlocks
- Dl-Check tool has been implemented
  - <https://github.com/Devexperts/dlcheck>
- Further work:
  - Contracts to describe lock acquisition rules

# Thank you for attention!

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