

## How to Calculate 'Algorithm Time Complexity' ?

✓ Iterative Algorithm

✓ Recursive Algorithm

## Example#1: Time Complexity of 'Iterative Algo'



5	18	3	54	26	...	55	41	...	19	1	10
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*FindBiggestNumber* (int arr[]):

biggestNumber = arr[0]

loop: i = 1 to length(arr)-1

if arr[i] > biggestNumber

biggestNumber = arr[i]

return biggestNumber

## Example#1: Time Complexity of 'Iterative Algo'(continued)

FindBiggestNumber (int arr[]):

biggestNumber = arr[0] -----  $O(1)$  ✓  
loop: i = 1 to length(arr)-1 -----  $O(n)$   
    if arr[i] > biggestNumber -----  $O(1)$   
        biggestNumber = arr[i] -----  $O(1)$   
return biggestNumber -----  $O(1)$

$$O(n-1) = O(n)$$
$$O(2 \times n) = O(n)$$

$$O(10) = O(1)$$

$$O(10000) = O(1)$$

$$13 + O(1)$$

$n$

$$\text{Time Complexity} = \underline{O(1)} + \underline{O(n)} + \underline{O(1)}$$
$$= \underline{O(n)} + \underline{O(2)}$$

## Example#2: Time Complexity of 'Recursive Algo'

5	18	3	54	26	...	55	41	...	19	1	10
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*FindBiggestNumber(A, n):*

*static highest = Integer.MIN*

*if n equals -1*

*return highest*

*else*

*if A[n] > highest*

*update highest*

*return FindBiggestNumber(A, n-1)*

## Example#2: Time Complexity of 'Recursive Algo' (Continued)

```
FindBiggestNumber(A,n): -----  $T(n)$ 
    static highest = Integer.Min -----  $O(1)$ 
    if (n equals -1) -----  $O(1)$  ✓
    return highest -----  $O(1)$  ✓
    else -----  $O(1)$ 
        if  $A[n] > \text{highest}$  -----  $O(1)$ 
        update highest -----  $O(1)$ 
    return FindBiggestNumber(A, n-1) -----  $T(n-1)$ 
```

### Back Substitution:

$$T(n) = O(1) + T(n-1) \text{ ----- Equation\#1} \quad \checkmark$$

$$T(-1) = O(1) \text{ ----- Base Condition} \quad \checkmark$$

$$T(n-1) = O(1) + T((n-1)-1) \text{ ----- Equation\#2} \quad \checkmark$$

$$T(n-2) = O(1) + T((n-2)-1) \text{ ----- Equation\#3} \quad \checkmark$$

## Example#2: Time Complexity of 'Recursive Algo' (continued)

### Back Substitution:

$$T(n) = O(1) + T(n-1) \text{ ----- Equation\#1 } \checkmark$$

$$T(-1) = O(1) \text{ ----- Base Condition}$$

$$T(n-1) = \underline{O(1)} + T((n-1)-1) \text{ ----- Equation\#2}$$

$$T(n-2) = \underline{O(1)} + T((n-2)-1) \text{ ----- Equation\#3}$$

$$T(n) = 1 + \underline{T(n-1)} \checkmark$$

$$= 1 + (1 + T((n-1)-1))$$

$$= 2 + T(n-2)$$

$$= 2 + 1 + T((n-2)-1)$$

$$= 3 + T(n-3)$$

$$= k + T(n-k)$$

## Example#2: Time Complexity of 'Recursive Algo' (continued)

### Back Substitution:

$$T(n) = O(1) + T(n-1) \text{ ----- Equation\#1}$$

$$T(-1) = O(1) \text{ ----- Base Condition}$$

$$T(n-1) = O(1) + T((n-1)-1) \text{ ----- Equation\#2}$$

$$T(n-2) = O(1) + T((n-2)-1) \text{ ----- Equation\#3}$$

$$T(n) = 1 + T(n-1)$$

$$= 1 + (1 + T((n-1)-1))$$

$$= 2 + T(n-2)$$

$$= 2 + 1 + T((n-2)-1)$$

$$= 3 + T(n-3)$$

$$= k + T(n-k)$$

$$= (n+1) + T(n-(n+1))$$

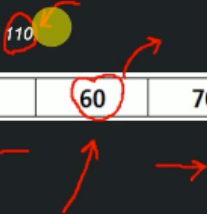
$$= n+1 + T(-1)$$

$$= n+1 + 1$$

$$= O(n)$$

### Example#3: Time Complexity of 'Recursive Algo'

✓ Problem Statement: Given a sorted array of 11 numbers, find number 110



10	20	30	40	50	60	70	80	90	100	110
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✓ Solution:



### Example#3: Time Complexity of 'Recursive Algo' (continued)

*BinarySearch*(int findNumber, int arr[], start, end): .....  $T(n)$  <sup>100</sup> <sup>10</sup>

if (start equals end) .....  $O(1)$  ✓

if (arr[start] equals findNumber) .....  $O(1)$  ✓

return start .....  $O(1)$  ✓

else return error message that number does not exists in the array .....  $O(1)$  ✓

mid = FindMid(arr[], start, end) .....  $O(1)$  ✓

if mid > findNumber .....  $O(1)$  ✓

*BinarySearch*(int findNumber, int arr[], start, mid) .....  $T(n/2)$  ✓ → SO ⇒  $\frac{10}{2} = 5$

else if mid < findNumber .....  $O(1)$  ✓

*BinarySearch*(int findNumber, int arr[], mid, end) .....  $T(n/2)$  ✓

else if mid = findNumber .....  $O(1)$  ✓

return mid .....  $O(1)$  ✓

Time Complexity:

$$T(n) = O(1) + T(n/2)$$

⊕  $T(n/2)$

### Example#3: Time Complexity of 'Recursive Algo' (continued):

#### Back Substitution:

$$T(n) = T(n/2) + 1 \text{ ----- Equation\#1}$$

$$T(1) = 1 \text{ ----- Base Condition}$$

$$T(n/2) = T(n/4) + 1 \text{ ----- Equation\#2}$$

$$T(n/4) = T(n/8) + 1 \text{ ----- Equation\#3}$$

$T(1) = 1$ , Hence  
so the for what value of  $K \rightarrow T(n/2^K)$   
becomes  $T(1)$

$$T(n) = T(n/2^K) + 1$$

$$= (T(n/4) + 1) + 1$$

$$= T(n/4) + 2$$

$$= (T(n/8) + 1) + 2$$

$$= T(n/8) + 3$$

$$= T(n/2^K) + k$$

$$= T(1) + \log n$$

$$= 1 + \log n$$

$$= \log n$$

$1 - n$

$$\frac{n}{2^K} = 1$$
$$n = 2^K$$
$$K = \log n$$