

Graph 4. Node Embedding

- Laplacian Eigenmaps (NIPS 2002)
 - Random Walk (KDD 2014)
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1

Lectures on Graph-based ML

- ✓ Graph 1-3. Social network analysis (e.g., HITS & PageRank)
- **Graph 4. Node embedding (Laplacian eigenmaps vs. random walk methods)**
- Graph 5-6. Graph neural networks (GCN, GAT, GIN, etc.)
- Graph 7-8. Knowledge graph embedding
- Graph 9-11. Neural solvers for combinatorial optimization (AR, NAR, LLM's)
- Graph 12-13. Graph-based learning for recommender systems (invited talks)
- Graph 14-15. Learning with heterogeneous graphs

2

Motivation for Graph Node Embedding

- GloVe co-occurrence graph → **word embedding** → NLP
- Hyperlinked websites → **page embedding** → websites classification
- Citation graphs → **article embedding** → literature classification
- Co-author graphs → **author embedding** → community detection
- Molecular structure graph → **atom embedding** → AI for science
- ...

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3

3

Unified View

- **Nodes** can be any objects (words, documents, authors, atoms, etc.)
- **Links** represent the interactions or dependencies among nodes.
- **Embedding Vectors**
 - Capturing the latent features of nodes **based on graph structures**
 - Supporting down-stream prediction tasks (node/graph classification, community detection, dense retrieval, etc.)

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4

4

Node Embedding Methods

- Based on ~~linked structures only~~ (this lecture)
 - Matrix-factorization based methods (e.g., Laplacian Eigenmaps)
 - Random-walk based methods
- Based on both linked structures and node-specific features (next lecture)

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5

5

Input and Output

- Input Graph $G = (V, A)$
 - V is the set of n nodes in the graph;
 - $A \in \mathbb{R}^{n \times n}$ is the adjacency matrix with $A_{ij} \geq 0$ and $A_{ii} = 0$ (zero at the diagonal).



Specifically, we focus on **undirected** graphs with $A_{ij} = A_{ji}$ (**symmetric**)

- Output Matrix $Z \in \mathbb{R}^{n \times d}$ ($d < n$)
 - Each row z_i is the embedding of a node;
 - Each column z_j is a feature (latent factor) of the embedding space.

Matrix Z

$$\begin{matrix}
 & \begin{matrix} z_{11} & z_{12} & \cdots & z_{1d} \end{matrix} \\
 \begin{matrix} z_i \\ z_{21} \\ \vdots \\ z_{n1} \end{matrix} & \begin{bmatrix} z_{12} & \cdots & z_{1d} \\ z_{22} & \cdots & z_{2d} \\ \vdots & \ddots & \vdots \\ z_{n2} & \cdots & z_{nd} \end{bmatrix}
 \end{matrix}$$

z_j

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6

6

Graph Laplacian Matrices

- Default Version

$$L \stackrel{\text{def}}{=} D - A \quad \text{with } D \stackrel{\text{def}}{=} \text{diag}(D_{11}, \dots, D_{nn}) \text{ and } D_{ii} = \sum_j A_{ij}$$

- Symmetric Version

$$L_{\text{sym}} \stackrel{\text{def}}{=} D^{-\frac{1}{2}} L D^{-\frac{1}{2}} = \overbrace{D^{-\frac{1}{2}} D D^{-\frac{1}{2}}}^{I_{n \times n}} - \overbrace{D^{-\frac{1}{2}} A D^{-\frac{1}{2}}}^{A_{\text{sym}}} \quad \left(A_{\text{sym}}[i, j] = \frac{A_{ij}}{\sqrt{D_{ii}} \sqrt{D_{jj}}} \right)$$

- Random-walk Version

$$L_{rw} \stackrel{\text{def}}{=} D^{-1} L = \overbrace{D^{-1} D}^{I_{n \times n}} - \overbrace{D^{-1} A}^{A_{rw}} \quad \left(A_{rw}[i, j] = \frac{A_{ij}}{D_{ii}}, \sum_j A_{rw}[i, j] = 1 \right)$$

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7

7

Properties of Graph Laplacian L


(Von Luxburg, 2007, <https://arxiv.org/pdf/0711.0189.pdf>)

- When A is symmetric, $L = D - A$ is also symmetric (obvious).

- L allows us to reinforce the smoothness of node embedding. 

- L is positive semi-definite. $z^T L z$

- L has n real-valued non-negative eigenvalues.

-  $\lambda_1 = 0$ is the smallest eigenvalue of L , and $\mathbf{1}_n$ as the corresponding eigenvector.

- L has a complete set of n eigenvectors (proof omitted).

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8

8

Laplacian eigenmaps and spectral techniques for embedding and clustering [M. Belkin and P. Niyogi. NIPS 2002]

- **Task:** Given graph $G = (V, A)$, find the optimal solution

$$Z^* = \arg \min_{Z \in \mathbb{R}^{n \times d}} |z_i - z_j|^2 A_{i,j} \quad \leftarrow \text{Smoothness Penalty}$$

where the row vectors are node embeddings (d-dimensional vectors).

- **Optimal Solution**

$$Z^* = (u_2, u_3, \dots, u_{d+1})$$



where u_2, u_3, \dots, u_{d+1} are the **eigenvectors of the graph Laplacian** (whichever the version of L, L_{sym} or L_{rm}) sorted by the eigenvalues $0 < \lambda_2 \leq \lambda_3 \leq \dots \leq \lambda_{k+1}$.

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9

9

Smoothness Penalty (when $d = 1$)

- Let us consider $d = 1$ for now, i.e., $z_i, z_j \in \mathbb{R}$ (scalars) and $Z \in \mathbb{R}^n$ (vector)

$$\begin{aligned} \sum_{i,j} |z_i - z_j|^2 A_{i,j} &= \sum_{i,j} (z_i - z_j)^2 A_{i,j} \\ &= \sum_{i,j} (z_i^2 + z_j^2 - 2z_i z_j) A_{i,j} \\ &= \sum_i z_i^2 \underbrace{\sum_j A_{i,j}}_{D_{ii}} + \sum_j z_j^2 \underbrace{\sum_i A_{i,j}}_{D_{jj}} - \sum_{i,j} 2z_i z_j A_{i,j} \\ &= 2\sum_i z_i^2 D_{ii} - 2\sum_{i,j} z_i z_j A_{i,j} = 2Z^T L Z \end{aligned}$$

- Also, we have

$$Z^T L Z = \sum_{i,j} z_i z_j L_{ij} = \sum_{i,j} z_i z_j (D_{ij} - A_{ij}) = \sum_i z_i^2 D_{ii} - \sum_{i,j} z_i z_j A_{i,j}$$

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10

10

Objective for Optimization (with $d = 1$)

$$Z^T L Z = \frac{1}{2} \sum_{i,j} |z_i - z_j|^2 A_{i,j} \geq 0 \quad (Z \in \mathbb{R}^n)$$

$$\min_{Z \in \mathbb{R}^n} \sum_{i,j} |z_i - z_j|^2 A_{i,j} = \min_{Z \in \mathbb{R}^n} Z^T L Z$$

Smoothness Penalty

L is positive-semidefinite

($\because Z^T L Z \geq 0$ for any $Z \in \mathbb{R}^n$).

Being positive-semidefinite (PSD) of L ...

1) For any $x \in \mathbb{R}^n$, we have $x^T L x \geq 0$, according to the definition of a PSD matrix.

2) For any eigenvector $u_i \in \mathbb{R}^n$ of graph Laplacian L , we have $u_i^T L u_i \geq 0$.

3) $L u_i = \lambda_i u_i \xrightarrow{\text{multiplying } u_i^T \text{ on both sides}} u_i^T L u_i = \lambda_i \xrightarrow{\text{imply}} \lambda_i \geq 0, \forall i$ (by Item 2 above)



4) $\lambda_1 = 0$ is the smallest eigenvalue of L .

5) $u_1 = 1_n$ is the eigenvector corresponding to $\lambda_1 = 0$ (next slide).

6) Vector $1_n = Z^* = \arg \min_{Z \in \mathbb{R}^n} Z^T L Z$ is **naively optimal** or 1D embedding.

7) We want to modify our objective as $Z^* = \arg \min_{Z \perp u_1} Z^T L Z$, which leads to u_2 with $\lambda_2 > 0$.

Eigenvector $u_1 = \mathbf{1}_n$ with $\lambda_1 = 0$

- Starting with $L = D - A$, we have

$$L\mathbf{1}_n = D\mathbf{1}_n - A\mathbf{1}_n \quad (\text{multiplying } \mathbf{1}_n \text{ on both sides})$$

$$= \begin{pmatrix} D_{11} \\ D_{22} \\ \vdots \end{pmatrix} - \begin{pmatrix} \sum_{j=1}^n A_{1j} \\ \sum_{j=1}^n A_{2j} \\ \vdots \end{pmatrix}$$

$$= \begin{pmatrix} 0 \\ 0 \\ \vdots \end{pmatrix} = 0 \cdot \mathbf{1}_n$$

$$\begin{aligned} L\mathbf{1}_n &= (D - A)\mathbf{1}_n \\ L\mathbf{1}_n &= D\mathbf{1}_n - A\mathbf{1}_n \end{aligned}$$

- Thus, we have $L\mathbf{1}_n = 0 \cdot \mathbf{1}_n$ which means $u_1 = \mathbf{1}_n$ and $\lambda_1 = 0$.

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13

13

Smoothness penalty when $d > 1$

$$\sum_{i,j} |z_i - z_j|^2 A_{i,j} = \sum_{i,j} \sum_{k=1}^d (z_i^{(k)} - z_j^{(k)})^2 A_{i,j} \quad (z_i, z_j \in \mathbb{R}^d)$$

$$= \sum_{k=1}^d \sum_{i,j} (z_i^{(k)} - z_j^{(k)})^2 A_{i,j}$$

$$= \sum_{k=1}^d 2Z_k^T LZ_k$$

$$\neq 2\text{tr}(Z^T LZ)$$

(using the proof for $d = 1$)
(next slide)



$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j} |z_i - z_j|^2 A_{i,j} = \min_{Z \in \mathbb{R}^{n \times d}} \text{tr}(Z^T LZ) \quad (d \geq 1)$$

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14

14

What is the trace of matrix $Z^T LZ$?

- Let (Z_1, Z_2, \dots, Z_d) be the column vectors of matrix $Z \in \mathbb{R}^{n \times d}$.

$$Z^T LZ = \begin{bmatrix} Z_1^T \\ Z_2^T \\ Z_3^T \end{bmatrix} L \begin{bmatrix} Z_1 \\ Z_2 \\ Z_3 \end{bmatrix} = \begin{bmatrix} Z_1^T LZ_1 & Z_1^T LZ_2 & Z_1^T LZ_3 \\ Z_2^T LZ_1 & Z_2^T LZ_2 & Z_2^T LZ_3 \\ Z_3^T LZ_1 & Z_3^T LZ_2 & Z_3^T LZ_3 \end{bmatrix}$$

$tr(Z^T LZ) = \sum_{k=1}^d Z_k^T LZ_k$

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15

15

Optimal Embedding

- We want: $Z^* = \underset{Z_k \perp 1_n}{\operatorname{argmin}} \sum_{k=1}^d Z_k^T LZ_k$

- When $d = 1$, $Z^* = \underset{Z_1 \perp 1_n}{\operatorname{argmin}} Z_1^T LZ_1 = u_2$

- When $d > 1$, for $k = 1$ to $d+1$, find

$$Z_{k+1}^* = \underset{Z_{k+1} \perp \{u_1, \dots, u_k\}}{\operatorname{argmin}} Z_k^T LZ_k = u_{k+1}$$

return $Z^* = (u_2, \dots, u_{d+1})$

where $\operatorname{tr}(Z^{*T} LZ^*) = \sum_{k=2}^{d+1} u_k^T L u_k = \lambda_2 + \lambda_3 + \dots + \lambda_{d+1}$

- How do we choose d ?

it's a hyperparameter.

? we are basically trying to learn the eigen vector u_2 that is perpendicular to u_1 to avoid trivially minimizing $Z^T LZ$.

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16

16

Generalized Eigenvector Problem

[M. Belkin and P. Niyogi. NIPS 2002]

- **Standard eigenvector problem** is defined as to find all the vectors satisfying

$$Lu = \lambda u \quad \text{or} \quad u^T Lu = \lambda$$

- Solution: eigenvectors u_1, u_2, \dots, u_n ordered by $\lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$.

- **Generalized eigenvector problem** is defined as to find all the vectors satisfying

$$Ly = \lambda Dy \quad \text{or} \quad y^T Ly = \lambda \quad (D \text{ is a diagonal matrix.})$$

- Solution: **D-orthonormal** y_1, y_2, \dots, y_n (when D is splittable)

$$\forall_{i,j} y_i^T D y_j = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases} \quad \text{or} \quad Y^T D Y = I, Y = (y_1, y_2, \dots, y_n)$$

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17

17 This is because the dot product of a unit vector with itself is 1.

Other matrix-factorization based methods

- **Laplacian Eigenmaps** (M. Belkin and P. Niyogi. NIPS 2002)

$$Z^* = \arg \min_{Z_{n \times d}} |z_i - z_j|^2 A_{i,j}.$$

- **Graph Factorization** (A. Ahmed et al., WWW 2013)

$$Z^* = \arg \min_{Z_{n \times d}} \sum_{i,j} |z_i^T z_j - A_{i,j}|$$

- **GraRep** (S. Cao et al., CIKM 2015)

$$\min_{\text{orthonormal } Z^k} \sum_{i,j} \left\| z_i^T z_j - \log \left(\frac{A_{i,j}^k}{\sum_l A_{l,j}^k} \right) + \log \frac{\lambda}{n} \right\|_2^2$$

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18

Issue with the computational cost



- Eigen-decomposition of the full matrix is expensive when n is large.
- A cheaper alternative is the **random walk** approach (next).

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19

19

Node Embedding via Random Walk over a Graph (e.g., [DeepWalk by Perozzi et al., KDD 2014](#))

- Input graphs $G = (V, E)$, for example
 - **BlogCatalog**: a network of social relationships among blogger authors (**10k**)
 - **Flickr**: a network of the contacts between users (**80k**) on a photo sharing website
 - **YouTube**: a social network among users (**1,139k**) of a popular video sharing website
- Random walk for generating “word”/context pairs
 - From each node, randomly walk for t steps to obtain a sequence of nodes (“words”)
 - At each step, uniformly sample the next node from the neighbors of the current node
 - Apply a sliding window over the node sequences to obtain “word”/context pairs
 - Train a word2vec method (SkipGram) on the above training pairs to obtain the embedding of nodes

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20

20

Other Variants of Random Walk Methods

- DeepWalk used a uniform sampling strategy, but other teleportation strategies are also possible (A Grover & J Leskovec, KDD 2016)
 - To reduce the probability of turning back to each node
 - To reduce the link weights for the nodes which have many links
 - To sample the context of each node via beam search (sampling multiple nodes per step instead sampling one node per step)
- DeepWalk uses SkipGram as the word embedding algorithm, but other choices (CBOW, GloVe, etc.) are also possible.

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21

21

Concluding Remarks

- Matrix-factorization methods focus on the global smoothness of the entire graph (the eigen-decomposition of the graph Laplacian matrix).
- Random-walk methods focus on the local neighborhood of each node.
- The former could be too expensive for large graphs.
- The latter could be too simple for good performance.
- Both methods are task-agnostic, which is a limitation as the embeddings of nodes cannot be adapted to down-stream tasks.
- Both do not leverage node-specific features (even if available) as another weakness.
- Graph Neural Networks (GNNs) overcome both kinds of the limitations (next lecture).

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22

22

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23

23

Graph construction based on given node features

(M. Belkin and P. Niyogi. NIPS 2002)

- Let $x_1, x_2, \dots, x_n \in \mathbb{R}^l$ be the given feature vectors of nodes, we can construct a graph based on
 - ε -neighborhoods ($\varepsilon > 0$): connect nodes i and j if $\|x_i - x_j\|^2 < \varepsilon$;
 - k -nearest neighbors ($k \in \mathbb{N}$): connect nodes i to j if i is among the k -nearest neighbors of j or if j is among the k -nearest neighbors of i .
- Choices of link weights in adjacency matrices
 - Simple-minded: $A_{ij} = 1$ if and only if nodes i and j are connected;
 - Heat kernel (with parameter $t > 0$): $A_{ij} = e^{-\frac{\|x_i - x_j\|^2}{t}}$ if nodes i and j are connected.

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24

24