Model Based Testing

Dr Michael Foster Based on material from Professor Rob Hierons

Formal models of systems (FSMs)

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- \cdot Testing from finite state machines

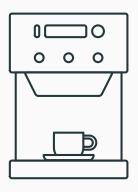
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 - · Unique I/O sequences
 - · The W method

Motivating Example - Simple Drinks Machine



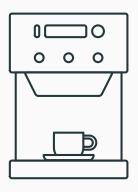
Select a drink

Motivating Example - Simple Drinks Machine



- · Select a drink
- Insert coins

Motivating Example - Simple Drinks Machine



- · Select a drink
- · Insert coins
- Press "vend" to dispense drink

Unit testing

 $\boldsymbol{\cdot}$ Generate tests according to code components

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- · Aim to achieve some level of code coverage

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There are lots of different modelling notations (Z, B, state machines) We will introduce MBT with state machines

An FSM (Mealy machine) is a six tuple (S, s_0 , X, Y, δ , λ) where

• S is a finite set of states

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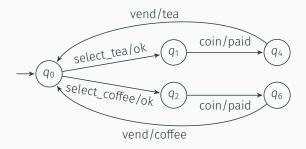
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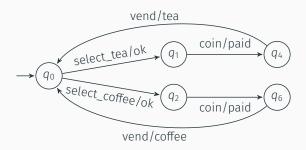
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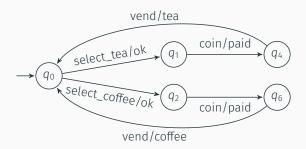




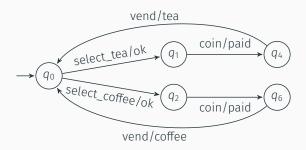
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 - Machine follows path $q_0 o q_1 o q_4 o q_0$

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We will focus on FSM testing with deterministic, minimal, and completely specified models

• An FSM is *deterministic* if there is at **most** one transition from each state for each input (implicit since δ and λ are functions)

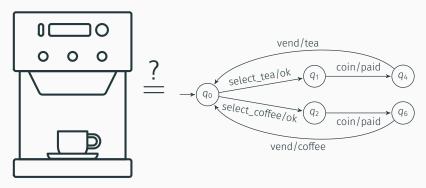
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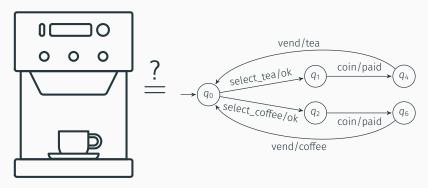
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- An FSM is minimal if there is no trace equivalent FSM with fewer states
 - We can minimise any FSM by following the algorithm

Testing from an FSM



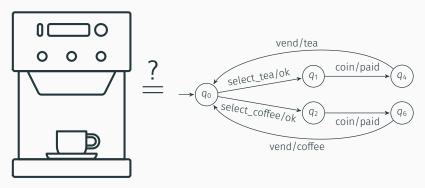
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Testing from an FSM



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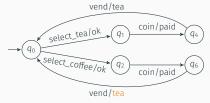
Testing from an FSM



- · Assume the software behaves like an FSM model
- · Submit inputs to the FSM and software in parallel
- Observe and compare the outputs

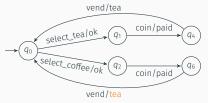
Faults

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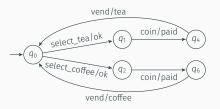


State transfer faults: A transition goes to the wrong state



Transition Tour

Assume no state transfer faults, then we can test by just executing every transition



The input sequence select_tea, coin, vend, select_coffee, coin, vend will do that for us

We validate that the output sequence is ok, paid, tea, ok, paid, coffee

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- If there are state transfer faults, a transition tour may not find them
- · Output faults may also be masked
- We want to explicitly check for state transfer faults

To test from an FSM, we want to check every transition $(q_i, q_j, x/y)$

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Challenges

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Challenges

Controllability: How do we get the FSM to q_i ?

Observability: How do we know the FSM is in q_j ?

Controllability

Find a sequence that gets the *specification* to the desired state.

Observability

Characterise states in terms of the I/O actions they can perform:

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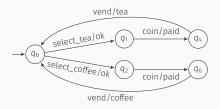
Characterise states in terms of the I/O actions they can perform:

- · Distinguishing sequences
- Unique I/O sequences
- · Characterising set

Distinguishing Sequences

A distinguishing sequence is an input sequence that produces a different output for **each state**.

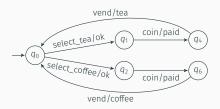
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Yes! (assuming we know when an input has been refused)

select_tea, coin, vend

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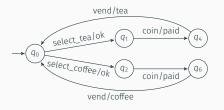
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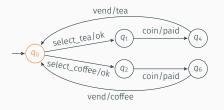
A state cover V is a set of input sequences such that each state of an FSM M is reached from s_0 by a sequence from V, and V also contains the empty input sequence.

The W Method for Drinks



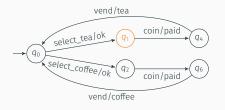
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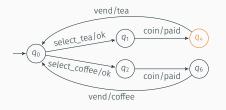
$$W = \{[select_tea, coin, vend]\}\$$

 $V = \{\epsilon\}$

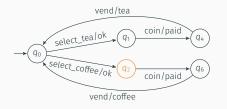


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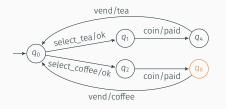
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W = \{[select\_tea, coin, vend]\} V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}
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For FSM with n+1 states, the W method produces the test set $VW \cup VXW$

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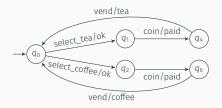
For n+m states, we get $VW \cup VXW \cup ... \cup VX^mW$

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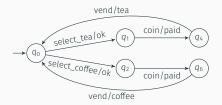
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Make sure you reset before each sequence!



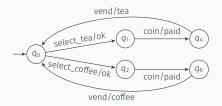
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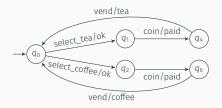
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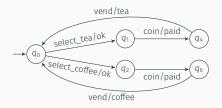
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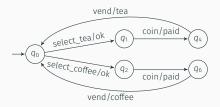
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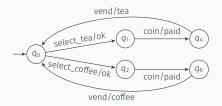
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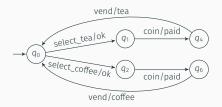
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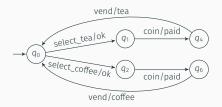
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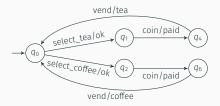
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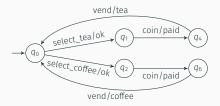
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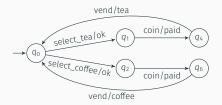
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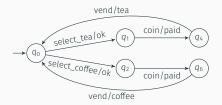
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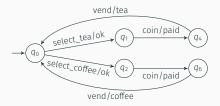
W = \{[select\_tea, coin, vend]\}
```



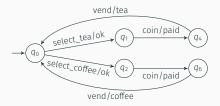
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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}

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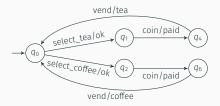
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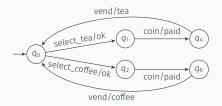
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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}
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W = \{[select\_tea, coin, vend]\}
```



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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}

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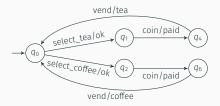
W = \{[select\_tea, coin, vend]\}
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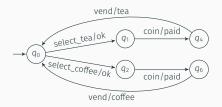
```
V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}

X = \{[select\_tea], [select\_coffee], [coin], [vend]\}

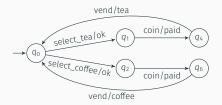
W = \{[select\_tea, coin, vend]\}
```



```
V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}
X = \{[select\_tea], [select\_coffee], [coin], [vend]\}
W = \{[select\_tea, coin, vend]\}
```



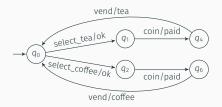
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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}
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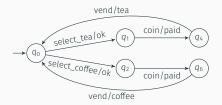
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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}

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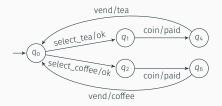
W = \{[select\_tea, coin, vend]\}
```



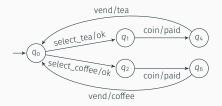
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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\}
X = \{[select\_tea], [select\_coffee], [coin], [vend]\}
W = \{[select\_tea, coin, vend]\}
```



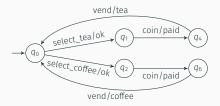
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V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\} X = \{[select\_tea], [select\_coffee], [coin], [vend]\} W = \{[select\_tea, coin, vend]\}
```



```
V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\} X = \{[select\_tea], [select\_coffee], [coin], [vend]\} W = \{[select\_tea, coin, vend]\}
```



```
V = \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\} X = \{[select\_tea], [select\_coffee], [coin], [vend]\} W = \{[select\_tea, coin, vend]\}
```



```
\label{eq:V} \begin{split} V &= \{\epsilon, [select\_tea], [select\_tea, coin], [select\_coffee], [select\_coffee, coin]\} \\ X &= \{ [select\_tea], [select\_coffee], [coin], [vend]\} \\ W &= \{ [select\_tea, coin, vend]\} \end{split}
```

• Where do models come from? Have you ever (voluntarily) drawn a model for your software?

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- We can test that the system matches an FSM specification using the W method