

# CSCI 270 - Spring 2023 - HW10

Due April 12, 2023

## Problem 1 (25pts)

Consider the partial satisfiability problem, denoted as 3-Sat( $\alpha$ ). We are given a collection of  $k$  clauses, each of which contains exactly three literals, and we are asked to determine whether there is an assignment of true/false values to the literals such that at least  $\alpha k$  clauses will be true. Note that 3-Sat(1) is exactly the 3-SAT problem from lecture.

Prove that 3-Sat(15/16) is **NP**-complete.

Hint: If  $x$ ,  $y$ , and  $z$  are literals, there are eight possible clauses containing them:  $(x \vee y \vee z)$ ,  $(!x \vee y \vee z)$ ,  $(x \vee !y \vee z)$ ,  $(x \vee y \vee !z)$ ,  $(!x \vee !y \vee z)$ ,  $(!x \vee y \vee !z)$ ,  $(x \vee !y \vee !z)$ ,  $(!x \vee !y \vee !z)$

## Problem 2

[25 pts.] Consider a graph  $G = (V, E)$  and two integers  $k, m$ .

### 2a

A **k-clique** is a subset of nodes  $u_i \in G, i = 1, \dots, k$  such that there is an edge connecting every pair of distinct vertices  $u_i, u_j$ . In other words, the **k-clique** is a complete sub-graph of  $G$ . Prove that finding a clique of size  $k$  is NP-Complete. [15 pts.]

### 2b

The **Dense Subgraph** problem is to find a subset  $V'$  of  $V$ , whose size is at most  $k$  and are connected by at least  $m$  edges. Prove that the **Dense Subgraph** problem is NP-Complete. [10 pts.]

### Problem 3 (25 pts)

Consider a modified SAT problem, SAT' in which given a CNF formula having  $m$  clauses and  $n$  variables  $x_1, x_2, \dots, x_n$ , the output is YES if there is an assignment to the variables such that exactly  $m - 2$  clauses are satisfied, and NO otherwise. Prove that SAT' is NP-Complete.

#### **Problem 4 (25 pts)**

Show that Vertex Cover is still NP-complete even when all vertices in the graph are restricted to have even degree.

## Practice Problems

### Problem 5 (25 pts)

(Kleinberg and Tardos, Chapter 8, Exercise 5)

Consider a set  $A = \{a_1, \dots, a_n\}$  and a collection  $B_1, B_2, \dots, B_m$  of subsets of  $A$  (i.e.,  $B_i \subseteq A$  for each  $i$ ).

We say that a set  $H \subseteq A$  is a *hitting set* for the collection  $B_1, B_2, \dots, B_m$  if  $H$  contains at least one element from each  $B_i$ —that is, if  $H \cap B_i$  is not empty for each  $i$  (so  $H$  “hits” all the sets  $B_i$ ).

We now define the *Hitting Set Problem* as follows. We are given a set  $A = \{a_1, \dots, a_n\}$ , a collection  $B_1, B_2, \dots, B_m$  of subsets of  $A$ , and a number  $k$ . We are asked: Is there a hitting set  $H \subseteq A$  for  $B_1, B_2, \dots, B_m$  so that the size of  $H$  is at most  $k$ ?

Prove that Hitting Set is NP-complete.