3.10 Poth properties of r.w. 1. M. /a,b) = # of poths from a b b in a steps = $\left(\begin{array}{c} h \rightarrow b - e \\ 2 \end{array}\right)$ $P(S_n = b)S_0 = a) = \left(\frac{n+b-a}{2}\right)P^{\frac{n+b-a}{2}}$ 2. For a, b>0, Nn°(a,b) = # of paths from a to b in a skys via tero = Nn (-a, b) 3. For b \$0, N, (0, b) = # of poths from 0 to 6 in a steps without zero visits = 1b/ Na (0, b) Cloim 1. Let So=0, b + 0. Then $P(S_1,...,S_n \neq 0, S_n = b) = \frac{1bl}{n} |^2(S_n = b), or$ $P(S_1...S_n \neq 0 \mid S_n = b) = \frac{1bl}{n}$ Some exercises Ix2. Consider simple r.w., So= 6, O<6 < N. Let IL = unis /42/1: Su=0 or Su= NS: IL is hime needed to reach boundary from le. a) Find E (Ta) Answer. 1. 1st step endyrin gives a system of your for Du = E(Tu): (1) $\int_{1}^{1} D_{k} = P D_{k+1} + y D_{k-1} + 1, k = 1, 2, ..., W-1$ becouse Du = E(T4 | X,=1) P(X,=1) + E(T4 | X,=-1) P(X,=-1) $= (1 + D_{k+1}) + (1 + D_{k-1}) = g D_k + y D_k + 1.$

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Given X, = 1, Zh = 1 + Zh+1: E(Zh+1) = 1 + E(Zh+1) = 1+ Dh+1
          X_{1} = -1, \quad \tau_{k} = 1 + \tau_{k-1} : E(\tau_{k} | X_{1} = -1) = 1 + E(\tau_{k-1}) - 1 + D_{k-1}
 Now (1) is non homopeneous lines equ. General
 Lo lu Lon: Du = A + B (4) + La, when La is
porticular es lection do (1), and A, B to be found
  Lu is of the form
Lu = \left( \begin{array}{cc} c & k \\ c & k^2 \end{array} \right), \quad g \neq f  with c \text{ found } foun(1).
A, B ove found from Do = DN = O.
  Result;
D_{k} = \begin{cases} k(N-h), & p = y = \frac{1}{2} \\ (q-p)(1-N), & \frac{1-(4)^{k}}{1-(4)^{N}}, & p \neq f. \end{cases}
  b) find Du = lim Ph
 Answer \widehat{D}_4 = \left\{ \begin{array}{c} + & \text{if } p > p \text{ (includes } p = p : \frac{1}{2}). \\ \frac{k}{q-p} & \text{if } p > p \end{array} \right.
 Remorte 1. Let h=0, Jo= h, Th= min /n>,1: Sn=0f.
Then Dh = E(Ta) is expected ruin time.
  Recoll for p=y=2, P(Th 20)=1,
  but E(\overline{z}_{k}) = + \infty.
\pm 43. Let S_0 = 0. We found for b > 0,
 \mathcal{N}_{n}^{\circ}(0,b) = \mathcal{N}_{n-1}^{\circ}(1,b) = \mathcal{N}_{n-1}^{\circ}(1,b) - \mathcal{N}_{n-1}^{\circ}(1,b) =
 = N_{n-1}((1,b) - N_{n-1}(-1,b) = N_{n-1}(0,b-1) - N_{n-1}(0,b+1).
Show that
= N_{n-1}(0,b-1) - N_{n-1}(0,b+1).
Rzn = # of paths in { S,... Szn + 0} = Nzn (0,0).
A. ...... R. = 2, 5 / S, ... Szn + 0, Szn = 26} delescopping sum
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Answer Ran = 2 = 1 { S. ... San ≠ 0, San = 26 } delescoping sum $=2\sum_{k=1}^{\infty}N_{2n}^{o}(0,2k)=2\sum_{k=1}^{\infty}N_{2n-1}(0,2k-1)-N_{2n-1}(0,2k+1)$ $=2 N_{2n-1}(0,1)-2 N_{2n-1}(2+2n+1)=N_{2n-1}(0,1)+N_{2n-1}(0,-1)$ = Nzn (0,0). Thm 4. Let So=0, p=1=2. Then $P(S_1...S_{2n}\neq 0) = P(\overline{S}_{2n}\neq 0) = (\frac{2n}{n})z^{-2n}$ where to = min / n >, 1: Sn = 0/ Proof. P(5,..., 52 n ≠ 0) = R2 n · 2 -2 n = N2 n (0, 0) 2-2 n $= \begin{pmatrix} 2n \\ 2n \end{pmatrix} 2^{-2n}$ Remorh 1. Recall our homework: P(Szn = D) So = O) $= \left(\begin{array}{c} 2n \\ n \end{array}\right) 2^{-2n} \approx \left(\begin{array}{c} 1 \\ 1 \end{array}\right)$ Corollary. Let So = 0, p = q = = 1. Then P(To 2 00) = 1. why? P([060) = P(U (70 = 2n)) = = lim P(c. = 2n) = lim [1-P(to>2n)] = 1. Remark 2. For 0>0, p=1=2, $P(t, > 2n | S_0 = a) = \sum_{j=1-a} P(S_{2n} = j | S_0 = 0)$ Hence P(To < 0 | So= e) = 1.

Moment of the last return Example. Fair win is tomed represhedly, $S_{n} = \mathcal{H}_{n} - \mathcal{T}_{n} \quad n > 1, \quad S_{o} = 0.$ $\# \text{ of } H_{s} \quad \# \text{ of } T_{s}'.$ By Bernoulli than, $\frac{S_n}{n} = \frac{Il_n}{n} - \frac{\Gamma_n}{n} \approx 0$ for largey. Question: How feequent on zew visits! Time moment of the lost return. Let So = 0, 9=9=1/2. At time moment n, consider on = mox { le s n: Sk = 0}, Def. r.v. on in collect moment of the last return (return because S. = 0) Range of on = 40,1,..., n3. Cloim 1. Let So=O, To = min {n = 1: Sn = 0} & moment of the first return. Then for j = 0,1,..., on, P(on=j) = P(S; =0) P(To>n-j) Proof. Cose j'=n: $P(\sigma_n = n) = P(S_n = 0) P(\tau_0 > 0) = P(S_n = 0)$ Cose j Z h; $P(\sigma_{n} = j) = P(S_{j} = 0, S_{j+1} - S_{n-1} S_{n} \neq 0)$

= P(S;+,---S, +0|S;=0) P(S;=0) = P(S,---Sn-j =0)So=n) P(S;=0) = P(to > n-j | So=0) P(s;=0). Corollary Let So=0, p=y=1/2. Then for j=0,1,..,n, [? ([2n = 2j) = P([2; =0) P([a> 2(n-j))) $= P(S_{2i} = 0) P(S_{2(n-j)} = 0) \approx \frac{1}{(\pi_{i})} \frac{1}{(\pi_{i})}$ j and n-j one large.