

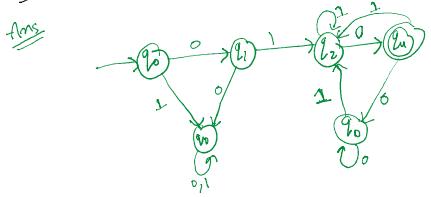
Practice Set - TOC

08 February 2024 00:55

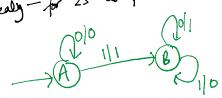
TOC Solution - Vaibhav Singh (Section-11)
(cse-cse)

Q1 DFA → Even No. of a's or Even No. of b's
Ans: $\{ \epsilon, aa, bb, aab, aba, abb, bab, bba, aabb, \dots \}$

Q2 DFA → Starting with 01 and end with 10
Ans:



Ans: Mealy - for 2's complement



Ans: Mealy - for 3's complement



Q3 $S \rightarrow aSS | aSaS | aSab | b$
find left factoring for the given grammar.

Q4 find derivation of string $id_1 + id_2 * id_3$ for the grammar where, $E \rightarrow E+E | E * E | id$

Ans: (i) Left derivation

$$\begin{aligned} E &\rightarrow E+E \\ &\rightarrow id + E \\ &\rightarrow id + E * E \\ &\rightarrow id + id_2 * E \\ &\rightarrow id_1 + id_2 * id_3 \end{aligned}$$

Right derivation

$$\begin{aligned} E &\rightarrow E+E \\ &\rightarrow E+E * E \\ &\rightarrow E+E * id_3 \\ &\rightarrow E+id_2 * id_3 \\ &\rightarrow id_1 + id_2 * id_3 \end{aligned}$$

Parse Derivation



(ii) Left Derivation

$$\begin{aligned} E &\rightarrow E * E \\ &\rightarrow E+E * E \\ &\rightarrow id_2 + E * E \\ &\rightarrow id_1 + id_2 * E \end{aligned}$$

Right Derivation

$$\begin{aligned} E &\rightarrow E * E \\ &\rightarrow E+E * E \\ &\rightarrow E+E * id_3 \\ &\rightarrow E+id_2 * id_3 \\ &\rightarrow id_1 + id_2 * id_3 \end{aligned}$$

Parse Derivation



$$\begin{array}{l}
 \rightarrow \Sigma^* \Sigma^* \Sigma^* \\
 \rightarrow id_1 + \Sigma^* \Sigma^* \\
 \rightarrow id_1 + id_2 * \Sigma^* \Sigma^* \\
 \rightarrow id_1 + id_2 * id_3
 \end{array}
 \quad
 \begin{array}{l}
 \rightarrow \Sigma^* \Sigma^* \Sigma^* \\
 \rightarrow \Sigma + id_2 * id_3 \\
 \rightarrow id_1 + id_2 * id_3
 \end{array}
 \quad
 \begin{array}{l}
 \downarrow \Sigma + \Sigma \\
 \downarrow id_1 id_2
 \end{array}
 \quad
 \text{Ans}$$

Q8 CFG to GNF

$$\begin{array}{l}
 S \rightarrow CB | AB \\
 A \rightarrow a | AA \\
 B \rightarrow b \\
 C \rightarrow d
 \end{array}$$

Step 1: Unit Productions ✓
Null Productions ✓

Step 2: CNF ✓

$$\begin{array}{ll}
 \text{Step 3:} & \begin{array}{l}
 S \rightarrow CB | AB \\
 A \rightarrow a | AA \\
 B \rightarrow b \\
 C \rightarrow d
 \end{array} \quad \begin{array}{l}
 S \text{ with } A_1 \\
 C \text{ with } A_2 \\
 B \text{ with } A_3 \\
 A \text{ with } A_4
 \end{array}
 \end{array}$$

$$\begin{array}{l}
 A_1 \rightarrow A_2 A_3 / A_4 A_3 \\
 A_4 \rightarrow a / A_4 A_4 \\
 A_3 \rightarrow b \\
 A_2 \rightarrow d
 \end{array}$$

$$\begin{array}{l}
 \text{Step 4: (i)} \quad A_1 \rightarrow A_2 A_3 / A_4 A_3 \checkmark \\
 \text{(ii)} \quad A_4 \rightarrow a / A_4 A_4 \times
 \end{array}$$

$$\begin{array}{l}
 \text{Step 5: (i)} \quad A_4 \rightarrow a / A_4 A_4 \\
 z = A_4 / A_4 z \\
 A_4 \rightarrow a / az
 \end{array}$$

Now, The grammar is—

$$\begin{array}{l}
 A_1 \rightarrow A_2 A_3 / A_4 A_3 \\
 A_4 \rightarrow a / az \\
 A_3 \rightarrow b \\
 A_2 \rightarrow d \\
 z \rightarrow A_4 / A_4 z
 \end{array}$$

$$\begin{array}{l}
 \text{(ii)} \quad A_1 \rightarrow A_2 A_3 / A_4 A_3 \\
 A_1 \rightarrow d A_3 / a A_3 / a z A_3 \\
 z \rightarrow A_4 / A_4 z \\
 z \rightarrow a / az / az / a z z
 \end{array}$$

$$\begin{array}{l}
 \rightarrow \text{Now, The grammar is—} \\
 A_1 \rightarrow a A_2 / a A_3 / a z A_3 \\
 A_2 \rightarrow d \\
 A_3 \rightarrow b \\
 A_4 \rightarrow a / az \\
 z \rightarrow a / az / a z z
 \end{array}$$

Q9 CFG for—

$$\text{(i) } L = \{ a^n b^n \mid n \geq 1 \}$$

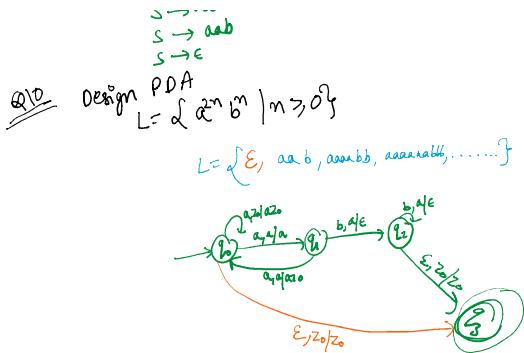
$$\text{Ans: } L = \{ ab, aabb, aaabbb, \dots \}$$

$$\begin{array}{l}
 S \rightarrow aSb \\
 S \rightarrow ab
 \end{array}$$

$$\begin{array}{l}
 \text{(ii) } L = \{ a^n b^n \mid n = 2m, m \geq 0 \} \\
 L = \{ a^{2n} b^{2n} \} \\
 L = \{ \epsilon, ab, aaaaabbb, \dots \}
 \end{array}$$

$$\begin{array}{l}
 S \rightarrow aab \\
 S \rightarrow ab \\
 S \rightarrow \epsilon
 \end{array}$$

Q10 Design PDA
 $r: 2n:m \mid m \geq n \}$



Q11 CFG \Rightarrow CNF

$$\begin{aligned} S &\rightarrow ASB \\ A &\rightarrow aAS | a|\epsilon \\ B &\rightarrow SbS | A | bb \end{aligned}$$

Step 1: (i) $\begin{aligned} p: \quad S &\rightarrow ASB \\ A &\rightarrow aAS | a | \epsilon \\ B &\rightarrow SbS | A | bb \\ S' &\rightarrow S \end{aligned}$

Step 2: (i) $\begin{aligned} A &\rightarrow aAS | a | \epsilon \\ S' &\rightarrow S \\ S &\rightarrow ASB | SB \\ A &\rightarrow aAS | a | aS \\ B &\rightarrow SbS | A | bb \\ S' &\rightarrow S \end{aligned}$

(ii) $\begin{aligned} B &\rightarrow SbS | A | bb | \epsilon \\ S &\rightarrow ASB | SB | AS | S \\ A &\rightarrow aAS | a | aS \\ B &\rightarrow SbS | A | bb \\ S' &\rightarrow S \end{aligned}$

Step 3: (i) $\begin{aligned} S &\rightarrow S \\ S &\rightarrow ASB | SB | AS \\ A &\rightarrow aAS | a | aS \\ B &\rightarrow SbS | A | bb \\ S' &\rightarrow S \end{aligned}$

(ii) $\begin{aligned} B &\rightarrow A \\ S &\rightarrow ASB | SB | AS \\ A &\rightarrow aAS | a | aS \\ B &\rightarrow aAS | a | aS \\ S' &\rightarrow S \end{aligned}$

(iii) $\begin{aligned} S' &\rightarrow S \\ S &\rightarrow ASB | SB | AS \\ A &\rightarrow aAS | a | aS \\ B &\rightarrow aAS | a | aS \\ S' &\rightarrow ASB | SB | AS \end{aligned}$

Step 4: $\begin{aligned} S &\rightarrow ASB | SB | AS \\ A &\rightarrow aAS | a | aS \\ B &\rightarrow aAS | a | aS \\ S' &\rightarrow ASB | SB | AS \end{aligned}$

$X \rightarrow 'AS'$

$\begin{aligned} S &\rightarrow Xb | SB | AS \\ A &\rightarrow aX | a | aS \\ B &\rightarrow aX | a | aS \\ S' &\rightarrow Xb | SB | AS \end{aligned}$

Step 5: $\begin{aligned} S &\rightarrow Xb | SB | AS \\ A &\rightarrow aX | a | aS \\ B &\rightarrow aX | a | aS \end{aligned}$

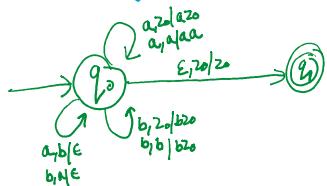
~~Step~~

$$\begin{array}{l}
 S \rightarrow XB|SB|AS \\
 A \rightarrow \alpha X|\alpha|AS \\
 B \rightarrow \alpha X|\alpha|AS \\
 S \rightarrow XB|SB|AS \\
 X \rightarrow SA \\
 \hline
 A \rightarrow \alpha X|\alpha S \\
 B \rightarrow \alpha X|\alpha S \\
 \hline
 A \rightarrow YX|YS \\
 B \rightarrow YX|YS
 \end{array}$$

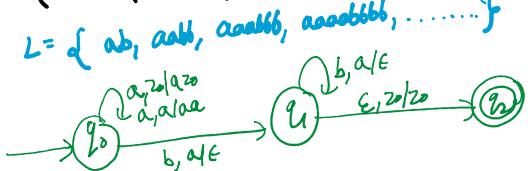
Now, This is CNF -

$$\begin{array}{l}
 A \rightarrow YX|YS|\alpha \\
 B \rightarrow YX|YS|\alpha \\
 S \rightarrow XB|SB|AS \\
 S \rightarrow XB|SB|AS \\
 X \rightarrow SA \\
 Y \rightarrow \alpha
 \end{array}$$

Q12 $L = \{ab, ba, aabb, abab, bbaa, \dots\}$

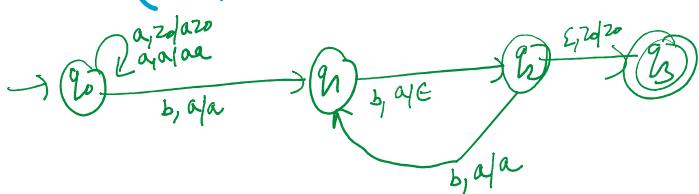


Q13 $L = \{a^n b^n \mid n \geq 1\}$



Q15 PDA - $L = \{a^n b^{2n} \mid n \geq 1\}$

$$L = \{abb, aabb, aaabbb, \dots\}$$

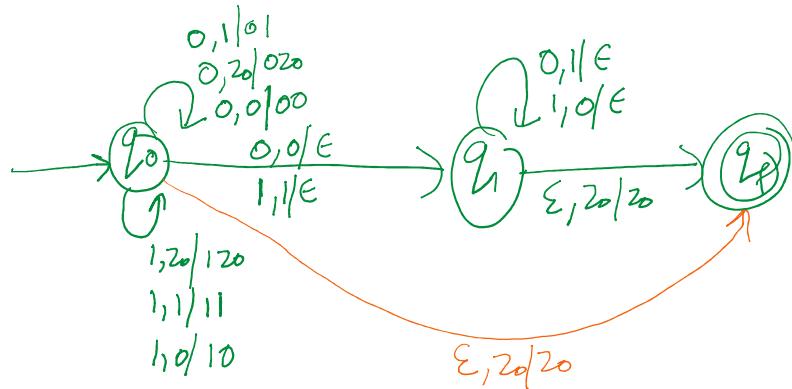


Q16 Similar to Q11

Q17 PDA $L = \{ww^R \mid w \in \{0,1\}^*\}$

$$L = \{\epsilon, 00, 11, 1001, 0110, \dots\}$$

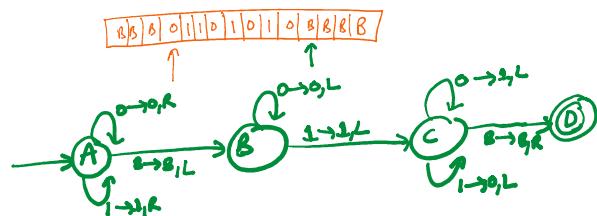
$$L = \{ \epsilon, 00, 11, 1001, 0110, \dots \}$$



Q1b $L = \{ w \mid w \in (0,1)^+ \}$

$$L = \{ 0<0, 1<1, 10<10, 000<000, \dots \}$$

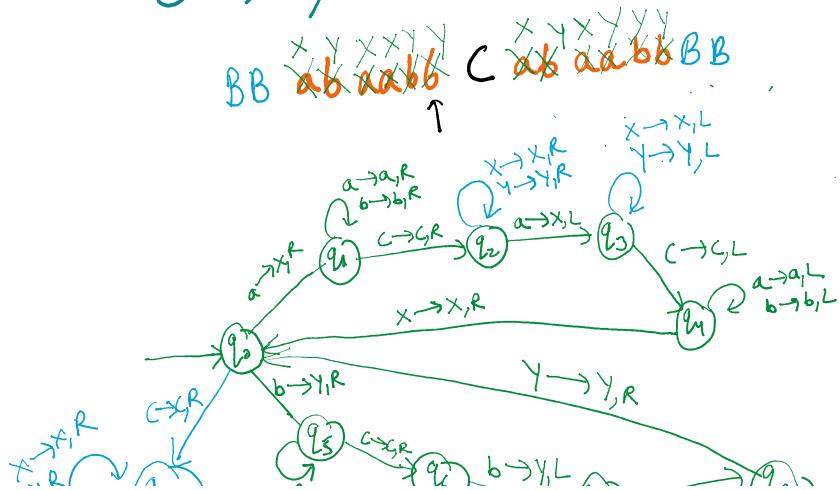
Q2a TM - Binary Value to 2's Complement

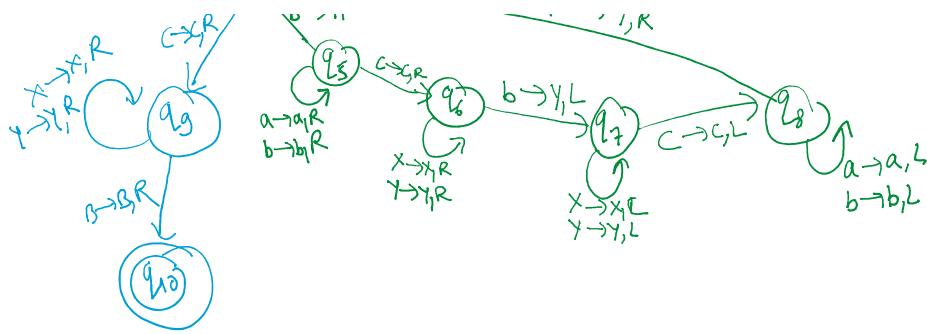


Q2b $L = \{ w \mid w \in (0,1)^{*} \}$

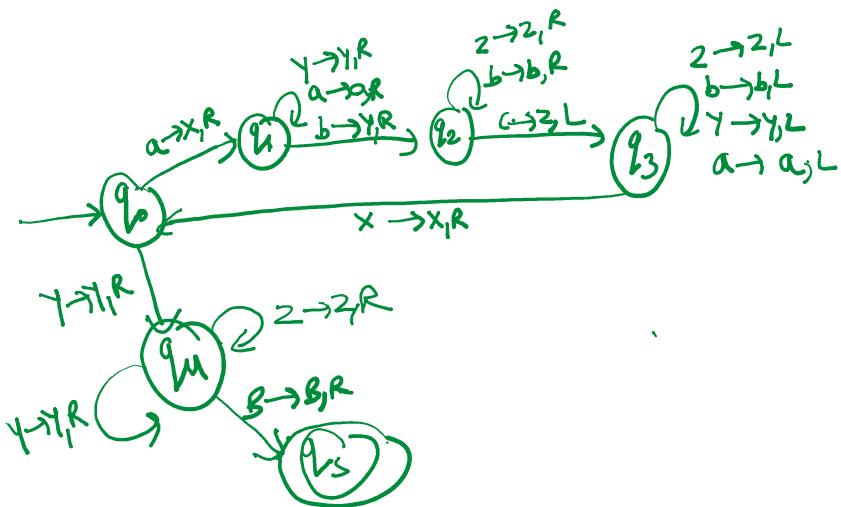
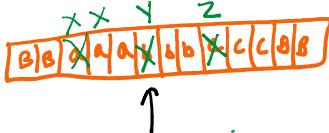
$$L = \{ \epsilon, 00, 11, 0000, 1010, \dots \}$$

$0 \rightarrow a, 1 \rightarrow b$



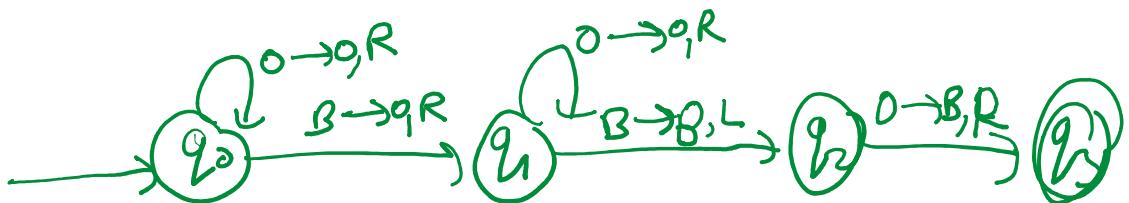
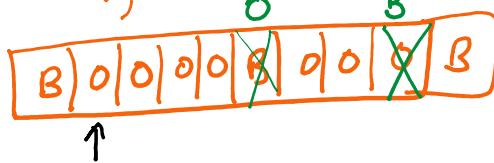


Q23 $L = \{ a^n b^m c^n \mid n \geq 0 \}$
 $L = \{ \epsilon, abc, aabbcc, aaabbbcc, \dots \}$

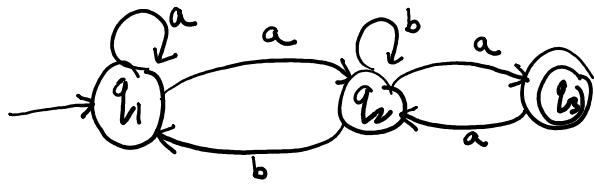


Q24 TM - Adder for Unary Value

$$a=4, b=3$$



Q28



$$q_1 = q_1 a + q_2 b + \epsilon \quad \text{--- (i)}$$

$$q_2 = q_2 b + q_1 a + q_3 a \quad \text{--- (ii)}$$

$$q_3 = q_2 a \quad \text{--- (iii)}$$

$$R = Q + RP$$

$$R = QP^*$$

$$\begin{aligned} q_2 &= q_2 b + q_1 a + q_3 a \\ &= q_2 b + q_1 a + q_2 aa \\ &= q_1 a + q_2 b + q_2 aa \\ &= \underbrace{q_1 a}_{a} + \underbrace{q_2}_{R} \underbrace{(b+aa)}_{P} \end{aligned}$$

$$= q_1 a (b+aa)^* \quad \text{--- (iv)}$$

$$\begin{aligned} h &= q_1 a + q_2 b + \epsilon \\ &= q_1 a + q_1 a (b+aa)^* b + \epsilon \\ &= \epsilon + q_1 a + q_1 a (b+aa)^* b \\ &= \underbrace{\epsilon}_{a} + \underbrace{q_1}_{R} \underbrace{(a+a(b+aa)^*)}_{P} b \\ &= \epsilon \cdot (a+a(b+aa)^*) b^* \\ &= (a+a(b+aa)^*) b^* \end{aligned}$$

$$q_3 = q_2 a$$

$$= (q_1 a (b+aa)^*) a$$

$$= (a+a(b+aa)^*) b^* a (b+aa)^* a$$

Ans

~~Q30~~ Same as Q23

~~Q31~~ Same as Q24