

NAME	Ms.NEERJA DOSHI
UID	2021300029
DIV	SE – CE -A
DAA EXPT 2	
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AIM :	Experiment Based on Divide and Conquer Techniques 1) MERGE SORT 2) QUICK SORT
ALGORITHM :	MERGE SORT step 1: start step 2: declare array and left, right, mid variable step 3: perform merge function. if left > right return mid= (left+right)/2 mergesort(array, left, mid) mergesort(array, mid+1, right) merge(array, left, mid, right) step 4: Stop QUICK SORT FOR QUICK SORT : QUICKSORT (array A, start, end) 1. { 2. 1 if (start < end) 3. 2 { 4. 3 p = partition(A, start, end) 5. 4 QUICKSORT (A, start, p - 1) 6. 5 QUICKSORT (A, p + 1, end) 7. 6} 8. } FOR PARTIONING ALGORITHM PARTITION (array A, start, end) 1. { 2. 1 pivot ? A[end] 3. 2 i ? start-1 4. 3 for j ? start to end -1 {

	<pre>5. 4 do if (A[j] < pivot) { 6. 5 then i = i + 1 7. 6 swap A[i] with A[j] 8. 7 } 9. 8 swap A[i+1] with A[end] 10. 9 return i+1 11. }</pre>
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CODE :

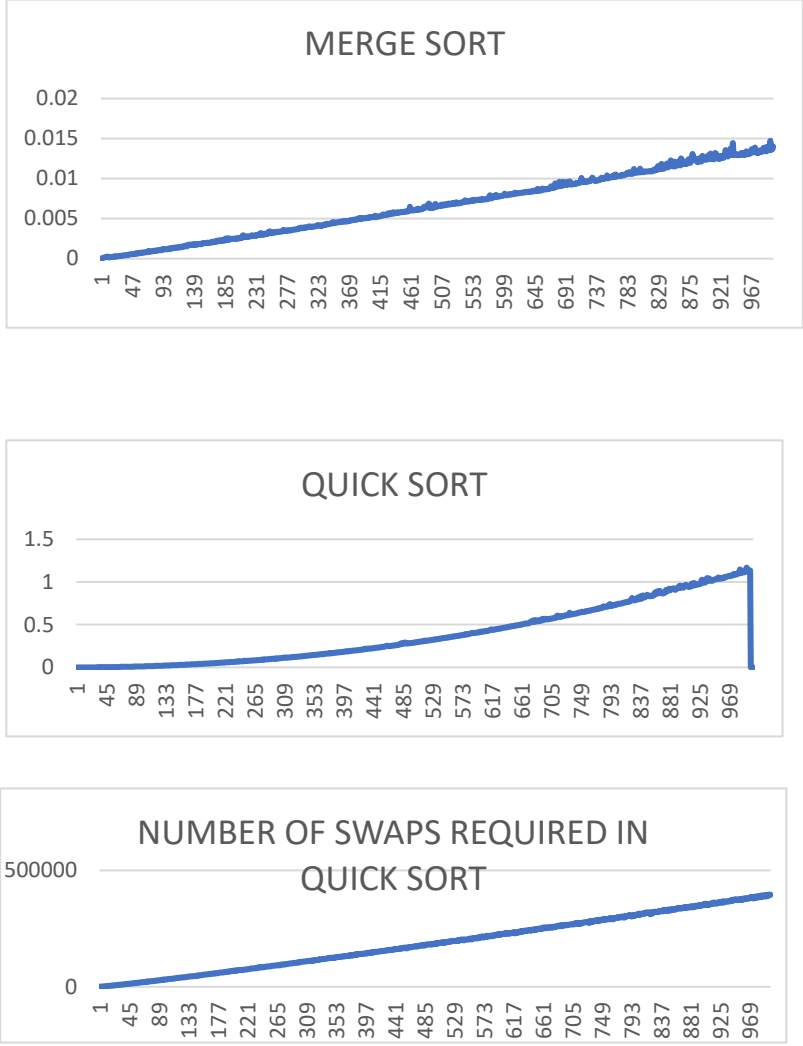
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#include <stdio.h>
#include <stdlib.h>
#include <time.h>
void merge(int a[], int beg, int mid, int end)
{
    int i, j, k;
    int n1 = mid - beg + 1;
    int n2 = end - mid;
    int LeftArray[n1], RightArray[n2];
    for (int i = 0; i < n1; i++)
        LeftArray[i] = a[beg + i];
    for (int j = 0; j < n2; j++)
        RightArray[j] = a[mid + 1 + j];
    i = 0,
    j = 0;
    k = beg;
    while (i < n1 && j < n2)
    {
        if (LeftArray[i] <= RightArray[j])
        {
            a[k] = LeftArray[i];
            i++;
        }
        else
        {
            a[k] = RightArray[j];
            j++;
        }
        k++;
    }
    while (i < n1)
    {
        a[k] = LeftArray[i];
        i++;
        k++;
    }
    while (j < n2)
    {
        a[k] = RightArray[j];
        j++;
        k++;
    }
}
void mergeSort(int a[], int beg, int end)
{
    if (beg < end)
    {
        int mid = (beg + end) / 2;
        mergeSort(a, beg, mid);
        mergeSort(a, mid + 1, end);
        merge(a, beg, mid, end);
    }
}
```

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    }
}
void printArray(int a[], int n)
{
    int i;
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    printf("\n");
}
int partition(int a[], int start, int end)
{
    int pivot = a[end];
    int i = (start - 1);
    for (int j = start; j <= end - 1; j++)
    {
        if (a[j] < pivot)
        {
            i++;
            int t = a[i];
            a[i] = a[j];
            a[j] = t;
        }
    }
    int t = a[i + 1];
    a[i + 1] = a[end];
    a[end] = t;
    return (i + 1);
}
void quick(int a[], int start, int end)
{
    if (start < end)
    {
        int p = partition(a, start, end);
        quick(a, start, p - 1);
        quick(a, p + 1, end);
    }
}
void printArr(int a[], int n)
{
    int i;
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
}
void main()
{
    int n = 0;
    for (int k = 0; k < (100000 / 100); k++)
    {
        n = n + 100;
        int num[n];
        int quicksort[n];
    }
}

```

```
int merge[n];
int j, min;
clock_t start_t, end_t;
double total_t;
printf("%d\t", n);
for (int i = 0; i < n; i++)
{
    num[i] = rand() % 10;
    merge[i] = num[i];
    quicksort[i] = num[i];
}
start_t = clock();
mergeSort(merge, 0, n - 1);
end_t = clock();
total_t = (double)(end_t - start_t) / CLOCKS_PER_SEC;
printf("%f\n", total_t);
start_t = clock();
quick(quicksort, 0, n - 1);
end_t = clock();
total_t = (double)(end_t - start_t) / CLOCKS_PER_SEC;
printf("%f\n", total_t);
}
}
```

<p>GRAPHICAL REPRESENTATION :</p>	 <p>The figure contains three line graphs. The first graph, titled 'MERGE SORT', shows a linear increase in time complexity from 0 to approximately 0.015 as the number of elements increases from 1 to 967. The second graph, titled 'QUICK SORT', shows a parabolic increase in time complexity from 0 to approximately 1.2 as the number of elements increases from 1 to 969. The third graph, titled 'NUMBER OF SWAPS REQUIRED IN QUICK SORT', shows a linear increase in the number of swaps from 0 to approximately 400,000 as the number of elements increases from 1 to 969.</p>
<p>OBSERVATIONS :</p>	<p>1) Merge sort follows a linear graph 2) Quick sort has a parabolic graph 3) Approximately 3 lakh swaps are required to sort the given Input of 1 lakh number in Quick Sort method</p>
<p>CONCLUSION :</p>	<p>Merge sort is more efficient as its worst case time complexity is $O(\log n)$ while in case of quick sort, it remains constant throughout all operations as we can see from its graph which is linear in nature.</p>