

# 1      **RESTARTING THE NON-SYMMETRIC LANCZOS ALGORITHM** 2      **VIA THE IMPLICITLY SHIFTED LR ALGORITHM\***

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4      **Abstract.** The shifted QR iteration is used as a restart procedure for the Arnoldi method  
5      for the calculation of a few eigenvalues of a large matrix. We show that the underlying idea can  
6      be utilized in much the same manner via the shifted LR iteration to create a restart procedure for  
7      the non-symmetric Lanczos algorithm for eigenvalue calculations. Additionally, we show that the  
8      (shifted) LR iteration can be performed implicitly in a manner similar to the Francis' algorithm,  
9      resulting in a bulge-chase type procedure which does not require the explicit construction of the full  
10     lower and upper triangular matrices.

11     **Key words.** LR algorithm, unsymmetric Lanczos, implicit restart

12     **MSC codes.** 68Q25, 68R10, 68U05

13     **1. Introduction.** The Arnoldi iteration [1] is a popular Krylov space method  
14     for calculating a few eigenvalues of a large matrix. The method relies on the gen-  
15     eration of a sequence of Krylov vectors which determine the subspace within which  
16     approximations of the eigenvalue-eigenvector pairs are obtained. Depending on the  
17     accuracy and number of eigenpair approximations needed, the Krylov space size can  
18     become exceedingly large so that the quality of the results may be limited by the  
19     available memory. Sorensen [9] introduced an elegant procedure for restarting the  
20     Arnoldi factorization based on polynomial filters, which are applied through the im-  
21     plicitly shifted QR iterations on the reduced Hessenberg matrix obtained through the  
22     Arnoldi method. In particular, the use of exact shifts was shown to be successful in the  
23     convergence process of the eigenspace[9] of the specified eigenvalues. The method has  
24     subsequently found widespread application through the ARPACK library [4]. The use  
25     of QR iterations ensures that the reduced matrix preserves its Hessenberg structure  
26     through the transforms that make up the restart process. If the underlying matrix is  
27     symmetric, the Arnoldi iteration reduces to the Lanczos algorithm and, the Hessen-  
28     berg matrix reduces to a symmetric tridiagonal matrix. The QR iteration preserves  
29     the symmetric tridiagonal structure as well and, as pointed out by Sorensen [9], the  
30     implicit restart process applies equally well for the Lanczos method for symmetric  
31     matrices.

32     One would like to extend this procedure to the case of the non-symmetric Lanczos  
33     method. However, the reduced matrix that one obtains is a non-symmetric tridiago-  
34     nal matrix, with the tridiagonal structure being the result of the recurrence relations  
35     of the Lanczos algorithm [8]. Since the QR iterations do not preserve the banded  
36     structure of non-symmetric matrices, a straightforward application of the restart pro-  
37     cedure propounded by Sorensen will lead to a loss of this tridiagonal structure of the  
38     reduced matrix (the Hessenberg structure will still be preserved). This loss of struc-  
39     ture can be circumvented if one looks to the predecessor of the QR algorithm namely,  
40     the LR algorithm proposed by Rutishauser [6, 7, 5], which has the attractive property  
41     of preserving the band structure of a matrix. This property was already pointed out  
42     by Rutishauser in [6] where the banded matrices were referred to as striped matrices.

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As we will show in the next section, shifted LR transforms is the appropriate generalization of the restart procedure to the case of non-symmetric Lanczos iteration. The process would necessarily require refining both the right as well as the left Krylov spaces simultaneously.

The rest of the paper is organized as follows. In the next section we start with the introduction of the non-symmetric Lanczos iteration and then develop the restart procedure. We also show that the LR algorithm can be implemented as a bulge-chase method, similar to the Francis' algorithm. In section 3 we apply the restart process to the Grcar matrix, and make some concluding remarks in section 4.

**2. Non-symmetric Lanczos.** Lanczos first introduced his algorithm in [3] as a method for tridiagonalizing a matrix, but also realized that the method could be used iteratively to find eigenvalues. For an arbitrary matrix  $A$ , the method generates a pair of Krylov subspaces  $\{v_1, \dots, v_j\}$  and  $\{w_1, \dots, w_j\}$ , through repeated action of  $A$  and  $A^H$  respectively. We refer to these as the right and left Krylov spaces respectively and they satisfy the biorthogonality relation  $w_i^H v_j = \delta_{ij}$ . The two subspaces are generated through the following recurrence relations:

$$(2.1a) \quad \delta_{j+1} v_{j+1} = A v_j - \alpha_j v_j - \beta_j v_{j-1}$$

$$(2.1b) \quad \beta_{j+1} w_{j+1} = A^H w_j - \alpha_j w_j - \delta_j w_{j-1}$$

which in matrix form can be written as

$$(2.2a) \quad A V_m = V_m T_m + v_{m+1} e_m^T$$

$$(2.2b) \quad A^H W_m = W_m T_m^H + w_{m+1} e_m^T$$

**3. Main results.** We interleave text filler with some example theorems and theorem-like items.

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Here we state our main result as Theorem 3.1; the proof is deferred to ??.

**THEOREM 3.1** (*LDL<sup>T</sup> Factorization [2]*). *If  $A \in \mathbb{R}^{n \times n}$  is symmetric and the principal submatrix  $A(1 : k, 1 : k)$  is nonsingular for  $k = 1 : n - 1$ , then there exists a unit lower triangular matrix  $L$  and a diagonal matrix*

$$D = \text{diag}(d_1, \dots, d_n)$$

such that  $A = LDL^T$ . The factorization is unique.

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THEOREM 3.2 (Mean Value Theorem). Suppose  $f$  is a function that is continuous on the closed interval  $[a, b]$ . and differentiable on the open interval  $(a, b)$ . Then there exists a number  $c$  such that  $a < c < b$  and

$$f'(c) = \frac{f(b) - f(a)}{b - a}.$$

In other words,

$$f(b) - f(a) = f'(c)(b - a).$$

Observe that Theorems 3.1 and 3.2 and Corollary 3.3 correctly mix references to multiple labels.

COROLLARY 3.3. Let  $f(x)$  be continuous and differentiable everywhere. If  $f(x)$  has at least two roots, then  $f'(x)$  must have at least one root.

*Proof.* Let  $a$  and  $b$  be two distinct roots of  $f$ . By Theorem 3.2, there exists a number  $c$  such that

$$f'(c) = \frac{f(b) - f(a)}{b - a} = \frac{0 - 0}{b - a} = 0.$$

□

Note that it may require two L<sup>A</sup>T<sub>E</sub>X compilations for the proof marks to show.

Display matrices can be rendered using environments from `amsmath`:

$$(3.1) \quad S = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \quad \text{and} \quad C = \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}.$$

Equation (3.1) shows some example matrices.

We calculate the Fréchet derivative of  $F$  as follows:

$$(3.2a) \quad F'(U, V)(H, K) = \langle R(U, V), H\Sigma V^T + U\Sigma K^T - P(H\Sigma V^T + U\Sigma K^T) \rangle \\ = \langle R(U, V), H\Sigma V^T + U\Sigma K^T \rangle$$

$$(3.2b) \quad = \langle R(U, V)V\Sigma^T, H \rangle + \langle \Sigma^T U^T R(U, V), K^T \rangle.$$

Equation (3.2a) is the first line, and (3.2b) is the last line.

**4. Algorithm.** Sed gravida lectus ut purus. Morbi laoreet magna. Pellentesque eu wisi. Proin turpis. Integer sollicitudin augue nec dui. Fusce lectus. Vivamus faucibus nulla nec lacus. Integer diam. Pellentesque sodales, enim feugiat cursus volutpat, sem mauris dignissim mauris, quis consequat sem est fermentum ligula. Nullam justo lectus, condimentum sit amet, posuere a, fringilla mollis, felis. Morbi nulla nibh, pellentesque at, nonummy eu, sollicitudin nec, ipsum. Cras neque. Nunc augue. Nullam vitae quam id quam pulvinar blandit. Nunc sit amet orci. Aliquam erat elit, pharetra nec, aliquet a, gravida in, mi. Quisque urna enim, viverra quis, suscipit quis, tincidunt ut, sapien. Cras placerat consequat sem. Curabitur ac diam. Curabitur diam tortor, mollis et, viverra ac, tempus vel, metus.

Our analysis leads to the algorithm in Algorithm 4.1.

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**Algorithm 4.1** Build tree

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Define  $P := T := \{\{1\}, \dots, \{d\}\}$ 
while  $\#P > 1$  do
  Choose  $C' \in \mathcal{C}_p(P)$  with  $C' := \operatorname{argmin}_{C \in \mathcal{C}_p(P)} \varrho(C)$ 
  Find an optimal partition tree  $T_{C'}$ 
  Update  $P := (P \setminus C') \cup \{\bigcup_{t \in C'} t\}$ 
  Update  $T := T \cup \{\bigcup_{t \in \tau} t : \tau \in T_{C'} \setminus \mathcal{L}(T_{C'})\}$ 
end while
return  $T$ 

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**5. Experimental results.** Quisque facilisis auctor sapien. Pellentesque gravida hendrerit lectus. Mauris rutrum sodales sapien. Fusce hendrerit sem vel lorem. Integer pellentesque massa vel augue. Integer elit tortor, feugiat quis, sagittis et, ornare non, lacus. Vestibulum posuere pellentesque eros. Quisque venenatis ipsum dictum nulla. Aliquam quis quam non metus eleifend interdum. Nam eget sapien ac mauris malesuada adipiscing. Etiam eleifend neque sed quam. Nulla facilisi. Proin a ligula. Sed id dui eu nibh egestas tincidunt. Suspendisse arcu.

Figure 1 shows some example results. Additional results are available in the supplement in Table 1.

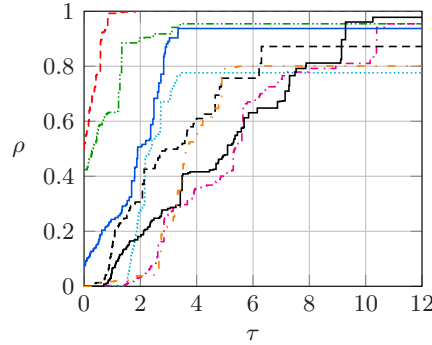


FIG. 1. Example figure using external image files.

Table 1 shows additional supporting evidence.

TABLE 1  
Example table.

Species	Mean	Std. Dev.
1	3.4	1.2
2	5.4	0.6
3	7.4	2.4
4	9.4	1.8

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erat. Sed tristique. Duis pede leo, fermentum quis, consecetur eget, vulputate sit amet, erat.

**6. Discussion of  $Z = X \cup Y$ .** Curabitur nunc magna, posuere eget, venenatis eu, vehicula ac, velit. Aenean ornare, massa a accumsan pulvinar, quam lorem laoreet purus, eu sodales magna risus molestie lorem. Nunc erat velit, hendrerit quis, malesuada ut, aliquam vitae, wisi. Sed posuere. Suspendisse ipsum arcu, scelerisque nec, aliquam eu, molestie tincidunt, justo. Phasellus iaculis. Sed posuere lorem non ipsum. Pellentesque dapibus. Suspendisse quam libero, laoreet a, tincidunt eget, consequat at, est. Nullam ut lectus non enim consequat facilisis. Mauris leo. Quisque pede ligula, auctor vel, pellentesque vel, posuere id, turpis. Cras ipsum sem, cursus et, facilisis ut, tempus euismod, quam. Suspendisse tristique dolor eu orci. Mauris mattis. Aenean semper. Vivamus tortor magna, facilisis id, varius mattis, hendrerit in, justo. Integer purus.

**7. Conclusions.** Some conclusions here.

**Appendix A. An example appendix.** Aenean tincidunt laoreet dui. Vestibulum ante ipsum primis in faucibus orci luctus et ultrices posuere cubilia Curae; Integer ipsum lectus, fermentum ac, malesuada in, eleifend ut, lorem. Vivamus ipsum turpis, elementum vel, hendrerit ut, semper at, metus. Vivamus sapien tortor, eleifend id, dapibus in, egestas et, pede. Pellentesque faucibus. Praesent lorem neque, dignissim in, facilisis nec, hendrerit vel, odio. Nam at diam ac neque aliquet viverra. Morbi dapibus ligula sagittis magna. In lobortis. Donec aliquet ultricies libero. Nunc dictum vulputate purus. Morbi varius. Lorem ipsum dolor sit amet, consecetur adipiscing elit. In tempor. Phasellus commodo porttitor magna. Curabitur vehicula odio vel dolor.

LEMMA A.1. *Test Lemma.*

**Acknowledgments.** We would like to acknowledge the assistance of volunteers in putting together this example manuscript and supplement.

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