

Tutorial 3

Q1 Construct predictive parsing table for following grammar and verify (give moves of the parser) for the string "id + id * id".

$$E \rightarrow T E'$$

$$E' \rightarrow + T E' \mid \epsilon$$

$$T \rightarrow F T'$$

$$T' \rightarrow * F T' \mid \epsilon$$

$$F \rightarrow (E) \mid \text{id}$$

Q2 Check the grammar is LL(1) or not:

1) $S \rightarrow A a A b \mid B b B a$

$$A \rightarrow \epsilon$$

$$B \rightarrow \epsilon$$

2) $S \rightarrow i E + S S'$

$$S' \rightarrow \epsilon S \mid \epsilon$$

$$E \rightarrow b$$

Q3 Given the grammar

$$S \rightarrow a \mid (L)$$

$$L \rightarrow L, S \mid S$$

i) Is the above grammar LL(1)? Justify your answer.

ii) What changes are necessary to make it suitable for LL(1) parser?

iii) Show the moves made by the LL(1) predictive parser on the input (a, (a, a)).

Q 4 Consider a grammar G as follows:

$$S \rightarrow W$$

$$W \rightarrow ZXY \mid XY$$

$$Y \rightarrow c \mid \epsilon$$

$$Z \rightarrow a \mid d$$

$$X \rightarrow Xb \mid \epsilon$$

Draw the $LL(1)$ parsing table for the given grammar.

Q 5 Consider the following grammar G :

$$S \rightarrow AC$$

$$A \rightarrow CbD$$

$$C \rightarrow BC \mid \epsilon$$

$$B \rightarrow cdD \mid acdD$$

$$D \rightarrow c \mid \epsilon$$

In $LL(1)$ parse table of above grammar G , How many cells are having multiple entries?

1) $E \rightarrow TE'$

$$E' \rightarrow +TE' \mid \epsilon$$

$$T \rightarrow FT'$$

$$T' \rightarrow *FT' \mid \epsilon$$

$$F \rightarrow (E) \mid id$$

FIRST (F_i)FOLLOW (F_o)

E	(, id	\$,)
E'	+, ε	\$,)
T	(, id	\$, +,)
T'	*, ε	\$, +,)
F	(, id	\$, +, *,)

$$F_i(E) = F_i(TE')$$

$$= F_i(T)$$

$$= F_i(FT')$$

$$= F_i(F)$$

$$= F_i((E)) + F_i(id)$$

$$= (, id$$

$$F_i(E') = F_i(+TE') + F_i(ε)$$

$$= +, ε$$

$$F_i(T) = F_i(FT')$$

$$= (, id$$

$$F_i(T') = F_i(*FT') + F_i(ε)$$

$$= *, ε$$

$$F_i(F) = F_i((E)) + F_i(id)$$

$$= (, id$$

- ① $E \rightarrow TE'$
- ② $E' \rightarrow + TE'$
- ③ $E' \rightarrow \epsilon$
- ④ $T \rightarrow FT'$
- ⑤ $T' \rightarrow * FT'$
- ⑥ $T' \rightarrow \epsilon$
- ⑦ $F \rightarrow (E)$
- ⑧ $F \rightarrow id$

$$F_0(E) = \emptyset, \$,)$$

$$F_0(E') = F_0(E) \cup \{ (, \$,) \}$$

$$= \$,)$$

$$F_0(T) = F_0(E') \cup \{ + \}$$

$$= +, \$,)$$

$$F_0(T') = F_0(T) \cup \{ * \}$$

$$= +, \$,)$$

$$F_0(F) = F_0(T') \cup \{ (,) \}$$

$$= F_0(T') \cup \{ (,) \}$$

$$= +, \$,) \cup \{ (,) \}$$

$$= +, \$,)$$

$$(27) \quad ① \quad E \rightarrow TE'$$

$$② \quad E' \rightarrow + TE'$$

$$③ \quad E' \rightarrow \epsilon$$

$$④ \quad T \rightarrow FT'$$

$$⑤ \quad T' \rightarrow * FT'$$

$$⑥ \quad T' \rightarrow \epsilon$$

$$⑦ \quad F \rightarrow (E)$$

$$⑧ \quad F \rightarrow id$$

$$F_0(E) = \$,)$$

$$F_0(E') = F_0(E) \quad \text{[scribbled out]} \\ = \$,)$$

$$F_0(T) = F_1(E') \quad \text{[scribbled out]} \\ = + , \epsilon \rightarrow + \cup F_0(E) \cup F_0(E) \\ = + , \$,)$$

$$F_0(T') = F_0(T) \\ = + , \$,)$$

$$F_0(F) = F_1(T') \\ = F_1(* , \epsilon) \\ = * , F_0(T) , F_0(T') \\ = * , + , \$,)$$

	id	+	*	()	\$
E	(1)			(1)		
E'		(2)			(3)	(3)
T	(4)			(4)		
T'		(6)	(5)		(6)	(6)
F	(8)			(7)		

2) 1) $S \rightarrow AaAb \mid BbBa$
 $A \rightarrow \epsilon$
 $B \rightarrow \epsilon$

No Left recursion. ~~not~~

$$F_i(S) = F_i(AaAb) \cup F_i(BbBa)$$

$$= a, b$$

$$F_i(A) = \epsilon$$

$$F_i(B) = \epsilon$$

$$F_0(S) = \$$$

$$F_0(A) = a, b$$

$$F_0(B) = a, b$$

$$S \rightarrow AaAb$$

$$S \rightarrow BbBa$$

$$A \rightarrow \epsilon$$

$$B \rightarrow \epsilon$$

	\$	a	b
S	(1)(2)	(1)(2)	(1)(2)
A		(3)	(3)
B		(3)	(3)

Not a LL(1) grammar

2) $S \rightarrow iEtSS'$ (1)
 $S' \rightarrow eS / \epsilon$ (2) (3)
 $E \rightarrow b$ (4)

	<u>First</u>	<u>Follow</u>
S	i	e, ϵ
E	b	t
E'	e, ϵ	e, ϵ

~~SSSSSSSS~~

	i	t	c	b	ϵ
S	(1)		(3)		
E				(4)	
S'		(2)	(3)		

Conflicting grammar, not LL1

3) $S \rightarrow a | (L)$
 $L \rightarrow L, S | S$

- (i) Left recursion exists, so not LL1
(ii) Removing LR,
 $L \rightarrow ,SL'$
 $L' \rightarrow ,S2' / \epsilon$

iii)

	<u>First</u>	<u>Follow</u>
S	(, a	\$, ,, a
L	(, a	,
L'	,, €	,

	a	,	()	\$
S	②		①		
L	③		③		
L'		④			

accept

∴ LL1 grammar

4) $S \rightarrow W$
 $W \rightarrow ZXY \mid XY$
 $Y \rightarrow c \mid \epsilon$
 $Z \rightarrow a \mid d$
 $X \rightarrow Xb \mid \epsilon$

Remove LR, $X \rightarrow X'$
 $X' \rightarrow bX' \mid \epsilon$

	<u>First</u>	<u>Follow</u>
S	a, c, d, €	\$
W	a, b, c, d, €	\$
X	b, €	c, b, \$
Y	c, €	ⓐ \$
Z	a, d	b, c, \$

	a	b	c	d	\$
S	(1)	(1)	(1)	(1)	(1)
W	(2)	(3)	(3)	(2)	(3)
X		(5)(9)	(9)		
Y			(4)		
Z	(6)			(7)	

Ambiguous grammar, not LL1

- 5)
- $S \rightarrow Ae$
 $A \rightarrow CbD$
 $C \rightarrow BC \mid E$
 $B \rightarrow cdD \mid acdD$
 $D \rightarrow c \mid E$

	<u>First</u>	<u>Follow</u>
S	abc	\$
A	abc	e
B	ca	abc
C	cab	b
D	cf	abce

	a	b	c	d	e	\$
S	(1)		(1)			
A	(2)	(2)	(2)			
B	(6)		(5)			
C	(4)	(5)	(4)			
D	(8)	(8)	(8)(7)	(8)		

Conflicting rules, not LL1