		Data	Structure Midterm Examinat	ion
	Class:	ID	Name	上課班級:□甲□乙
Pai	rt I Choose the best answ	ver (20%)		
(Qu	uestion 1~5)On a byte-ad	dressable machine, A	is a 2-dimensional array A[m	[n], the location of A [30][20] is 3271660, and
A [1	10][40] is 3273220. Assu	me that each element	occupies 2 bytes.	
1.	The memory allo	ction of array A is (a)r	row-major (b)column-major (c	e)first-in-first-out (d)last-in-first-out (e)none of
	above.			
2.	The row size m is	? (a)10 (b)20 (c)30 (d	1)40 (e)can't determine.	
3.	The column size n	is? (a)10 (b)20 (c)30) (d)40 (e) can't determine.	
4.	What is the address	ss of array beginning?	(a)3270000 (b)3270010 (c)32	270020 (d)3270030 (e) can't determine.
5.	What is the location	on of A [15][5]? (a)32°	70410 (b)3270420 (c)3270430) (d)3270440 (e) can't determine.
(Qu	uestion 6~10) The Knuth-	-Morris-Pratt string-s	earching algorithm (or KMP a	algorithm) searches for occurrences of a pattern
P =	p ₀ p ₁ p ₂ p ₃ p ₄ within a m	nain "text string" S =s ₀	s ₁ s ₂ s ₃ s ₄ by employing the	observation that when a mismatch occurs (failue
fun	ction as below), the patte	ern P itself embodies s	sufficient information to deter	mine where the next match could begin, thus
byp	bassing re-examination of	f previously matched	characters. If a partial match is	s found such that s_{i-j} $s_{i-1} = p_0 p_1 p_{j-1}$ and $s_i \neq$
p_j t	hen matching may be res	umed by comparing s	and $p_{f(j-1)+1}$ if $j \neq 0$. If $j = 0$, t	then we may continue to compare s_{i+1} and p_0 . If
pati	tern P is aabbccaabbca a	\mathbf{b} and text \mathbf{S} is aabbec	caabbcab	
	largest k such	that $k < j$ and p_0	$p_1p_k = p_{:k}p_{:k+1}p_{:k}$ if su	$ch \ a \ k \ge 0$ exists
	$f(j) = \begin{cases} -1 \end{cases}$		$p_1p_k = p_{j-k}p_{j-k+1}p_j$ if su otherwise.	
6.	f(4)=?(a)-1(b)0	(c)1 (d)2 (e)3		
	f(7)=? (a)-1 (b)0 (
	f(12)=? (a)-1 (b)0			
9.			, and next comaparsion is s_x	and p_y . What is the x=? (a)11 (b)12 (c)13 (d)14
	(e)15	(1	
10.	What is the y=? (a	a)0 (b)1 (c)2 (d)3 (e)4		
	、	, , , , , , , , ,		
Par	rt II Blanks Filling(20%))		
1.	. ,		e, i.e., each byte is accessed vi	ia an address. Let the locations of an array be
	allocated in a row-majo	or manner, and each el	ement takes one byte. What is	is the address of the element $A[5][15][25]$ in an
	=			lress of the first byte of the array is 2222. If the
	•		he address of the element abo	•
2.		-		and satisfy the following criteria, zero or more
	quantities input,			
3.				cle) are requirements,,
			_,	-
Pai	rt III Essay Questions (t	wo quesitions are sel	ected)	
1.	(30%)(a) One way to re	epresent a polynomial	is to use an array to hold the	non-zero terms of polynomial by format
	(coefficient, exponent).	Please use the forma	t to represent the polynomials	$A(x)=2x^5+5x^3-2x$ and $B(x)=7x^8+5x^3-3x^2+2x$
	-			(poly B) to caculate $C(x) = A(x) - B(x)$, when
	polynomials are represe			
((b)Design an algorithm fo	or polynomial addition	n when polynomials are imple	mented by linked list.

- (c)How to delete a linked list in O(1)?
- (d)Please describe the steps of deletion a node X (e.g. $5x^3$ in B(x)) from a duoble linked list and insertion a node Y after node Z (e.g. insert $4x^2$ after $5x^3$ in A(x)).
- 2. (20%)(a) Write the infix and prefix form of the following postfix expression, 5 35 7 / 2 * + (b) Places describe the data structure and algorithm to transfer infix expression into postfix form. You may
 - (b) Please describe the data structure and algorithm to transfer infix expression into postfix form. You must demonstrate your algorithm using expression in question (a).

- (c) Please describe the data structure and algorithm to evaluate postfix expression. You must demonstrate your algorithm using expression in question (a).
- 3. (30%) Sparse matrix by triple <row,col,value> style is adopted for computation efficiency. (a)

 Show the example according to right side matrix.(Generate array **a** []).

 (b) Please describe the sparse matrix addition algorithm. Time complexity of the algorithm must be

O(m+n) that m,n are non-zero item number in each sparse matrix.

(c) The time complexity of Fast_transpose algorithm for sparse matrix transposing is O(terms+columns) (hint: using row_terms and starting_pos array). Please show the Fast_transpose algorithm and analyze the time complexity. You must demonstrate your algorithm by using expression in question (a).

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- 4. (20%) (a)Describe the principles of popular data structures 'Stack' and 'Queue'.(b)Implement a Stack and a Queue using array. Do implementation again using linked list. Compare the advantages in two implementation methods.
- 5. (20%) (a) What are the definitions of asymptotic notations, O (big oh), Ω (Omega) and and Θ (Theta), which are used for time complexity comparison?
 - (b)Please derive the corresponding time complexity (Big-Oh) for each of the following five program segments.

(1)	(2)Hanoi(from,to,temp,n){	(3)	(4)	(5)
k=0;	if(n<=1){ cout<<"1"< <from<<"->"<<to<<endl;< td=""><td>k=0;</td><td>float rsum (float *a, int n) {</td><td>int ft(int n){</td></to<<endl;<></from<<"->	k=0;	float rsum (float *a, int n) {	int ft(int n){
for(i=0;i< N;i++)	return;}	for(i=1;i <n;i*=2)< td=""><td>if $(n \le 0)$ return 0;</td><td>int ac=$1,c=0$;</td></n;i*=2)<>	if $(n \le 0)$ return 0 ;	int ac= $1,c=0$;
for(j=0;j< i;j++)	else{Hanoi(from,temp,to,n-1);	for(j=0;j< i;j++)	else	while(ac <n){< td=""></n){<>
k++;	cout< <n<<from<<"->"<<to<<endl;< td=""><td>k++;</td><td>return (rsum(a, n-1) +</td><td>ac*=2; c++;</td></to<<endl;<></n<<from<<"->	k++;	return (rsum(a, n-1) +	ac*=2; c++;
	Hanoi(temp,to,from,n-1);		a[n-1]);	}
	}		}	return c;
				}