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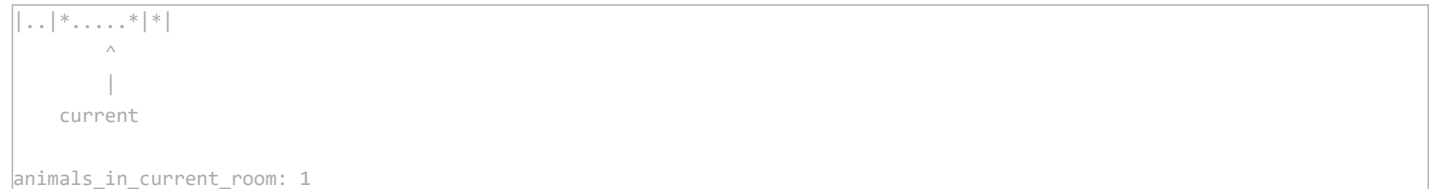
# Animal separation

Consider a sequence **S** of characters describing an animal shelter. The characters are **|** (viewed as “wall”), **.** (viewed as “nothing”), and **\*** (viewed as “animal”). The first and last characters of **S** are always a wall.

I want to check that the shelter is “safe”, i.e., all animals are separated by walls. For instance, **|..|\*|\*|\*|\*|..|** is “safe”, but **|\*\*|** is not, and neither is **|..|\*.....\*|\*|**.

Assume the sequence **S** is given as a linked list; each **Node** in the linked list has two instance variables: **c** is the character stored at this node (so **c** is one of **|**, **.**, **\***), and **next** is a reference to the following node.

To solve the safety problem, my program uses two nested **while**-loops; resetting an animal counter each time it “starts a new room”. A pointer **current** to the current Node will advance to the right, a counter **animals\_in\_current\_room** counts the animals encountered since the last **|**. In the middle of an execution, the situation could look like this:



Here is a high-level sketch, it assumes that **current** begins by pointing to the first (leftmost) Node:

```

animals_in_current_room = 0
while True:
    while current.c != '|':
        if current.c == '*':
            animals_in_current_room += 1
        if animals_in_current_room == 2:
            exit("unsafe!")
        current = current.next
    animals_in_current_room = 0
    if current.next != null:
        current = current.next
    else:
        break
print("safe!")
  
```

You are strongly encouraged to understand this on the level of pseudocode. As a service of *highly* questionable usefulness, here are also two minimal implementations in Java and Python that make the above idea concrete:



```
public class Animals
{
    static class Node {
        char c;
        Node next;
    }

    static Node build(String C) {
        Node prev = null;
        for (int i = C.length() - 1; i >= 0; --i) {
            Node fresh = new Node();
            fresh.next = prev;
            fresh.c = C.charAt(i);
            prev = fresh;
        }
        return prev;
    }

    public static void main(String[] args) {
        Node S = build("|**...|..*|*..|..|");

        Node current = S;
        int animals_in_current_room = 0;
        while (true) {
            while (current.c != '|') {
                if (current.c == '*')
                    animals_in_current_room += 1;
                if (animals_in_current_room == 2) {
                    System.out.println("unsafe!");
                    return;
                }
            }
            current = current.next;
        }
        animals_in_current_room = 0;
        if (current.next != null)
            current = current.next;
        else
            break;
    }
    System.out.println("safe!");
}
```



```

class Node:
    slots = ['c', 'succ']
    def __init__(self, c, next):
        self.c = c
        self.next = next

def build(chars) -> Node:
    prev = None
    for c in reversed(chars):
        prev = Node(c, prev)
    return prev

S = build("|...|..*|*..|..|")

current = S
animals_in_current_room = 0
while True:
    while current.c != '|':
        if current.c == '*':
            animals_in_current_room += 1
        if animals_in_current_room == 2:
            print("unsafe!")
            exit(0)
        current = current.next
    animals_in_current_room = 0
    if current.next:
        current = current.next
    else:
        break
print("safe!")

```

## Question 17

Not yet answered

Marked out of 1.00

Professor Ynot thinks this is a silly implementation and suggests the following idea:

"First, transform **S** into a string or an array of characters. That takes linear time and allows constant-time access to the positions using **S[i]** or **S.charAt(i)**. And then I can just iterate over all pairs of **\*** and check that there is a **|** between them."

Explain to the good professor the error of their ways. Mark all the correct responses.

Select one or more:

- ☐ a. Huh? That's basically the same solution; I'm just using lists instead of arrays. There is no difference in running time.
- ☐ b. There's no way you can transform the given linked list into an array in linear time, silly man! You'd be wasting all your time in the very first step.
- ☒ c. It sounds to me like your solution uses cubic time (in the length of ``s``) in the worst case, unless you also use a symbol table or something. That's really slow.
- ☐ d. My answer is faster because linked lists *also* have constant-time access and use memory much more efficiently.
- ☒ e. Wasteful cretin! That takes quadratic time *at least*! If there are **k** many **\***s in the input then you'd have to check  $\binom{k}{2}$  pairs of animals, and since **k** can be linear in **n**, the value of  $\binom{k}{2} \sim \frac{1}{2}k^2$  is quadratic in **n**.

## Question 18

Not yet answered

Marked out of 1.00

What would be an appropriate *cost model* to analyse Ynot's algorithm?

Select one or more:

- ☒ a. Number of comparisons (including `!=` and `==`, on integers or characters)
- ☐ b. Number of `print` statements
- ☒ c. Number of assignments
- ☐ d. Number of multiplications
- ☒ e. Number of array accesses
- ☐ f. Number of times a pointer (reference) is followed.
- ☐ g. Number of function calls

## Question 19

Not yet answered

Marked out of 1.00

What is the largest number of iterations of a single execution of the inner ``while``-loop in the worst case, and how can such a worst case input look?

Select one or more:

- ☐ a. Constant
- ☐ b. Logarithmic
- ☒ c. Linear
- ☐ d. Linearithmic
- ☐ e. Quadratic
- ☐ f. `|.....|`
- ☐ g. `|*|*|*|*|...|*|`
- ☐ h. `|*****...**|`



## Question 20

Not yet answered

Marked out of 1.00

What is the worst-case running time for our algorithm, and how can such a worst case input look?

Select one or more:

- ☐ a. Constant
- ☐ b. Logarithmic
- ☐ c. Linear
- ☐ d. Linearithmic
- ☒ e. Quadratic
- ☐ f. | ..... |
- ☐ g. | \* | \* | \* | \* | ... | \* |
- ☐ h. | \* \* \* \* \* ... \* \* |

## Question 21

Not yet answered

Marked out of 1.00

What is the best-case running time for our algorithm, and how can such a best-case input look?

Select one or more:

- ☒ a. Constant
- ☐ b. Logarithmic
- ☐ c. Linear
- ☐ d. Linearithmic
- ☐ e. Quadratic
- ☐ f. | ..... |
- ☐ g. | \* | \* | \* | \* | ... | \* |
- ☐ h. | \* \* \* \* \* ... \* \* |



## Question 22

Not yet answered

Marked out of 1.00

Which of the following arguments for the worst-case running time of our algorithm are both true and relevant?

Select one:

- ☐ a. We can speed up the running time by sorting the input.
- ☐ b. The number of nodes is halved in every iteration. This can happen only a logarithmic number of times.
- ☐ c. The number of pairs of animals in the input is at most quadratic.
- ☐ d. The inner loop takes at most linear time each time it is executed. Two nested loops with  $t_1$  and  $t_2$  may iterations, respectively, take at most  $\min(t_1, t_2)$  many iterations in total. Therefore, the whole programme takes linear time.
- ☒ e. In each iteration of the inner loop, **current** advances by one node, and **current** is never reset. Since there are exactly as many nodes as the input size, this can happen at most a linear number of times.
- ☐ f. The **break** statement must be executed sooner or later, otherwise the outer loop would run indefinitely.
- ☐ g. There can be at most a linear number of animals in the input.

[Clear my choice](#)

