## CSC 225 SUMMER 2016 ALGORITHMS AND DATA STRUCTURES I ASSIGNMENT 3 UNIVERSITY OF VICTORIA

- 1. Consider a version of the deterministic quick-sort algorithm that uses the element at rank  $\lfloor n/2 \rfloor$  as the pivot for a sequence on n elements.
  - (a) What is the running time of this version on a sequence that is already sorted?
  - (b) Describe the kind of sequence that would cause this version of quick-sort to run in  $\Theta(n^2)$  time.
- 2. Illustrate the performance of the heap-sort algorithm on the following input sequence, *S*, by drawing the heap tree after each insert() and removeMin() call. That is, there should be 20 trees, each the final result of the given operation after bubbling is complete.

$$S = (2, 5, 16, 4, 10, 23, 39, 18, 26, 15)$$

- 3. Design algorithms for the following operations for a node v in a binary tree T:
  - (a) preorderNext( $\nu$ ): return the node visited after  $\nu$  in a preorder traversal of T.
  - (b) inorderNext( $\nu$ ): return the node visited after  $\nu$  in an inorder traversal of T.
  - (c) postorderNext( $\nu$ ): return the node visited after  $\nu$  in a postorder traversal of T.

What are the worst-case running times of your algorithms?

- 4. Define the *internal path length*, I(T), of a tree T to be the sum of the depths of all the internal nodes in T. Likewise, define the *external path length*, E(T), of a tree T to be the sum of the depths of all the external nodes in T. Show that if T is a proper binary tree with n internal nodes, then E(T) = I(T) + 2n.
- 5. Suppose we are given a sequence S of n elements, each of which is an integer in the range  $[0, n^3 1]$ . Describe a simple method for sorting S in O(n) time.

**Hint:** Think of an alternate way of representing the elements.