

CSC 225 SUMMER 2016
ALGORITHMS AND DATA STRUCTURES I
ASSIGNMENT 1
UNIVERSITY OF VICTORIA

1. Order the following list of functions by their big-Oh notation. Group together (for example, by underlining) those functions that are big-Theta of one another. (No justification needed).

Note: $\log n = \log_2 n$ unless otherwise stated.

$$\begin{array}{ccccc}
 6n \log n & 2^{100} & \log \log n & \log^2 n & 2^{\log n} \\
 2^{2^n} & \sqrt{n} & n^{0.01} & 1/n & 4n^{3/2} \\
 3n^{0.5} & 5n & 2n \log^2 n & 2^n & n \log_4 n \\
 4^n & n^3 & n^2 \log n & 4^{\log n} & \sqrt{\log n}
 \end{array}$$

Hint: When in doubt about two functions $f(n)$ and $g(n)$, consider $\log f(n)$ and $\log g(n)$ or $2^{f(n)}$ and $2^{g(n)}$.

2. Justify the fact that if $d(n)$ is $O(f(n))$ and $e(n)$ is $O(g(n))$, then the product $d(n)e(n)$ is $O(f(n)g(n))$.

3. Show that $\log_b f(n)$ is $\Theta(\log_2 f(n))$ if $b > 1$ is a constant.

Hint: One of the properties of logarithms from the slide I gave you is very useful here.

4. Consider the recurrence equation,

$$T(n) = \begin{cases} 1 & \text{if } n = 0 \\ T(n-1) + 2^n & \text{otherwise.} \end{cases}$$

Show, by induction, that $T(n) = 2^{n+1} - 1$.

5. Consider the Algorithm `arrayFind`, given below, which searches an array A for an element x .

Algorithm `arrayFind`(x, A):

Input: An element x and an n -element array, A .

Output: The index i such that $x = A[i]$ or -1 if no element of A is equal to x .

```

i ← 0
while i < n do
  if x = A[i] then
    return i
  else
    i ← i + 1
return -1

```

- (a) Counting assignments, comparisons, and returns only, calculate the worst-case, $T(n)$, and best-case, $T_b(n)$, running times of `arrayFind`.

- (b) Prove by induction (loop invariants) that `arrayFind` is correct.