

Problem 2858: Minimum Edge Reversals So Every Node Is Reachable

Problem Information

Difficulty: **Hard**

Acceptance Rate: 0.00%

Paid Only: No

Problem Description

There is a

simple directed graph

with

n

nodes labeled from

0

to

$n - 1$

. The graph would form a

tree

if its edges were bi-directional.

You are given an integer

n

and a

2D

integer array

edges

, where

$\text{edges}[i] = [u$

i

, v

i

$]$

represents a

directed edge

going from node

u

i

to node

v

i

.

An

edge reversal

changes the direction of an edge, i.e., a directed edge going from node

u

i

to node

v

i

becomes a directed edge going from node

v

i

to node

u

i

.

For every node

i

in the range

$[0, n - 1]$

, your task is to

independently

calculate the

minimum

number of

edge reversals

required so it is possible to reach any other node starting from node

i

through a

sequence

of

directed edges

.

Return

an integer array

answer

, where

answer[i]

is the

minimum

number of

edge reversals

required so it is possible to reach any other node starting from node

i

through a

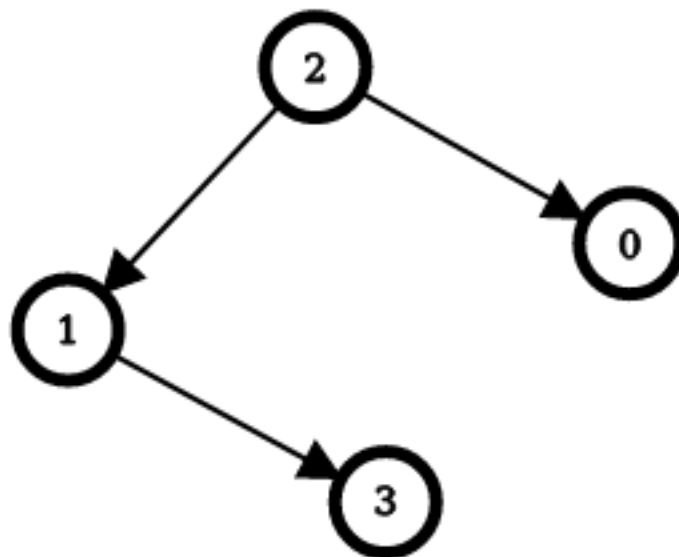
sequence

of

directed edges

.

Example 1:



Input:

$n = 4$, edges = $[[2,0],[2,1],[1,3]]$

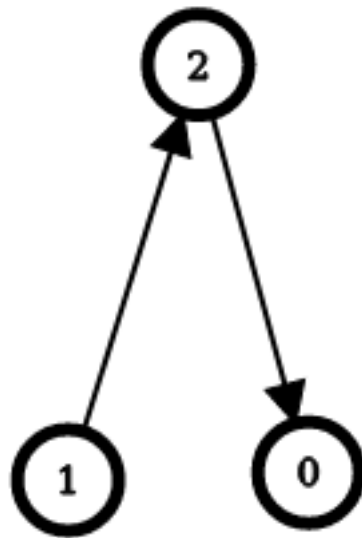
Output:

[1,1,0,2]

Explanation:

The image above shows the graph formed by the edges. For node 0: after reversing the edge [2,0], it is possible to reach any other node starting from node 0. So, $\text{answer}[0] = 1$. For node 1: after reversing the edge [2,1], it is possible to reach any other node starting from node 1. So, $\text{answer}[1] = 1$. For node 2: it is already possible to reach any other node starting from node 2. So, $\text{answer}[2] = 0$. For node 3: after reversing the edges [1,3] and [2,1], it is possible to reach any other node starting from node 3. So, $\text{answer}[3] = 2$.

Example 2:



Input:

$n = 3$, edges = [[1,2],[2,0]]

Output:

[2,0,1]

Explanation:

The image above shows the graph formed by the edges. For node 0: after reversing the edges [2,0] and [1,2], it is possible to reach any other node starting from node 0. So, $\text{answer}[0] = 2$. For node 1: it is already possible to reach any other node starting from node 1. So, $\text{answer}[1] = 0$. For node 2: after reversing the edge [1, 2], it is possible to reach any other node starting from node 2. So, $\text{answer}[2] = 1$.

Constraints:

$2 \leq n \leq 10$

5

$\text{edges.length} == n - 1$

$\text{edges}[i].\text{length} == 2$

$0 \leq u$

i

$== \text{edges}[i][0] < n$

$0 \leq v$

i

$== \text{edges}[i][1] < n$

u

i

$!= v$

i

The input is generated such that if the edges were bi-directional, the graph would be a tree.

Code Snippets

C++:

```
class Solution {
public:
    vector<int> minEdgeReversals(int n, vector<vector<int>>& edges) {

    }
};
```

Java:

```
class Solution {
    public int[] minEdgeReversals(int n, int[][] edges) {

    }
}
```

Python3:

```
class Solution:
    def minEdgeReversals(self, n: int, edges: List[List[int]]) -> List[int]:
```

Python:

```
class Solution(object):
    def minEdgeReversals(self, n, edges):
        """
        :type n: int
        :type edges: List[List[int]]
        :rtype: List[int]
        """
```

JavaScript:

```
/**
 * @param {number} n
 * @param {number[][]} edges
 * @return {number[]}
 */
var minEdgeReversals = function(n, edges) {
```



```
};
```

TypeScript:

```
function minEdgeReversals(n: number, edges: number[][]): number[] {  
  
};
```

C#:

```
public class Solution {  
    public int[] MinEdgeReversals(int n, int[][] edges) {  
  
    }  
}
```

C:

```
/**  
 * Note: The returned array must be malloced, assume caller calls free().  
 */  
int* minEdgeReversals(int n, int** edges, int edgesSize, int* edgesColSize,  
int* returnSize) {  
  
}
```

Go:

```
func minEdgeReversals(n int, edges [][]int) []int {  
  
}
```

Kotlin:

```
class Solution {  
    fun minEdgeReversals(n: Int, edges: Array<IntArray>): IntArray {  
  
    }  
}
```

Swift:

```

class Solution {
  func minEdgeReversals(_ n: Int, _ edges: [[Int]]) -> [Int] {

  }
}

```

Rust:

```

impl Solution {
  pub fn min_edge_reversals(n: i32, edges: Vec<Vec<i32>>) -> Vec<i32> {

  }
}

```

Ruby:

```

# @param {Integer} n
# @param {Integer[][]} edges
# @return {Integer[]}
def min_edge_reversals(n, edges)

end

```

PHP:

```

class Solution {

  /**
   * @param Integer $n
   * @param Integer[][] $edges
   * @return Integer[]
   */
  function minEdgeReversals($n, $edges) {

  }
}

```

Dart:

```

class Solution {
  List<int> minEdgeReversals(int n, List<List<int>> edges) {

  }
}

```

```
}
```

Scala:

```
object Solution {  
  def minEdgeReversals(n: Int, edges: Array[Array[Int]]): Array[Int] = {  
  
  }  
}
```

Elixir:

```
defmodule Solution do  
  @spec min_edge_reversals(n :: integer, edges :: [[integer]]) :: [integer]  
  def min_edge_reversals(n, edges) do  
  
  end  
end
```

Erlang:

```
-spec min_edge_reversals(N :: integer(), Edges :: [[integer()]]) ->  
  [integer()].  
min_edge_reversals(N, Edges) ->  
  .
```

Racket:

```
(define/contract (min-edge-reversals n edges)  
  (-> exact-integer? (listof (listof exact-integer?)) (listof exact-integer?))  
  )
```

Solutions

C++ Solution:

```
/*  
 * Problem: Minimum Edge Reversals So Every Node Is Reachable  
 * Difficulty: Hard  
 * Tags: array, tree, graph, dp, search
```

```

*
* Approach: Use two pointers or sliding window technique
* Time Complexity: O(n) or O(n log n)
* Space Complexity: O(n) or O(n * m) for DP table
*/

class Solution {
public:
vector<int> minEdgeReversals(int n, vector<vector<int>>& edges) {

}
};

```

Java Solution:

```

/**
* Problem: Minimum Edge Reversals So Every Node Is Reachable
* Difficulty: Hard
* Tags: array, tree, graph, dp, search
*
* Approach: Use two pointers or sliding window technique
* Time Complexity: O(n) or O(n log n)
* Space Complexity: O(n) or O(n * m) for DP table
*/

class Solution {
public int[] minEdgeReversals(int n, int[][] edges) {

}
}

```

Python3 Solution:

```

"""
Problem: Minimum Edge Reversals So Every Node Is Reachable
Difficulty: Hard
Tags: array, tree, graph, dp, search

Approach: Use two pointers or sliding window technique
Time Complexity: O(n) or O(n log n)
Space Complexity: O(n) or O(n * m) for DP table

```

```

"""

class Solution:
    def minEdgeReversals(self, n: int, edges: List[List[int]]) -> List[int]:
        # TODO: Implement optimized solution
        pass

```

Python Solution:

```

class Solution(object):
    def minEdgeReversals(self, n, edges):
        """
        :type n: int
        :type edges: List[List[int]]
        :rtype: List[int]
        """

```

JavaScript Solution:

```

/**
 * Problem: Minimum Edge Reversals So Every Node Is Reachable
 * Difficulty: Hard
 * Tags: array, tree, graph, dp, search
 *
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 * Time Complexity: O(n) or O(n log n)
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 */

/**
 * @param {number} n
 * @param {number[][]} edges
 * @return {number[]}
 */
var minEdgeReversals = function(n, edges) {

};

```

TypeScript Solution:

```

/**
 * Problem: Minimum Edge Reversals So Every Node Is Reachable
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 * Tags: array, tree, graph, dp, search
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 * Approach: Use two pointers or sliding window technique
 * Time Complexity: O(n) or O(n log n)
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 */

function minEdgeReversals(n: number, edges: number[][]): number[] {

};

```

C# Solution:

```

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 * Approach: Use two pointers or sliding window technique
 * Time Complexity: O(n) or O(n log n)
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 */

public class Solution {
    public int[] MinEdgeReversals(int n, int[][] edges) {

    }
}

```

C Solution:

```

/*
 * Problem: Minimum Edge Reversals So Every Node Is Reachable
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```

*/

/**
 * Note: The returned array must be malloced, assume caller calls free().
 */
int* minEdgeReversals(int n, int** edges, int edgesSize, int* edgesColSize,
int* returnSize) {

}

```

Go Solution:

```

// Problem: Minimum Edge Reversals So Every Node Is Reachable
// Difficulty: Hard
// Tags: array, tree, graph, dp, search
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// Approach: Use two pointers or sliding window technique
// Time Complexity: O(n) or O(n log n)
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func minEdgeReversals(n int, edges [][]int) []int {

}

```

Kotlin Solution:

```

class Solution {
    fun minEdgeReversals(n: Int, edges: Array<IntArray>): IntArray {

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class Solution {
    func minEdgeReversals(_ n: Int, _ edges: [[Int]]) -> [Int] {

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}

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Rust Solution:

```

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// Difficulty: Hard
// Tags: array, tree, graph, dp, search
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// Approach: Use two pointers or sliding window technique
// Time Complexity: O(n) or O(n log n)
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impl Solution {
pub fn min_edge_reversals(n: i32, edges: Vec<Vec<i32>>) -> Vec<i32> {

}
}

```

Ruby Solution:

```

# @param {Integer} n
# @param {Integer[][]} edges
# @return {Integer[]}
def min_edge_reversals(n, edges)

end

```

PHP Solution:

```

class Solution {

/**
 * @param Integer $n
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 * @return Integer[]
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function minEdgeReversals($n, $edges) {

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Dart Solution:

```

class Solution {
List<int> minEdgeReversals(int n, List<List<int>> edges) {

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}  
}
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object Solution {  
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-spec min_edge_reversals(N :: integer(), Edges :: [[integer()]]) ->  
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