

Problem 2642: Design Graph With Shortest Path Calculator

Problem Information

Difficulty: **Hard**

Acceptance Rate: 64.46%

Paid Only: No

Tags: Graph, Design, Heap (Priority Queue), Shortest Path

Problem Description

There is a **directed weighted** graph that consists of n nodes numbered from 0 to $n - 1$. The edges of the graph are initially represented by the given array `edges` where `edges[i] = [fromi, toi, edgeCosti]` meaning that there is an edge from `fromi` to `toi` with the cost `edgeCosti`.

Implement the `Graph` class:

`Graph(int n, int[][] edges)` initializes the object with n nodes and the given edges. `addEdge(int[] edge)` adds an edge to the list of edges where `edge = [from, to, edgeCost]`. It is guaranteed that there is no edge between the two nodes before adding this one. `shortestPath(int node1, int node2)` returns the **minimum** cost of a path from `node1` to `node2`. If no path exists, return `-1`. The cost of a path is the sum of the costs of the edges in the path.

Example 1:



Input `["Graph", "shortestPath", "shortestPath", "addEdge", "shortestPath"]` `[[4, [[0, 2, 5], [0, 1, 2], [1, 2, 1], [3, 0, 3]], [3, 2], [0, 3], [[1, 3, 4]], [0, 3]]` **Output** `[null, 6, -1, null, 6]`
Explanation `Graph g = new Graph(4, [[0, 2, 5], [0, 1, 2], [1, 2, 1], [3, 0, 3]]);`
`g.shortestPath(3, 2);` // return 6. The shortest path from 3 to 2 in the first diagram above is `3 -> 0 -> 1 -> 2` with a total cost of `3 + 2 + 1 = 6`. `g.shortestPath(0, 3);` // return -1. There is no path from 0 to 3. `g.addEdge([1, 3, 4]);` // We add an edge from node 1 to node 3, and we get the second diagram above. `g.shortestPath(0, 3);` // return 6. The shortest path from 0 to 3 now is 0

-> 1 -> 3 with a total cost of $2 + 4 = 6$.

****Constraints:****

* `1` $\leq n \leq 100$ * `0` $\leq \text{edges.length} \leq n * (n - 1)$ * `edges[i].length == edge.length == 3` *
`0` $\leq \text{fromi, toi, from, to, node1, node2} \leq n - 1$ * `1` $\leq \text{edgeCosti, edgeCost} \leq 106$ * There
are no repeated edges and no self-loops in the graph at any point. * At most `100` calls will be
made for `addEdge`. * At most `100` calls will be made for `shortestPath`.

Code Snippets

C++:

```
class Graph {
public:
    Graph(int n, vector<vector<int>>& edges) {

    }

    void addEdge(vector<int> edge) {

    }

    int shortestPath(int node1, int node2) {

    }
};

/**
 * Your Graph object will be instantiated and called as such:
 * Graph* obj = new Graph(n, edges);
 * obj->addEdge(edge);
 * int param_2 = obj->shortestPath(node1,node2);
 */
```

Java:

```
class Graph {

    public Graph(int n, int[][] edges) {
```

```

    }

    public void addEdge(int[] edge) {

    }

    public int shortestPath(int node1, int node2) {

    }
}

/**
 * Your Graph object will be instantiated and called as such:
 * Graph obj = new Graph(n, edges);
 * obj.addEdge(edge);
 * int param_2 = obj.shortestPath(node1,node2);
 */

```

Python3:

```

class Graph:

    def __init__(self, n: int, edges: List[List[int]]):

    def addEdge(self, edge: List[int]) -> None:

    def shortestPath(self, node1: int, node2: int) -> int:

    # Your Graph object will be instantiated and called as such:
    # obj = Graph(n, edges)
    # obj.addEdge(edge)
    # param_2 = obj.shortestPath(node1,node2)

```