

# 0-1 Knapsack Problem: BFS and DFS

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## 1 The problem

### 1.1 Representation of a solution

#### 1.1.1 Variables

The set of variables  $X = \{X_1, \dots, X_i, \dots, X_n\}$  where  $n$  is the number of total items and each variable represents an item to be put in the knapsack.

#### 1.1.2 Domain

The domain of the  $i$ -th variable  $D_i = \{true, false\}$  as an item can either be in the knapsack or not.

### 1.2 Objective function

To define the objective function, the parameter  $b_i$  needs to be defined as the benefit of the variable  $X_i$ . Hence:  $\max \sum_{i=1}^n X_i * b[i]$

### 1.3 Constraint

The total capacity of the knapsack  $K$  shall never be exceeded by the items selected as a solution. Each item has a weight  $w_i$ .

$$\sum_{i=1}^n X_i * w_i \leq K$$

## 2 Comparison of the algorithms

### 2.1 Time complexity

The theoretical time complexity of BFS is  $O(b^d)$  where  $b$  is the maximum branching factor and  $d$  is the depth of the shallowest solution, while for DFS it's  $O(b^m)$  where  $m$  is the maximum depth. In my case  $b = n = 11$  (since on the first expansion of the tree  $n$  children are generated) and  $d = 8$  (practical result from running the algorithm), while  $m = 11$  (the maximum depth is equal to  $n$  as the left-most node will be expanded from 1 to  $n$ ). The corresponding time

usages of the algorithms are, respectively: 0.06s and 0.09s which is in line with the theoretical outcome.

## **2.2 Space complexity**

The theoretical space complexity of BFS is  $O(b^d)$  while for DFS it's  $O(b \cdot m)$ . It's difficult to assess the actual space usage but theoretically it is lower for the latter:  $11^8 = 214358881$  for BFS and  $11 \cdot 11 = 121$  for DFS.