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[Quiz 4](#)

Started on Sunday, 26 March 2023, 5:50 PM

State Finished

Completed on Sunday, 26 March 2023, 6:15 PM

Time taken 25 mins 8 secs

Marks 5.00/7.00

Grade 7.14 out of 10.00 (71%)

Question 1

Incorrect

Mark 0.00 out of 1.00

Select the asymptotic upper and lower bounds for $T(n)$ in the following recurrence. Assume that $T(n)$ is constant for $n \leq 3$. Make your bounds as tight as possible.

$$T(n) = 7T(n/2) + n^3$$

Select one:

- ☐ a. $T(n) = \Theta(n^{7/3})$
- ☐ b. $T(n) = \Theta(n^3)$
- ☒ c. $T(n) = \Theta(n^{\log_2(7)})$
- ☐ d. $T(n) = \Theta(\log_7(2n))$

✗

Your answer is incorrect.

$$a=7, b=2, d=3$$

According to master thm $T(n) = \Theta(nd) = \Theta(n^3)$ because $(b^d = 2^3 = 8 > a)$

The correct answer is:

$$T(n) = \Theta(n^3)$$

Question 2

Correct

Mark 1.00 out of 1.00

What is/are the best data structure(s) to implement recursive function calls?

- ☐ a. Array
- ☒ b. Stack
- ☐ c. Binary Tree
- ☐ d. Linked List



The correct answer is: Stack

Question 3

Correct

Mark 1.00 out of 1.00

Solve the following Recursive Algorithm:

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ 2T(\frac{n}{2}) + F'(n) & \text{if } n>1 \end{cases}$$

Note: $F'(n)$ function is in the order of $O(n)$

- ☐ a. $T(n) = O(n)$
- ☐ b. $T(n) = O(\log(n))$
- ☒ c. $T(n) = O(n \log(n))$
- ☐ d. $T(n) = O(n^2)$



The correct answer is: $T(n) = O(n \log(n))$

Question 4

Correct

Mark 1.00 out of 1.00

For the following recurrence, select the correct expression for run time $T(n)$ if the recurrence can be solved using Master Theorem, Otherwise, indicate that the Master Theorem does not apply.

$$T(n) = 3T(n/3) + \sqrt{n}$$

- ☐ a. Master Theorem does not apply.
- ☐ b. $T(n) = \Theta(n \log n)$
- ☐ c. $T(n) = \Theta(n^2)$
- ☒ d. $T(n) = \Theta(n)$



The correct answer is: $T(n) = \Theta(n)$

Question 5

Correct

Mark 1.00 out of 1.00

Given a set 'S' of n integers and another integer x, an algorithm should determine whether or not there exists two elements in S whose sum is exactly x. A possible algorithm for this task is described below.

- 1) Sort the elements in S using any efficient sorting algorithm.
- 2) Remove the last element from S. Let y be the value of the removed element.
- 3) If S is non-empty, look whether an element z exist in S where $z = x - y$
- 4) If S contains such an element z, then stop, since we have found y and z such that $x = y + z$; otherwise repeat Step 2.
- 5) If S is empty, then no two elements in S sum to x.

Select the correct statement(s) regarding above approach.

- ☐ a. There are algorithms which can solve this task with better time complexity than above described algorithm
- ☒ b. Step 1 can be achieved through merge sort with $\Theta(n \lg n)$ time complexity.
- ☐ c. Best time complexity to do Step 3 is $\Theta(n)$.
- ☒ d. Time complexity of this algorithm is $\Theta(n \lg n)$.



The correct answers are: Step 1 can be achieved through merge sort with $\Theta(n \lg n)$ time complexity., Time complexity of this algorithm is $\Theta(n \lg n)$.

Question 6

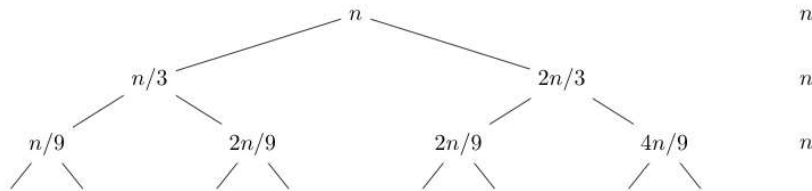
Correct

Mark 1.00 out of 1.00

Recursion Tree is one way to analyze recursive functions. Consider a function with following time complexity.

$$T(n) = T(n/3) + T(2n/3) + n$$

Following figure shows the first 3 levels of the recursion tree.



What is/are the number(s) which can be appear in the next (4th) level in this recursion tree?

- ☒ a. $n/27$
- ☐ b. $16n/27$
- ☒ c. $2n/27$
- ☒ d. $8n/27$



The correct answers are: $n/27, 2n/27, 8n/27$

Question 7

Incorrect

Mark 0.00 out of 1.00

A function is defined as following.

```
Function process(int n){
    If (n == 0) return;
    mark(n);
    process(n/3);
    process(n/3);
    process(n/3);
    process(n/3);
}
```

Assuming `mark(n)` function has the asymptotic bound of $\Theta(n^2)$, what is the asymptotic bound of process function.

Select one:

- ☐ a. $T(n) = \Theta(n^3 \log(n))$
- ☐ b. $T(n) = \Theta(n^2)$
- ☐ c. $T(n) = \Theta(n^3)$
- ☒ d. $T(n) = \Theta(n^{\log_3(4)})$

✖

Your answer is incorrect.

The recurrence of the process function is $T(n) = 4T(n/3) + \Theta(n^2)$

According to master theorem, time complexity is $\Theta(n^2)$

($a=4, b=3, d=2$ so $b^d > a$)

The correct answer is:

$T(n) = \Theta(n^2)$

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