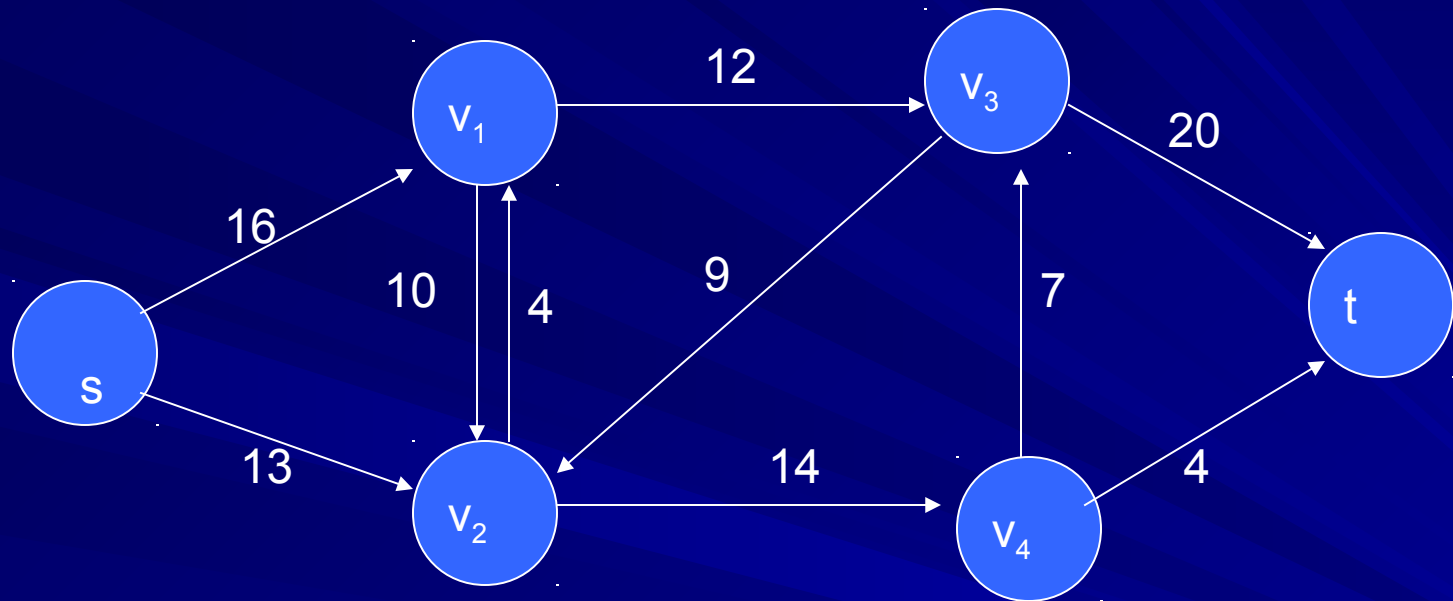


# Max-Flow Min Cut

## Algorithms

# Network



# Max Flow Applications

- Data mining.
- Open-pit mining.
- Project selection.
- Airline scheduling.
- Bipartite matching.
- Baseball elimination.
- Image segmentation.
- Network connectivity.
- Network reliability.
- Distributed computing.
- Egalitarian stable matching.
- Security of statistical data.
- Network intrusion detection.
- Multi-camera scene reconstruction

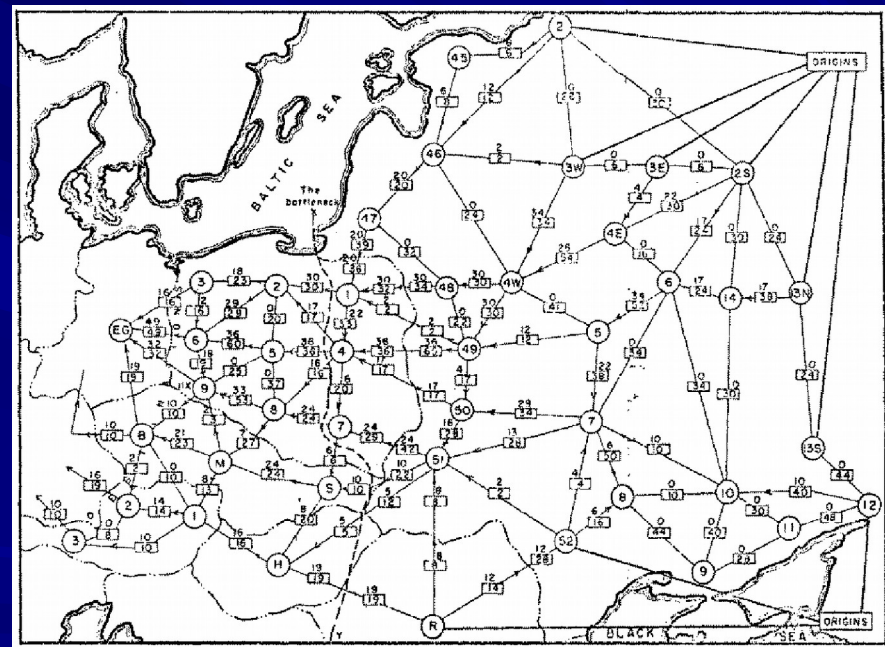
# Soviet Rail Network, 1955

"The Soviet rail system also roused the interest of the Americans, and again it inspired fundamental research in optimization."

-- Schrijver

G. Danzig\* 1951...First soln...

Again formulated by  
Harris in 1955 for the  
US Airforce  
("unclassified in 1999")  
What were they looking for?



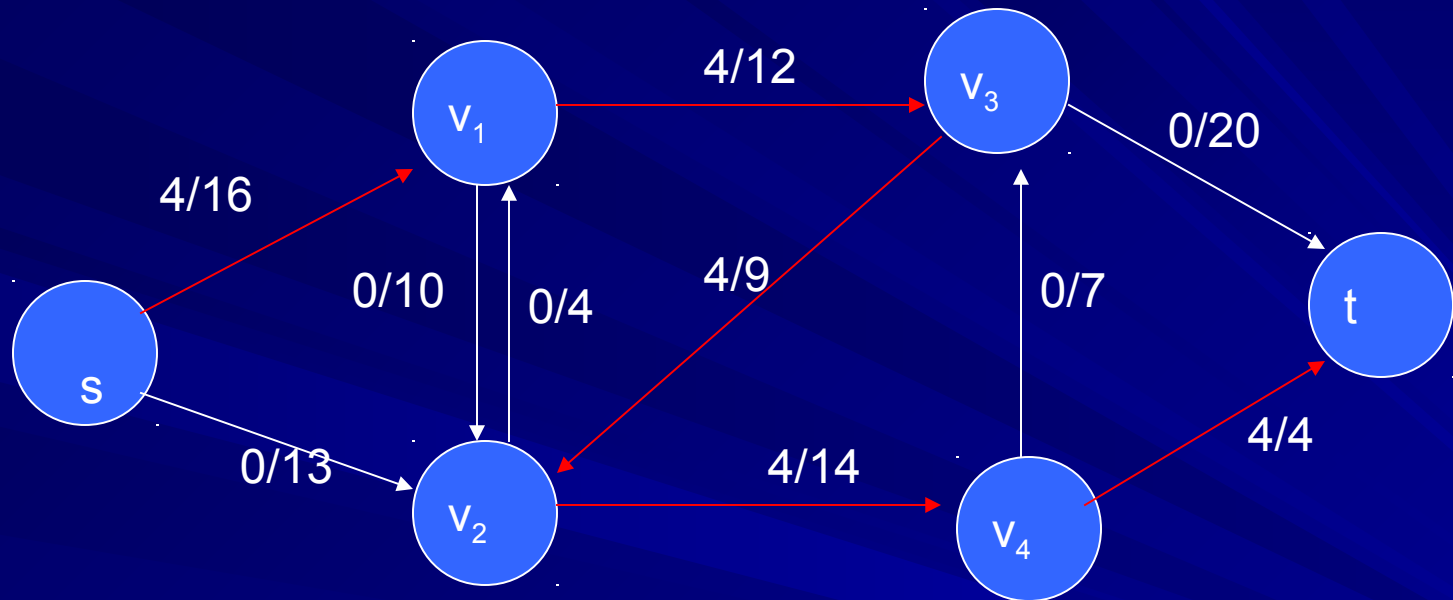
Reference: *On the history of the transportation and maximum flow problems.*  
Alexander Schrijver in Math Programming, 91: 3, 2002.

year	discoverer(s)	bound
1951	Dantzig [11]	$O(n^2mU)$
1956	Ford & Fulkerson [17]	$O(nmU)$
1970	Dinitz [13] Edmonds & Karp [15]	$O(nm^2)$
1970	Dinitz [13]	$O(n^2m)$
1972	Edmonds & Karp [15] Dinitz [14]	$O(m^2 \log U)$
1973	Dinitz [14] Gabow [19]	$O(nm \log U)$
1974	Karzanov [36]	$O(n^3)$
1977	Cherkassky [9]	$O(n^2m^{1/2})$
1980	Galil & Naamad [20]	$O(nm \log^2 n)$
1983	Sleator & Tarjan [46]	$O(nm \log n)$
1986	Goldberg & Tarjan [26]	$O(nm \log(n^2/m))$
1987	Ahuja & Orlin [2]	$O(nm + n^2 \log U)$
1987	Ahuja et al. [3]	$O(nm \log(n\sqrt{\log U}/m))$
1989	Cheriyān & Hagerup [7]	$E(nm + n^2 \log^2 n)$
1990	Cheriyān et al. [8]	$O(n^3 / \log n)$
1990	Alon [4]	$O(nm + n^{8/3} \log n)$
1992	King et al. [37]	$O(nm + n^{2+\epsilon})$
1993	Phillips & Westbrook [44]	$O(nm(\log_{m/n} n + \log^{2+\epsilon} n))$
1994	King et al. [38]	$O(nm \log_{m/(n \log n)} n)$
1997	Goldberg & Rao [24]	$O(\min(n^{2/3}, m^{1/2})m \log(n^2/m) \log U)$

# Ford-Fulkerson( $G, s, t$ )

- $F = 0$
- While there is augmenting path from  $s$  to  $t$ ,  $P$ 
  - $F = f + f_p$
  - Subtract  $f_p$  from all capacities of edges on  $P$

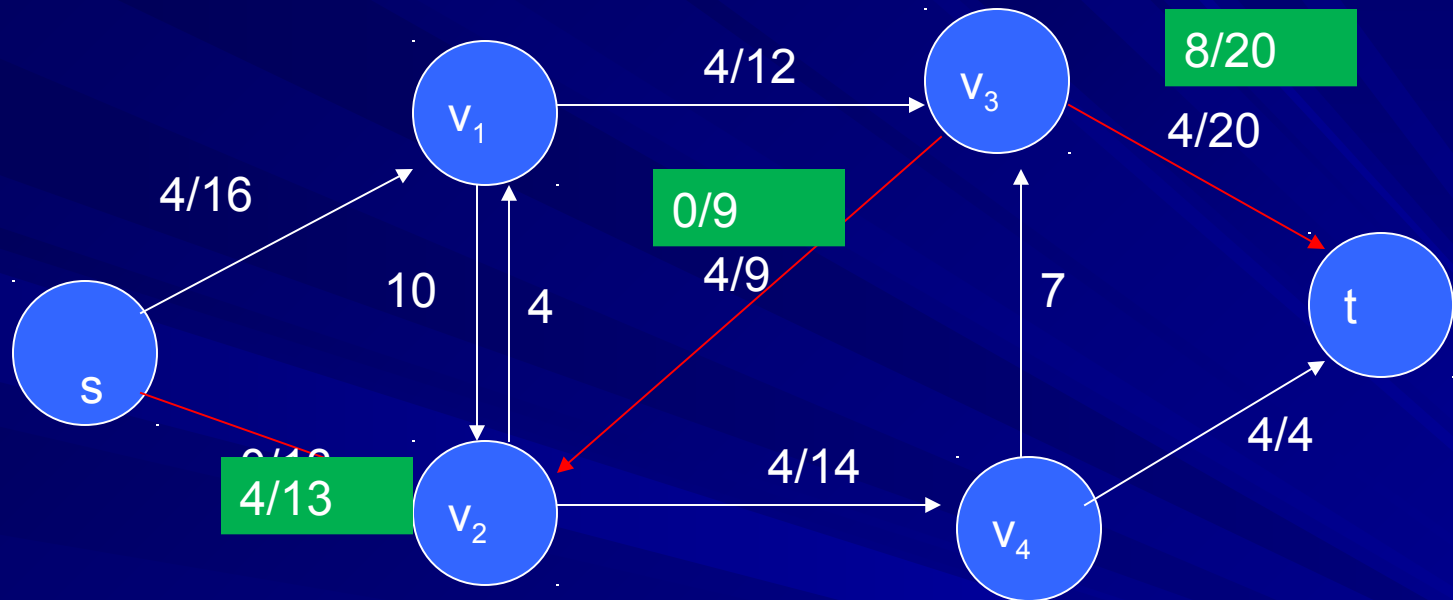
# Network



Flow of path is 4

$$F=4$$

# Network

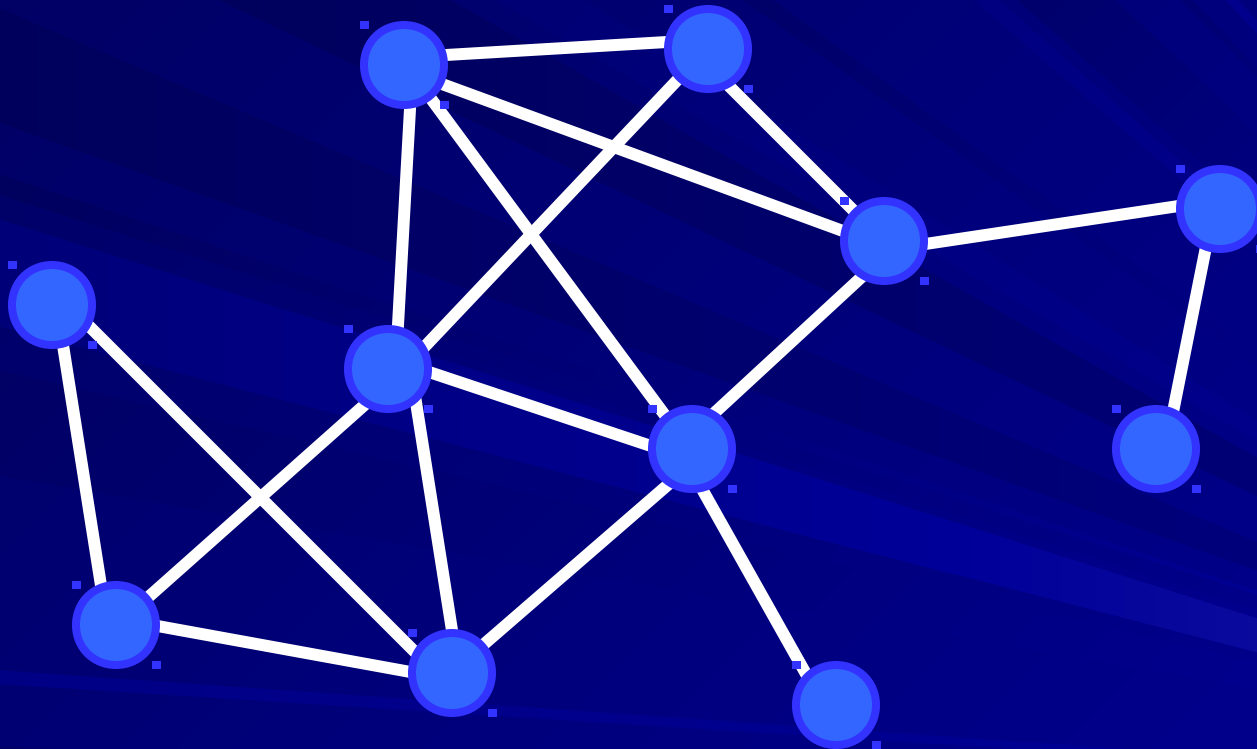


Flow of path is 4,  
Total = 8

$$F=4$$



# Matching



# Maximal Matching

## ■ Maximal Matching

- $k$  edges sharing no common vertices.
- No other edges from  $G$  can be added