Introduction to NP

Analysis of Algorithms
CSCI 5620

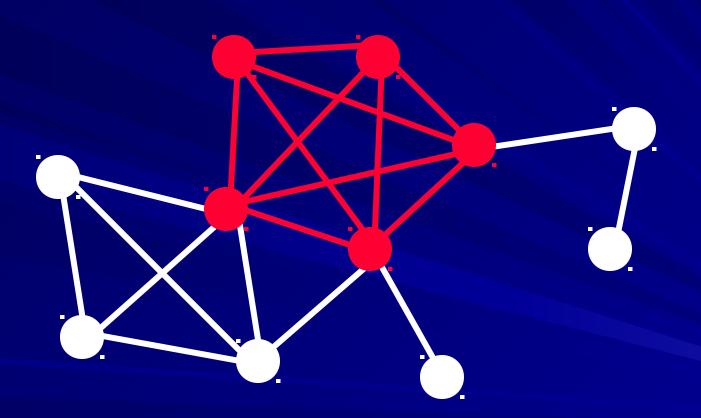
Decision Problem

- Decision problem: One for which the output can is a simple yes or no answer for each input.
- Example:
 - The input is a positive integer C and n positive integers s1, ... sn.
 - Is there a subset of the integers which add up to C?

Decision Problem

- Decision Problem
 - Does the graph have a Spanning Tree with total weight K?
- Optimization Problem
 - What is the MST weight?
- Decision Problem
 - For a given graph, is there a path from A to B of length k?
- Optimization Problem
 - What is length of a shortest path from A to B?
- Optimization Problems are at least as hard as Decision Problems.
 - So if the optimization problem is easy, then the decision problem is easy.

Clique Size



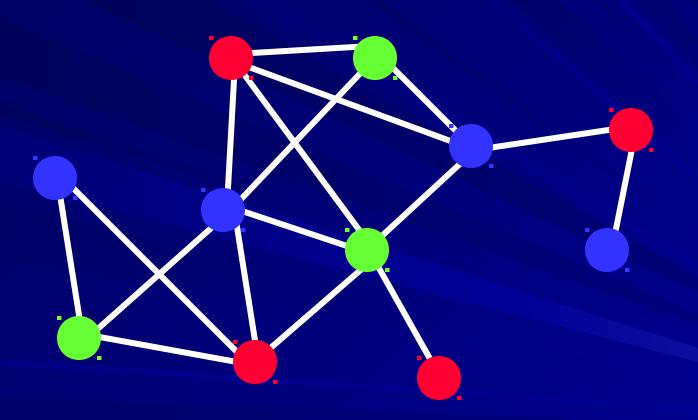
Clique Formulation

$$\max \sum_{i=1}^{n} w_i x_i,$$

s.t.
$$x_i + x_j \le 1, \ \forall \ (i, j) \in \overline{E},$$

$$x_i \in \{0, 1\}, i = 1, \dots, n.$$

Coloring



Coloring Formulation

Given a set of n colors $A = \{1, 2, ..., n\}$, the graph coloring problem associated with a graph G = (V, E) can be formulated as the following integer program:

(P) min
$$\sum_{k=1}^{n} y_k$$
,

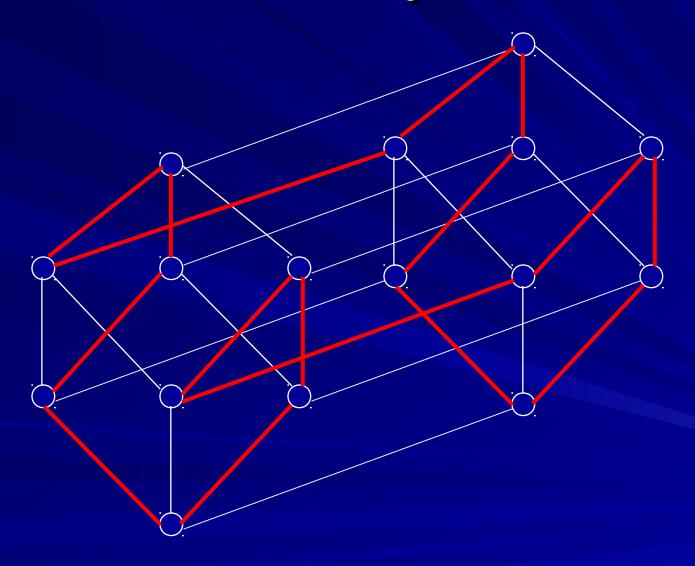
s.t.
$$\sum_{k=1}^{n} x_{ik} = 1, \qquad \forall i \in V, \tag{1}$$

$$x_{ik} + x_{jk} \le 1, \qquad \forall (i, j) \in E, \tag{2}$$

$$y_k \ge x_{ik}, \qquad \forall i \in V, k = 1, \dots, n,$$
 (3)

$$x_{ik} \in \{0, 1\}, \qquad \forall i \in V, k = 1, \dots, n.$$
 (4)

Hamiltonian Cycle



NP

- NP "informally" is the class of decision problems whose answer "yes/no" can be verified in polynomial time given a "guess" or "certificate".
- Example Does G contain a Hamiltonian cycle?
 - What is a guess?
 - Can it be verified in polynomial time?
 - How many total guesses?
 - How many does it take to prove G contains a cycle?
- Example: Is n a composite (not prime)?
 - Does d divide n? (can be checked in polynomial time)
- Example: Can graph G be colored with k colors?
 - Chromatic Number, smallest number of colors in which G can be colored (optimization problem)

Strings

- Catenation
 - If s_1 =abac and s_2 =dade, then $s_1 \cdot s_2$ or s_1s_2 =abacdade.
- Empty string, λ="
 - $-s_1 \lambda = \lambda s_1 = abac$

Alphabet / Words

- For a set A, the set A* consists of all finite sequences of elements of A.
 - $Ex A = \{0,1\}$
 - λ € A*, 0 € A*, 1 € A*, 01 € A*, 10 € A*
 - -00000000... ∉ A*
- A is called an alphabet, elements of A* are called words.
- Ex. A={a, b, c,..., z}, A* contains all words, such as "discrete", but also "ababccad".

Example

- Let A={ab, bc, ba}
 - − Is ababab € A*?
 - Is abc € A*?
 - Is abba € A*?
 - − Is abbcbaba € A*?
 - Is bcabbab € A*?
 - Is abbbcba € A*?

Regular Expressions

- A <u>regular expression over A</u>, is a string constructed from elements of A and symbols (,), v, *, λ.
 - First λ is a regular expression.
 - If x E A, then symbol x is a regular expression.
 - If α and β are regular expression then $(\alpha \vee \beta)$ is regular.
 - If α is a regular expression than (α)* is a regular expression.

Regular Expressions

- $Ex. Let A = \{0,1\}$
- 0* is a regular expression.
- (01)* is a regular expression
- (0v1)* is a regular expression
- 1(0v1)* is a regular expression
- (011)* is a regular expression
- 0 v 1* is a regular expression
- (0 is not a regular expression

Regular Expressions

Give the regular expression for the set of strings over {a,b} in which all the a's precede the b's which in turn precedes the c's.

Regular Set

- For each regular expression, there is a corresponding subset of A*, called a <u>regular set</u>.
- Example Let A={a, b}
 - For a*, the regular set Y contains
 - ■λ, a, aa, aaa, aaaa, aaaaa
 - $\P Y = \{s, s = aaa...a (repeated n times for <math>n \ge 0\}$
 - ■aaaaaaaa... ∉ Y

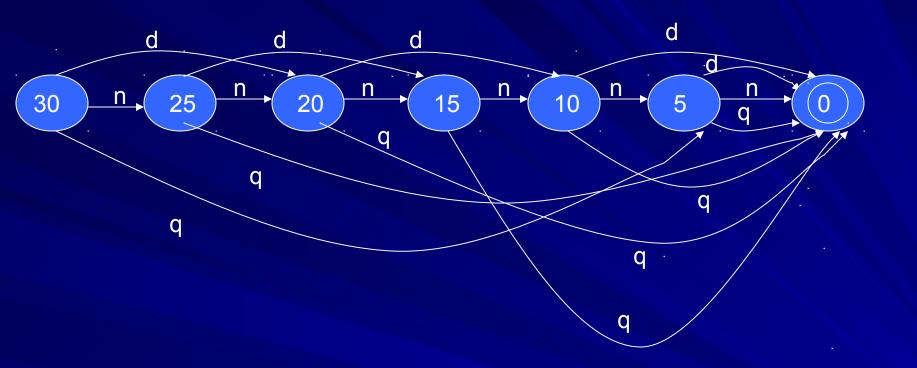
Language

- Let Σ be an alphabet.
- \blacksquare A **language** over Σ is a subset of Σ^* .
- Computer Language
 - Contains words (if, then, else, etc)
- English Language

 - Subset of Σ^*

Deterministic Finite Automata (DFA)

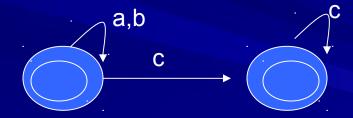
- Finite State Machine
- Vending machine
 - Simple newspaper vending machine
 - <u>- 30 cents</u>
 - No change given
 - Takes Nickels, dimes, and quarters
 - No memory used



Regular expressions

- A={a,b,c}
- a*b

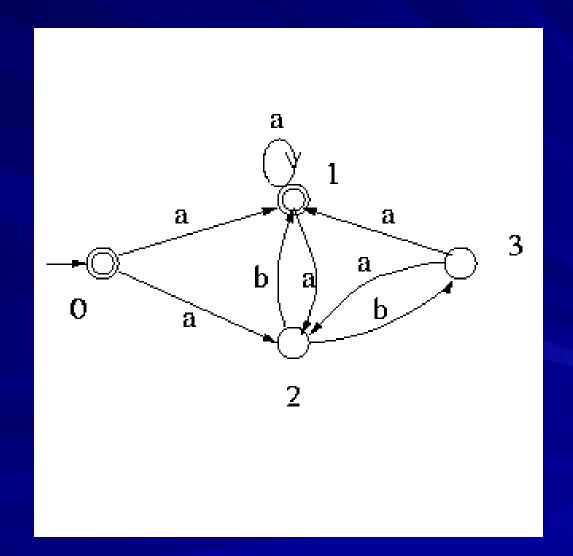
(aub)*c*



Regular Language

- Equivalent to A regular set
- Accepted by a DFA (deterministic finite automation), NFA (nondeterministic finite automation).

NFA example



DFA example

