

Translating Addresses



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Goals of Today's Lecture

Three different kinds of addresses

- Host names (e.g., www.cnn.com)

- IP addresses (e.g., 64.236.16.20)

- MAC addresses (e.g., 00-15-C5-49-04-A9)

Protocols for translating between addresses

- Domain Name System (DNS)

- Dynamic Host Configuration Protocol (DHCP)

- Address Resolution Protocol (ARP)

Two main topics

- Decentralized management of the name space

- Boot-strapping an end host that attaches to the 'net



Host Names vs. IP addresses

Host names

- Mnemonic name appreciated by humans

- Variable length, alpha-numeric characters

- Provide little (if any) information about location

- Examples: `www.cnn.com` and `ftp.eurocom.fr`

IP addresses

- Numerical address appreciated by routers

- Fixed length, binary number

- Hierarchical, related to host location

- Examples: `64.236.16.20` and `193.30.227.161`



Separating Naming and Addressing

Names are easier to remember

`www.cnn.com` vs. `64.236.16.20`

Addresses can change underneath

Move `www.cnn.com` to `173.15.201.39`

E.g., renumbering when changing providers

Name could map to multiple IP addresses

`www.cnn.com` to multiple replicas of the Web site



Separating Naming and Addressing

Map to different addresses in different places

Address of a nearby copy of the Web site

E.g., to reduce latency, or return different content

Multiple names for the same address

E.g., aliases like ee.mit.edu and cs.mit.edu



Solution #1: Local File

Original name to address mapping

- Flat namespace

- /etc/hosts

- Downloaded regularly

Count of hosts was increasing: moving from a machine per domain to machine per user

- Many more downloads

- Many more updates



Solution #2: Central Server

Central server

One place where all mappings are stored

All queries go to the central server

Need a distributed,
hierarchical collection of
servers

Many practical problems

Single point of failure

High traffic volume

Distant centralized database

Does not scale



Domain Name System (DNS)

Properties of DNS

- Hierarchical name space divided into zones

- Distributed over a collection of DNS servers

Hierarchy of DNS servers

- Root servers (13 clusters, labelled A-M)

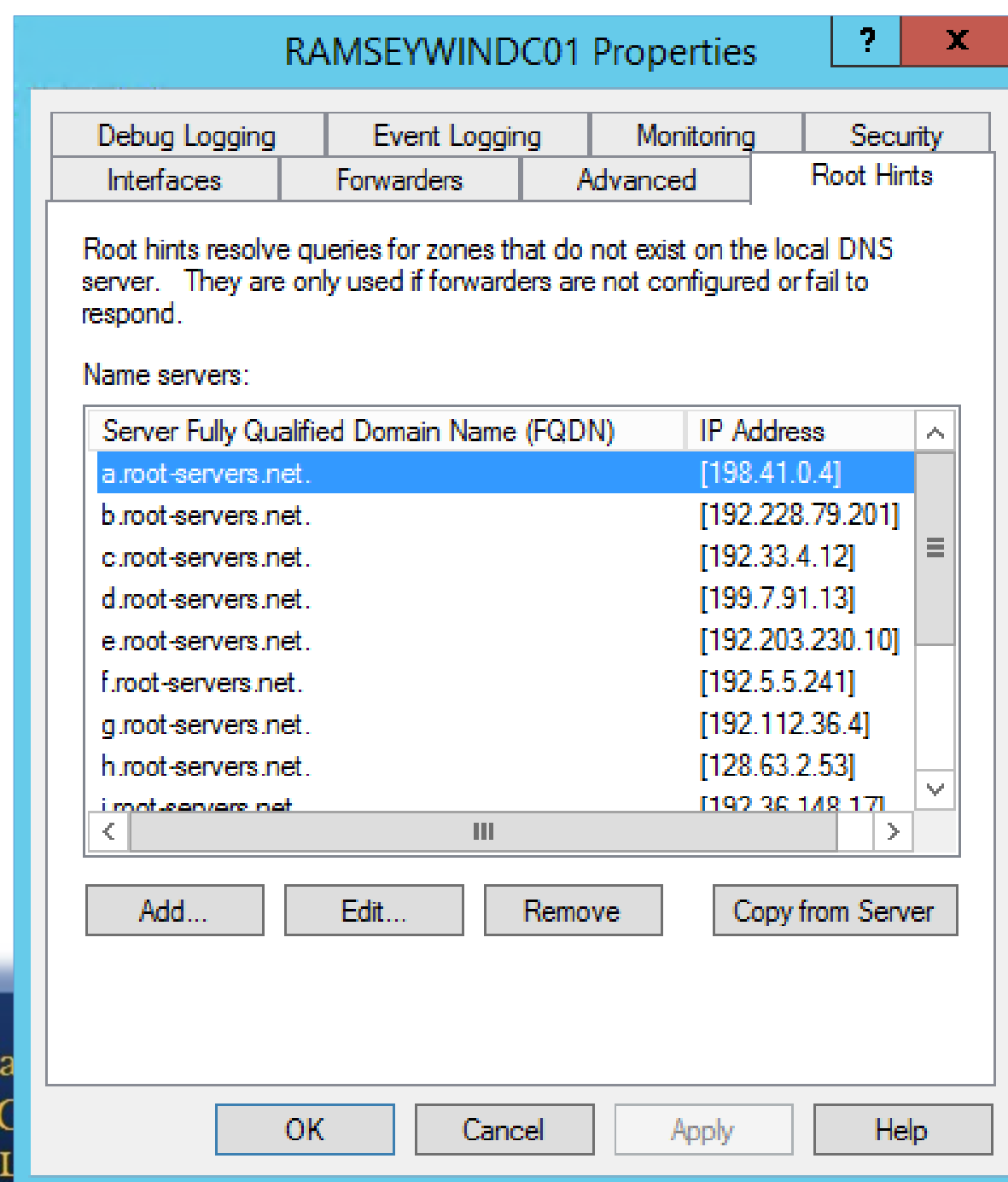
- Top-level domain (TLD) servers

- Authoritative DNS servers





DNS Root Servers



TLD and Authoritative DNS Servers

Top-level domain (TLD) servers

Generic domains (e.g., com, org, edu)

Country domains (e.g., uk, fr, ca, jp)

Typically managed professionally

Network Solutions maintains servers for “com”

Educause maintains servers for “edu”



TLD and Authoritative DNS Servers

Authoritative DNS servers (e.g., RAMSEYDC01)

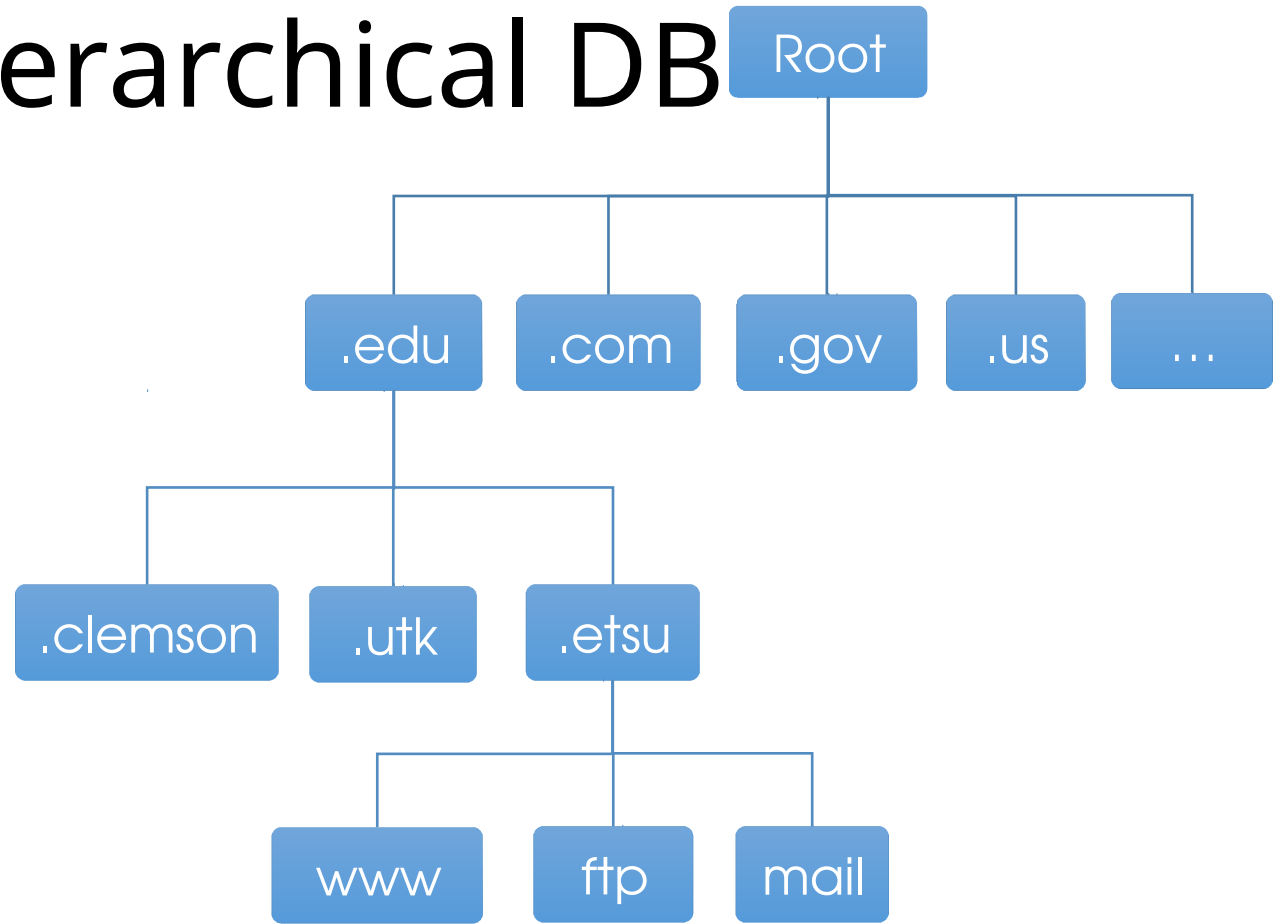
- Provide public records for hosts at an organization

- For the organization's servers (e.g., Web and mail)

- Can be maintained locally or by a service provider



Distributed Hierarchical DB



Using DNS

Local DNS server (“default name server”)

Usually (geographically) near the end hosts who use it

Local hosts configured with local server (e.g.,
/etc/resolv.conf) or learn the server via DHCP

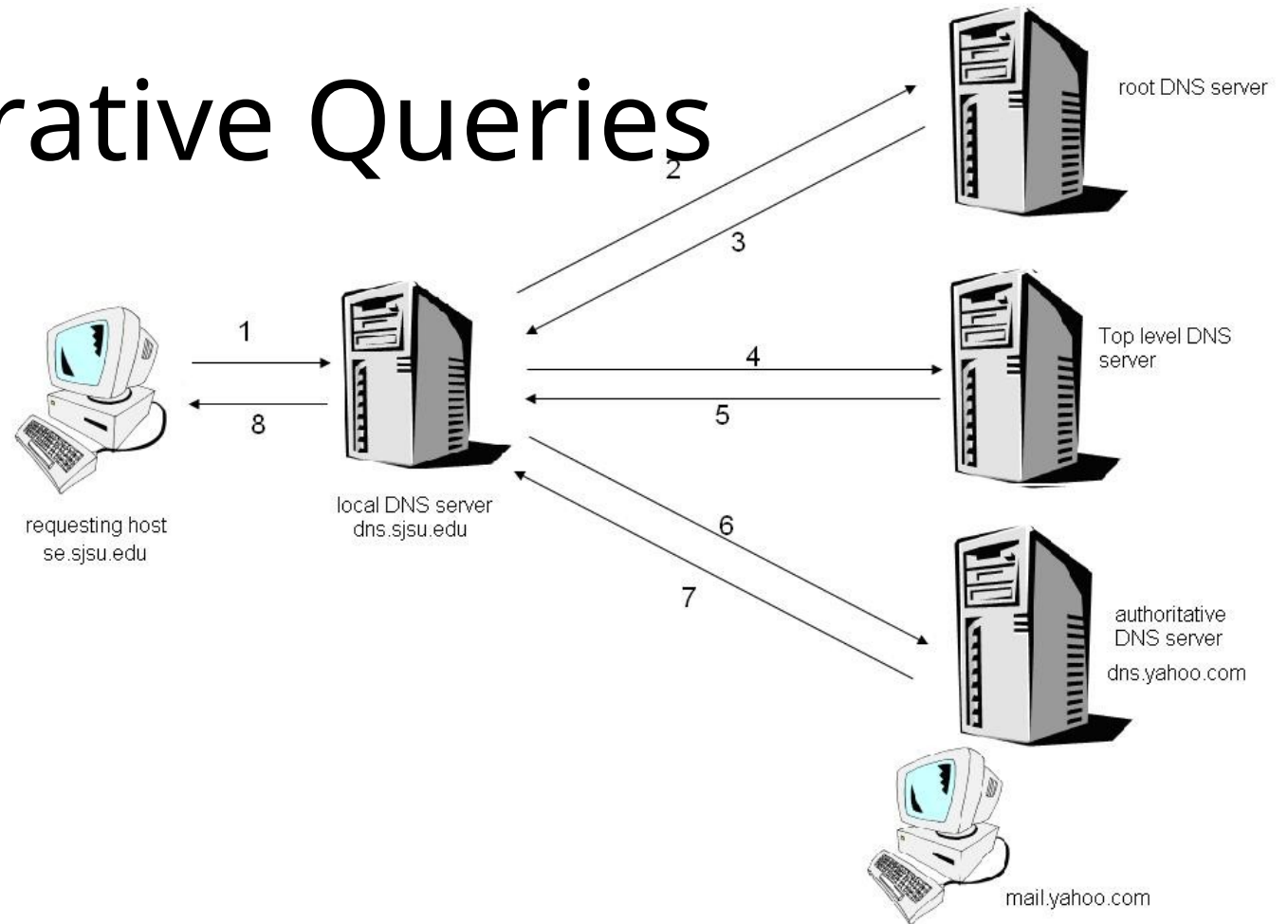


Recursive vs. Iterative Queries

Iterative query

Ask server who
to ask next

E.g., all other request-
response pairs

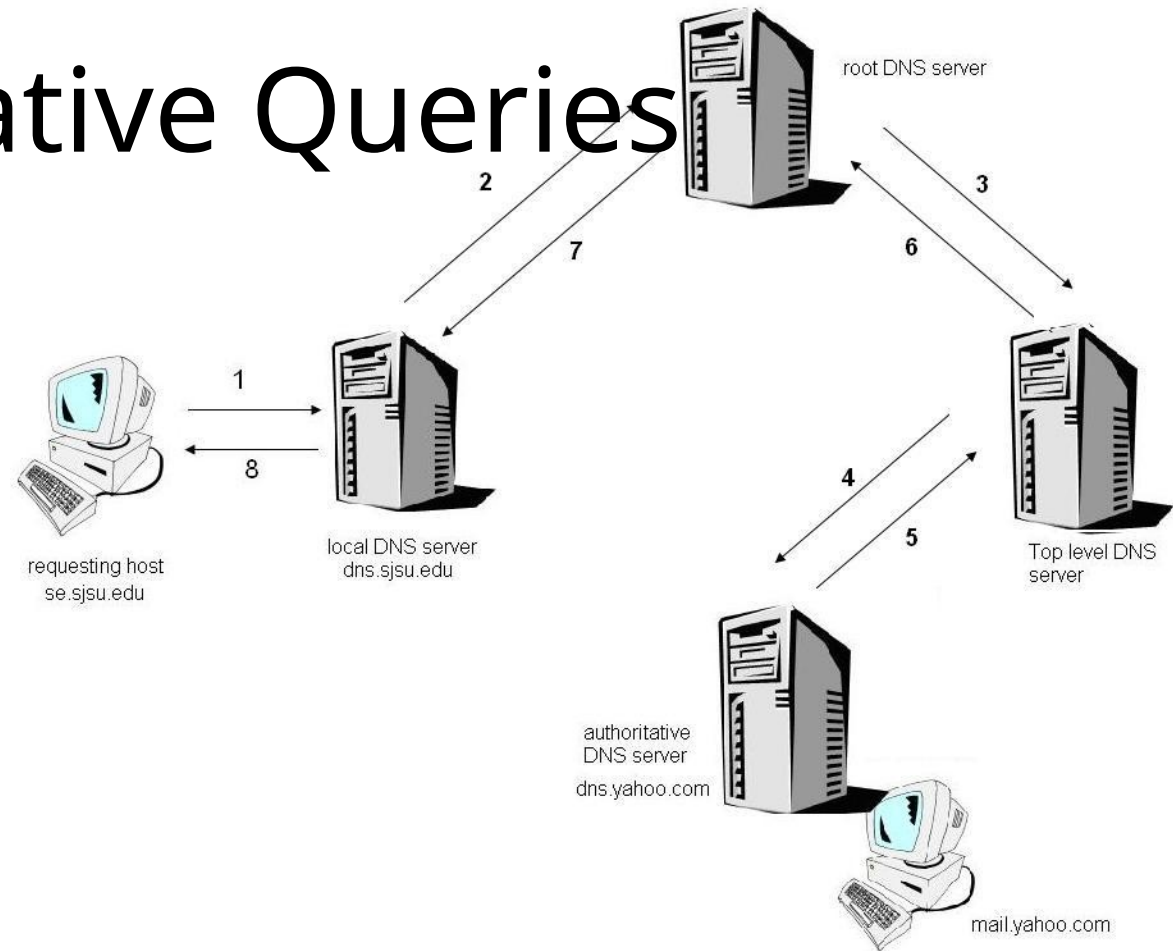


Recursive vs. Iterative Queries

Recursive query

Ask server to get answer for you

E.g., request 1 and response 8



DNS Caching

Performing all these queries take time

And all this before the actual communication takes place

E.g., 1-second latency before starting Web download

Caching can substantially reduce overhead

The top-level servers very rarely change

Popular sites (e.g., www.cnn.com) visited often

Local DNS server often has the information cached



DNS Caching

How DNS caching works

DNS servers cache responses to queries

Responses include a “time to live” (TTL) field

Server deletes the cached entry after TTL expires



Negative Caching

Remember things that don't work

Misspellings like www.cnn.comm and www.cnnn.com

These can take a long time to fail the first time

Good to remember that they don't work

... so the failure takes less time the next time around



DNS Resource Records

DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

Type=A

- name is hostname
- value is IP address

Type=NS

name is domain (e.g. foo.com)

value is hostname of authoritative name server for this domain

DNS Resource Records

DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

Type=CNAME

- name is alias name for some “canonical” (the real) name
`www.ibm.com` is really
`servereast.backup2.ibm.com`
- value is canonical name

Type=MX

- value is name of mailserver associated with name



Reliability

DNS servers are replicated

- Name service available if at least one replica is up

- Queries can be load balanced between replicas

UDP used for queries

- Need reliability: must implement this on top of UDP

Try alternate servers on timeout

Same identifier for all queries

- Don't care which server responds



Inserting Resource Records into DNS

Example: just created startup “FooBar”

Register foobar.com at Network Solutions

Provide registrar with names and IP addresses of your authoritative name server (primary and secondary)

Registrar inserts two RRs into the com TLD server:

(foobar.com, dns1.foobar.com, NS)

(dns1.foobar.com, 212.212.212.1, A)



Inserting Resource Records into DNS

Put in authoritative server dns1.foobar.com

Type A record for www.foobar.com

Type MX record for foobar.com



Playing With Dig on UNIX

Dig program

“domain information grouper”

Allows querying of DNS system

Use flags to find name server (NS)

Disable recursion so that operates one step at a time



```
; <<>> DiG 9.9.5-3-Ubuntu <<>> +recurse etsu.edu
;; global options: +cmd
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 34482
;; flags: qr rd ra; QUERY: 1, ANSWER: 4, AUTHORITY: 0, ADDITIONAL: 1

;; OPT PSEUDOSECTION:
; EDNS: version: 0, flags:; MBZ: 0005 , udp: 4000
;; QUESTION SECTION:
;etsu.edu.                IN      A

;; ANSWER SECTION:
etsu.edu.                 5       IN      A      151.141.8.110
etsu.edu.                 5       IN      A      151.141.8.101
etsu.edu.                 5       IN      A      151.141.5.62
etsu.edu.                 5       IN      A      151.141.8.100

;; Query time: 5 msec
;; SERVER: 127.0.1.1#53(127.0.1.1)
;; WHEN: Thu Mar 19 12:09:45 EDT 2015
;; MSG SIZE rcvd: 101
```

jack@azure:~\$

ipconfig

Windows Utility

Can display client network information



```
C:\Windows\System32\cmd.exe

C:\Windows\System32>ipconfig /all

Windows IP Configuration

    Host Name . . . . . : Hurin
    Primary Dns Suffix . . . . . :
    Node Type . . . . . : Hybrid
    IP Routing Enabled. . . . . : No
    WINS Proxy Enabled. . . . . : No
    DNS Suffix Search List. . . . . : wifi.etsu.edu

Wireless LAN adapter Local Area Connection* 3:

    Media State . . . . . : Media disconnected
    Connection-specific DNS Suffix . :
    Description . . . . . : Microsoft Wi-Fi Direct Virtual Adapter
    Physical Address. . . . . : 80-86-F2-7E-AB-13
    DHCP Enabled. . . . . : Yes
    Autoconfiguration Enabled . . . . : Yes

Ethernet adapter Bluetooth Network Connection:

    Media State . . . . . : Media disconnected
    Connection-specific DNS Suffix . :
    Description . . . . . : Bluetooth Device (Personal Area Network)
    Physical Address. . . . . : 80-86-F2-7E-AB-16
    DHCP Enabled. . . . . : Yes
    Autoconfiguration Enabled . . . . : Yes

Wireless LAN adapter Wi-Fi:

    Connection-specific DNS Suffix . : wifi.etsu.edu
    Description . . . . . : Intel(R) Dual Band Wireless-N 7260
    Physical Address. . . . . : 80-86-F2-7E-AB-12
    DHCP Enabled. . . . . : Yes
    Autoconfiguration Enabled . . . . : Yes
    Link-local IPv6 Address . . . . . : fe80::f5e6:25f3:638a:1922%4(Preferred)
    IPv4 Address. . . . . : 151.141.227.37(Preferred)
    Subnet Mask . . . . . : 255.255.248.0
    Lease Obtained. . . . . : Tuesday, October 14, 2014 8:01:55 AM
    Lease Expires . . . . . : Tuesday, October 14, 2014 7:01:44 PM
    Default Gateway . . . . . : 151.141.224.1
```



ipconfig

Windows Utility

Can display client network information

Can also display/flush cached DNS information



```
C:\Windows\System32>ipconfig /displaydns
```

Windows IP Configuration

i.creativecommons.org

```
-----  
Record Name . . . . . : i.creativecommons.org  
Record Type . . . . . : 1  
Time To Live . . . . . : 885  
Data Length . . . . . : 4  
Section . . . . . : Answer  
A (Host) Record . . . : 213.138.111.164
```

api.mixpanel.com

```
-----  
Record Name . . . . . : api.mixpanel.com  
Record Type . . . . . : 1  
Time To Live . . . . . : 75242  
Data Length . . . . . : 4  
Section . . . . . : Answer  
A (Host) Record . . . : 174.36.240.85
```

```
Record Name . . . . . : api.mixpanel.com  
Record Type . . . . . : 1  
Time To Live . . . . . : 75242  
Data Length . . . . . : 4  
Section . . . . . : Answer  
A (Host) Record . . . : 174.36.240.86
```

```
Record Name . . . . . : api.mixpanel.com  
Record Type . . . . . : 1  
Time To Live . . . . . : 75242  
Data Length . . . . . : 4  
Section . . . . . : Answer  
A (Host) Record . . . : 174.36.240.84
```

Boot-Strapping an End Host

DHCP and ARP



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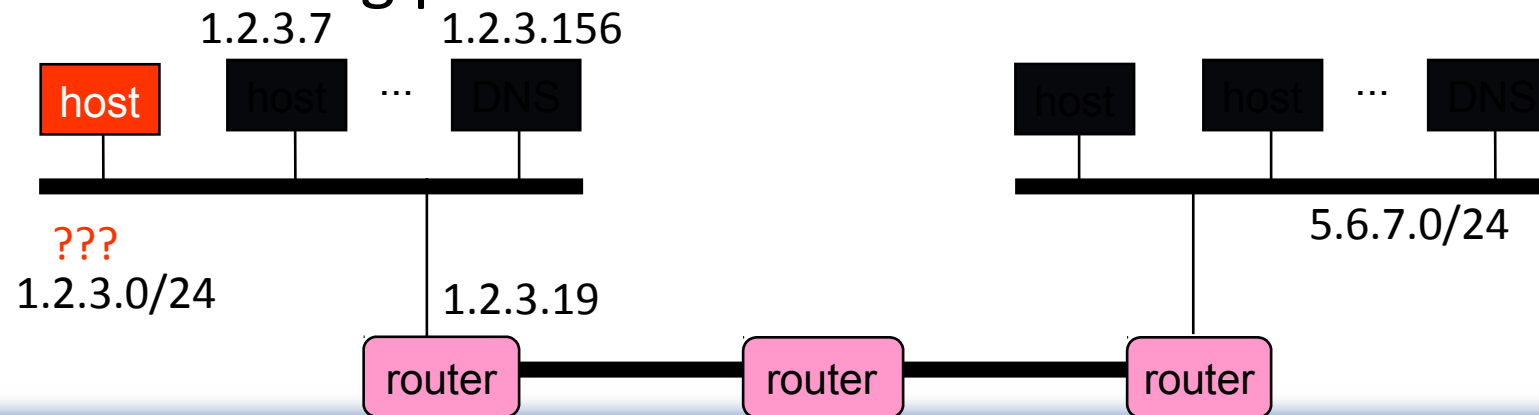
How To Bootstrap an End Host?

What local Domain Name System server to use?

What IP address the host should use?

How to send packets to remote destinations?

How to ensure incoming packets arrive?



Avoiding Manual Configuration

Dynamic Host Configuration Protocol (DHCP)

- End host learns how to send packets

- Learn IP address, DNS servers, and gateway

Address Resolution Protocol (ARP)

- Others learn how to send packets to the end host

- Learn mapping between IP address & interface address



Key Ideas in Both Protocols

Broadcasting: when in doubt, shout!

Broadcast query to all hosts in the local-area-network

... when you don't know how to identify the right one

Caching: remember the past for a while

Store the information you learn to reduce overhead

Remember your own address & other host's addresses



Key Ideas in Both Protocols

Soft state: ... but eventually forget the past

Associate a time-to-live field with the information

... and either refresh or discard the information

Key for robustness in the face of unpredictable change



Need Yet Another Kind of Identity

LANs are designed for arbitrary network protocols

Not just for IP and the Internet

Using IP address would require reconfiguration

Every time the adapter was moved or powered up

Broadcasting all data to all adapters is expensive

Requires every host on the LAN to inspect each packet

Motivates separate Media Access Control (MAC) addresses



MAC Address vs. IP Address

MAC addresses

Hard-coded in read-only memory when adaptor is built

Like a social security number

Flat name space of 48 bits (e.g., 00-0E-9B-6E-49-76)

Portable, and can stay the same as the host moves

Used to get packet between interfaces on same network



MAC Address vs. IP Address

IP addresses

Configured, or learned dynamically

Like a postal mailing address

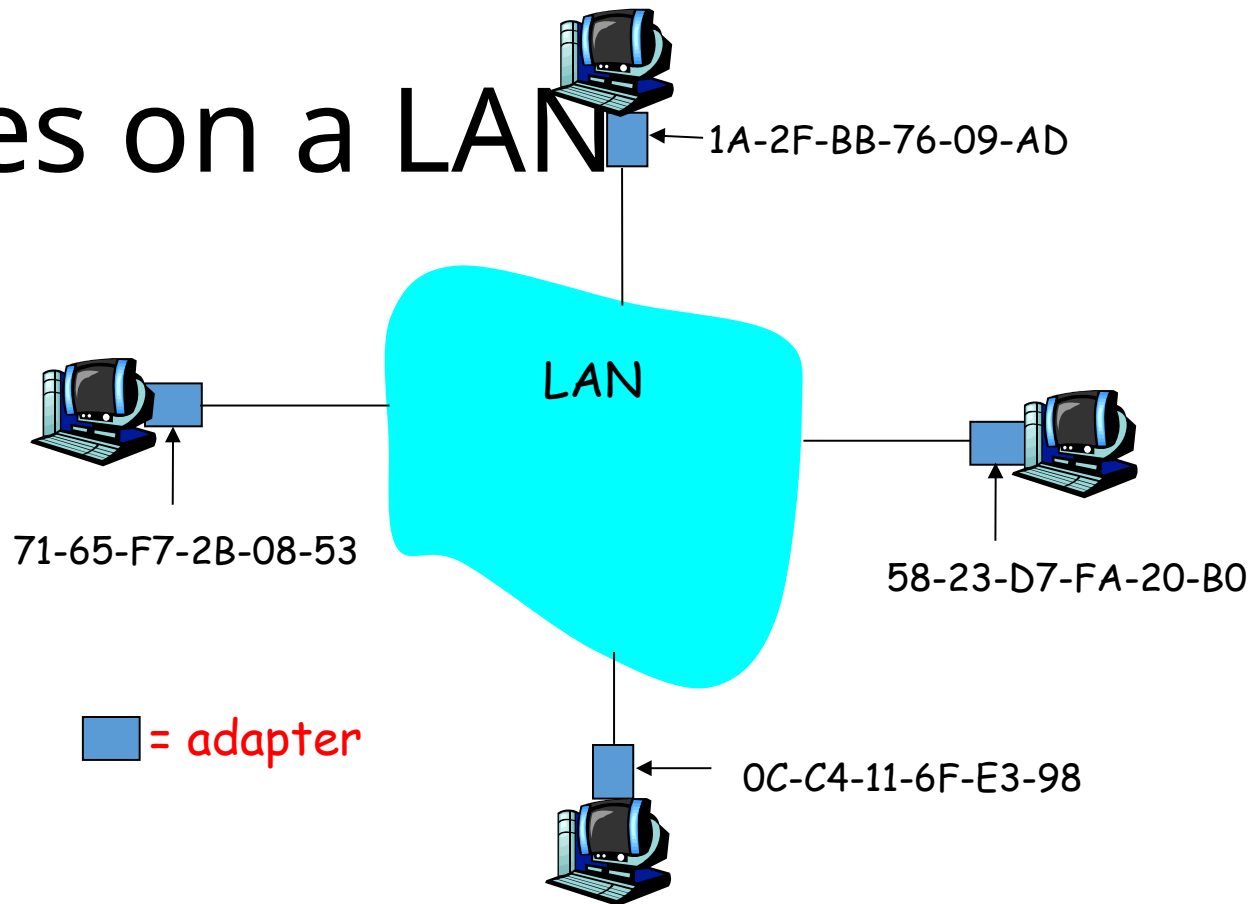
Hierarchical name space of 32 bits (e.g., 12.178.66.9)

Not portable, and depends on where the host is attached

Used to get a packet to destination IP subnet



MAC Addresses on a LAN



Bootstrapping Problem

Host doesn't have an IP address yet

So, host doesn't know what source address to use

Host doesn't know who to ask for an IP address

So, host doesn't know what destination address to use

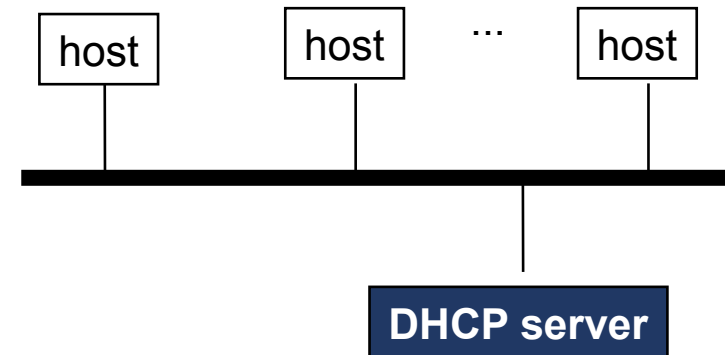


Bootstrapping Problem

Solution: shout to discover a server who can help

Broadcast a server-discovery message

Server sends a reply offering an address



Broadcasting

Broadcasting: sending to everyone

Special destination address: FF-FF-FF-FF-FF-FF

All adapters on the LAN receive the packet

Delivering a broadcast packet

Easy on a “shared media”

Like shouting in a room – everyone can hear you



Response from the DHCP Server

DHCP “offer message” from the server

Configuration parameters (proposed IP address, mask, gateway router, DNS server, ...)

Lease time (the time the information remains valid)

Multiple servers may respond

Multiple servers on the same broadcast media

Each may respond with an offer

The client can decide which offer to accept



Response from the DHCP Server

Accepting one of the offers

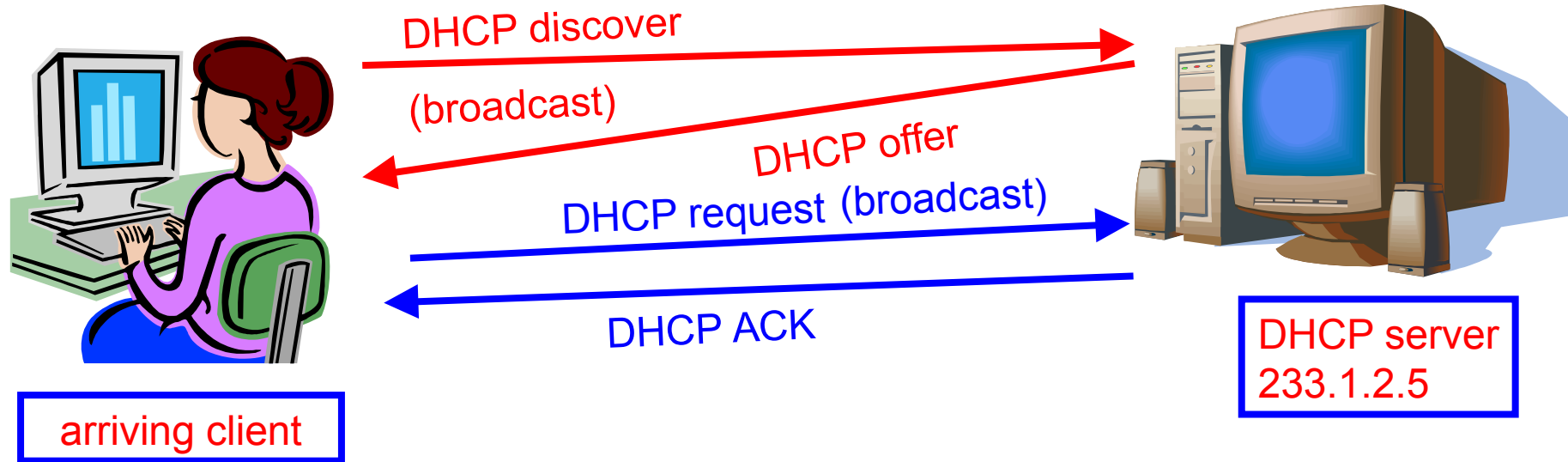
Client sends a DHCP request echoing the parameters

The DHCP server responds with an ACK to confirm

... and the other servers see they were not chosen



Dynamic Host Configuration Protocol



Deciding What IP Address to Offer

Server as centralized configuration database

All parameters are statically configured in the server

E.g., a dedicated IP address for each MAC address

Avoids complexity of configuring hosts directly

... while still having a permanent IP address per host



Deciding What IP Address to Offer

Or, dynamic assignment of IP addresses

Server maintains a pool of available addresses

... and assigns them to hosts on demand

Leads to less configuration complexity

... and more efficient use of the pool of addresses

Though, it is harder to track the same host over time



Soft State: Refresh or Forget

Why is a lease time necessary?

Client can release the IP address (DHCP RELEASE)

E.g., “ipconfig /release” at the DOS prompt

E.g., clean shutdown of the computer

But, the host might not release the address

E.g., the host crashes

E.g., buggy client software

And you don't want the address to be allocated forever



Soft State: Refresh or Forget

Performance trade-offs

Short lease time: returns inactive addresses quickly

Long lease time: avoids overhead of frequent renewals



So, Now the Host Knows Things

IP address

Mask

Gateway router

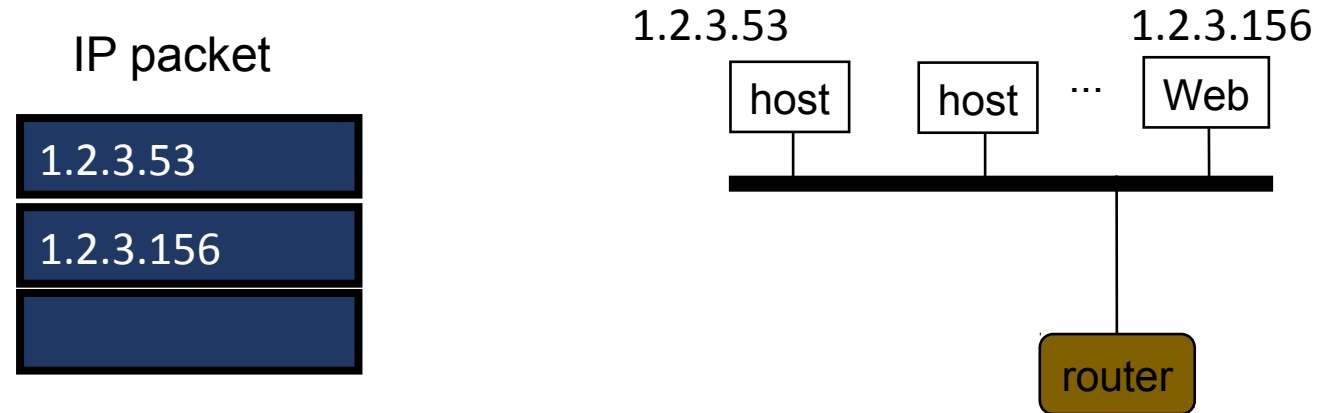
DNS server,

And can send packets to other IP addresses

But, how to learn the MAC address of the destination?



Sending Packets Over a Link



Adaptors only understand MAC addresses

Translate the destination IP address to MAC address

Encapsulate the IP packet inside a link-level frame

Address Resolution Protocol Table

Every node maintains an ARP table

(IP address, MAC address) pair

Consult the table when sending a packet

Map destination IP address to destination MAC address

Encapsulate and transmit the data packet



Address Resolution Protocol Table

But, what if the IP address is not in the table?

Sender broadcasts: “Who has IP address 1.2.3.156?”

Receiver responds: “MAC address 58-23-D7-FA-20-B0”

Sender caches the result in its ARP table

No need for network administrator to get involved



```
C:\Windows\System32\cmd.exe

C:\Windows\System32>arp -a

Interface: 151.141.227.37 --- 0x4
  Internet Address      Physical Address      Type
  151.141.224.1          58-bf-ea-be-c2-80     dynamic
  151.141.231.255        ff-ff-ff-ff-ff-ff     static
  224.0.0.22             01-00-5e-00-00-16     static
  224.0.0.251            01-00-5e-00-00-fb     static
  224.0.0.252            01-00-5e-00-00-fc     static
  239.255.255.250        01-00-5e-7f-ff-fa     static
  255.255.255.255        ff-ff-ff-ff-ff-ff     static

Interface: 192.168.92.1 --- 0x17
  Internet Address      Physical Address      Type
  192.168.92.255        ff-ff-ff-ff-ff-ff     static
  224.0.0.22             01-00-5e-00-00-16     static
  224.0.0.251            01-00-5e-00-00-fb     static
  224.0.0.252            01-00-5e-00-00-fc     static
  239.255.255.250        01-00-5e-7f-ff-fa     static

Interface: 192.168.244.1 --- 0x18
  Internet Address      Physical Address      Type
  192.168.244.255        ff-ff-ff-ff-ff-ff     static
  224.0.0.22             01-00-5e-00-00-16     static
  224.0.0.251            01-00-5e-00-00-fb     static
  224.0.0.252            01-00-5e-00-00-fc     static
  239.255.255.250        01-00-5e-7f-ff-fa     static
  255.255.255.255        ff-ff-ff-ff-ff-ff     static

C:\Windows\System32>
```

What happens if there's no DHCP?

APIPA

“Automatic Private IP Addressing”

Microsoft creation allows a client to self-configure its IP and subnet when no DHCP server available

Reserved address range: 169.254.0.1 through 169.254.255.254

Continues to check for DHCP server, obtains network IP address – replacing APIPA address – if DHCP is found

Conclusion

Domain Name System

- Distributed, hierarchical database

- Distributed collection of servers

- Caching to improve performance

Bootstrapping an end host

- Dynamic Host Configuration Protocol (DHCP)

- Address Resolution Protocol (ARP)



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