Sequential Decision-Making

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Informed Search

- We are informed about the goal
- We need an **estimation** about **how close** a state is to **goal** → heuristic function
- Manhattan distance and Euclidean distance in path finding

Best First Search

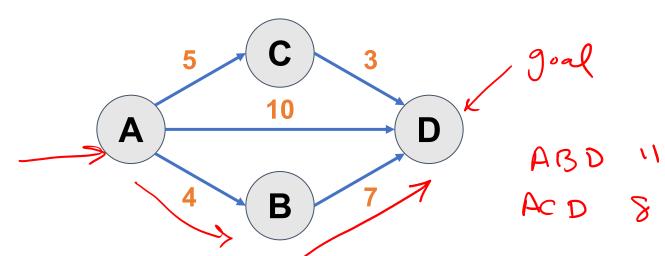
Strategy:

- Use heuristic for each node as an estimation to its "desirability"
- Expand the most desirable unexpanded node

Greedy Search

Strategy:

- Expand a node which has the best (smallest) heuristic
- Heuristic: estimate of distance to nearest goal for each state
- Common case: closest distance to the goal is not always the optimal solution
- Not optimal
- It is not complete (might stock in loops)



Combining UCS and Greedy

h=7

h=6

Uniform-cost: orders by path cost g(n) SabcGreedy: orders by goal proximity h(n) Saed G A* search: orders by the sum of f(n) = g(n) + h(n)SadG 8 h=1 h=5 h=0 h=2b

Admissible Heuristic

A heuristic h is admissible (optimistic) if:

$$0 \le h(n) \le h^*(n)$$

- Where $h^*(n)$ is the true cost to the nearest goal
- Is it possible using the actual cost as a heuristic?

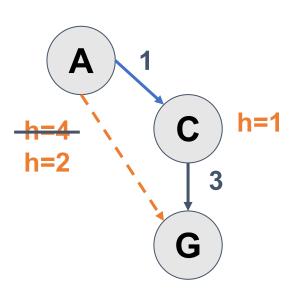
As heuristics get closer to the true cost, you will expand fewer nodes but usually do more work per node to compute the heuristic itself

Consistency

Heuristic "arc" cost is less than or equal to the actual cost for each arc

$$h(A) - h(c) \le cost(AC)$$

- The f value along a path never decreases
- A* graph search is optimal



Optimality of A* Graph Search

- Assumption: A* with a consistent heuristic
 - A* expands nodes in increasing total f value (f-contours)
 - For every state s, nodes that reach s optimally are expanded before s is expanded
- Uniform Cost Search (UCS) is optimal since h=0 is a consistent heuristic

$$0 \le h(n)$$
 $h(n) = 0$

Summary

- Informed search
- Best first search
- Greedy search
- A* search
- Admissibility
- Consistency
- A* optimality