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1 DataStruct

1.1 基础-单调队列

```
1 <u>int</u> main() {//luogu1440
    rep(i,1,n){
3
       if(que.empty()) ans[i]=0;
       else ans[i]=a[que.front()];
       if(!que.empty()&&i-que.front()>=m) que.pop_front();
       while(!que.empty()&&a[i]<a[que.back()]) que.pop_back();</pre>
       que.push_back(i);
     }//min(ai-m,...,ai-1)
9
10 <u>int</u> main(){//luogu2852
    rep(i,1,n){
11
      while(!que.empty()&&h[i]<h[que.back()]) que.pop_back();</pre>
12
13
      que.push_back(i);
       while(!que.empty()&&i-que.front()>=m) que.pop_front();
14
      if (que.empty()) ans[i]=0;
15
       else ans[i]=h[que.front()];
16
17
     }//min(ai-m+1,...,ai)
18|}
```

1.2 基础-单调栈

```
int main() {//hdu1506
     cin>>n;
3
     rep(i,1,n) <u>cin</u>>>a[i];
     11 ans=a[1];
     while(!stk.empty()) stk.pop();
     rep(i,1,n){
       while(!stk.empty()&&a[stk.top()]>=a[i]) stk.pop();
       <u>if</u>(!stk.empty()) l[i]=stk.top()+1;
8
9
       <u>else</u> l[i]=1;
10
       stk.push(i);
     }
11
12
     while(!stk.empty()) stk.pop();
13
     per(i,1,n){
       while(!stk.empty()&&a[stk.top()]>=a[i]) stk.pop();
14
15
       \underline{if}(!stk.empty()) r[i] = stk.top()-1;
16
       else r[i]=n;
17
       stk.push(i);
18
19
     rep(i,1,n) ans=max(ans,a[i]*(r[i]-l[i]+1));
20
     cout<<ans<<endl;</pre>
21| \}
```

L.3 基础-拉链法哈希表

```
1 const int maxsz=3e5+7;//maxsz素数表
 2|//1e7+19,2e7+3,3e7+23,4e5+9 maxsz最好为素数
3 //1e6+3,2e6+3,3e6+7,4e6+9,1e5+3,2e5+3,3e5+7
41//因为是vector不需要限制操作次数了
5 //count操作不增加新节点
6 template typename key, typename val>
  class hash_map{public:
    struct node{key u;val v;int next;};
    vector<node> e;
10
    int head[maxsz], nume, numk, id[maxsz];
|11|
    int geths(key &u){return u;}
12| //geths是把key映射到[0,maxsz-1]的函数
13
    bool count(key &u){
      int hs=geths(u);
14
15
      for(int i=head[hs];i;i=e[i].next)
16
        <u>if</u>(e[i].u==u) <u>return</u> 1;
      return 0;
17
    }
18
    val& <u>operator</u>[](key &u){//视情况加引用,可能引起CE
19
20
      int hs=geths(u);
21
      for(int i=head[hs];i;i=e[i].next)if(e[i].u==u) return e[i].v;
22
      if(!head[hs])id[++numk]=hs;
23
      if(++nume>=e.size())e.resize(nume<<1);</pre>
24
      return e[nume] = (node) {u,0,head[hs]},head[hs] = nume,e[nume].v;
25
    void clear(){
```

```
rep(i,0,numk)head[id[i]]=0;
28
      numk=nume=0;
29
    }
30 };
```

基础-可删堆

```
1 // 可删堆
  //保证remove元素被包含在x中,不能多删
3 template < typename T > class re_heap { public :
    priority_queue<T> x,y;
5
    void push(T a){x.push(a);}
    void remove(T a){y.push(a);}
7
      while (y.size()&&x.top()==y.top())
9
      x.pop(),y.pop();return x.top();
10
11
    int size(){return x.size()-y.size();}
12
    void pop(){
13
      \underline{\text{while}}(y.size()\&\&x.top()==y.top())
14
      x.pop(),y.pop();
15
      x.pop();
16
    T sectop(){
17
18
      T a=top();pop();
19
      T b=top();push(a);
20
      return b;
21
22 };
23 re_heap<pii> heap;
```

并查集-非递归+按秩合并

```
1 int pre[maxn], rk[maxn];
 2 int find(int a){
    if(pre[a] == -1) return a;
 4
    int t,rt=a;
5
    while(pre[rt]!=-1) rt=pre[rt];
6
    while(a!=rt)t=pre[a],pre[a]=rt,a=t;
7
    return rt;
8
9 #define same(a,b) (find(a)==find(b))
10 <u>void</u> unite(<u>int</u> a, <u>int</u> b){
     a=find(a),b=find(b);
11
12
     if(a==b) return;
13
     if(rk[a]>rk[b])pre[b]=a;
|14|
       pre[a]=b;
15
       <u>if</u>(rk[a]==rk[b])rk[b]++;
16
17
18|}
19
  class ufs{public:
20
     void init(int n){rep(i,0,n) pre[i]=i;}
     int find(int a){return pre[a] == a?a:pre[a] = find(pre[a]);}
21
    bool same(int a, int b) {return find(a)==find(b);}
23
    void unite(int a,int b){
24
       a=find(a),b=find(b);
25
       if(a==b) return;
26
       pre[b]=a;
    }
27
28 }dsu;
```

并查集-可删除

```
1| #define same(a,b) (find(a)==find(b))
   int pre[maxn],id[maxn];
   int numid;
 4
   \underline{int} fd(\underline{int} a) {
       return pre[a] == a?a:pre[a] = find(pre[a]);
5
6
7
   int find(int a){return fd(id[a])};
8 void unite(int a, int b){
       a=find(id[a]),b=find(id[b]);
10
       pre[b]=a;
11|}
```

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```
12 void remove(int now){
13    int pos=find(id[now]);
14    id[now]=++numid;
15    pre[numid]=numid;
16 }
```

1.7 并查集-可撤销

```
class ufs{public:
    int fa[maxn],sz[maxn];
    stack<pii> stk;
    void init(int tot){
      rep(i,1,tot) fa[i]=i,sz[i]=1;
5
6
7
    int find(int now){
8
      if(fa[now] == now) return now;
9
       return find(fa[now]);
10
11
    bool unite(int a,int b){
12
      a=find(a),b=find(b);
13
       if(a==b)return 0;
14
       <u>if</u>(sz[a]<sz[b]) fa[a]=b,sz[b]+=sz[a],stk.push(make_pair(a,b));
       else fa[b]=a,sz[a]+=sz[b],stk.push(make_pair(b,a));
15
16
       <u>return</u> 1;
17
    void undo(){
18
      if(!stk.empty()){
19
        auto now=stk.top();
20
21
         stk.pop();
|22|
        fa[now.fi]=now.fi;
23
         sz[now.se] -=sz[now.fi];
24
25
26| }dsu;
```

1.8 区间信息-数列分块

```
const int maxn=1e5+10, maxm=sqrt(maxn+1)+10;
  class sqblock{public:
    int id[maxn],sz,cnt;
    struct block{int l,r;ll sum,tag;}node[maxm];
5
    void init(int n){
6
      sz=sqrt(n+0.5);
7
      rep(i,1,n) id[i]=(i-1)/sz+1;cnt=id[n];
8
      rep(i,1,cnt) node[i].l=(i-1)*sz+1;
9
      rep(i,1,cnt-1) node[i].r=i*sz;node[cnt].r=n;
10
      rep(i,1,n) node[id[i]].sum+=arr[i];
11
12
    void update(int s,int t,ll x){
      int p1=id[s],p2=id[t];
13
      <u>if</u>(p1==p2)rep(i,s,t)node[p1].sum+=x,arr[i]+=x;
|14|
15
16
        rep(i,s,node[p1].r) arr[i]+=x,node[p1].sum+=x;
17
        rep(i,node[p2].1,t) arr[i]+=x,node[p2].sum+=x;
18
        rep(i,p1+1,p2-1) node[i].tag+=x;
19
20
21
    11 query(<u>int</u> s,<u>int</u> t){
22
       <u>int</u> p1=id[s],p2=id[t];ll ans=0;
23
       \underline{if}(p1==p2) rep(i,s,t) ans+=arr[i]+node[p1].tag;
24
       else {
25
        rep(i,s,node[p1].r) ans+=arr[i]+node[p1].tag;
26
        rep(i,node[p2].1,t) ans+=arr[i]+node[p2].tag;
27
        rep(i,p1+1,p2-1) ans+=node[i].sum+node[i].tag*sz;
28
29
      return ans;
    }
30
31 }square;
```

1.9 区间信息-ST表

```
1 const int maxp=21;
2 class stable{public:
3 int logn[maxn],dp[maxp][maxn];
4 int *a;
```

```
void init(int n=maxn-1){
6
       logn[2]=1;
       rep(i,3,n) logn[i]=logn[i>>1]+1;
8
9
     void cal(int *_a,int n){//init(n)
10
       rep(i,1,n) dp[0][i]=a[i],pos[0][i]=i;
11
12
       for(int j=1;(1<<j)<=n;j++) for(int i=1;i+(1<<j)-1<=n;++i)</pre>
13
         dp[j][i]=min(dp[j-1][i],dp[j-1][i+(1<<(j-1))]);
       for(int j=1;(1<<j)<=n;j++) for(int i=1;i+(1<<j)-1<=n;++i)</pre>
14
         <u>if</u>(a[pos[j-1][i]]<=a[pos[j-1][i+(1<<(j-1))]]) pos[j][i]=pos[j-1][i];
15
16
         <u>else</u> pos[j][i]=pos[j-1][i+(1<<(j-1))];
17
18
     inline int query(int 1,int r){
19
       int lg=logn[r-l+1];
20
       <u>return</u> min(dp[lg][l],dp[lg][r-(1<<lg)+1]);
21
22
     <u>inline</u> <u>int</u> query_pos(<u>int</u> 1, <u>int</u> r){
23
       int lg=logn[r-l+1];
24
       <u>if</u>(a[pos[lg][1]]<=a[pos[lg][r-(1<<lg)+1]])<u>return</u> pos[lg][1];
25
       else return pos[lg][r-(1<<lg)+1];</pre>
26
27 }st;
```

区间信息-树状数组 1.10

```
1 //树状数组区间求和,单点修改
  template<typename T> class bit{public:
    T val[maxn];
     \underline{int} lb(\underline{int} x) {\underline{return} x&(-x);}
     int len;
5
    void init(int _n){
       len=_n;
8
       fill_n(val+1,len+1,INF);
9
10
    void update(int pos,T x){
       \underline{for}(;pos>0\&\&pos<=n;pos+=lb(pos)) val[pos]+=x;
11
12
13
     T psum(<u>int</u> pos){
14
       T sum=0;
       for(;pos>0;pos-=lb(pos)) sum+=val[pos];
15
16
       return sum;
17
     T query(int 1,int r){return psum(r)-psum(l-1);}
18
19
  }tree;
20| //树状数组前缀最小值,支持单点的取min覆盖操作
21
  template<typename T> class bit{public:
22
     T val[2*maxn];
23
     int lb(int x) {return x&-x;}
     int len;
24
     void init(int _n){
25
       len=_n;
26
27
      fill_n(val+1,len+1,INF);
28
29
     void update(int pos,T x){
30
       for(;pos<=len;pos+=lb(pos)) val[pos]=min(val[pos],x);</pre>
31
32
    T query(<u>int</u> pos){
33
       T res=INF;
34
       for(;pos;pos-=lb(pos)) res=min(val[pos],res);
35
       return res;
36
37
38 bit<<u>int</u>> tree;
```

区间信息-扫描线算法 1.11

```
1 double dis[maxn];
 struct node{
   double x,y1,y2;int tag;
   bool operator <(node &other) const{return x<other.x;}</pre>
5| \seg[maxn];
6 class segtree {public:
7 #define nd node[now]
```

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```
#define ndl node[now<<1]</pre>
   #define ndr node[now<<1|1]</pre>
9
10
     struct segnode {
       int l,r,tag;double dis;
11
       inline int mid(){return (r+1)>>1;}
12
13
       inline int len(){return r-l+1;}
     node[maxn << 2|3];
|14|
15
     inline void update(int now){
16
       if(nd.tag) nd.dis=dis[nd.r+1]-dis[nd.l];
17
       else if(nd.len()==1) nd.dis=0;
18
       else nd.dis=ndl.dis+ndr.dis;
19
20
     void maketree(int s,int t,int now=1){
21
       nd={s,t,0,0};
22
       if(s==t) return ;
23
       maketree(s,nd.mid(),now<<1);</pre>
       maketree(nd.mid()+1,t,now<<1|1);
24
25
26
     void update(int s,int t,int x,int now=1){
27
       \underline{if}(s \le nd.1\&\&t \ge nd.r) {
28
       nd.tag+=x;update(now);
29
       return;
30
31
       if(s<=ndl.r) update(s,t,x,now<<1);</pre>
32
       if(t>ndl.r) update(s,t,x,now<<1|1);</pre>
33
       update(now);
34
35
   }tree;
   int main() {I0;cout<<fixed<<setprecision(2);</pre>
37
     \underline{\text{while}}((\underline{\text{cin}} >> n) \&\& n) \{
38
     m=0;
39
     rep(i,1,n){
40
       <u>double</u> a,b,c,d;<u>cin</u>>>a>>b>>c>>d;
41
       dis[++m]=b; seg[m]={a,b,d,1};
42
       dis[++m]=d; seg[m]=\{c,b,d,-1\};
43
     sort(dis+1,dis+m+1);
44
45
     sort(seg+1,seg+m+1);
46
     int cnt=unique(dis+1,dis+1+m)-dis-1;
47
     tree.maketree(1,cnt);
     double ans=0;
|48|
49
     rep(i,1,m-1){
50
       int l=lower_bound(dis+1,dis+1+cnt,seg[i].y1)-dis;
51
       int r=lower_bound(dis+1,dis+1+cnt,seg[i].y2)-dis;
52
53
       if(l<=r) tree.update(l,r,seg[i].tag);</pre>
54
       ans+=tree.node[1].dis*(seg[i+1].x-seg[i].x);
55
     cout<<"Test case #"<<++casn<<endl;</pre>
56
57
     cout<<"Total explored area: "<<ans<<endl<<endl;</pre>
58
59 }
```

1.12 区间信息-动态开点线段树

```
1 int root;
  const 11 rangel=0,ranger=1e9+10;
  class dsegtree{public:
4 #define nd node[now]
5 #define ndl node[node[now].son[0]]
 6 #define ndr node[node[now].son[1]]
    struct dsegnode{
       int son[2],mx,tag;
      void update(int x){mx+=x,tag+=x;}
10
    node[maxn*50];
11
12
    void pushup(int now){nd.mx=max(ndl.mx,ndr.mx);}
13
    void pushdown(int now){
14
      \underline{if}(nd.tag){
15
         <u>if</u>(!nd.son[0]) nd.son[0]=++cnt;
16
         <u>if</u>(!nd.son[1]) nd.son[1]=++cnt;
17
        ndl.update(nd.tag),ndr.update(nd.tag);
18
        nd.tag=0;
19
```

```
20
21
     void update(ll s,ll t,int x,ll l=rangel,ll r=ranger,int &now=root){
22
       if(!now) {now=++cnt;nd={0,0,0,0};}
23
       \underline{if}(s \le 1\&\&t \ge r){
24
         nd.update(x);
25
         return ;
26
27
       pushdown(now);
28
       if(s <= ((1+r) >> 1)) update(s,t,x,1,(1+r) >> 1,nd.son[0]);
29
       <u>if</u>(t>((l+r)>>1)) update(s,t,x,1+((l+r)>>1),r,nd.son[1]);
30
       pushup(now);
31
32| }tree;
```

区间信息-线段树优化建图 1.13

```
1 int cnt;
  class segtree{public:
  #define nd node[now]
  #define ndl node[now<<1]</pre>
  #define ndr node[now<<1|1]</pre>
     int flag;//1==intree,0==outtree
     struct segnode {
8
       int l,r,id;
9
       inline int mid(){return (r+1)>>1;}
10
       inline int len(){return r-l+1;}
     };
|11|
12
     segnode node[maxn<<2|3];
13
     vector<int> v;
|14|
     void init(int n, int flag){
15
       this->flag=flag;
16
       maketree(1,n);
17
18
     void pushup(int now){
       if(!flag){
19
         g.add(nd.id,ndl.id,0);
20
21
         g.add(nd.id,ndr.id,0);
22
       }else {
23
         g.add(ndl.id,nd.id,0);
24
         g.add(ndr.id,nd.id,0);
25
     }
26
27
     void maketree(int s,int t,int now=1){
28
       nd={s,t,++cnt};
       \underline{if}(s==t){}
29
30
         if(flag) g.add(s,nd.id,0);
31
         else g.add(nd.id,s,0);
32
         <u>return</u>;
33
34
       maketree(s,nd.mid(),now<<1);</pre>
35
       maketree(nd.mid()+1,t,now<<1|1);</pre>
36
       pushup(now);
37
38
     vector<int > query(int s,int t){v.clear();find(s,t);return v;}
39
     void find(int s,int t,int now=1){
40
       \underline{if}(s \le nd.l\&t \ge nd.r) {
41
         v.emplace_back(nd.id);
42
         return ;
43
44
       if(s<=ndl.r) find(s,t,now<<1);</pre>
45
       if(t>ndl.r) find(s,t,now<<1|1);</pre>
46
47 }intree, outree;
```

区间信息-线段树合并

```
1 int a[maxn], cnt, root[maxn];
2 vector<int> g[maxn];
3 \mid 11 \text{ ans}[maxn];
4 <u>class</u> dsegtree{<u>public</u>:
5 #define nd node[now]
6 #define ndl node[node[now].son[0]]
7 #define ndr node[node[now].son[1]]
    struct dsegnode {
```

```
int son[2],cnt,id,ans;
10
     node[maxn*50];
     void pushup(int now){
11
12
       if(ndl.cnt>ndr.cnt){
13
         nd.cnt=ndl.cnt;
14
         nd.id=ndl.id;
15
         nd.ans=ndl.ans;
16
       }else if(ndr.cnt>ndl.cnt){
17
         nd.cnt=ndr.cnt;
18
         nd.id=ndr.id;
19
         nd.ans=ndr.ans;
20
       }<u>else</u> {
21
         nd.cnt=ndr.cnt;
|22|
         nd.id=ndr.id;
23
         nd.ans=ndr.ans+ndl.ans;
24
25
     }
26
     void update(int 1,int r,int pos,int &now){
27
       if(!now) now=++cnt;
28
       \underline{if}(l==r){}
29
         nd.id=nd.ans=1;
30
         nd.cnt+=1;
31
         <u>return</u>;
32
33
       <u>int</u> mid=(l+r)>>1;
34
       if(pos<=mid) update(1,mid,pos,nd.son[0]);</pre>
35
       else update(mid+1,r,pos,nd.son[1]);
36
       pushup(now);
37
38
     int merge(int now,int b,int l,int r){
39
       if(!now||!b) return now^b;
       \underline{if}(l==r){}
40
         nd.id=nd.ans=1;
41
42
         nd.cnt+=node[b].cnt;
43
         return now;
44
45
       nd.son[0] = merge(nd.son[0], node[b].son[0], 1, (1+r)/2);
       nd.son[1] = merge(nd.son[1], node[b].son[1], (1+r)/2+1,r);
46
47
       pushup(now);
48
       return now;
     }
49
50 }tree;
51
   void dfs(int now,int fa){
52
     for(int to:g[now]){
53
       if(to==fa) continue;
54
       dfs(to,now);
55
       tree.merge(root[now],root[to],1,1e5);
56
57
     tree.update(1,1e5,a[now],root[now]);
58
     ans[now] = tree.node[root[now]].ans;
59|}
60 <u>int</u> main() {
61
     cin >> n;
62
     rep(i,1,n){
        <u>cin</u>>>a[i];
63
64
        root[i]=i;
     }
65
66
     cnt=n;
     rep(i,2,n){
67
68
       <u>int</u> a,b;<u>cin</u>>>a>>b;
69
       g[a].push_back(b);
70
       g[b].push_back(a);
71
72
     dfs(1,0);
73
     rep(i,1,n) <u>cout</u><<ans[i]<<' ';
```

1.15 区间信息-二维线段树

```
1 class sstree{public:
2 #define nd node[nowx][nowy]
3     struct segnode {int val;};
4     int n,m,x1,x2,y1,y2,x,nowx;int ans;
5     segnode node[maxn][maxn];
```

```
void init(int nn,int mm) {
|7|
           n=nn, m=mm;
8
           memset(node ,0,sizeof node);
9
10
       void update(int xx1,int xx2,int yy1,int yy2){
11
           x1=xx1, y1=yy1, x2=xx2, y2=yy2;
12
           updatex(1,n);
13
       void updatey(int 1,int r,int nowy=1){
14
15
           if(y1>r||y2<1) return;
           <u>if</u>(y1<=1&&y2>=r){nd.val^=1;<u>return</u>;}
16
           updatey(1,(1+r)>>1,nowy<<1); updatey(((1+r)>>1)+1,r,nowy<<1|1);
17
18
19
       void updatex(int 1,int r,int now=1){
20
           \underline{if}(x1>r||x2<1) \underline{return};
21
           if(x1<=1&&x2>=r) {nowx=now;updatey(1,m);return ;}
22
           updatex(1,(1+r)>>1,now<<1); updatex(((1+r)>>1)+1,r,now<<1|1);
23
       int query(int xx,int yy){
24
25
           x1=xx,y1=yy;ans=0;
26
           queryx(1,n);
27
           <u>return</u> ans;
28
29
       void queryy(int l,int r,int nowy=1){
           <u>if</u>(y1>r||y1<1) <u>return</u>;
30
31
           ans^=nd.val;
32
           if(l==r) return ;
           queryy(1,(1+r)>>1,nowy<<1);queryy(((1+r)>>1)+1,r,nowy<<1|1);
33
34
35
       void queryx(int 1,int r,int now=1){
36
           \underline{if}(x1>r||x1<1) \underline{return};
37
       nowx=now;queryy(1,m);
38
           <u>if</u>(l==r) <u>return</u>;
39
           queryx(1,(1+r)>>1,now<<1); queryx(((1+r)>>1)+1,r,now<<1|1);
40
41| }tree;
```

1.16 可持久化数据结构-01字典树

```
int rt[maxn];
  class ptrie{public:
    int node[maxn*40][2],top;
    void init(){rt[0]=node[0][1]=node[0][0]=top=0;}
5
    int add(int pre,int val,int bit=31){
      int now=++top;
 6
      if(bit<0) return now;</pre>
      int t=val>>bit&1;
9
      node[now][t]=add(node[pre][t],val,bit-1);
10
      node[now][t^1]=node[pre][t^1];
      return now;
11
12
    }
13
    int query(int now,int pre,int val,int bit=31,int ans=0){
14
      if(bit<0) return ans;</pre>
15
      int t=val>>bit&1;
16
      if(node[now][t^1]-node[pre][t^1]>0) return query(node[now][t^1],node[pre][t^1],val,bit-1,ans|(1<<bit</pre>
      return query(node[now][t],node[pre][t],val,bit-1,ans);
17
19
20 }tree;
```

1.17 可持久化数据结构-主席树

```
1 int rt[maxn];//树根
2 class ptree{public:
3 #define nd node[now]
4 #define mid (s+t)/2
6 int cnt;
7 struct segnode{int l,r,sum;}node[maxn*30];
8 void maketree(int s,int t,int &now=rt[0]){
9 now=++cnt;nd={s,t,0};
10 if(s==t) return;
11 maketree(s,mid,nd.l); maketree(mid+1,t,nd.r);
```

```
12
13
     void update(int pos,int val,int s,int t,int &now,int pre){
14
      now=++cnt;nd=ndp;nd.sum+=val;
15
       <u>if</u>(s==t) <u>return</u> ;
16
       if(pos<=mid) update(pos,val,s,mid,nd.1,ndp.1);</pre>
17
       else update(pos,val,mid+1,t,nd.r,ndp.r);
18
19
     11 query(int 1,int r,int s,int t,int now,int pre){
20
       if(l<=s&&r>=t)return nd.sum-ndp.sum;
21
      11 sum=0;
22
       if(l<=mid) sum+=query(l,r,s,mid,nd.l,ndp.l);</pre>
23
       if(r>mid) sum+=query(l,r,mid+1,t,nd.r,ndp.r);
24
       return sum;
25
    }
26
  #undef mid
27| }tree;
```

1.18 可持久化数据结构-主席树区间第k大

```
#include<bits/stdc++.h>
   #define rep(ii,a,b) for(int ii=a;ii<=b;++ii)</pre>
 3 #define all(x) x.begin(),x.end()
 4 using namespace std;//head
 5 \mid \underline{\text{const}} \mid \underline{\text{int}} \mid \text{maxn} = 2e5 + 10, \\ \underline{\text{maxm}} = 2e6 + 10;
 6 int casn,n,m,k;
 7 <u>int</u> rt[maxn];//树根
 8 <u>class</u> ptree{<u>public</u>:
 9 #define nd node[now]
10 #define ndp node[pre]
11 \frac{\text{#define}}{\text{mid}} \text{ mid } (s+t)/2
12
     <u>int</u> cnt;
     struct segnode{int 1,r,sum;}node[maxn*20];
13
14
     void init(){cnt=0;}
15
     void update(int pre,int &now,int pos,int s=1,int t=k){
16
        now=++cnt;nd=ndp;nd.sum++;
17
        if(s==t) return ;
18
        if(pos<=mid) update(ndp.1,nd.1,pos,s,mid);</pre>
19
        else update(ndp.r,nd.r,pos,mid+1,t);
20
21
     int query(int pre,int now,int pos,int s=1,int t=k){
|22|
        if(s==t) return s;
23
        int sum=node[nd.1].sum-node[ndp.1].sum;
|24|
        if(pos<=sum) return query(ndp.1,nd.1,pos,s,mid);</pre>
25
        else return query(ndp.r,nd.r,pos-sum,mid+1,t);
     }
26
27 #undef mid
28 }tree;
29 int a [maxn];
30 vector<int>pos;
31 <u>int</u> main() {IO;
32
     <u>cin</u>>>n>>m;
33
     rep(i,1,n) <u>cin</u>>>a[i];
34
     rep(i,1,n) pos.push_back(a[i]);
35
     sort(all(pos));
36
     pos.erase(unique(all(pos)),pos.end());
37
     k=pos.size();
38
     rep(i,1,n){
39
        int id=lower_bound(all(pos),a[i])-pos.begin();
40
        tree.update(rt[i-1],rt[i],id+1);
41
42
     \underline{\text{while}}(m--) {
43
        <u>int</u> a,b,c;<u>cin</u>>>a>>b>>c;
        \underline{\text{cout}} << \text{pos}[\text{tree.query}(\text{rt}[a-1],\text{rt}[b],c)-1] << '\n';
44
45
     return 0;
46
47| \}
```

1.19 可持久化数据结构-bit套主席树

```
1    namespace dsegtree{
2    #define nd node[now]
3    #define ndl node[node[now].son[0]]
4    #define ndr node[node[now].son[1]]
5    struct dsegnode{int son[2],sum;}node[maxn*200];
```

```
int cnt;
 |7|
     int pos,s,t,x;
     void update(int 1,int r,int &now){
 8
 9
       if(!now) now=++cnt;
10
       nd.sum+=x;
11
       if(l==r) return ;
12
        <u>if</u>(pos<=((l+r)>>1)) update(l,(l+r)>>1,nd.son[0]);
13
       else update(1+((1+r)>>1),r,nd.son[1]);
14
15
     void update_1(int _pos,int _x,int &root){
16
       pos=_pos,x=_x;
17
        update(1,n,root);
18
19
     int query(int 1,int r,int now){
20
       if(!now||nd.sum==0)return 0;
21
        if(s<=l&&t>=r) return nd.sum;
22
        int mid=(l+r)>>1;
23
24
        if (nd.son[0] &&s<=mid) sum+=query(1,mid,nd.son[0]);
25
        \underline{if}(nd.son[1]\&\&t>mid) sum+=query(mid+1,r,nd.son[1]);
26
       return sum;
27
28
     int query_1(int _s,int _t,int root){
29
        s=_s,t=_t;
30
       return query(1,n,root);
31
32| \}
33
   namespace bit{
34
     int node[maxn];
35
     inline void update(int pos,int y,int val){
36
        \underline{\text{for}}(;\text{pos}\leq n;\text{pos}+=\text{pos}\&(-\text{pos}))
37
          dsegtree::update_1(y,val,node[pos]);
38
39
     <u>inline</u> <u>int</u> ask(<u>int</u> pos, <u>int</u> x, <u>int</u> y){
40
        int sum=0;
41
        for(;pos;pos==pos&(-pos))
42
          sum+=dsegtree::query_1(x,y,node[pos]);
43
        return sum;
     }
44
45
     <u>inline</u> <u>int</u> query(<u>int</u> 1, <u>int</u> r, <u>int</u> x, <u>int</u> y){
46
        return ask(r,x,y)-ask(l-1,x,y);
47
48 }
   int aa[maxn];
49
50 int main() {
     \underline{cin} >> n >> m;
51
52
     rep(i,1,n) {
53
       cin>>aa[i];
        <u>if</u>(aa[i]!=aa[i-1]) bit::update(i,aa[i],1);
54
55
56
     register int x,y,b,c,d,e,a;
57
     \underline{\text{while}}(m--)
58
        cin>>a;
59
        if(a==1){
60
          <u>cin</u>>>x>>y;
          if(aa[x]==y) continue;
61
62
          <u>if</u>(aa[x]!=aa[x-1]) bit::update(x,aa[x],-1);
63
          \underline{if}(x+1 < n\&\&aa[x+1]!=aa[x]) bit::update(x+1,aa[x+1],-1);
          aa[x]=y
64
65
          \underline{if}(aa[x]!=aa[x-1]) bit::update(x,aa[x],1);
66
          \underline{if}(x+1<n\&\&aa[x+1]!=aa[x]) bit::update(x+1,aa[x+1],1);
67
68
          cin>>b>>c>>d>>e;
69
          int ans=bit::query(b,c,d,e);
70
          \underline{if}(aa[b] == aa[b-1] \&\&aa[b] >= d\&\&aa[b] <= e) ans++;
71
          cout<<ans<<end1</pre>
72
73
     }
```

1.20 可持久化数据结构-数组

```
2| <u>int</u> a0 [maxn];//初始数组
  <u>int</u> rt[maxn];//树根
   <u>template<typename</u> T><u>class</u> parray{<u>public</u>:
   #define nd node[now]
6
   #define ndp node[pre]
7
   #define mid (s+t)/2
8
     <u>int</u> cnt;
     <u>static</u> <u>int</u> 10,r0;//数组区间
9
     struct segnode{int l,r;T val;}node[maxn*20];
10
11
     void init(int s,int t,T a[]){
12
       10=s,r0=t,cnt=0;
13
       makearray(a);
14
15
     void makearray(T a[], int &now=rt[0], int s=10, int t=r0){
16
       now=++cnt;
17
       if(s==t){nd.val=a[s];return;}
18
       makearray(a,nd.l,s,mid);makearray(a,nd.r,mid+1,t);
19
20
     void update(int pre,int &now,int pos,T val,int s=10,int t=r0){
21
       now=++cnt;nd=ndp;
22
       if(s==t){nd.val=val; return;}
23
       if(pos<=mid) update(ndp.1,nd.1,pos,val,s,mid);</pre>
24
       else update(ndp.r,nd.r,pos,val,mid+1,t);
25
26
     T query(<u>int</u> now,<u>int</u> pos,<u>int</u> s=10,<u>int</u> t=r0){
27
       if(s==t) return nd.val;
28
       if(pos<=mid) return query(nd.1,pos,s,mid);</pre>
29
       else return query(nd.r,pos,mid+1,t);
30
31 <u>#undef</u> mid
32 };
33 parray<<u>int</u>>arr;
```

1.21 可持久化数据结构-并查集

```
1 //可持久化并查集,主席树实现
 2 int rt[maxn];
3 int 10,r0;
 4 class pdsu{public:
5 #define nd node[now]
6 #define ndp node[pre]
7
  #define mid (s+t)/2
8
    <u>int</u> cnt;
9
    struct segnode{int 1,r,fa,dep;}node[maxn*30];
10
    void init(int n){
11
      10=1, r0=n, cnt=0;
12
      makearray();
13
14
    void makearray(int &now=rt[0],int s=10,int t=r0){
15
      now=++cnt;nd=\{0,0,s,0\};
16
      if(s==t)return;
17
      makearray(nd.1,s,mid);makearray(nd.r,mid+1,t);
18
19
    void merge(int pre,int &now,int pos,int fa,int s=10,int t=r0){
20
      now=++cnt;nd=ndp;
21
      if(s==t){nd.fa=fa;return;}
22
      if(pos<=mid) merge(ndp.1,nd.1,pos,fa,s,mid);</pre>
23
      else merge(ndp.r,nd.r,pos,fa,mid+1,t);
24
25
    void update(int now,int pos,int s=10,int t=r0){
26
      <u>if</u>(s==t) {nd.dep++;<u>return</u>;}
27
      if(pos<=mid) update(nd.1,pos,s,mid);</pre>
28
      else update(nd.r,pos,mid+1,t);
29
30
    int query(int now,int pos,int s=10,int t=r0){
31
      if(s==t) return now;
32
      if(pos<=mid) return query(nd.1,pos,s,mid);</pre>
33
      else return query(nd.r,pos,mid+1,t);
34
  //找到第ver个版本的集合根:
35
36
    int find(int ver,int pos){
37
      int now=query(ver,pos);
38
      if(nd.fa==pos) return now;
```

```
return find(ver,nd.fa);
|40|
  //在ver1的基础上,合并a,b集合,得到ver2:
41
42
    void unite(int ver1,int ver2,int a,int b){
43
        rt[ver2]=rt[ver1];
        int fa=find(rt[ver2],a),fb=find(rt[ver2],b);
44
        if(node[fa].fa==node[fb].fa) return;
45
        if(node[fa].dep>node[fb].dep) swap(fa,fb);
46
        merge(rt[ver1],rt[ver2],node[fa].fa,node[fb].fa);
47
48
        if(node[fa].dep==node[fb].dep)
49
         update(rt[ver2],node[fb].fa);
50
  //复制ver1的状态到ver2,并查询a,b是否为同集合
51
    bool same(int ver1,int ver2,int a,int b){
52
      rt[ver2]=rt[ver1];
53
54
      int fa=find(rt[ver2],a),fb=find(rt[ver2],b);
55
      return node[fa].fa==node[fb].fa;
56
  //令ver2的状态回到ver1:
57
    void popback(int ver1,int ver2){
58
59
      rt[ver2]=rt[ver1];
60
  #undef mid
61
62| }dsu;
```

1.22 平衡树-老司机树

```
class odt{public:
    struct segnode{
      int l,r;mutable int val;
      bool operator<(const segnode &b)const {return 1<b.1;}</pre>
    };
6
    set < segnode > nd;
7
    void init(int n=maxn-5){nd.insert(1,n,0);}
    #define iter set<segnode>::iterator
8
9
    auto split(int pos){
10
      auto it=nd.lower_bound({pos,pos,0});
      if(it!=nd.end()&&it->l==pos) return it;
11
12
13
      int l=it->1,r=it->r,val=it->val;
|14|
      nd.erase(it);nd.insert({1,pos-1,val});
15
      return nd.insert({pos,r,val}).fi;
16
17
    void update(int l,int r,int val){
18
      auto itr=split(r+1);auto itl=split(l);
19
      nd.erase(itl,itr);nd.insert({1,r,val});
20
21
    int query(int pos){
22
          auto it=nd.lower_bound({pos,pos,0});
23
           if(it!=nd.end()&&it->l==pos) return it->val;
24
           it--; return it->val;
25
  }tree;
^{26}
```

1.23 平衡树-Treap

```
class splaytree{public:
   #define nd node[now]
   #define ndl node[node[now].son[0]]
   #define ndr node[node[now].son[1]]
       struct splaynode{
6
           int son[2],fa,val,size;
           splaynode(){size=1,fa=son[0]=son[1]=0;}
7
9
       int cnt,root;
10
       vector<splaynode> node;
11
       <u>inline</u> <u>void</u> pushup(<u>int</u> now){nd.size=ndl.size+ndr.size+1;}
12
       <u>inline</u> <u>void</u> pushdown(<u>int</u> now){}
13
       <u>inline</u> <u>int</u> wh(<u>int</u> now){<u>return</u> node[nd.fa].son[1]==now;}
14
       void rotate(int now){
           int fa=nd.fa,gf=node[fa].fa,c=wh(now);
15
16
           pushdown(fa);pushdown(now);
           if(gf) node[gf].son[wh(fa)]=now;
17
18
           nd.fa=gf;
```

```
node[fa].son[c]=nd.son[c^1];
          node[node[fa].son[c]].fa=fa;nd.son[c^1]=fa;node[fa].fa=now;
20
21
          pushup(fa);pushup(now);
22
23
      void splay(int now,int dst=0){
          for(;nd.fa!=dst;rotate(now))
24
25
              if(node[nd.fa].fa!=dst)rotate(wh(now)==wh(nd.fa)?nd.fa:now);
26
              if(!dst) root=now;
27
28
      void insert(int pos){
29
          int now=root,fa=0,val=node[pos].val;
30
          while(now) fa=now,now=val<nd.val?nd.son[0]:nd.son[1];</pre>
31
32
          node[fa].son[val>node[fa].val]=now;
33
          nd.fa=fa;
34
          splay(now);
35
36
      void order(int now){
37
          int l=nd.son[0],r=nd.son[1];
38
          nd.son[0]=nd.son[1]=nd.fa=0;
39
          nd.size=1;
40
          <u>if</u>(1) order(1);
41
          insert(now);
42
          if(r) order(r);
      }
43
      void merge(int a,int b){
44
          if(a==b) return ;
45
46
          splay(a);splay(b);
          if(node[a].size>node[b].size) swap(a,b);
47
48
          pre[a]=b;root=b;
          order(a);
49
50
51
      int kth(int now,int k){
          splay(now); int lsize=0;
52
53
          while(now){
54
              int lsum=lsize+ndl.size;
55
              if(k<=lsum) now=nd.son[0];</pre>
56
              else if(k==lsum+1) return now;
57
              else lsize=lsum+1,now=nd.son[1];
58
59
          return -1;
60
61
      splaytree(int n){
62
          node.resize(n+7,splaynode());
63
          rep(i,1,n) node[i].val=val[i];
64
          node[0].size=0;
65
          root=0,cnt=0;
66
67| };
```

1.24 平衡树-Splay

```
1| <u>int</u> root;//树根
  class splaytree{public:
    int fa[maxn], son[maxn][2], sz[maxn], val[maxn], cnt[maxn];
    int tot;//根权值重复次数子树大小
    inline void pushup(int now){
6
      sz[now]=sz[son[now][0]]+sz[son[now][1]]+cnt[now];
    }//更新当前节点信息
8
    <u>inline</u> <u>bool</u> getson(<u>int</u> now) {<u>return</u> now==son[fa[now]][1];}//真为右儿子
    inline void clear(int now){
9
      fa[now]=son[now][0]=son[now][1]=sz[now]=val[now]=cnt[now]=0;
10
11
12
    inline void rotate(int now){
      int f=fa[now],gf=fa[fa[now]],flag=getson(now);
13
14
      son[f][flag]=son[now][flag^1];
      fa[son[now][flag^1]]=f;
15
      son[now][flag^1]=f;
16
17
      fa[f]=now;fa[now]=gf;
18
      if(gf) son[gf][f==son[gf][1]]=now;
19
      pushup(now);pushup(f);
20
    }//旋转一层
    void splay(int now){
```

```
for(int f=fa[now];f=fa[now],f;rotate(now))
23
         if(fa[f]) rotate(getson(now)==getson(f)?f:now);
24
      root=now:
     }//旋转到根
25
     void insert(int x,int now=root,int f=0){
26
27
       \underline{if}(!now){}
28
         now=++tot;
29
         val [now] =x, cnt [now] ++;
30
         if(!root)root=now;
         pushup(now);
31
         if(f){}
32
33
           fa[now]=f;
34
           son[f][val[f]<x]=now;
35
          pushup(f);splay(now);
36
       }else if(val[now] == x){
37
38
         ++cnt[now];
39
         pushup(now);pushup(f);
40
         splay(now);
       }else insert(x,son[now][val[now]<x],now);</pre>
41
     }//插入新点
42
     int getrank(int x,int now=root,int ans=0){
43
44
      \underline{\text{while}}(1)
45
         if(x<val[now]) now=son[now][0];</pre>
46
         else {
47
           ans+=sz[son[now][0]];
48
           if(x==val[now]) return splay(now),ans+1;
49
           ans+=cnt[now];now=son[now][1];
50
      }
51
52
     }//多少个元素小于x
53
     int get(int k,int now=root){
54
      \underline{\text{while}}(1)
         \underline{if}(son[now][0]\&\&k \le sz[son[now][0]]) now=son[now][0];
55
56
57
           k=cnt[now]+sz[son[now][0]];
58
           if(k<=0) return val[now];</pre>
59
           now=son[now][1];
60
      }
61
     }//查询元素
62
63
     int pre(){
       int now=son[root][0];
64
65
       while(son[now][1]) now=son[now][1];
66
       return now;
     }//前驱
67
68
     int nxt(){
69
       int now=son[root][1];
       while(son[now][0]) now=son[now][0];
70
71
       return now;
72
     }//后缀
     int lower(int x){
73
74
       insert(x);
75
       int ans=val[pre()];
       erase(x); return ans;
76
77
     }//查询前驱
78
     int upper(int x){
79
       insert(x);
80
       int ans=val[nxt()];
81
       erase(x);return ans;
     }//查询后缀
82
83
     void erase(int x){
       getrank(x);
84
85
       if(cnt[root]>1){
86
         --cnt[root];pushup(root);
       }<u>else</u> <u>if</u>(!son[root][0]&&!son[root][1]){
87
88
         clear(root);root=0;
       }else if(!son[root][0]||!son[root][1]){
89
90
         int t=root;
         if(!son[root][0]) root=son[root][1];
91
         else root=son[root][0];
92|
```

```
fa[root]=0;clear(t);
94
95
           int now=pre(),t=root;
96
           splay(now);
           fa[son[t][1]]=now;
97
98
           son[now][1]=son[t][1];
99
           clear(t);pushup(root);
100|
101
      }//删除元素
102 }tree;
103 | \underline{int}  main() {
104
      cin>>n;
105
      \underline{\text{while}}(n--)
106
         int opt,x;cin>>opt>>x;
107
         if(opt==1) tree.insert(x);
108
         if(opt==2) tree.erase(x);
109
         <u>if</u>(opt==3) <u>cout</u><<tree.getrank(x)<<<u>endl</u>;
        if(opt==4) cout<<tree.get(x)<<endl;</pre>
110|
111
         <u>if</u>(opt==5) <u>cout</u><<tree.lower(x)<<<u>endl</u>;
112
         <u>if</u>(opt==6) <u>cout</u><<tree.upper(x)<<<u>endl</u>;
113
114
      return 0;
115|}
```

2 Graph

2.1 基础-链式前向星

```
class graph{public:
    struct node{int to,next,cost;};
    node e[maxm]; int head[maxn],nume;
    void init(int n=maxn-5){nume=0;fill_n(head,n+1,0);}
    void add(int a,int b,int c){e[++nume]={b,head[a],c};head[a]=nume;}
}
```

2.2 最大流-isap算法

```
1 template < typename T > class mxf {public:
     struct node{int to,next;T cap;}e[maxm<<1];</pre>
     int cur[maxn],head[maxn],dis[maxn],gap[maxn];
     int nume=1,s,t,tot;
     void init(int n){
6
      rep(i,0,n) head[i]=gap[i]=dis[i]=0;
      nume=1;
     }
8
9
     void add(int a,int b,T c){
10
       e[++nume]={b,head[a],c};head[a]=nume;
|11|
       e[++nume]={a,head[b],0};head[b]=nume;
12
13
     T dfs(int now, T flow=INF){
       if (now==t||!flow) return flow;
14
15
       T use=0,tmp;
16
       int d=dis[now]-1,to;
       for (int &i=cur[now];i;i=e[i].next) {
17
         <u>if</u>(dis[to=e[i].to]==d&&(tmp=e[i].cap)){
18
19
           e[i].cap-=(tmp=dfs(to,min(flow-use,tmp)));
20
           e[i^1].cap+=tmp;
21
           if((use+=tmp)==flow) return use;
22
23
24
       <u>if</u> (!--gap[dis[now]]) dis[s]=tot+1;
25
       ++gap[++dis[now]];
26
       cur[now] = head[now];
27
       return use;
28
29
    T getflow(<u>int</u> ss,<u>int</u> tt,<u>int</u> n,T ans=0){
30
       tot=n; s=ss; t=tt; gap[0]=tot;
31
       memcpy(cur,head,(tot+1)<<2);</pre>
32
       while(dis[s]<=tot) ans+=dfs(s);</pre>
33
       return ans;
34
    }
35 };
36 mxf<int> net;
```

2.3 最大流-dinic算法

```
1 template<typename T>class mxf{public:
     struct node{int to,next;T cap;}e[maxm<<1];</pre>
3
     int cur[maxn],head[maxn],que[maxn],dis[maxn];
    int nume=1,s,t,tot,tp,ed;
5
    inline void adde(int a,int b,T c){e[++nume]={b,head[a],c};head[a]=nume;}
    inline void add(int a,int b,T c){adde(a,b,c);adde(b,a,0);}
    void init(int n=maxn-1){memset(head,0,(n+1)<<2);nume=1;}</pre>
     bool bfs(){
9
      rep(i,0,ed) dis[que[i]]=0;
10
      dis[t]=1,que[0]=t;
11
      tp=0,ed=1;
12
      cur[t]=head[t];
13
      int now,to;
      while(tp!=ed) for(int i=head[now=que[tp++]];i;i=e[i].next)
14
15
        <u>if</u>(dis[to=e[i].to]==0&&e[i^1].cap>0){
16
          cur[to]=head[to];
17
          dis[to]=dis[now]+1;
18
          if((que[ed++]=to)==s) return true;
19
20
      <u>return</u> <u>false</u>;
21
|22|
    T dfs(<u>int</u> now,T flow=INF){
23
      if(now==t||flow==0) return flow;
24
      int to,d=dis[now]-1;
25
      T use=0,tmp;
26
      for(int &i=cur[now];i;i=e[i].next){
27
        if(dis[to=e[i].to]!=d||!(tmp=e[i].cap))continue;
28
        e[i].cap-=(tmp=dfs(to,min(tmp,flow-use)));
29
        e[i^1].cap+=tmp,use+=tmp;
30
        if(use==flow) return use;
31
32
      if(use==0) dis[now]=-1;
      return use;
33
34
35
    T getflow(int ss,int tt,int n,T ans=0){
36
      s=ss,t=tt,tot=n;
37
      while(bfs())ans+=dfs(s);
38
      return ans;
39
40 };
41 mxf<int> net;
```

网络流-上下界可行流-dinic

```
template<typename T>class mxf{public:
     struct node{int id,to,rev;T cap;}w[maxm<<1],e[maxm<<1];</pre>
     int cur[maxn],head[maxn],dis[maxn],que[maxn];
5
     T up [maxm<<1], in [maxn], sum, ans [maxm<<1];
6
     int num[maxn], numv, nume, s, t, tot, tp, ed, fr[maxm<<1];</pre>
7
     bool bfs(){
8
      rep(i,0,ed) dis[que[i]]=0;
9
      dis[t]=1,que[0]=t;
10
      tp=0,ed=1;
11
       cur[t]=head[t];
12
      int now, to;
13
       while(tp!=ed) for(int i=head[now=que[tp++]];i<=num[now];++i)</pre>
14
         <u>if</u>(dis[to=e[i].to]==0&&e[e[i].rev].cap>0){
15
           cur[to] = head[to];
           dis[to]=dis[now]+1;
16
           if((que[ed++]=to)==s) return true;
17
18
19
      return false;
20
21
     T dfs(<u>int</u> now,T flow=0x3f3f3f3f){
      if(now==t||flow==0) return flow;
22
23
       int to,d=dis[now]-1;
|24|
      T use=0,tmp;
25
       for(int &i=cur[now];i<=num[now];++i){</pre>
26
         if(dis[to=e[i].to]!=d||!(tmp=e[i].cap))continue;
27
         e[i].cap-=(tmp=dfs(to,min(tmp,flow-use)));
28
         e[e[i].rev].cap+=tmp,use+=tmp;
29
         if(use==flow) return use;
```

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```
if(use==0) dis[now]=-1;
31
32
       return use;
33
34
    T getflow(<u>int</u> ss,<u>int</u> tt,<u>int</u> n,T ans=0){
35
       s=ss,t=tt,tot=n;
       while(bfs())ans+=dfs(s);
36
37
       return ans;
38
39
     void init(int n){
      rep(i,0,n) num[i]=in[i]=head[i]=dis[i]=0;
40
41
      nume=0;tot=n;sum=0;
42
43
     void add(int a,int b,T c,int id){
44
      w[++nume]=(node){id,b,0,c};++num[a],fr[nume]=a;
45
       w[++nume]=(node)\{0,a,0,0\}; ++num[b],fr[nume]=b;
46
47
     void addbound(int a, int b, T c, T d, int id){
48
       add(a,b,d-c,id);
49
       up[id]=d,in[b]+=c,in[a]-=c;
     }
50
     bool fesbflow(int n){
51
52
       s=n+1, t=n+2; numv=n; tot=t;
53
       rep(i,1,numv){
         <u>if</u>(in[i]>0) add(s,i,in[i],0),sum+=in[i];
54
55
        <u>if</u>(in[i]<0) add(i,t,-in[i],0);
56
57
       head[1]=1;
       rep(i,2,tot) head[i]=head[i-1]+num[i-1];
58
59
       rep(i,1,tot-1) num[i]=head[i+1]-1;
60
       num[tot] = nume;
61
       rep(i,1,nume){
        e[head[fr[i]]+cur[fr[i]]++]=w[i];
62
63
         <u>if</u>(!(i%2)){
64
           e[head[fr[i]]+cur[fr[i]]-1].rev=head[w[i].to]+cur[w[i].to]-1;
           e[head[w[i].to]+cur[w[i].to]-1].rev=head[fr[i]]+cur[fr[i]]-1;
65
66
67
68
       T flow=getflow(s,t,t);
69
       if(flow<sum) return 0;</pre>
       rep(i,1,nume){
70
71
        node &x=e[i];
72
         if(x.id) ans[x.id]=up[x.id]-x.cap;
73
74
       return 1;
75
    }
76|};
77 mxf<int> net;
```

2.5 网络流-上下界网络流-isap

```
1 template < typename T > class mxf {public:
     struct node{int to,rev;T cap;}w[maxm<<1],e[maxm<<1];</pre>
     int cur[maxn],head[maxn],dis[maxn],gap[maxn];
     int num[maxn],numv,nume,s,t,tot,last,fr[maxm<<1];</pre>
     T in[maxn],sum;
     T dfs(int now, T flow=INF){
       <u>if</u> (now==t||!flow) <u>return</u> flow;
       T use=0,tmp;
8
9
       int d=dis[now]-1,to;
10
       for (int &i=cur[now];i<=num[now];++i) {</pre>
11
         <u>if</u>(dis[to=e[i].to]==d&&(tmp=e[i].cap)){
12
           e[i].cap-=(tmp=dfs(to,min(flow-use,tmp)));
           e[e[i].rev].cap+=tmp;
13
           if((use+=tmp)==flow) return use;
|14|
15
16
17
       <u>if</u> (!--gap[dis[now]]) dis[s]=tot+1;
       ++gap[++dis[now]];
18
19
       cur[now] = head[now];
20
       return use;
21
22
     T getflow(<u>int</u> ss,<u>int</u> tt,<u>int</u> n,T ans=0){
```

```
rep(i,0,n)dis[i]=gap[i]=0;
24
       tot=n;s=ss;t=tt;gap[0]=tot;
25
       memcpy(cur,head,(tot+1)<<2);</pre>
26
       while(dis[s]<=tot) ans+=dfs(s);</pre>
27
       return ans;
28
29
     void init(int n){
30
       rep(i,0,n) num[i]=in[i]=head[i]=dis[i]=0;
31
       nume=0;tot=n;sum=0;
32
33
     void add(int a,int b,T c){
34
       w[++nume]=(node){b,0,c};++num[a],fr[nume]=a;
35
       w[++nume] = (node)\{a,0,0\}; ++num[b],fr[nume] = b;
36
37
     void addbound(int a,int b,T c,T d){
       add(a,b,d-c);
38
       in[b]+=c,in[a]-=c;
39
     }
40
41
     void makeflow(int n){
42
       s=n+1, t=n+2; numv=n; tot=t;
       rep(i,1,numv){
43
         <u>if</u>(in[i]>0) add(s,i,in[i]),sum+=in[i];
44
45
        <u>if</u>(in[i]<0) add(i,t,-in[i]);
46
47
      head[1]=1;
48
      rep(i,2,tot) head[i]=head[i-1]+num[i-1];
49
       rep(i,1,tot-1) num[i]=head[i+1]-1;
50
       num[tot] = nume;
51
       rep(i,1,nume){
52
        e[head[fr[i]]+cur[fr[i]]++]=w[i];
53
         if(!(i%2)){
54
           e[head[fr[i]]+cur[fr[i]]-1].rev=head[w[i].to]+cur[w[i].to]-1;
           e[head[w[i].to]+cur[w[i].to]-1].rev=head[fr[i]]+cur[fr[i]]-1;
55
56
       }
57
58
     T fesbflow(int n){
59
60
       makeflow(n);
61
       T flow=getflow(s,t,t);
62
       if(flow!=sum) return -1;
63
       return flow;
|64|
     T fesbflow(<u>int</u> ss, <u>int</u> tt, <u>int</u> n){
65
66
       add(tt,ss,INF);
67
       makeflow(n);
68
       rep(i,head[tt],num[tt])
69
         <u>if</u>(e[i].to==ss&&e[i].cap==INF) {
70
          last=i;
71
           break;
72
73
       T flow=getflow(s,t,t);
74
       if(flow!=sum) return -1;
75
       return flow;
76
77
    T maxflow(int ss,int tt,int n){
78
       if(fesbflow(ss,tt,n)==-1) return -1;
79
       return getflow(ss,tt,n+2);
80
81
     T minflow(int ss,int tt,int n){
82
       if(fesbflow(ss,tt,n)==-1) return -1;
83
       node &x=e[last];
84
       T ans=INF-x.cap;
85
       x.cap=e[x.rev].cap=0;
       return ans-getflow(tt,ss,n+2);
86
87
88|};
89 mxf<<u>int</u>> net;
```

2.6 网络流-SW全局最小割算法

```
1 //全局最小割
2 #include <cstdio>
3 #include <iostream>
```

```
4 #include <cstring>
  #include <algorithm>
6 #include <queue>
7 #include <numeric>
8 typedef long long LL;
 9 | \underline{\text{const}} | \underline{\text{int}} | \underline{\text{MAXV}} = 3010;
10 | \underline{\text{const}} \underline{\text{int}} | \text{MAXE} = 100010 * 2;
12 int head [MAXV], val [MAXV], ecnt;
13 int to [MAXE], next [MAXE], weight [MAXE];
14 bool vis[MAXV];
15 int fa[MAXV], link[MAXV];
16 int n, m;
17 void init() {
18
    memset(head + 1, -1, \underline{\text{sizeof}}(\underline{\text{int}}) * n);
    memset(link + 1, -1, \underline{\text{sizeof}}(\underline{\text{int}}) * n);
19
20
    <u>for</u> (<u>int</u> i = 1; i <= n; ++i)
21
      fa[i] = i;
     ecnt = 0;
22
23 }
24 void add_edge(int u, int v, int w) {
    to[ecnt] = v; weight[ecnt] = w; next[ecnt] = head[u]; head[u] = ecnt++;
    to[ecnt] = u; weight[ecnt] = w; next[ecnt] = head[v]; head[v] = ecnt++;
27| \}
28 int findset(int u) {
    return u == fa[u] ? u : fa[u] = findset(fa[u]);
29
30|}
31 void merge(int u, int v) {
32
     \underline{int} p = u;
     while ("link[p]) p = link[p];
33
34
     link[p] = v;
35
    fa[v] = u;
|36| }
37 int MinimumCutPhase(int cnt, int &s, int &t) {
    38
39
40
     std::priority_queue<std::pair<int, int>> que;
     t = 1;
41
42
     while (--cnt) {
43
       vis[s = t] = true;
       for (int u = s; ~u; u = link[u]) {
44
         for (int p = head[u]; ~p; p = next[p]) {
45
46
           int v = findset(to[p]);
47
           if (!vis[v])
             que.push(std::make_pair(val[v] += weight[p], v));
48
49
       }
50
       t = 0;
51
52
       while (!t) {
53
         if (que.empty()) return 0;
54
         auto pa = que.top(); que.pop();
55
         if (val[pa.second] == pa.first) t = pa.second;
56
     }
57
58
     return val[t];
59|}
60 int StoerWagner() {
     int res = INF;
61
62
     for (int i = n, s, t; i > 1; --i) {
63
       res = std::min(res, MinimumCutPhase(i, s, t));
64
       if (res == 0)
         break;
65
66
       merge(s, t);
67
68
     return res;
69 }
70 \mid \underline{int} \mid main()  {
     while (scanf("dd', &n, &m) != EOF) {
71
72
       init();
73
       for (int i = 0, u, v, w; i < m; ++i) {</pre>
         scanf("%d%d%d", &u, &v, &w);
74
75
         add_edge(u, v, w);
76
```

```
77 printf("%d\n", StoerWagner());
78 }
79 }
```

2.7 网络流-最大密度子图

```
const double eps=1e-8;
  template<typename T>class mxf{public:
    struct node{int to,next;T cap;}e[maxm<<1];</pre>
    int cur[maxn],head[maxn],que[maxn],dis[maxn],nume=1,s,t;
5
    <u>inline void</u> adde(<u>int</u> a,<u>int</u> b,T c){
6
       e[++nume]={b,head[a],c};head[a]=nume;
8
    inline void add(int a, int b,T c){adde(a,b,c);adde(b,a,0);}
9
    void init(int n=maxn-1){memset(head,0,(n+1)<<2);nume=1;}</pre>
10
    bool bfs(){
11
      memset(dis,-1,(t+1)<<2);
12
      dis[t]=0,que[0]=t;
13
       int tp=0,ed=1;
14
       while(tp!=ed){
15
        int now=que[tp++]; if(tp==maxn) tp=0;
16
        for(int i=head[now];i;i=e[i].next){
17
           int to=e[i].to;
          if(dis[to] == -1 \&\&e[i^1].cap>0){
18
19
             dis[to]=dis[now]+1;
             if(to==s) return true;
20
21
            que [ed++]=to;
22
             if(ed==maxn) ed=0;
23
        }
24
25
26
       <u>return</u> <u>false</u>;
27
28
    T dfs(<u>int</u> now,T flow=1e9){
29
       if(now==t||flow==0) return flow;
30
      T use=0;
31
       for(int &i=head[now];i&&use!=flow;i=e[i].next){
32
         int to=e[i].to;
33
         if(dis[to]+1!=dis[now])continue;
34
        T tmp=dfs(to,min(e[i].cap,flow-use));
35
        e[i].cap-=tmp,e[i^1].cap+=tmp,use+=tmp;
36
37
       if(use==0) dis[now]=-1;
38
      return use;
39
40
    T getflow(int ss, int tt){
41
       s=ss,t=tt;T ans=0;
42
      memcpy(cur, head, (t+1) << 2);
43
       while(bfs()){
44
        ans+=dfs(s);
45
        memcpy(head, cur, (t+1) << 2);
46
47
      <u>return</u> ans;
48
    }
49|};
50 mxf<double> net;
51 const int maxn2=500;
  int mt[maxn2] [maxn2];
52
53 double d[maxn2];
54 int val[maxn2], tag[maxn2];
55 void init(int n){rep(i,1,n)rep(j,i+1,n) mt[i][j]=mt[j][i]=0;}
56 void adde(int a, int b, int v){mt[a][b]=mt[b][a]=v;}
  <u>const double</u> all=400*2200;//点权和+边权和*2
57
58 bool check(double mid){
59
    int s=n+1, t=n+2;
    net.init(n+3);
60
61
    double f=0;
62
    rep(i,1,n){
63
      d[i]=0.0;
64
      rep(j,1,n){
65
        if(i==j||!mt[i][j])continue;
66
        d[i]+=mt[i][j];
67 //如果公式计算出来,边权跟mid有关,就要加上相应的mid
```

```
net.add(i,j,mt[i][j]);
69
     }
70
71
     rep(i,1,n){
72
        <u>if</u>(tag[i]){
73
          f+=all+2*mid-d[i];
74
          net.add(s,i,INF);
75
76
          net.add(s,i,all);
77
          net.add(i,t,all+2*mid-d[i]);
   //有点权的话,这个2*mid还要再乘那个点权
78
79
     }
80
81
     double x=net.getflow(s,t);
82
     \underline{\text{double}} ans=(all*n-f-x)*0.5;
83
     return ans>eps;
84|}
85
   int main(){IO;
86
     cin>>casn;
87
     cout << fixed << setprecision(10);</pre>
88
     rep(kase,1,casn){
89
        cin>>n;
90
        init(n);
91
        rep(i,1,n)<u>cin</u>>>val[i];
92|
        rep(i,1,n)rep(j,i+1,n)
93
          <u>if</u>(val[j]<val[i]) adde(i,j,1);
94
          rep(i,1,n) tag[i]=0;//是否必须用
95
        double l=0,r=n,mid;
96
        \underline{\text{while}}(r-1) = \exp s
97
          mid=(1+r)*0.5;
98
          if(check(mid)) l=mid;
99
          else r=mid;
100
        cout<<"Case #"<<kase<<": "<<(1+r)*0.5<<'\n';</pre>
101
102
103
     return 0;
104 }
```

费用流-原始对偶算法

```
1 //原始对偶算法,dijkstra寻找增广路,单路增广
 2 | //无优化空间
3 template < typename T1, typename T2 > class mcf {public:
  #define pdi pair<T2,int>
    priority_queue<pdi, vector<pdi>, greater<pdi>>que;
     struct node{int to;T1 cap;T2 cost;int rev;};
 6
     int prev[maxn],pree[maxn],numv;
8
     T2 dis[maxn],h[maxn];
9
     vector<node> g[maxn];
10
     void init(int n=maxn-2){
11
       numv=n;
12
       rep(i,0,n) g[i].clear();
13
14
     <u>inline</u> <u>void</u> add(<u>int</u> from, <u>int</u> to, T1 cap, T2 cost){
       g[from].push_back({to,cap,cost,(<u>int</u>)g[to].size()});
15
16
       g[to].push_back({from,0,-cost,(<u>int</u>)g[from].size()-1});
17
18
     pair<T1,T2>getcost(<u>int</u> s,<u>int</u> t,<u>int</u> n){
19
       numv=n:
20
       T1 flow=0; T2 cost=0;
21
       fill_n(h,numv+1,0);
22
       while(1){
23
           fill_n(dis,numv+1,INF);
24
           dis[s]=0;que.push(make_pair(0,s));
25
           while(!que.empty()){
26
             auto now=que.top();que.pop();
27
             if(dis[now.second] < now.first) continue;</pre>
             int x=now.second;
28
29
             int cnt=0;
30
             for(auto &i:g[x])
31
               if(i.cap>0\&\&dis[i.to]>dis[x]+h[x]-h[i.to]+i.cost){
                 dis[i.to]=dis[x]+i.cost+h[x]-h[i.to];
32
33
                 prev[i.to]=x;
```

```
pree[i.to]=cnt++;
35
                que.push(make_pair(dis[i.to],i.to));
36
               }else cnt++;
37
          if(dis[t] == INF) break;
38
39
          rep(i,0,numv) h[i]+=dis[i];
40
          T1 d=INF;
41
           for(int now=t;now!=s;now=prev[now])
42
             d=min(d,g[prev[now]][pree[now]].cap);
43
           if(d==INF)break;
44
           flow+=d;cost+=d*h[t];
45
           for(int now=t;now!=s;now=prev[now]){
46
            node &e=g[prev[now]][pree[now]];
47
             e.cap-=d,g[now][e.rev].cap+=d;
48
49
      return make_pair(flow,cost);
50
51
52 };
53 mcf<<u>int</u>,<u>int</u>>net;
```

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费用流-EK算法

```
1 //zkw费用流,单路增广
  //可将单路增广改为多路增广+当前弧优化,但提升不大
  template<typename T1,typename T2>class mcf{public:
     int nume=1,s,t,numv,head[maxn],pre[maxn];
     bool vis[maxn];
 6
     queue<<u>int</u>>q;
     T1 flow[maxn], mflow;
     T2 dis[maxn], mcost;
     struct node{int to,next;T1 cap;T2 cost;}e[maxm<<1];</pre>
10
     void init(int n=maxn-10){
11
       numv=n;
|12|
       fill(head,head+n+2,0);nume=1,mflow=mcost=0;
13
|14|
     inline void add(int from,int to,int cap,T2 cost){
15
       e[++nume] = {to,head[from],cap,cost};head[from] = nume;
16
       e[++nume]={from,head[to],0,-cost};head[to]=nume;
17
18
     bool spfa(){
19
       fill(dis,dis+2+numv,INF);
20
       fill(vis, vis+2+numv, false);
21
       dis[s]=0;flow[s]=INF;q.push(s);
22
       while (!q.empty()){
23
         int now=q.front();q.pop();
24
        vis[now] = false;
25
         for (int i=head[now];i;i=e[i].next){
26
           int to=e[i].to;
27
          T2 cost=e[i].cost;
28
           if (e[i].cap&&dis[now]+cost<dis[to]){</pre>
29
             dis[to]=dis[now]+cost;
30
            flow[to]=min(flow[now],e[i].cap);
31
            pre[to]=i;
             <u>if</u> (!vis[to]){
32
33
              vis[to]=true;
              q.push(to);
34
35
36
        }
37
38
39
       return dis[t] < INF;</pre>
40
41
     void dfs(){
42
       int x=t;
43
       \underline{\text{while}} (x!=s){
44
        int i=pre[x];
        e[i].cap-=flow[t];
45
46
        e[i^1].cap+=flow[t];
        x=e[i^1].to;
47
48
49
      mflow+=flow[t];
      mcost+=(T2)flow[t]*dis[t];
50
```

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```
51  }
52  pair<T1,T2> getcost(int ss,int tt){
53    s=ss,t=tt;
54    while (spfa())dfs();
55    return make_pair(mflow,mcost);
56  }
57  };
58  mcf<int,int>net;
```

```
强连通-求桥
  2.10
 1| struct node{int to,nx;}e[maxm],e2[maxm];
  int head[maxn],head2[maxn],nume,nume2;
3 bool bg[maxm]; int dfn[maxn], low[maxn];
  int numc,cnt,vis1[maxn],belong[maxn];
 5 void add(int a, int b){
     e[++nume]={b,head[a]};head[a]=nume;
8 void add2(int a, int b){
9
     e2[++nume2]={b,head2[a]};head2[a]=nume2;
10 }
  void tdfs(int now,int in){
11
12
     dfn[now] = low[now] = ++cnt;
     for(int i=head[now];i;i=e[i].nx){
13
14
       int to=e[i].to;
       <u>if</u>(!dfn[to]) {
15
16
         tdfs(to,i);
         low[now] = min(low[now], low[to]);
17
18
         <u>if</u>(low[to]>dfn[now]) bg[i]=bg[i^1]=1;
19
       }else if(i!=(in^1))low[now]=min(low[now],dfn[to]);
20
21|
22
  void dfs(int now){
23
     belong[now]=numc;
     for(int i=head[now];i;i=e[i].nx){
24
25
       int to=e[i].to;
26
       if(belong[to]||bg[i]) continue;
27
       dfs(to);
28
29
  }
30
  pii dfs2(<u>int</u> now, <u>int</u> fa, <u>int</u> d=0){
31
     vis1[now]=1;pii x={d,now};
     for(int i=head2[now];i;i=e2[i].nx){
32
33
       int to=e2[i].to;
34
       if(to==fa) continue;
35
       pii t=dfs2(to,now,d+1);
36
       \underline{if}(x < t) x = t;
37
38
     return x;
39 }
40 int main() {
41
     cin>>casn;
42
     while(casn--){
43
       cin >> n >> m;
44
       rep(i,1,n){
45
         belong[i]=head[i]=head2[i]=0;
46
         low[i]=vis1[i]=dfn[i]=0;
47
48
       rep(i,1,m*2+2)bg[i]=0;
49
       cnt=numc=0,nume=nume2=1;
50
       \underline{\text{while}}(m--) {
         <u>int</u> a,b;<u>cin</u>>>a>>b;
51
52
         add(a,b);add(b,a);
53
54
       rep(i,1,n) <u>if</u>(!dfn[i]) tdfs(i,0);
55
       rep(i,1,n)
         if(!belong[i]) {
56
57
           numc++;dfs(i);
58
59
       k=0;
60
       <u>for</u>(<u>int</u> i=2;i<=nume;i+=2){
         <u>int</u> a=e[i].to,b=e[i^1].to;
61
62
         if(belong[a]!=belong[b]){
```

```
k++;
           add2(belong[a],belong[b]);
64
65
           add2(belong[b],belong[a]);
66
       }
67
       int c=0;
68
69
      rep(i,1,numc){
70
         if(vis1[i]) continue;
71
         int t=dfs2(i,i).se;
72
         c=max(c,dfs2(t,t).fi);
73
74
       cout<<k-c<<endl;</pre>
75
76 }
```

2.11 强连通-求割顶

```
1 void init(int n){
     for(int i=0;i<=n;i++){if(i!=0) f[i]=i;}</pre>
|3|
     memset(cut, false, n << 2); memset(low, 0, n << 2);
4
     memset(dfn,0,n<<2); memset(head,0,n<<2);
5
     nume=cnt=0;
6
7
   void tarjan(int u,int p){
     low[u]=dfn[u]=++cnt;
     int son=0;
     for(int i=head[u];i;i=e[i].next){
11
       int v=e[i].to;
12
       if(v==p)continue;
13
       \underline{if}(!dfn[v]){
14
         son++;
15
         tarjan(v,u);
         low[u]=min(low[u],low[v]);
16
17
         if(u!=p\&\&low[v]>=dfn[u]){
18
           cut[u]=<u>true</u>;
19
20
       } else low[u]=min(low[u],dfn[v]);
21
22
     \underline{if}(u==p\&\&son>1) cut[u]=\underline{true};
|23| }
```

2.12 强连通-缩点

```
int casn,n,m,k;
  struct node {int to,next;}e[maxm];int head[maxn],nume;
  inline void add(int a,int b){e[++nume]=(node){b,head[a]};head[a]=nume;}
  namespace tarjan{
    int stk[maxn],top,cnt,dfn[maxn],low[maxn],numc,belong[maxn],vis[maxn];
6
    vector<int>g[maxn];
    void tdfs(int now,int fa){
8
      dfn[now] = low[now] = ++cnt;
9
      stk[top++]=now,vis[now]=1;
10
      for(int i=head[now];i;i=e[i].next){
11
        int to=e[i].to;
12
        if(!dfn[to]){tdfs(to,now);low[now]=min(low[now],low[to]);}
13
        else if(vis[to]) low[now]=min(low[now],dfn[to]);
14
15
        if(to==fa) continue;
16
        if(!dfn[to]){tdfs(to,now);low[now]=min(low[now],low[to]);}
17
        else low[now] = min(low[now], dfn[to]);
18
      }
19
20
      if(low[now] ==dfn[now]){
21
        numc++;
22
        int to;
23
        do{to=stk[--top];
          belong[to]=numc;
25
          vis[to]=0;
26
        }while(to!=now);
27
    }
28
29
    void makegraph(int n){
30
      for(int i=1;i<=n;i++) if(!dfn[i]) tdfs(i,i);</pre>
31
      rep(i,1,n){
```

```
int a=belong[i];
         for(int j=head[i];j;j=e[j].next){
33
34
           int b=belong[e[j].to];
35
           \underline{if}(a!=b){
36
             g[a].emplace_back(b);
             g[b].emplace_back(a);
37
38
39
40
     }
41
42|}
```

2.13 强连通-无向图双连通

```
1 int stk[maxn],top,cnt,dfn[maxn],low[maxn],numc,belong[maxn],vis[maxn],sz[maxn];
  struct node {int to,next;}e[maxm];int head[maxn],nume;
  inline void add(int a,int b){e[++nume]=(node){b,head[a]};head[a]=nume;}
  void tdfs(int now,int pre){
    dfn[now] = low[now] = ++cnt;
    stk[top++]=now,vis[now]=1;
    for(int i=head[now];i;i=e[i].next){
8
      int to=e[i].to;
9
      if(to==pre) continue;
10
      if(!dfn[to]){tdfs(to,now);low[now]=min(low[now],low[to]);}
11
      else low[now] = min(low[now], dfn[to]);
12
    if(low[now] == dfn[now]){
13
14
      numc++;
15
      int to;
16
      do{to=stk[--top];
17
        belong[to]=numc;
18
        sz[numc]++;
19
      }while(to!=now);
20
21 }
```

2.14 最短路-Dijkstra算法

```
1 namespace dij{
    struct road{
      int now; ll dis;
|4|
      road(<u>int</u> a=0,ll _dis=0):now(a),dis(_dis){}
5
      bool operator<(const road &x)const {return dis>x.dis};
6
7
    11 dis[maxn];
8
    bool vis[maxn];
9
    priority_queue<road>que;
    void cal(int st,int n=maxn-5){
10
11
      fill_n(dis,n+1,1e18);fill_n(vis,n+1,0);
12
      que.emplace(st,0);dis[st]=0;
13
      while(!que.empty()){
|14|
        road t=que.top();que.pop();
15
        for(auto e:g[t.now]){
          11 cost=t.dis+e.cost;
16
17
          <u>if</u>(cost<dis[e.to]){</pre>
18
            dis[e.to]=cost;
19
            que.emplace(e.to,cost);
20
21
22
23
24
  template<typename T> class shortpath{public:
25
26
    T dis[maxn];
27
    bool vis[maxn];
    #define pdi pair<T, int>
28
29
    priority_queue<pdi, vector<pdi>, greater<pdi> >que;
30
    void cal(int st,int n,vector<pdi> g[]){
31
      fill_n(dis,n+1,INF);fill_n(vis,n+1,0);
32
      que.emplace(0,st);dis[st]=0;
33
      while(!que.empty()){
34
        pdi t=que.top();que.pop();
35
        for(auto e:g[t.se]){
          T cost=t.fi+e.fi;
36
```

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```
最短路-floyd求最小环
  2.15
 1 const int inf=500;
   const int N=105;
3 <u>int</u> n,m,ans,num;
   int dis[N][N],mapp[N][N],pre[N][N],path[N];
  void floyd() {
6
     ans=inf;
     for(int k=1; k<=n; k++) {</pre>
       for(int i=1; i<=n; i++) {</pre>
8
9
         <u>for</u>(<u>int</u> j=1; j<=n; j++)
            <u>if(i==j||k==i||k==j)</u> {
10
11
              continue;
12
           <u>if</u>(dis[i][j]+mapp[i][k]+mapp[k][j]<ans) {
13
              ans=dis[i][j]+mapp[i][k]+mapp[k][j];
14
15
              path[num++]=k;
16
17
              int temp=j;
              while(temp!=i) {
18
19
                path[num++]=temp;
20
                temp=pre[i][temp];
21
             path[num++]=i;
22
23
24
         }
25
26
       for(int i=1; i<=n; i++) {</pre>
         for(int j=1; j<=n; j++)
if(i==j||k==i||k==j) {</pre>
27
28
29
              continue;
30
            <u>if</u>(dis[i][j]>dis[i][k]+dis[k][j]) {
31
              dis[i][j]=dis[i][k]+dis[k][j];
32
              pre[i][j]=pre[k][j];
33
34
35
36
     }
37
38|}
39
   int main() {
40
     \underline{\operatorname{cin}} >> n >> m;
41
     for(int i=1; i<=n; i++) {</pre>
42
       for(int j=1; j<=n; j++) {</pre>
         pre[i][j]=i;
43
44
         <u>if</u>(i==j)
           mapp[i][j]=dis[i][j]=0;
45
46
         } else {
47
           mapp[i][j]=dis[i][j]=inf;
48
       }
49
50
     }
51
     <u>int</u> u,v,w;
     for(int i=1; i<=m; i++) {</pre>
52
53
       cin>>u>>v>>w;
       mapp[u][v]=mapp[v][u]=dis[u][v]=dis[v][u]=min(w,mapp[u][v]);
54
55
56
     floyd();
     if(ans==inf) {
57
       printf("No solution.\n");
58
59
     } <u>else</u> {
60
       for(int i=num-1; i>=0; i--) {
         printf("%d",path[i]);
61
         if(i==0) {
62|
```

2.16 树-Prim算法

```
1 //最小生成树prim算法
  int head[30],next[200],point[200],val[200],size,dist[30];
3 <u>bool</u> vis[30];
4 void add (int a, int b, int v){
5
    int i;
6
    for(i=head[a];~i;i=next[i]){
       \underline{if}(point[i]==b){
8
        if(val[i]>v)val[i]=v;
9
        return;
10
    }
11
    point[size]=b;
12
13
    val[size]=v;
|14|
    next[size] = head[a];
15
    head[a]=size++;
16|}
17
  struct cmp{
    bool operator()(pii a,pii b){
18
19
      return a.first>b.first;
20
21 };
22
  void prim(int s){
23
    int i,ans=0;
24
    memset(dist,-1,sizeof(dist));
25
    memset(vis,0,sizeof(vis));
26
    priority_queue<pii, vector<pii>, cmp>q;
27
    for (i=head[s];~i;i=next[i]){
28
      dist[point[i]]=val[i];
29
      q.push(make_pair(dist[point[i]],point[i]));
30
31
    dist[s]=0;
32
    vis[s]=1;
33
    while(!q.empty()){
34
      pii u=q.top();
35
      q.pop();
36
      if(vis[u.second])continue;
37
      vis[u.second]=1;
38
      ans+=u.first;
      for(i=head[u.second];~i;i=next[i]){
39
40
        int j=point[i];
41
        <u>if</u>(!vis[j]&&(dist[j]>val[i]||dist[j]==-1)){
42
          dist[j]=val[i];
43
           q.push(make_pair(dist[j],j));
44
45
46
    }
    printf("%d\n",ans);
47
48|}
```

2.17 树-灭绝树算法

```
1 int cntin[maxn];
2 struct node{int to,next;};
3 class graph{public:
4    node e[maxn]; int head[maxn], nume;
5    void init(int n=maxn-5){nume=0;fill_n(head,n+1,0);}
6    void add(int a, int b){e[++nume]={b,head[a]};head[a]=nume;}
7 }inv,nxt,dom;
8 class domtree{public://DAG
9    int deep[maxn],anc[maxn][maxp],que[maxn];
10    vector<int>edge;
11    void init(int n=maxn-5){
12    inv.init(n),nxt.init(n),dom.init(n);
```

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```
edge.clear();
    }
14
15
     void bit(int &x,int h){
       for(int i=0;h>0;++i){
16
17
         \underline{if}(h\&1) x=anc[x][i];
18
         h>>=1;
19
20
     }
21
     void add(int a,int b){nxt.add(a,b);inv.add(b,a);}
22
     int lca(int a,int b){
23
       if(deep[a] < deep[b]) swap(a,b);</pre>
24
       bit(a,deep[a]-deep[b]);
25
       \underline{if}(a==b) \underline{return} a;
26
       per(i,0,maxp-1) if(anc[a][i]!=anc[b][i])
27
         a=anc[a][i],b=anc[b][i];
28
       return anc[a][0];
29
30
     void tpsort(int n){
       int tp=0,ed=0;
31
32
       rep(i,1,n) {
         if(!cntin[i]) {
33
34
           que[ed++]=i;
35
           inv.add(0,i);
36
           nxt.add(i,0);
37
           edge.push_back(i);
38
       }
39
40
       while (ed!=tp){
         int now=que[tp++];
41
42
         forn(i,now,inv.head,inv.e){
43
           int to=inv.e[i].to;
44
           cntin[to]--
45
           <u>if</u>(!cntin[to]) que[ed++]=to,edge.push_back(to);
46
47
48
     }
     void maketree(int n){
49
50
       tpsort(n);
51
       for(auto i:edge){
52
         int fa=-1;
53
         forn(j,i,nxt.head,nxt.e){
54
           int to=nxt.e[j].to;
           if(fa==-1) fa=to;
55
56
           else fa=lca(fa,to);
         }fa=fa==-1?0:fa;
57
         deep[i] = deep[fa] + 1, anc[i][0] = fa;
58
59
         rep(j,1,maxp-1) anc[i][j]=anc[anc[i][j-1]][j-1];
60
         dom.add(fa,i);
61
    }
62
63
    int ans[maxn];
64
     int cal(int now){
65
       ans [now] = 1;
       forn(i,now,dom.head,dom.e) ans[now]+=cal(dom.e[i].to);
66
67
       return ans[now];
68
    }
69| }tree;
```

2.18 树-支配树算法

```
1 int root;
  <u>class</u> domtree{<u>public</u>://dom为最终的支配树,root为根,cnt为每个点的支配点的数量
    int dfn[maxn],rev[maxn],anc[maxn];
    int semi[maxn],idom[maxn];
    int fa[maxn],mi[maxn],clo;
6
    struct node{int to,next;};
    struct graph{
8
      node e[maxn]; int head[maxn], nume;
9
      void init(int n=maxn-5){nume=0;fill_n(head,n+1,0);}
      void add(int a,int b){e[++nume]={b,head[a]};head[a]=nume;}
10
11
    }inv,nxt,dom;
12
    void init(int n=maxn-5){
13
      clo=0;
```

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```
rep(i,1,n)fa[i]=mi[i]=semi[i]=i,rev[i]=dfn[i]=anc[i]=idom[i]=0;
15
      nxt.init(n),inv.init(n),dom.init(n);
16
    void add(int a,int b){inv.add(b,a),nxt.add(a,b);}
17
    int find(int now){
18
      if(fa[now] == now) return now;
19
20
      int fx=fa[now],y=find(fa[now]);
21
      if(dfn[semi[mi[fx]]]<dfn[semi[mi[now]]])</pre>
22
        mi[now]=mi[fx];
      return fa[now] = y;
23
    }
24
25
    void tdfs(int now){
26
      dfn[now]=++clo;rev[clo]=now;
27
      forn(i,now,nxt.head,nxt.e) if (!dfn[nxt.e[i].to])
28
        anc[nxt.e[i].to]=now,tdfs(nxt.e[i].to);
29
30
    void maketree(int root,int n=maxn-5){
31
      tdfs(root);
32
      per(i,2,n){
33
        int now=rev[i],tmp=n;
34
        forn(i,now,inv.head,inv.e){
35
          int to=inv.e[i].to;if(!dfn[to]) continue;
          if(dfn[to]<dfn[now]) tmp=min(tmp,dfn[to])</pre>
36
37
          else find(to),tmp=min(tmp,dfn[semi[mi[to]]]);
38
39
        semi[now] = rev[tmp]; fa[now] = anc[now];
40
        dom.add(semi[now],now);
        now=rev[i-1];
41
        forn(i,now,dom.head,dom.e){
42
          int to=dom.e[i].to;find(to);
43
          if(semi[mi[to]] == now) idom[to] = now;
44
45
          else idom[to]=mi[to];
46
47
      rep(i,2,n){
|48|
49
        int to=rev[i];
        if(idom[to]!=semi[to]) idom[to]=idom[idom[to]];
50
51
52
      dom.init(n);
53
      rep(i,1,n) if(i!=root)dom.add(idom[i],i);
54
55| }tree;
```

2.19 树-最小树形图

```
1 //定根最小树形图
  struct node{int a,b,c;}e[maxm];
   int in[maxn],pre[maxn],vis[maxn],id[maxn];
4
  11 mdst(){
5
    11 ans=0; int cnt=0,a,b,laz;
6
     \underline{\text{while}}(1){
7
       rep(i,1,n) in[i]=INF,id[i]=vis[i]=0;
       rep(i,1,m) if(e[i].a^e[i].b\&\&e[i].c < in[e[i].b])
8
         pre[e[i].b]=e[i].a,in[e[i].b]=e[i].c;
10
       in[k]=0;
11
       rep(i,1,n){
12
         if(in[i] == INF) return -1;
13
|14|
         <u>for</u>(a=i;a^k&&vis[a]^i&&!id[a];a=pre[a])vis[a]=i;
         if(a^k&&!id[a]){
15
16
           id[a]=++cnt;
17
           <u>for</u>(b=pre[a];a^b;b=pre[b])id[b]=cnt;
18
19
20
       if(!cnt) return ans;
       rep(i,1,n) <u>if</u>(!id[i]) id[i]=++cnt;
21
22
       rep(i,1,m) {
23
         laz=in[e[i].b];
24
         <u>if</u>((e[i].a=id[e[i].a])^(e[i].b=id[e[i].b]))
25
           e[i].c-=laz;
26
27
       n=cnt; k=id[k], cnt=0;
    }
28
```

```
29 }
30 int main() {IO;
31 cin>>n>m>>k;
32 rep(i,1,m)cin>>e[i].a>>e[i].b>>e[i].c;
33 cout<<mdst()<<endl;
}
}
```

2.20 树-RMQ求LCA+树上链交

```
1|//rmq求lca+快速求树上链交
  const int maxp=18;
3 <u>class</u> graph{public:
     struct node{int to,next;}e[maxn<<1];</pre>
     int head[maxn], nume, dfn[maxn], deep[maxn];
6
     int logn[maxn],pos[maxp][maxn],cnt;
     inline void add(int a, int b){e[++nume]={b,head[a]};head[a]=nume;}
8
     void init(int n){rep(i,1,n) head[i]=0;cnt=0,nume=1;}
9
     void cal(int n){
10
       logn[2]=1;
       rep(i,3,n) logn[i]=logn[i>>1]+1;
11
12
       <u>for(int</u> j=1;(1<<j)<=n;j++) <u>for(int</u> i=1;i+(1<<j)-1<=n;++i){
13
         <u>int</u> r=i+(1<<(j-1));
         <u>if</u>(deep[pos[j-1][i]]<deep[pos[j-1][r]]) pos[j][i]=pos[j-1][i];
14
15
         <u>else</u> pos[j][i]=pos[j-1][r];
16
     }
17
18
     void dfs(int now=root,int fa=root,int d=1){
19
       dfn[now]=++cnt;deep[now]=d;pos[0][cnt]=now;
20
       forn(i,now){
21
         if(e[i].to==fa) continue;
22
         dfs(e[i].to,now,d+1);pos[0][++cnt]=now;
23
     }
24
25
     inline int lca(int l,int r){
26
       l=dfn[l],r=dfn[r];<u>if</u>(l>r) swap(l,r);
27
       int lg=logn[r-l+1];
       <u>if</u>(deep[pos[lg][1]]<deep[pos[lg][r-(1<<lg)+1]])<u>return</u> pos[lg][1];
28
29
       else return pos[lg][r-(1<<lg)+1];</pre>
30
31
     inline int getdis(int a,int b){return deep[a]+deep[b]-2*deep[lca(a,b)];}
32
     void getlca(){dfs();cal(cnt);}
33
     <u>inline</u> <u>bool</u> check(<u>int</u> a, <u>int</u> b){return lca(a,b)==a;}
34
     int getans(int a1,int a2,int b1,int b2){
35
       int ra=lca(a1,a2);
36
       bool f1=check(ra,b1),f2=check(ra,b2);
37
       <u>if</u>(!f1&&!f2) <u>return</u> 0;
       <u>if</u>(f1&&f2){
38
39
         int rb=lca(b1,b2);
         if(!( check(rb,a1)||check(rb,a2)))return 0;
40
41
         <u>int</u> r1=lca(a1,b1),r2=lca(a1,b2);
42
         <u>int</u> r3=lca(a2,b1),r4=lca(a2,b2);
43
         <u>if</u>(r1==r3&&r2==r4) <u>return</u> 1;
         <u>return</u> getdis(r1==ra?r3:r1,r2==ra?r4:r2)+1;
44
45
46
       if(!f1)swap(b1,b2);
47
       int r1=lca(a1,b1),r3=lca(a2,b1);
48
       return getdis(r1==ra?r3:r1,ra)+1;
49
50 }g;
```

2.21 树-点分治

```
1 namespace graph{
    vector<int>g[maxn];
    int all,sz[maxn],root,maxt;
    bool vis[maxn];
5
    int dfs_root(int now,int fa){
6
      int cnt=1;
7
      for(auto to:g[now])if(to!=fa&&!vis[to])
8
        cnt+=dfs_root(to,now);
9
      int tmp=max(cnt-1,all-cnt);
10
      if(maxt>tmp) maxt=tmp,root=now;
      return sz[now] = cnt;
11
    }//基础部分
12
```

```
int ans[maxn];
    void dfs_col(int now,int fa,int c){
14
15
      ans[now]=c;
16
      for(auto to:g[now])if(to!=fa&&!vis[to])
17
        dfs_col(to,now,c);
18
19
    void dfs_dv(int now,int d=0){
20
      vis[now]=1;dfs_col(now,now,d);
      for(auto to:g[now]){
21
22
        if(vis[to]) continue;
23
        maxt=root=n+1;all=sz[to];
24
        dfs_root(to,now);dfs_dv(root,d+1);
25
    }
26
27
    void solve(int n){
28
      all=maxt=root=n+1;
      dfs_root(1,1);
29
30
      all-=maxt;
31
      dfs_dv(root);
32
33 }
```

2.22 树-轻重链剖分

```
1 class graph{//按边
     struct node{int from,to,cost,next;}e[maxn<<1];</pre>
     int head[maxn], nume, cnt2;
     <u>inline</u> <u>void</u> add(<u>int</u> a, <u>int</u> b, <u>int</u> c){
5
       e[++nume]={a,b,c,head[a]};head[a]=nume;
6
7
     int fa[maxn],sz[maxn],top[maxn],remp[maxn],ans[maxn];
8
     int son[maxn],in[maxn],cnt,deep[maxn];
     void dfs1(int now,int pre,int d){
10
       deep[now]=d;sz[now]=1;fa[now]=pre;
11
       for(int i=head[now];i;i=e[i].next){
12
         if(e[i].to==pre) continue;
13
         dfs1(e[i].to,now,d+1);
|14|
         sz[now] += sz[e[i].to];
15
         if(sz[son[now]] < sz[e[i].to]) son[now] = e[i].to;</pre>
16
     }
17
18
     void dfs2(int now,int pre,int st){
19
       top[now]=st;in[now]=++cnt;remp[cnt]=now;
20
       if(son[now]) dfs2(son[now],now,st);
21
       for(int i=head[now];i;i=e[i].next)
22
         <u>if</u>(e[i].to!=pre&&e[i].to!=son[now])
23
           dfs2(e[i].to,now,e[i].to);
24
25
     int query(int a,int b){
26
       int sum=0;
27
       while(top[a]!=top[b]){
         if(deep[top[a]] < deep[top[b]]) swap(a,b);</pre>
28
29
         sum+=tree.query(in[top[a]],in[a]);
30
        a=fa[top[a]];
31
32
       if(a==b)return sum;
33
       if(deep[a]>deep[b]) swap(a,b);
34
       sum+=tree.query(in[son[a]],in[b]);
35
       return sum;
36
37
     void getchain(){dfs1(1,1,0);dfs2(1,1,1);}
38| }g;
39 <u>int</u> root=1;
40 class graph{public://按点
     struct node{int to,next;}e[maxn<<1];</pre>
41
42
     int head[maxn], nume, mp[maxn];
43
     inline void add(int a,int b){
44
       e[++nume]={b,head[a]};
45
      head[a]=nume;
    }
46
47
     int ltop[maxn],fa[maxn],deep[maxn];
     int sz[maxn],remp[maxn];
48
49
     int son[maxn],cnt;
```

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```
void init(int n){rep(i,1,n) head[i]=0;cnt=0,nume=1;}
51
    void dfs1(int now=root,int pre=root,int d=0){
52
       deep[now]=d,fa[now]=pre,sz[now]=1,son[now]=0;
53
       forn(i,now){
        int to=e[i].to;
54
        if(to!=pre) {
55
56
          dfs1(to,now,d+1);
57
          sz[now] += sz[to];
58
           <u>if</u>(sz[to]>sz[son[now]]) son[now]=to;
59
      }
60
    }
61
62
    void dfs2(int now=root,int pre=root,int sp=root){
63
      ltop[now] = sp; mp[now] = ++cnt; remp[cnt] = now;
64
        if(son[now]) dfs2(son[now],now,sp);
65
        forn(i,now){
          int to=e[i].to;
66
67
           <u>if</u>(to!=son[now]&&to!=pre) dfs2(to,now,to);
68
69
70
    void getchain(){dfs1();dfs2();}
71
    int lca(int x,int y){
       for(;ltop[x]!=ltop[y];deep[ltop[x]]>deep[ltop[y]]?x=fa[ltop[x]]:y=fa[ltop[y]]);
72
73
       return deep[x] < deep[y] ?x:y;</pre>
74
    inline int getdis(int a,int b){return deep[a]+deep[b]-2*deep[lca(a,b)];}
75
76
    inline bool check(int a,int b){return dfn[a]<=dfn[b]&&dfn[a]+sz[a]-1>=dfn[b]+sz[b]-1;}
     //基础部分
77
78
    void update(int a,int b,int val){
      while(ltop[a]!=ltop[b]){
79
           \underline{if}(deep[ltop[a]] < deep[ltop[b]])swap(a,b);
80
          tree.update(mp[ltop[a]],mp[a],val);
81
82
          a=fa[ltop[a]];
83
        if (deep[a] > deep[b]) swap(a,b);
84
85
        tree.update(mp[a],mp[b],val);
86
87
    int query(int a,int b,int k){
88
       int sum=0;
89
       while(ltop[a]!=ltop[b]){
90
           <u>if</u>(deep[ltop[a]]<deep[ltop[b]])swap(a,b);
91
           sum+=tree.query(mp[ltop[a]],mp[a],k);
92
          a=fa[ltop[a]];
93
94
        if(deep[a]>deep[b])swap(a,b);
95
        sum+=tree.query(mp[a],mp[b],k);
96
        return sum;
97
98|}g;
```

2.23 杂项-欧拉路径

```
1 int vis[maxn];
 2 int cnt;
3 struct node {
4
    int to,flag,id,next;
   int head[maxn], nume, deg[maxn];
   <u>inline void</u> _add(<u>int</u> a,<u>int</u> b,<u>int</u> c){
     e[++nume]=(node){b,1,c,head[a]};
8
9
     head[a]=nume;
10 }
11 <u>inline void</u> add(<u>int</u> a, <u>int</u> b, <u>int</u> c){
     _add(a,b,c);_add(b,a,-c);
12
13|}
14 vector < int > ans [maxn];
15 void dfs(int now){
16
     vis[now]=1;
     for(int i=head[now];i;i=e[i].next){
17
       if(!e[i].flag) continue;
18
       e[i].flag=e[i^1].flag=0;
19
20
       dfs(e[i].to);
21
       ans[cnt].push_backe(-e[i].id);
```

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```
22
23 }
24 void solve(){
25
     rep(i,1,n){
26
       <u>if</u>(!vis[i]&&deg[i]&1) {
27
         cnt++;
         dfs(i);
28
29
     }
30
     rep(i,1,n){
31
       <u>if</u>(!vis[i]&&deg[i]){
32
33
         cnt++;
34
         dfs(i);
35
36
     }
37|}
```

2.24 杂项-三元环计数

```
1 vector<pii>g[maxn];
   int deg[maxn],a[maxn],b[maxn],cnt[maxn],pos[maxn],v[maxn];
   int main() {
     \underline{\text{while}}(\underline{\text{cin}}>>n>>m)
5
       rep(i,1,n){
6
         g[i].clear();
7
         v[i]=deg[i]=pos[i]=0;
8
       rep(i,1,m){
10
         cin>>a[i]>>b[i];
|11|
         deg[a[i]]++,deg[b[i]]++;
12
13
       rep(i,1,m){
14
         cnt[i]=0;
15
         <u>if</u>(deg[a[i]] < deg[b[i]])g[a[i]].emplace_back(b[i],i);
16
         else if(deg[a[i]]>deg[b[i]])g[b[i]].emplace_back(a[i],i);
17
         else {
18
           <u>if</u>(a[i] < b[i]) g[a[i]].emplace_back(b[i],i);
19
           else g[b[i]].emplace_back(a[i],i);
20
21
22
       rep(i,1,m){
23
         <u>int</u> u=a[i],to=b[i];
         for(auto j:g[u]) pos[j.fi]=j.se,v[j.fi]=i+1;
24
         for(auto j:g[to]){
  int t=j.fi;
25
26
           <u>if(</u>v[t]==i+1){
27
28
              cnt[i]++;
29
              cnt[pos[t]]++;
30
              cnt[j.se]++;
31
         }
32
33
34
       ll ans=0;
35
       rep(i,1,m) ans+=1ll*cnt[i]*(cnt[i]-1)/2;
36
       cout << ans << endl;</pre>
37
38|}
```

3 Geometry

3.1 几何类-点类与基础

```
\underline{\text{return}} \text{ cmp}(r1,12)!=-1 \&\& \text{ cmp}(r2,11)!=-1;
15 }
16 struct point {
    db x, y;
point(){}
17
18
     point(db k1, db k2){ x = k1, y = k2; }
19
     //向量加法、点+向量=点:*/
20
     point operator + (const point &k1) const { return point(x+k1.x, y+k1.y); }
21
     //向量减法、点-点=向量:*/
|22|
23
     point operator - (const point &k1) const { return point(x-k1.x, y-k1.y); }
     //向量数乘:*/
24
25
     point operator * (db k1) const { return (point){x*k1, y*k1}; }
26
     //向量数除:*/
27
     point operator / (db k1) const { return (point){x/k1, y/k1}; }
28
     //比较两个点(向量)是否相同:*/
29
     bool operator == (const point &k1) const {
30
       \underline{\text{return}} \text{ cmp}(x,k1.x) == 0 \&\& \text{ cmp}(y,k1.y) == 0;
31
32
     //逆时针旋转:*/
33
    point turn(db k1){
34
       \underline{\text{return}} \text{ (point)} \{x*\cos(k1)-y*\sin(k1), x*\sin(k1)+y*\cos(k1)\};
35
     //逆时针旋转90度:*/
36
    point turn90(){return (point){-y, x};}
37
     //比较两个点(向量)的大小:
38
     //x越小则点越小,若x相等,则y越小点越小.可以实现按点的坐标排序*/
39
40
     bool operator < (const point k1) const{</pre>
       int a = cmp(x, k1.x);
41
       \underline{if}(a == -1) \underline{return} 1;
42
43
       else if(a == 1) return 0;
44
       else return cmp(y,k1.y)==-1;
45
46
     //向量模长:
47
    db len(){ return sqrt(x*x+y*y); }
     //向量模长的平方:
48
    db len2(){ return x*x+y*y; }
49
     //单位向量:
50
     point unit(){ return (*this)/(*this).len(); }
51
     //向量的极角:
52
     db angle() { return atan2(y, x); }
53
     //将点放入第一象限:
54
55
     //当横坐标为负时,或横坐标为0纵坐标为负时,将点按原点做对称角度是[-/2,/2]
    point getdel(){
56
       \underline{if} (sign(x)==-1||(sign(x)==0&&sign(y)==-1)) \underline{return} (*\underline{this})*(-1);
57
       else return (*this);
58
59
60
     //判断点是否在1 2象限,或者在x的负半轴上角度是(0, π]
     \underline{bool} \text{ getp()} \underline{const} \{\underline{return} \text{ sign(y)==1 } || (sign(y)==0\&\&sign(x)==-1); \}
61
62
     void scan(){cin>>x>>y;}
     void print(){cout<<x<<' '<<y<<'\n'; }</pre>
63
64 };
65 //判断k1 在 [k2,k3] 内:
66 bool inmid(point k1, point k2, point k3){
    <u>return</u> inmid(k1.x,k2.x,k3.x) && inmid(k1.y,k2.y,k3.y);
67
68|}
69| //得到两点中点:
70 point midpo(point k1, point k2){ return (k1+k2)/2; }
71 //两点距离的平方
72 db dis2(point k1, point k2){
    \underline{\text{return}} \ (k1.x-k2.x)*(k1.x-k2.x) + (k1.y-k2.y)*(k1.y-k2.y);
73
74| \}
75 db dis(point k1, point k2){ return sqrt(dis2(k1, k2)); }
76 // 叉乘:
77 db cross(point k1, point k2) { return k1.x*k2.y - k1.y*k2.x; }
78| //点乘:
79 db dot(point k1, point k2) { \underline{return} k1.x*k2.x + k1.y*k2.y; }
80 //向量夹角:
81 db rad(point k1, point k2){
    return acos(dot(k1,k2)/k1.len()/k2.len());
83 //return atan2(cross(k1,k2), dot(k1,k2));
```

```
84| }
85| //极角排序,[-π, π]:
86| bool compareangle (point k1,point k2){
87| return k1.getp()<k2.getp() ||
88| (k1.getp()==k2.getp() && sign(cross(k1,k2))>0);
89| }
90| //k1 k2 k3 逆时针1 顺时针-1 否则0:
91| int clockwise(point k1,point k2,point k3){return sign(cross(k2-k1,k3-k1));}
```

```
几何类-直线类与线段类
 1| //直线与线段
2 //直线类
3 struct line{
 4 //方向为p[0]->p[1]
5
    point p[2];
 6
    line(){}
    line(db x1,db y1,db x2,db y2)\{p[0]=point(x1,y1),p[1]=point(x2,y2);\}
7
    line(point k1,point k2){p[0]=k1; p[1]=k2;}
    point& operator [] (int k){return p[k];}
10 //点在直线左侧的判定:
11| //沿着p0->p1的左侧为1,右侧为0
12
    bool include(point k){
13
      return sign(cross(p[0]-k,p[1]-k))>0;
|14|
  //方向向量:
15
    point dir(){return p[1]-p[0];}
16
17
  //向外(左)平移eps
18
    line push(){
19
      point delta=(p[1]-p[0]).turn90().unit()*eps;
20
      return {p[0]-delta, p[1]-delta};
21
22|\ \};
23 //线段类:
24 struct segment{
25
    point p[2];
26
    segment(){}
27
    segment(db x1,db y1,db x2,db y2){p[0]=point(x1,y1),p[1]=point(x2,y2);}
28
    segment(point a, point b) { p[0] = a, p[1] = b; }
29
    point dir(){return p[1]-p[0];}
30
    point& operator [] (int k){ return p[k]; }
31| };
32| //q 到直线 k1,k2 的投影:
33
  point proj(point q, point k1, point k2){
    point k=k2-k1;
34
35
    \underline{\text{return}} \text{ k1+k*(dot(q-k1,k)/k.len2())};
36|}
37 //q 关于直线 k1, k2 的对称点:
  point reflect(point q, point k1, point k2){
39
    return proj(q,k1,k2)*2-q;
40| \}
41 //点在线段上的判定:
42 bool checkons(point q,point k1,point k2){
43
    \underline{\text{return}} \text{ inmid}(q,k1,k2) && \text{sign}(\text{cross}(k1-q, k2-k1)) == 0;
44|}
45 //点在直线上的判定:
46 bool checkonl(point q,point k1,point k2){
47
    return sign(cross(k1-q, k2-k1))==0;
48| \}
49| //点在射线k1->k2上的判定:
50 bool checkonr(point q, point k1, point k2){
    \underline{\text{return}} \text{ sign}(\text{cross}(q-k1, k2-k1)) == 0 \&\& \text{ sign}(\text{dot}(q-k1, k2-k1)) >= 0;
511
52|}
53 //直线平行判定,可以重合:
54 bool parallel(line k1,line k2){ return sign(cross(k1.dir(),k2.dir()))==0; }
55| //直线同向判定:
56 bool samedir(line k1, line k2){
    return parallel(k1,k2)&&sign(dot(k1.dir(),k2.dir()))==1;
57
58| }
59| //直线的比较,极角排序,范围是[-π, π]:
60 bool operator < (line k1, line k2) {
```

```
if (samedir(k1,k2)) return k2.include(k1[0]);
|62|
     return compareangle(k1.dir(),k2.dir());
63 }
64 //直线相交判定:
65|//叉积计算面积,两直线不平行必相交(除去重合的情况),平行时,三角形面积相等:
66 bool checkll(point k1, point k2, point k3, point k4){
67
     return cmp(cross(k3-k1,k4-k1),cross(k3-k2,k4-k2))!=0;
68
69 //直线相交判定:
70 bool checkll(line k1,line k2){
     return checkll(k1[0],k1[1],k2[0],k2[1]);
71
72 }
73 //直线交点:
74 point getll(point k1, point k2, point k3, point k4) {
75
     db w1=cross(k1-k3,k4-k3), w2=cross(k4-k3,k2-k3);
76
     <u>return</u> (k1*w2+k2*w1)/(w1+w2);
77|}
78| //直线交点:
79 point getll(line k1, line k2){
80
    <u>return</u> getll(k1[0],k1[1],k2[0],k2[1]);
81|}
82| //直线与线段相交判定:
83| //线段的两端点在直线的两侧
84 bool checkls(point k1, point k2, point k3, point k4){
     \underline{\text{return}} \text{ sign}(\text{cross}(k1-k3, k2-k3)) * \text{sign}(\text{cross}(k1-k4, k2-k4)) <= 0;
85
86|}
87 // 线段相交判定:
88 bool checkss(point k1, point k2, point k3, point k4){
89
     return intersect(k1.x,k2.x,k3.x,k4.x)&&intersect(k1.y,k2.y,k3.y,k4.y) &&
90
           sign(cross(k3-k1,k4-k1))*sign(cross(k3-k2,k4-k2)) <= 0 \&\&
91
           sign(cross(k1-k3,k2-k3))*sign(cross(k1-k4,k2-k4)) <= 0;
92| \}
93 // 线段相交判定:
94 bool checkss(segment k1, segment k2){
95
     <u>return</u> checkss(k1[0], k1[1], k2[0], k2[1]);
96 }
97 //线段规范相交判定:
98| //端点相交不算
99 bool strictcheckss(point k1, point k2, point k3, point k4){
     <u>return</u> sign(cross(k3-k1,k4-k1))*sign(cross(k3-k2,k4-k2))<0 &&
101
           sign(cross(k1-k3,k2-k3))*sign(cross(k1-k4,k2-k4))<0;
102| \}
103| // 线段规范相交判定:
104 bool strictcheckss(segment k1, segment k2){
     return strictcheckss(k1[0], k1[1], k2[0], k2[1]);
105
106| \}
107 //点到直线的距离:
108 db displ(point q, point k1, point k2){
     if(k1 == k2) return dis(q, k1);
     return fabs(cross(k2-k1, q-k1)) / (k2-k1).len();
110
111|}
112| //点到直线的距离:
113 db displ(point q, line l){
     <u>return</u> displ(q, 1[0], 1[1]);
115| \}
116| //点到线段的距离:
117 db disps(point q,point k1,point k2){
118
    point k3 = proj(q,k1,k2);
119
     if (inmid(k3,k1,k2)) return dis(q, k3);
120|
     else return min(dis(q, k1),dis(q, k2));
121|}
122| //点到线段的距离:
123 db disps(point q, segment k1){
124
     <u>return</u> disps(q, k1[0], k1[0]);
125| \}
126| //线段到线段间的距离:
127 db disss(point k1, point k2, point k3, point k4) {
128
     if (checkss(k1,k2,k3,k4)) return 0;
129
     else return min(min(disps(k3,k1,k2),disps(k4,k1,k2)),
130
                    min(disps(k1,k3,k4),disps(k2,k3,k4)));
```

```
131|}
132| //线段到线段间的距离:
133 db disss(segment k1, segment k2){
    <u>return</u> disss(k1[0], k1[1], k2[0], k2[1]);
134
135|}
        几何类-圆类
  3.3
 1 // 圆类
```

```
2 struct circle{
3
    point o; db r;
4
    circle(){}
    circle(point _o, db _r){ o = _o, r = _r; }
 6
  //点在圆内判足:
7
    bool include(point k){ return cmp(dis(o, k), r) <= 0; }</pre>
8|};
9|//求直线与圆的交点沿着 k2->k3 方向给出
                                          , 相切给出两个:
10 <u>vector</u><point> getcl(circle k1,point k2,point k3){
    point k=proj(k1.o,k2,k3);
11
    db d=k1.r*k1.r-(k-k1.o).len2();
12
    if (sign(d)==-1) return {};
13
    point del=(k3-k2).unit()*sqrt(max((db)0.0, d));
14
    return {k-del, k+del};
15
16|}
17 // 返回两个圆的公切线数量:
18
  int checkposcc(circle k1,circle k2){
    <u>if</u> (cmp(k1.r,k2.r)==-1) swap(k1,k2);
19
    db d=dis(k1.o,k2.o); <u>int</u> w1=cmp(d,k1.r+k2.r), w2=cmp(d,k1.r-k2.r);
20
21
    <u>if</u> (w1>0) <u>return</u> 4; //相离:
    <u>else if</u> (w1==0) <u>return</u> 3; //相切:
22
    <u>else if</u> (w2>0) <u>return</u> 2; //相交:
23
24
    <u>else if (w2==0) return 1; //内切:</u>
    else return 0; //内含:
|26| }
27 //求两圆交点沿圆 k1 逆时针给出 , 相切给出两个:
28 vector <point > getcc(circle k1, circle k2){
    int pd=checkposcc(k1,k2);
30
    if(pd==0||pd==4) return {};
31
    db a=(k2.o-k1.o).len2();
    db cosA=(k1.r*k1.r+a-k2.r*k2.r)/(2*k1.r*sqrt(max(a,(db)0.0)));
32
33
    db b=k1.r*cosA;
34
    db c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
35
    point k=(k2.o-k1.o).unit(), m=k1.o+k*b, del=k.turn90()*c;
36
    return {m-del, m+del};
37|}
38| //过圆外一点作圆的切线的切点:
39 //沿圆 k1 逆时针给出
40 <u>vector</u><point> tangentcp(circle k1,point k2){
    db a=(k2-k1.o).len(),b=k1.r*k1.r/a,c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
41
42
    point k=(k2-k1.0).unit(),m=k1.0+k*b,del=k.turn90()*c;
43
    return {m-del, m+del};
44|}
45| //求两圆的外切线:
46
  vector<line> tangentoutcc(circle k1,circle k2){
47
    int pd=checkposcc(k1,k2); if (pd==0) return {};
  //内含,返回一条切线
48
49
    <u>if</u> (pd==1){
      point p1=getcc(k1,k2)[0]; point p2=p1+((p1-k1.o).turn90()/(p1-k1.o).len());
50
51
      <u>return</u> {(line){p1,p2}};
52
53
    if (cmp(k1.r,k2.r)==0){
      point del=(k2.o-k1.o).unit().turn90().getdel();
54
55
      <u>return</u> {(line){k1.o-del*k1.r,k2.o-del*k2.r},(line){k1.o+del*k1.r,k2.o+del*k2.r}};
56
    } else {
57
      point p=(k2.0*k1.r-k1.0*k2.r)/(k1.r-k2.r);
      vector<point>A=tangentcp(k1,p),B=tangentcp(k2,p);
58
59
      vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});</pre>
60
      <u>return</u> ans;
61
62|}
63 //求两圆的内切线:
64 vector < line > tangentincc (circle k1, circle k2) {
```

```
int pd=checkposcc(k1,k2); if (pd<=2) return {};</pre>
66
     if (pd==3){
67
       point p1=getcc(k1,k2)[0];    point p2=p1+((p1-k1.o).turn90()/(p1-k1.o).len());
68
       <u>return</u> {(line){p1,p2}};
69
70
     point p=(k2.0*k1.r+k1.0*k2.r)/(k1.r+k2.r);
71
     vector<point>A=tangentcp(k1,p),B=tangentcp(k2,p);
72
     vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});</pre>
73
     return ans;
74|}
75| //求两圆所有切线:
76 vector < line > tangentcc (circle k1, circle k2) {
     <u>int</u> flag=0; <u>if</u> (k1.r<k2.r) swap(k1,k2),flag=1;
     vector<line>A=tangentoutcc(k1,k2),B=tangentincc(k1,k2);
78
79
     for (line k:B) A.push_back(k);
80
     if (flag) for (line &k:A) swap(k[0],k[1]);
81
82|}
83| // 圆 k1 与三角形 k2 k3 k1.o 的有向面积交:
84 db circleinsarea(circle k1,point k2,point k3){
    point k=k1.o; k1.o=k1.o-k; k2=k2-k; k3=k3-k;
851
     int pd1=k1.include(k2),pd2=k1.include(k3);
861
87
     vector<point>A=getcl(k1,k2,k3);
88| //有一个点在圆内或圆上:
89
     if (pd1>=0){
90| //三角形整个落在圆内,返回三角形的有向面积
91|
       if (pd2>=0) return cross(k2,k3)/2;
92 //三角形的一个点落在圆内,一个点落在圆外
       else return k1.r*k1.r*rad(A[1],k3)/2+cross(k2,A[1])/2;
93
94 //三角形的一个点落在圆内,一个点落在圆外
     } else if (pd2>=0){
95
       return k1.r*k1.r*rad(k2,A[0])/2+cross(A[0],k3)/2;
96
97 //否则,三角形的两个点都落在圆外:
     }else {
98
99
       int pd=cmp(k1.r,disps(k1.o,k2,k3));
100| //返回一个扇形面积:
       <u>if</u> (pd<=0) <u>return</u> k1.r*k1.r*rad(k2,k3)/2;
101|
102 //返回两个扇形加一个三角形:
103
       else return cross(A[0],A[1])/2+k1.r*k1.r*(rad(k2,A[0])+rad(A[1],k3))/2;
104
105|}
106| //以k1k2为直径的圆:
107 circle getexcir2(point k1, point k2){
108
     point c = midpo(k1, k2);
     db r = dis(c, k1);
109
110|
     return circle(c, r);
111|}
112| //三角形的外接圆:
113 circle getexcir3(point k1, point k2, point k3){
     point c = getll( midpo(k1, k2)
114
115
                       (k1-k2).turn90()+midpo(k1, k2),
116
                      midpo(k1, k3),
117|
                       (k1-k3).turn90()+midpo(k1, k3));
118
     db r = dis(c, k1);
119
     return circle(c, r);
120|}
121 //最小圆覆盖:
122 circle mincircover(vector <point>A){
     random_shuffle(A.begin(), A.end()); int n = A.size();
123
     circle now = circle(A[0], 0);
124
     \underline{\text{for}}(\underline{\text{int}} \ i = 0; \ i < n; \ i++)\underline{\text{if}}(!\text{now.include}(A[i]))
125|
126
       now = circle(A[i], 0);
127
       \underline{\text{for}}(\underline{\text{int}} \ j = 0; \ j < i; \ j++)\underline{\text{if}}(!\text{now.include}(A[j]))\{
         now = getexcir2(A[i], A[j]);
128
129
         \underline{\text{for}}(\underline{\text{int}} \ k = 0; \ k < j; \ k++) \ \underline{\text{if}}(!\text{now.include}(A[k]))
130
           now = getexcir3(A[i], A[j], A[k]);
131
     }
132
133
     return now;
134|}
```

3.4 几何函数-多边形与半平面

```
1 //多边形函数
 2 //三角形面积:
 3 db tarea(point a, point b, point c) {
     \underline{\text{return}} \text{ fabs}((b.x-a.x)*(c.y-a.y)-(b.y-a.y)*(c.x-a.x))/2;
 5 }
 6 //多边形面积:
                                        , 逆时针
 7 //多边形用 vector<point> 表示
 8 db polyarea(<u>vector</u><point>A){
     db ans = 0;
10
     sort(all(A),compareangle);
     \underline{\text{for}}(\underline{\text{int}} \ \text{i=0;i<A.size();i++}) \ \text{ans += cross(A[i],A[(i+1)%A.size()]);}
11
12
     return fabs(ans/2);
13|}
14 //多边形周长:
15 db polyperimeter(<u>vector</u><point>&A){
     db ans = 0;
16
     \underline{\text{for}}(\underline{\text{int}} \ i = 0; \ i < A.size(); \ i++) \ \text{ans} += dis(A[i], A[(i+1)\%A.size()]);
17|
18
     return ans;
19|}
20 //多边形重心:
21 point polyfocus(vector<point>&A){
22
     \underline{int} n = A.size();
23
     db sumx= 0, sumy = 0, sumarea = 0, area;
     for(int i = 1; i+1 < n; i++){
25
       area = cross(A[i]-A[0], A[i+1]-A[0])/2.0;
26
       sumarea += area;
       sumx += (A[0].x+A[i].x+A[i+1].x)*area;
27
28
       sumy += (A[0].y+A[i].y+A[i+1].y)*area;
29
     return point(sumx/sumarea/3.0, sumy/sumarea/3.0);
30
31|}
32| //点与多边形的位置关系:
33| // 2 内部 1 边界 0 外部
34 int contain(vector<point>&A, point q){
35
     int pd=0; A.push_back(A[0])
36
     for (int i=1;i<A.size();i++){</pre>
37
       point u=A[i-1], v=A[i];
38
       if (checkons(q,u,v)) return 1;
39
       <u>if</u> (cmp(u.y,v.y)>0) swap(u,v);
40
       \underline{if} (cmp(u.y,q.y)>=0||cmp(v.y,q.y)<0) \underline{continue};
41
       if (sign(cross(u-v,q-v))<0) pd^=1;</pre>
42
43
     return pd<<1;</pre>
     \underline{int} wn = 0;
44
     \underline{int} n = A.size();
45
     for(int i = 0; i < n; i++){
46
       if(checkons(q, A[i], A[(i+1)%n])); return -1;//onside
47
       \underline{int} k = sign(cross(A[(i+1)\%n]-A[i], q-A[i]));
48
       int d1 = sign(A[i].y-q.y);
49
       \underline{int} d2 = sign(A[(i+1)\sqrt{n}].y-q.y);
50
51
       if(k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
       if(k < 0 \&\& d2 <= 0 \&\& d1 > 0) wn--;
52
53
     if(wn != 0) return 1;//inside
55
     return 0;//outside
56|}
57 //逆时针凸包判定:
58 int checkconvex(vector<point>&A){
     <u>int</u> n=A.size(); A.pb(A[0]); A.pb(A[1]);
     \underline{\text{for }}(\underline{\text{int }}i=0;i\leq n;i++)\underline{\text{if }}(\operatorname{sign}(\operatorname{cross}(A[i+1]-A[i],A[i+2]-A[i]))=-1)\underline{\text{return }}0;
60
61
     return 1;
62 }
63| //求凸包:
64 //flag=0 不严格 flag=1 严格
65 vector <point > convexhull(vector <point > A, int flag=1){
     int n=A.size(); vector<point>ans(n*2);
67
     sort(A.begin(), A.end());
68
     int now=-1;
69 // 下凸壳
    <u>for(int</u> i=0;i<n;i++){
```

```
while(now>0 && sign(cross(ans[now]-ans[now-1], A[i]-ans[now-1])) < flag) now--;</pre>
72
       ans[++now]=A[i];
73
74
     int pre=now;
   //上凸壳
75
76
     for(int i=n-2;i>=0;i--){
       while(now>pre && sign(cross(ans[now]-ans[now-1], A[i]-ans[now-1])) < flag) now--;</pre>
77
       ans[++now]=A[i];
78
79
   //因为A[0]会被算两次,所以舍弃最后一次的A[0]
80
81
     ans.resize(now);
82
     <u>return</u> ans;
83|}
84 //切割凸包:
85 //保留直线左边的所有点
86 vector<point> convexcut(vector<point>A,point k1,point k2){
87
     int n=A.size(); A.push_back(A[0]); vector<point>ans;
88
     <u>for(int</u> i=0;i<n;i++){
       int w1=clockwise(k1,k2,A[i]), w2=clockwise(k1,k2,A[i+1]);
89
90
       if (w1>=0) ans.push_back(A[i]);
91
       \underline{if} (w1*w2<0) ans.push_back(getll(k1,k2,A[i],A[i+1]));
92
93
     return ans;
94 }
95 //凸包最近点对:
96 // 先要按照 x 坐标排序
97 bool _cmp(point k1,point k2){return k1.y<k2.y;}
98 db closestpoint(vector<point>&A, int 1, int r){
     if (r-1 <= 5){
99
100 //当点数小于等于5时,暴力计算:
101
       db ans=1e20;
102
       <u>for</u> (<u>int</u> i=1;i<=r;i++) <u>for</u> (<u>int</u> j=i+1;j<=r;j++) ans=min(ans,dis(A[i],A[j]));
103
       <u>return</u> ans;
104
105
     int mid=l+r>>1; db ans=min(closestpoint(A,1,mid),closestpoint(A,mid+1,r));
106
     vector<point>B; for (int i=1;i<=r;i++) if (abs(A[i].x-A[mid].x)<=ans) B.push_back(A[i]);</pre>
     sort(B.begin(),B.end(),_cmp);
107
108
     for (int i=0;i<B.size();i++) for (int j=i+1;j<B.size()&&B[j].y-B[i].y<ans;j++)
109
       ans=min(ans,dis(B[i],B[j]));
110
     return ans;
111|}
112 //凸包的直径(最远点对):
113| //旋转卡壳,得到的答案为最远距离的平方
114 db convexdiameter(vector <point>&A){
1151
     int n = A.size();
116
     int now = 1;
     db res = 0;
117
     for(int i = 0; i < n; i++){
118
119
       \underline{\text{while}}(1)
120
         db x=cross(A[i]-A[(i+1)%n],A[i]-A[(now+1)%n]);
121
         db y=cross(A[i]-A[(i+1)%n],A[i]-A[now]);
122
         if(x<y) break;</pre>
123
         now=(now+1)%n;
124
125
       res = max(res, dis2(A[now], A[i]));
     }
126
127
     return res;
128|}
129 //点集中的最大三角形:
130 db maxtriangle(<u>vector</u><point>&A){
131
     \underline{int} m = A.size();
     \underline{int} a = 1, b = 2;
132
     db res = 0;
133|
134
     for(int i = 0; i < m; i++){}
       \frac{\text{while}}{\text{(cross(A[a]-A[i], A[(b+1)\%m]-A[i])}} > \text{cross(A[a]-A[i], A[b]-A[i])})
135
136
         b = (b + 1) \% m;
137
       res = \max(\text{res}, \text{cross}(A[a]-A[i], A[b]-A[i]) / 2.0);
138
       \frac{\text{while}(\text{cross}(A[(a+1)\%m]-A[i], A[b]-A[i])}{\text{cross}(A[a]-A[i], A[b]-A[i])}
139|
         a = (a + 1) \% m;
140
       res = max(res, cross(A[a]-A[i], A[b]-A[i]) / 2.0);
141
```

```
142
     return res;
143|}
144 //凸包间的最小距离:
145 db mindisbetconvex(<u>vector</u><point>&A, <u>vector</u><point>&B){
|146|
      \underline{int} n = A.size(), m = B.size();
|147|
      if(n < 3 \&\& m < 3){
        \underline{if}(n == 1){
|148|
          <u>if</u>(m == 1) <u>return</u> dis(A[0], B[0]);
149
150
          else return disps(A[0], B[0], B[1]);
151
        else{
152
153
          \underline{if}(m == 1) \underline{return} disps(B[0], A[0], A[1]);
154
          else return disss(A[0], A[1], B[0], B[1]);
155
      }
156
157
      int ai = 0, bi = 0;
      \underline{for}(\underline{int} \ i = 0; \ i < n; \ i++) \ \underline{if}(A[i].y < A[ai].y) \{ ai = i; \}
158
159
      <u>for(int</u> i = 0; i < m; i++) <u>if</u>(B[i].y > A[bi].y){ bi = i; }
160
      db ans = 1e18;
      for(int i = 0; i < n; i++){
161
162
        db ck;
        \frac{\text{while}(\text{ck} = \text{sign}(\text{cross}(B[(\text{bi}+1)\%m]-B[\text{bi}], A[(\text{ai}+1)\%n]-A[\text{ai}])) < 0) \text{ bi } = (\text{bi}+1)\%m;}
163
164
        if(ck == 0) ans = min(ans, disss(A[(ai+1)%n], A[ai], B[(bi+1)%m], B[bi]));
165
        <u>else</u> ans = min(ans, disps(B[bi], A[(ai+1)\%n], A[ai]));
166
        ai = (ai+1)%n;
167
168
      <u>return</u> ans;
169 }
170| //最小正方形覆盖:
171 db minsquarecover(<u>vector</u><point>&A, db rad){
      db minx = inf, maxx = -inf, miny = inf, maxy = -inf;
172
      for(int i = 0; i < A.size(); i++){</pre>
173
174
        point p = A[i].turn(rad);
        minx = min(minx, p.x);
175
        miny = min(miny, p.y);
176
177
        maxx = max(maxx, p.x);
        maxy = max(maxy, p.y);
178
179
      return max(maxx-minx, maxy-miny);
180
181|}
182| //三分--最小正方形覆盖:
183 db t_divide(<u>vector</u><point>&A, db 1, db r){
184
      db m,rm,eps=1e-8;
185
      \underline{\text{while}}(r-1) = \exp s
186
        m=1+(r-1)/3;
187
        rm=r-(r-1)/3;
188
        if(minsquarecover(A,m)>minsquarecover(A,rm)) l=m;
189
        else r=rm;
190
191
      return minsquarecover(A, (m+rm)/2);
192|}
193| //求半平面交:
194 //半平面是逆时针方向 , 输出按照逆时针
195 vector<point> gethalf(vector<line> L){
196
      \underline{int} n = L.size();
      sort(L.begin(), L.end());
int first = 0, last = 0;
197
198
199| //双端队列指针
200 line *q = \underline{\text{new}} line[n];
201| //双端队列
202 point *p = \underline{\text{new}} point[n];
203| //p[i] 为1[i] 和1[i+1] 的交点
204 q[last] = L[0];
205| //初始化为一个半平面
206
      for(int i = 0; i < n; i++){
207
        while(first < last && !L[i].include(p[last-1])) last--;</pre>
208
        while(first < last && !L[i].include(p[first])) first++;</pre>
        q[++last] = L[i];
209
210
        if(samedir(q[last], q[last-1])) last--;
211
        \underline{if}(first < last) p[last-1] = getll(q[last], q[last-1]);
212
```

```
213
    while(first < last && !q[first].include(p[last-1])) last--;</pre>
214
    vector<point>ans;
    if(last - first <= 1) return ans;</pre>
215
216
    p[last] = getll(q[last], q[first]);
    for(int i = first; i <= last; i++) ans.pb(p[i]);</pre>
217
218
    return ans;
219 }
220 int checkpos(line k1,line k2,line k3) {return k3.include(getll(k1,k2));}
221| //求半平面交:
222| //半平面是逆时针方向 , 输出按照逆时针
223 vector<line> gethl(vector<line> L){
    sort(L.begin(),L.end()); deque<line> q;
224
    for (int i=0;i<(int)L.size();i++){</pre>
225
226
      if (i&&samedir(L[i],L[i-1])) continue;
      \underline{\text{while}} \ (\text{q.size()} > 1 \& \& ! \text{checkpos}(\text{q[q.size()} - 2], \text{q[q.size()} - 1], \text{L[i])}) \ \text{q.pop\_back()};
227
228
      while (q.size()>1&&!checkpos(q[1],q[0],L[i])) q.pop_front();
229
      q.push_back(L[i]);
230
231
    while (q.size()>2&\&!checkpos(q[q.size()-2],q[q.size()-1],q[0])) q.pop_back();
232|
    while (q.size()>2\&\&!checkpos(q[1],q[0],q[q.size()-1])) q.pop_front();
    vector<line>ans; for (int i=0;i<q.size();i++) ans.push_back(q[i]);</pre>
233
234
    <u>return</u> ans;
235|}
       几何函数-圆的反演
  3.5
 1 /*
     反演的概念
 3|设在平面内给定一点O和常数k(k不等于零),对于平面内任意一点A,
 4] 确定A',使A'为直线OA上一点,并且有向线段OA与OA'满足OAOA2 = k,我们称这种变换是以O为反演中心,
 5 以k为反演幂的反演变换,简称反演.称A′为A关于O(r)的互为反演点.
  二、作已知点的反演点的方法
  给出反演极O和反演幂k>0,作点A的反演点A'.
 8 令k = r^2,作出反演基圆\odot0(r),
 9| 1)若点A在⊙O(r)外,则过点A作圆的切线(两条),两个切点相连与OA连线交点就是点A′.
10 2) 若点A在⊙0(r)内,则把上述过程逆过来:
| 11 | 连结OA,过点A作直线垂直于OA,直线与OO(r)的交点处的切线的交点就是点A'.
12|3)若点A在⊙O(r)上,反演点A′就是点A自身.
13 4)0没有反演点
14| 三、圆的反演变换
15| 圆在不同情形下的反演成像:
      当圆不经过反演中心,它的反演图形仍旧是个不过反演中心的圆,并且反演中心为这两个互为反形的圆的位似中心:
16 1.
      当圆与反演圆相交,交点是保持不变的;
17|2.
      当圆在反演圆的外面的时候,反演成像位于圆的内部;反之,当圆位于反演圆的内部,反演成像位于圆的外部.
18 3.
      当圆经过反演中心,它的反演图形是一条直线.
19|4.
      反之,任意一条不过反演中心的直线,其反演成像是一个经过反演中心的圆.
20
      相切两圆反向任相切,且切点不变,若切点是反演中心,
21 5.
      则其反象是两条平行直线;两圆相切,若反演中心在某圆上,则为反形为相切的直线与圆;
22
|23| */
24| //c1关于c0的反演圆:
25 circle getinvertcir(circle c1, circle c0){
26
    circle c2;
27
    db x0 = c0.o.x, y0 = c0.o.y, r0 = c0.r,
28
    db x1 = c1.o.x, y1 = c1.o.y, r1 = c1.r;
29
    db \ d01 = dis(c0.o, c1.o);
    c2.r = 0.5*((1/(d01-r1))-(1/(d01+r1)))*r0*r0;
30
31
    db d02 = r0*r0/(d01+r1)+c2.r;
32
    //db _d02 = r0*r0/(d01-r1)-c2.r;
33
    c2.o.x = x0 + d02/d01*(x1-x0);
34
    c2.o.y = y0 + d02/d01*(y1-y0);
35
    return c2;
36|}
37| //直线k1关于c0的反演圆:
38 circle getinvertcir(line k1, circle c0){
39
    point a = proj(c0.0, k1[0], k1[1]);
    db oa = dis(c0.o, a);
|40|
    db ob = c0.r*c0.r/oa;
41
    point v = a-c0.o; v = v/v.len();
42
    point b = c0.o+v*ob;
43
44
    circle res;
45
    res.o = midpo(c0.o, b);
```

```
res.r = ob/2;
47
    return res;
48 }
  //c1关于c0的反演直线:
49
50 line getinvertline(circle c1, circle c0){
51
    point v = c1.0 - c0.0;
    v = v / v.len();
52
    db d = c0.r*c0.r / (2*c1.r);
53
54
    point k1 = v * d + c0.o;
55
    v = v.turn90();
56
    point k2 = k1 + v;
57
    return line(k1, k2);
58 }
59| //p关于c的反演点:
60 point getinvertpoint(point p, circle c){
    point v = (p-c.o).unit();
62|
    db len = c.r*c.r/dis(c.o, p);
63
    return c.o+v*len;
64|}
```

几何函数-圆上整点

```
1 struct point{
     11 x,y;//两圆上整数点对
     point(11 _x=0,11 _y=0){x=_x,y=_y;}
     void print2(){printf("%lld %lld\n",x,y);}
5
     void print1(){printf("%lld %lld ",x,y);}
6
     bool operator == (const point&other) const{
       return x==other.x&&y==other.y;
8
9
     bool operator<(const point&other)const{</pre>
10
       if(x==other.x) return y<other.y;</pre>
11
       return x<other.x;</pre>
12
     }
13 };
14 ll dis(point a, point b){
     \underline{\text{return}} (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
15
|16| }
17 bool check(ll n){ll x=sqrt(n);return x*x==n;}
18 <u>void</u> solve(ll rr,ll r,ll d,<u>vector</u><point>&A){
19
     for (ll i=1;i*i<=rr;++i){</pre>
20
       11 t=rr-i*i;
21
       if(!check(t))continue;
22
       11 j=sqrt(t);
23
       if(i>=j)break;
       \underline{if}(\underline{-gcd(i*i,t)}==1){
24
25
         11 x=i*j*d;
26
         11 y=sqrt(r*r-x*x);
         A.pb(point(x,y));
27
28
         A.pb(point(-x,y));
29
         A.pb(point(x,-y));
30
         A.pb(point(-x,-y));
31
32
     }
33 }
34 <u>void</u> deal(<u>vector</u><point>&A,ll r){
     A.pb(point(0,r));
35
     A.pb(point(0,-r));
36
37
     A.pb(point(-r,0));
38
     A.pb(point(r,0));
39
     r<<=1;
     for (ll d=1;d*d<=r;d++){</pre>
40
41
       if(r%d!=0) continue;
42
       solve(r/d,r/2,d,A);
       if(d*d==r) break;
43
44
       solve(d,r/2,r/d,A);
45
46|}
47 <u>vector</u><point>veca, vecb;
48 struct node{
|49|
     point A,B;
     node(point _A,point_B){A=_A;B=_B;}
50
51
     bool operator<(const node&other)const{</pre>
```

```
if(A==other.A)return B<other.B;</pre>
53
       else return A<other.A;</pre>
    }
54
55 };
56
  vector<node>ans;
  int main(){
57
58
    int T;ll a,b,c;
59
     scanf("%d",&T);
     while (T--){
60
61
       veca.clear();vecb.clear();ans.clear();
       scanf("%11d%11d%11d",&a,&b,&c);
62
63
       deal(veca,a);deal(vecb,b);
64
       for(int i=0;i<veca.size();i++)</pre>
65
         for(int j=0;j<vecb.size();j++)</pre>
66
           <u>if</u>(dis(veca[i],vecb[j])==c*c)
67
             ans.pb(node(veca[i],vecb[j]));
68
       sort(ans.begin(),ans.end());
69
       printf("%d\n",ans.size());
70
       for(int i=0;i<ans.size();++i){</pre>
71
         ans[i].A.print1();ans[i].B.print2();
72
73
74 }
```

3.7 几何函数-多边形面积并

```
多边形的交, 多边形的边一定是要按逆时针方向给出
   * 还要判断是凸包还是凹包,调用相应的函数
    * 面积并,只要和面积减去交即可
   */
 6
   #include <bits/stdc++.h>
   using namespace std;
   \underline{\text{const}} \underline{\text{int}} \max = 300;
   const double eps = 1e-8;
10 int dcmp(double x){
       \underline{if}(x > eps) \underline{return} 1;
11
       <u>return</u> x < -eps ? -1 : 0;
12
13|}
14 struct Point{
15
       double x, y;
16|};
17 double cross(Point a, Point b, Point c) ///叉积
18 {
19
       \underline{\text{return}} (a.x-c.x)*(b.y-c.y)-(b.x-c.x)*(a.y-c.y);
20 }
21 Point intersection(Point a, Point b, Point c, Point d) {
22
       Point p = a;
       \frac{\text{double}}{\text{double}} \text{ t = ((a.x-c.x)*(c.y-d.y)-(a.y-c.y)*(c.x-d.x))/((a.x-b.x)*(c.y-d.y)-(a.y-b.y)*(c.x-d.x));}
23
       p.x +=(b.x-a.x)*t;
24
25
       p.y +=(b.y-a.y)*t;
26
       return p;
27| \}
28 //计算多边形面积
29 double PolygonArea(Point p[], int n){
30
       \underline{if}(n < 3) \underline{return} 0.0;
31
       <u>double</u> s = p[0].y * (p[n - 1].x - p[1].x);
       p[n] = p[0];
32
       for(int i = 1; i < n; ++ i)
33
            s += p[i].y * (p[i - 1].x - p[i + 1].x);
34
35
       \underline{\text{return}} fabs(s * 0.5);
36|}
37
   double CPIA(Point a[], Point b[], int na, int nb)//ConvexPolygonIntersectArea
38| {
39
       Point p[20], tmp[20];
       int tn, sflag, eflag;
40
41
       a[na] = a[0], b[nb] = b[0];
42
       memcpy(p,b, \frac{sizeof}{(Point)*(nb + 1));
43
       \underline{for}(\underline{int} \ i = 0; \ i < na \&\& \ nb > 2; \ i++){}
            sflag = dcmp(cross(a[i + 1], p[0],a[i]));
44
            \underline{\text{for}}(\underline{\text{int}} \ j = \text{tn} = 0; \ j < \text{nb}; \ j++, \ \text{sflag} = \text{eflag})\{
45
                if(sflag>=0) tmp[tn++] = p[j];
46
47
                eflag = dcmp(cross(a[i + 1], p[j + 1],a[i]));
```

```
\underline{if}((sflag \cdot eflag) == -2)
49
                   tmp[tn++] = intersection(a[i], a[i + 1], p[j], p[j + 1]); //求交点
50
          memcpy(p, tmp, sizeof(Point) * tn);
51
52
          nb = tn, p[nb] = p[0];
53
       if(nb < 3) return 0.0;
54
55
       return PolygonArea(p, nb);
56|}
  double SPIA(Point a[], Point b[], int na, int nb)//SimplePolygonIntersectArea 调用此函数
57
58
59
       <u>int</u> i, j;
       Point t1[4], t2[4];
60
61
       double res = 0, num1, num2;
       a[na] = t1[0] = a[0], b[nb] = t2[0] = b[0];
62
       for(i = 2; i < na; i++){
63
          t1[1] = a[i-1], t1[2] = a[i];
64
65
          num1 = dcmp(cross(t1[1], t1[2], t1[0]));
66
           if(num1 < 0) swap(t1[1], t1[2]);
67
           for(j = 2; j < nb; j++)
68
           {
69
              t2[1] = b[j - 1], t2[2] = b[j];
70
              num2 = dcmp(cross(t2[1], t2[2], t2[0]));
71
               if(num2 < 0) swap(t2[1], t2[2]);
72
              res += CPIA(t1, t2, 3, 3) * num1 * num2;
73
74
75
       <u>return</u> res;
76
77 Point p1 [maxn], p2 [maxn];
78
  <u>int</u> n1, n2;
  int main(){
79
       freopen("in.txt", "r", stdin);
80
81
       while(cin>>n1>>n2){
           for(int i = 0; i < n1; i++) scanf("%lf%lf", &p1[i].x, &p1[i].y);</pre>
82
           \overline{\text{for}(int \ i = 0; \ i < n2; \ i++) \ scanf("%lf%lf", \&p2[i].x, \&p2[i].y);}
83
           double Area = SPIA(p1, p2, n1, n2);
84
85
           cout << Area << endl;</pre>
86
87
       return 0;
88|}
```

3.8 几何函数-辛普森积分

```
1 /* 自适应辛普森积分*/
2 db f(db x){
      /*积分表达式,或积分微元*/
3
4 }
5| db simpson(db 1,db r){return (r-1)*(f(1)+f(r)+4*f((1+r)/2))/6;}/*辛普森积分公式*/
6 db asr(db 1,db r,db eps,db s){
       double mid=(1+r)/2, ls=simpson(1,mid),rs=simpson(mid,r);
       \underline{if}(fabs(ls+rs-s) \le eps*15){
9
          <u>return</u> ls+rs+(ls+rs-s)/15.0;
10
11
      \underline{\text{return}} asr(1, mid, eps/2, ls) + ars(mid, r, eps/2, rs);
12|}
13 <u>int</u> main(){
14
      db ans = asr();
|15| }
```

3.9 几何函数-最小圆覆盖

```
1 #include <algorithm>
2 #include <iostream>
3 #include <cstring>
4 #include <cstdio>
5 #include <cmath>

7 using namespace std;

9 struct vec
10 {
11     double x, y;
12     vec (const double& x0 = 0, const double& y0 = 0) : x(x0), y(y0) {}
```

```
vec operator + (const vec& t) const {return vec(x+t.x, y+t.y);}
       vec operator - (const vec& t) const {return vec(x-t.x, y-t.y);}
14
       vec operator * (const double& t) const {return vec(x*t, y*t);}
15
16
       vec operator / (const double& t) const {return vec(x/t, y/t);}
       const double len2 () const {return x*x + y*y;}
17
       const double len () const {return sqrt(len2());}
18
19
       vec norm() const {return *this/len();}
20
       vec rotate_90_c () {return vec(y, -x);}
21|};
23 double dot(const vec& a, const vec& b) {return a.x*b.x + a.y*b.y;}
24 double crs(const vec& a, const vec& b) {return a.x*b.y - a.y*b.x;}
26 vec lin_lin_int(const vec& p0, const vec& v0, const vec& p1, const vec& v1)
27 {
28
       \underline{\text{double}} \ \text{t = crs(p1-p0, v1) / crs(v0, v1)};
       return p0 + v0 * t;
29
30| \}
32 vec circle(const vec& a, const vec& b, const vec& c)
33| {
       return lin_lin_int((a+b)/2, (b-a).rotate_90_c(), (a+c)/2, (c-a).rotate_90_c());
34
35|}
37 int n;
38 vec pot[100005];
40 <u>int</u> main()
41| {
       scanf("%d", &n);
42
43
       for(int i=1; i<=n; i++) scanf("%lf%lf", &pot[i].x, &pot[i].y);</pre>
44
       random_shuffle(pot+1, pot+n+1);
45
       vec o;
46
       \underline{\text{double}} r2 = 0;
47
       for(int i=1; i<=n; i++)</pre>
48
           <u>if</u>((pot[i]-o).len2() > r2)
49
50
51
               o = pot[i], r2 = 0;
52
               <u>for</u>(<u>int</u> j=1; j<i; j++)
53
54
                   \underline{if}((pot[j]-o).len2() > r2)
55
                       o = (pot[i]+pot[j])/2, r2 = (pot[j]-o).len2();
56
57
                       <u>for</u>(<u>int</u> k=1; k<j; k++)
58
59
                           \underline{if}((pot[k]-o).len2() > r2)
60
61
                               o = circle(pot[i], pot[j], pot[k]), r2 = (pot[k]-o).len2();
62
                      }
63
                   }
64
65
               }
           }
66
67
68
       printf("%.101f\n%.101f\n", sqrt(r2), o.x, o.y);
69
       return 0;
70 }
```

3.10 几何函数-最小球覆盖

```
1 \mid double cx=0, cy=0, cz=0;
 2 double x[maxn],y[maxn],z[maxn];
3 double dis(int now){
     return sqrt((x[now]-cx)*(x[now]-cx)+
5
           (y[now]-cy)*(y[now]-cy)+
6
           (z[now]-cz)*(z[now]-cz));
7 }
8
   int main() {
9
     cin>>n;
     rep(i,1,n) cin>>x[i]>>y[i]>>z[i];
10
11
     int pos=1;
12
     double dmax=1e4,ans=1e18;
13
     \underline{\text{while}}(\text{dmax}>1\text{e}-7)
```

```
rep(i,1,n) <u>if</u>(dis(i)>dis(pos)) pos=i;
15
       double d=dis(pos);
16
       ans=min(ans,d);
       cx+=(x[pos]-cx)/d*dmax;
17
       cy+=(y[pos]-cy)/d*dmax;
18
       cz + = (z[pos] - cz)/d*dmax;
19
       dmax*=0.98;
20
|22|
     printf("%.6f \n",ans);
23
     return 0;
24| \}
```

4 Game

4.1 SG函数

```
1 int sg[maxn] [maxn];
  int getmex(bool *vis,int n){
3
     int t=0;
     while(vis[t]) ++t;
4
5
     return t;
6 }
7
   int getsg(int n,int m){
     if(~sg[n][m]) return sg[n][m];
     if(n<=m) return sg[n][m]=1;</pre>
     bool vis[n+1]=\{0\};
10
11
     rep(i,1,m)vis[getsg(n-i,m)]=1;
     return sg[n][m]=getmex(vis,n);
13|}
14 <u>int</u> main() {
15
    memset(sg,-1,sizeof sg);
16
     cin>>casn;
     while(casn--){
17
18
       \underline{\operatorname{cin}} >> n >> m;
       if(getsg(n,m)) cout<<"first\n";</pre>
19
20
       else cout<<"second\n";</pre>
21
22
     return 0;
23 }
```

4.2 威佐夫游戏

```
int main(){
    double k=(1+sqrt(5.0))/2;
    while(scanf("%d%d",&n,&m)!=EOF) {
        if (n>m) swap(n,m);
        int t=m-n;
        if (n==(int)((double)t*k)) printf("0\n");
        else printf("1\n");
}
```

4.3 K倍取石子博弈

```
2 #include <iostream>
3 <u>#include</u><cstdio>
4 #include <cstring>
5 #include <algorithm>
 6 using namespace std;
7 <u>const</u> <u>int</u> maxn=2000000;
8 int a[maxn],b[maxn];
9 <u>int</u> main()
10 {
11
       int t,n,k;
       cin>>t;
12
       for(int cas=1;cas<=t;cas++)</pre>
13
|14|
15
           <u>cin</u>>>n>>k; //n个石子每次拿前一次的最多k倍
           printf("Case %d: ",cas);
16
           <u>if</u>(n<=k+1)
17
18
               printf("lose\n");
19
20
                continue;
21
           a[0]=b[0]=1;
|22|
```

```
<u>int</u> i=0, j=0;
24
             \underline{\text{while}}(a[i] < n)
25
26
                  i++;
                  a[i]=b[i-1]+1;
27
                  while(a[j+1]*k<a[i]) j++;</pre>
28
29
                  <u>if</u>(a[j]*k<a[i]) b[i]=b[j]+a[i];
30
                  else b[i]=a[i];
31
             \underline{if}(a[i]==n)printf("lose\n");
33
34
             <u>else</u>
35
                  int ans=0;
36
37
                  while(n)
38
                  {
39
                       if(n)=a[i]
40
41
                            n-=a[i]
42
                            ans=a[i];
43
                       i--;
44
45
                  }
                  cout<<ans<<endl;</pre>
46
47
48
49 }
```

5 Math

5.1 数论-基础函数

```
1 ll p=1e9+7;
  11 gcd(l1 a,l1 b) {return b?gcd(b,a%b):a;}
3 11 lcm(11 a, 11 b) \{return a*gcd(a,b)/b;\}
 4 | 11 exgcd(11 a,11 b,11 &x,11 &y) {
    \underline{if}(b==0) \underline{return} (x=1,y=0,a);
     <u>if</u>(a==0) <u>return</u> (x=0,y=1,b);
    11 r=exgcd(b,a%b,y,x);
8
    y=(a/b)*x;
    return r;
10 }
11 | 11 | 1cm_mod(11 a,11 b,11 c=p) {
    return (a/gcd(a,b)*b)%c;
12
13|}
14 | 11 pow_mod(ll a, ll b, ll c=p, ll ans=1) {
15
    while(b) {
16
       if(b&1) ans=(a*ans)%c;
17
      a=(a*a)%c,b>>=1;
18
19
    return ans;
20| \}
21 //long double 1e9以下表现良好,不会出现误差
22 11 mul_mod_2(11 a,11 b,11 m){
23
   11 c=a*b-(11)((long double)a*b/m+0.5)*m;
24
   return c<0?c+m:c;</pre>
25| \}
26 //不丢失精度的快速乘
27 | 11 mul_mod(11 a,11 b,11 c){return (__int128)a*b%c;}
28 | 11 pow_mul_mod(ll a, ll b, ll c=p, ll ans=1){
29
    while(b){
30
       if(b&1)res=mul_mod(res,a,c);
31
       a=mul_mod(a,a,c),b>>=1;
32
33
    return res;
34|}
35 | 11 inv_gcd(11 a,11 c=p){
36
    a%=c;
    if(a<0)a+=c;
37
    ll b=c,u=0,v=1;
38
39
    while(a) {
40
      11 t=b/a;b=t*a;
      swap(a,b);
41
42
      u-=t*v;
```

```
swap(u,v);
     }
|44|
45
      if(u<0)u+=c;
46
     <u>return</u> u;
47 }
```

数论-欧拉函数线性筛

```
1 / / \#x \text{ that } x \le n \&\& \gcd(x,n) == 1
  int euler_phi(int n) {
     \underline{int} m = (\underline{int})sqrt(n+0.5);
     int ans = n;
     for (int i = 2; i <= m; ++ i) if (n % i == 0) {</pre>
 5
        ans = ans / i * (i-1);
 6
        while (n\%i == 0) n /= i;
 7
 8
     <u>if</u> (n > 1) ans = ans / n * (n-1);
10
     return ans;
11|}
   int phi[maxn];
12
13 void phi_table(int n) {
      \underline{\text{for}} \ (\underline{\text{int}} \ i = 2; \ i \le n; ++ i) \ \text{phi}[i] = 0;
14
     phi[1] = 1;
15
16
     for (int i = 2; i <= n; ++ i) {
17
        <u>if</u> (!phi[i]) {
18
          <u>for</u> (<u>int</u> j = i; j <= n; j += i) {
19
            <u>if</u> (!phi[j]) phi[j] = j;
20
            phi[j] = phi[j] / i * (i-1);
21
22
23
        phi[i] += phi[i-1];
24
25|}
```

数论-莫比乌斯函数线性筛

```
1 int mu[maxn],prime[maxn],sum[maxn],nump;
 2 bool isp[maxn];
3 void getmu(){
    mu[1]=1,nump=0;
5
     int n=maxn-10;
6
    rep(i,2,n){
7
       <u>if</u>(!isp[i]) prime[++nump]=i,mu[i]=-1;
       for(int j=1;j<=nump&&prime[j]*i<=n;j++){</pre>
8
         isp[i*prime[j]]=1;
9
10
         <u>if</u>(i%prime[j]==0) mu[i*prime[j]]=0,j=nump+10;
         else mu[i*prime[j]]=-mu[i];
|11|
12
    }
13
14
    rep(i,1,n){
       sum[i]=sum[i-1]+mu[i];
15
16
17 }
```

数论-Miller Rabin素数判定

```
11 //两种米勒罗宾素数筛实现
2| //第一种速度稍快,但实现麻烦,使用前需要先初始化
3 //单次复杂度约为1-maxpe中的素数个数*logn,maxpe不要小于50
4 const int maxpe=100;
5 auto randint=bind(uniform_int_distribution<int>(1,1e9),mt19937(rand()));
6 | 11 mul_mod(11 a,11 b,11 c){return (__int128)a*(__int128)b%(__int128)c;}
7 int mul_mod(int a,int b,int c){return (11)a*(11)b%c;}
  template<typename T> int pow_mod(int a,int b,int c){
9
    int res=1;
10
    while(b){
11
      if(b&1)res=mul_mod(res,a,c);
12
      a=mul_mod(a,a,c);
13
      b>>=1;
    }
|14|
15
    return res;
16|}
17 template < typename T > class miller_rabin {public:
18
    T prime[maxpe],cntp;
    void init(int n=maxpe-5){
19
```

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```
cntp=0;
       bool vis[n+1];
21
|22|
       rep(i,2,n)vis[i]=1;
23
       rep(j,2,n)if(vis[j]==1)
24
         for(int m=2; j*m<=n; ++m) vis[j*m]=0;</pre>
25
       rep(i,2,n)<u>if</u>(vis[i]==1) prime[cntp++]=i;
26
27
     bool _test(T n,T a,T d) {
28
       if(n==2||n==a) return true;
29
       if((n&1)==0) return false;
30
       while(!(d&1))d>>=1;
31
       T t=pow_mod(a,d,n);
32
       \underline{\text{while}}(d!=n-1\&\&t!=1\&\&t!=n-1)
33
         t=mul_mod(t,t,n);
34
         d <<=1;
35
36
       \underline{\text{return}} (t==n-1||(d&1)==1);
37
38
     bool test(T n) {
39
       if(n<2||n\%2==0) return false;
40
       rep(i,0,cntp-1) <u>if(!_test(n,prime[i],n-1)) return false;</u>
41
       return true;
42
43 };
44 miller_rabin<int> miller;
45 //第二种实现
46 //速度稍慢,复杂度约为logn*time,time不要低于10
47
   const int test_time=20;
  template<typename T> bool miller_rabin(T n) {
48
     \underline{if}(n<3)\underline{return} n==2;
49
     T = n-1, b=0;
50
     <u>while</u> (a\%2==0) a/=2,++b;
51
52
     for(int i=1,j; i<=test_time;++i){</pre>
53
       T x=randint()\%(n-2)+2,v=pow_mod(x,a,n);
       if(v==1||v==n-1) continue;
54
55
       for(j=0;j<b;++j){
56
         v=mul_mod(v,v,n);
57
         if(v==n-1) break;
58
59
       if(j>=b)return 0;
60
61
     return 1;
62| \}
```

5.5 数论-Pollard Rho因子分解

```
1 //得到n的一个随机因子,包括自身和1
2|//复杂度O(n^{1/4})
3 //要保证randint的随机范围大于等于测试数字
4 | 11 mul_mod(11 a,11 b,11 c){return (__int128)a*b%c;}
5 int mul_mod(int a, int b, int c){return (11)a*(11)b%c;}
6 auto randint=bind(uniform_int_distribution<11>(1e9,1e18),mt19937(rand()));
  template < typename T> T pollard_rho(T n,T c=randint()) {
    T i=1,k=2,x=randint()\%(n-1)+1,y=x,d;
q
    \underline{\text{while}}(1){
10
      i++;
11
      x=(mul_mod(x,x,n)+c)%n;
12
      d=_gcd(n,y-x);
13
      if(d>1&&d<n)return d;</pre>
14
      if(y==x)return n;
15
      \underline{if}(i==k){}
16
        k <<=1;
17
        y=x;
18
19
    }
20|}
21 //分解因子,map中即为素因子从小到大,first为因子,second为次幂
22 //分解1000个1e18的数字约为600ms,10000个1e9的数字约为200ms
23 map<11, int > factor;
24 template typename T > void get_factor(T n,T c=randint()) {
25
    if(n==1) return;
26
    if(miller_rabin(n)) {
      factor[n]++;
27
```

```
return;
29
    T p=n;
30
31
    while(p>=n) p=pollard_rho(p,c--);
32
    get_factor(p,c);
33
    get_factor(n/p,c);
34 }
        多项式-拉格朗日插值
1 class polysum {public:
    ll a[maxn],f[maxn],g[maxn],p[maxn],p1[maxn],p2[maxn],b[maxn],h[maxn][2],C[maxn];
    ll calcn(\underline{int} d,ll *a,ll n) {//len=d get(an)
      if (n<=d) return a[n];</pre>
5
      p1[0]=p2[0]=1;
6
      rep(i,0,d) {
7
        11 t=(n-i+mod) mod;
8
        p1[i+1]=p1[i]*t%mod;
9
10
      rep(i,0,d) {
11
        11 t=(n-d+i+mod)\%mod;
12
        p2[i+1]=p2[i]*t\mod;
13
      ll ans=0;
14
15
      rep(i,0,d) {
        11 t=g[i]*g[d-i]%mod*p1[i]%mod*p2[d-i]%mod*a[i]%mod;
16
17
        if ((d-i)\&1) ans=(ans-t+mod)\mod;
        else ans=(ans+t)%mod;
18
19
20
      return ans;
21
|22|
    void init(int maxm) {//init
23
      f[0]=f[1]=g[0]=g[1]=1;
24
      rep(i,2,maxm+4) f[i]=f[i-1]*i\%mod;
25
      g[maxm+4] = pow_mod(f[maxm+4], mod-2);
26
      per(i,1,maxm+3) g[i]=g[i+1]*(i+1)%mod;
27
28
    ll polysum(ll n,ll *a,ll m){//a[i]会被修改
      // 初始化预处理阶乘和逆元(取模乘法)a[0].. a[m] \sum_{i=0}^{n-1} a[i]
29
30
      // len=m,psum_n
      a[m+1]=calcn(m,a,m+1);
31
      rep(i,1,m+1) a[i]=(a[i-1]+a[i]) mod;
32
33
      return calcn(m+1,a,n-1);
34
    ll qpolysum(ll R,ll n,ll *a,ll m) { // a[0].. a[m] \sum_{i=0}^{n-1} a[i]*R^i
35
36
      if (R==1) return polysum(n,a,m);
37
      a[m+1]=calcn(m,a,m+1);
      ll r=pow_mod(R,mod-2),p3=0,p4=0,c,ans; h[0][0]=0;
38
39
40
      h[0][1]=1;
41
      rep(i,1,m+1) {
        h[i][0]=(h[i-1][0]+a[i-1])*r\mod;
42
        h[i][1]=h[i-1][1]*r%mod;
43
44
45
      rep(i,0,m+1) {
46
        ll t=g[i]*g[m+1-i]%mod;
        <u>if</u> (i&1) p3=((p3-h[i][0]*t)%mod+mod)%mod,p4=((p4-h[i][1]*t)%mod+mod)%mod;
47
        else p3=(p3+h[i][0]*t)\( mod, p4=(p4+h[i][1]*t)\( mod; )
48
49
```

多项式-快速傅立叶变换

if (ans<0) ans+=mod;</pre>

return ans;

 $c=pow_mod(p4,mod-2)*(mod-p3)%mod;$

ans= $(calcn(m,C,n)*pow_mod(R,n)-c)\mbox{mod};$

rep(i,0,m+1) C[i]=h[i][0];

 $rep(i,0,m+1) h[i][0]=(h[i][0]+h[i][1]*c)\mod;$

50

51

52

53

54

55

56 57 }

```
1| const double pi=acos(-1.0);
 struct cp{double x,y;};
3 cp operator*(cp a,cp b){return {a.x*b.x-a.y*b.y,a.x*b.y+a.y*b.x};}
4 cp operator + (cp a, cp b) {return {a.x+b.x,a.y+b.y};}
```

```
5| cp operator (cp a,cp b){return {a.x-b.x,a.y-b.y};}
6 #define carr vector <cp>
7 <u>const</u> <u>int</u> maxl=6e4+10;//卷积单个数组的最大长度
8 class fourier{public:
9
     int rev[max1<<2],len,pw;</pre>
10
     vector<cp> wt;
     void init_0(int ml=maxl){
11
12
       <u>for</u>(<u>int</u> mid=1;mid<2*ml;mid<<=1){
13
         wt.resize(mid*2+1);
|14|
         cp wn={cos(pi/mid),sin(pi/mid)};
15
         wt[mid] = (cp) \{1,0\};
         rep(j,1,mid-1) wt [mid+j]=wt[mid+j-1]*wn;
16
17
     }
18
19
     void init(int n){
20
       len=1,pw=0;
21
       \underline{\text{while}}(\text{len} \leq n) \text{ len} \leq 1, ++pw; --pw;
22
       rep(i,0,len-1) rev[i]=rev[i>>1]>>1|(i&1)<<pw;
23
24
     void transform(carr &a, int flag){
25
       \underline{if}(a.size()!=len) a.resize(len,(cp){0,0});
26
       rep(i,0,len-1) \underline{if}(i < rev[i]) swap(a[i],a[rev[i]]);
27
       for(int mid=1;mid<len;mid<<=1){</pre>
28
         for(int r=mid<<1,j=0;j<len;j+=r){</pre>
29
           for(int k=0;k<mid;++k){</pre>
30
             cp wn=wt[mid+k];
             if(flag==-1) wn.y=-wn.y;
31
             cp y=wn*a[mid+j+k];
32
33
             a[j+k+mid]=a[j+k]-y,a[j+k]=a[j+k]+y;
           }
34
         }
35
36
37
       <u>if</u>(flag==-1) rep(i,0,len-1) a[i].x/=len;
     }//会破坏掉a,b数组,视情况可以去掉引用
38
39
     carr mul(carr &a,carr &b){
40
       int la=a.size(),lb=b.size();
       if((11)1a*1b<=1000){
41
42
         carr c(la+lb,(cp)\{0,0\});
43
         rep(i,0,la-1) rep(j,0,lb-1) c[i+j]=c[i+j]+a[i]*b[j];
         return c;
44
45
46
       init(la+lb);
       carr c(len,(cp){0,0});
47
48
       transform(a,1);transform(b,1);
       rep(i,0,len-1)c[i]=a[i]*b[i];
49
50
       transform(c,-1);
51
       return c;
52
53 }fft;
```

多项式-快速傅立叶变换-数组实现

```
const double pi=acos(-1.0);
 2 struct cp{double x,y;};
3 \mid \text{cp operator}*(\text{cp a,cp b})\{\frac{\text{return}}{\text{a.x*b.x-a.y*b.y,a.x*b.y+a.y*b.x}};\}
4 cp operator + (cp a, cp b) {return {a.x+b.x,a.y+b.y};}
5 cp operator (cp a,cp b) {return {a.x-b.x,a.y-b.y};}
6 const int max1=2e6+10;
   class fourier{public:
     int rev[maxl<<2],len,pw;</pre>
9
     void init(int n){
10
       len=1,pw=0;
       <u>while</u>(len<=n) len<<=1,++pw;--pw;
11
12
       rep(i,0,len-1) rev[i]=rev[i>>1]>>1|(i&1)<<pw;
13
14
     cp c1[max1<<2],c2[max1<<2];</pre>
15
     vector<cp> wt;
16
     void init_0(){
       for(int mid=1;mid<2*maxl;mid<<=1){</pre>
17
18
         wt.resize(mid*2+1);
19
         cp wn{cos(pi/mid),sin(pi/mid)};
20
         wt[mid] = (cp) \{1,0\};
21
         rep(j,1,mid-1) wt[mid+j]=wt[mid+j-1]*wn;
```

```
}
23
24
     void transform(cp*a, int flag){
      rep(i,0,len-1) if(i<rev[i]) swap(a[i],a[rev[i]]);
25
       for(int mid=1;mid<len;mid<<=1){</pre>
26
27
         for(int r=mid<<1,j=0;j<len;j+=r){</pre>
28
           for(int k=0;k<mid;++k){</pre>
29
             cp wn=wt[mid+k];
30
            if(flag==-1) wn.y=-wn.y;
31
             cp y=wn*a[mid+j+k];
32
             a[j+k+mid]=a[j+k]-y,a[j+k]=a[j+k]+y;
33
        }
34
35
36
       \underline{if}(flag==-1) rep(i,0,len-1) a[i].x/=len;
37
     }
38
     void fix(vector<int>&a){
39
       while(!a.empty()&&!a.back())a.pop_back();
     }//传入整数,进行卷积
40
41
     void mul(vector<int> &a, vector<int> &b){
42
       int la=a.size(),lb=b.size();
43
       init(la+lb);
44
       rep(i,0,la-1) c1[i]=(cp){a[i],0};
45
       rep(i,la,len-1) c1[i]=(cp){0,0};
46
       rep(i,0,1b-1) c2[i]=(cp){b[i],0};
47
       rep(i,lb,len-1) c2[i]=(cp){0,0};
48
       transform(c1,1);transform(c2,1)
49
       rep(i,0,len-1) c1[i]=c1[i]*c2[i];
50
       transform(c1,-1);
51
       a.resize(len)
       rep(i,0,len-1) a[i]=\underline{int}(c1[i].x+0.5)?1:0;
52
53
       fix(a);
54
55
    void sqr(vector<int> &a){
       int la=a.size(); init(2*la);
56
57
       rep(i,0,la-1)c1[i]=(cp){a[i],0};
58
       rep(i,la,len-1)c1[i]=(cp){0,0};
59
       transform(c1,1);
       rep(i,0,len-1) c1[i]=c1[i]*c1[i];
60
61
       transform(c1,-1);
62
       a.resize(len);
63
       rep(i,0,len-1) a[i]=<u>int</u>(c1[i].x+0.5)?1:0;
64
       fix(a);
65
66
     vector<int> pow(vector<int>a,int k){
67
       fix(a);
68
       vector<int>ret;
69
       \underline{\text{while}}(k) {
70
         if(k&1){
71
           if(ret.empty()) ret=a;
72
           else mul(ret,a);
73
74
        sqr(a);
        k/=2;
75
76
77
       return ret;
78
79|}
```

5.9 多项式-杜教BM

```
1 namespace bm{
     const int maxl=1e4+10;
     11 res[maxl], base[maxl], c[maxl], md[maxl];
     vector<11> md;
     ll inv(ll a,ll c=mod) {
6
       a\%=c; if(a<0)a+=c;
7
       ll b=c,u=0,v=1;
8
       \underline{\text{while}}(a)  {
9
         11 t=b/a;b=t*a;
10
         swap(a,b);u=t*v;
11
         swap(u,v);
12
```

```
if(u<0)u+=c;
14
       <u>return</u> u;
     }
15
16
     void mul(ll *a,ll *b,int k) {
       for(int i=0;i<k+k;i++) _c[i]=0;</pre>
17
18
       for(int i=0;i<k;i++) if (a[i])</pre>
19
         for(int j=0;j<k;j++) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;</pre>
20
       for (ll i=k+k-1;i>=k;i--) if (_c[i])
21
         for(int j=0;j<md.size();j++)</pre>
22
            _c[i-k+md[j]]=(_c[i-k+md[j]]-_c[i]*_md[md[j]])%mod;
23
       for(int i=0;i<k;i++) a[i]=_c[i];</pre>
     }
24
25
     int solve(ll n, vector<11> a, vector<11> b) {
     //a 系数 b 初值 b[n+1]=a[0]*b[n]+...
26
     //求出的是第n+1项
27
       ll ans=0,pnt=0;
28
29
       11 k=a.size();
30
       for(int i=0;i<k;i++) _md[k-1-i]=-a[i];_md[k]=1;</pre>
31
       md.clear();
       \underline{\text{for}}(\underline{\text{int}} \text{ i=0;i<k;i++}) \underline{\text{if}} (\underline{\text{md}}[i]!=0) \text{ md.push_back(i);}
32
33
       for(int i=0;i<k;i++) res[i]=base[i]=0;</pre>
34
       res[0]=1;
35
       <u>while</u> ((111<<pnt)<=n) pnt++;
36
       <u>for</u> (ll p=pnt;p>=0;p--) {
37
         mul(res,res,k);
38
         <u>if</u> ((n>>p)&1) {
39
           for (ll i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
40
           for(int j=0;j<md.size();j++)</pre>
41
             res[md[j]]=(res[md[j]]-res[k]*_md[md[j]])%mod;
42
       }
43
       for(int i=0;i<k;i++) ans=(ans+res[i]*b[i])%mod;</pre>
|44|
45
       if (ans<0) ans+=mod;</pre>
46
       <u>return</u> ans;
47
48
     vector<ll> init(vector<ll> s) {
49
       vector<ll> coe(1,1), base(1,1);
50
       int len=0, m=1, b=1;
51
       \underline{for}(\underline{int} n=0;n< s.size();n++) {
52
         11 d=0;
53
         for(int i=0;i<len+1;i++) d=(d+(l1)coe[i]*s[n-i])%mod;</pre>
54
         if (d==0) ++m;
         \underline{else} \underline{if} (2*len<=n) {
55
56
            vector<ll> tmp=coe;
57
            11 c=mod-d*inv(b)%mod;
58
           while (coe.size() < base.size() + m) coe.push_back(0);</pre>
59
            <u>for(int</u> i=0;i<base.size();i++) coe[i+m]=(coe[i+m]+c*base[i])%mod;
           len=n+1-len; base=tmp; b=d; m=1;
60
61
         } <u>else</u> {
           11 c=mod-d*inv(b)%mod;
62
63
            while (coe.size()<base.size()+m) coe.push_back(0);</pre>
64
           for(int i=0;i<base.size();i++) coe[i+m]=(coe[i+m]+c*base[i])%mod;</pre>
65
            ++m;
66
       }
67
68
       return coe;
69
70
     vector<11> c,a;
71
     void inita(vector<11> _a){
72
       a=_a;
73
       c=init(a);c.erase(c.begin());
74
       for(auto &i:c) i=(mod-i)%mod;
     }
75
76
     int get(ll n) {
77
       return solve(n,c,vector<11>(a.begin(),a.begin()+c.size()));
78
79
     int get(vector<11> a,11 n) {
80
       vector<1l> c=init(a);
81
       c.erase(c.begin());
82
       for(int i=0;i<c.size();i++) c[i]=(mod-c[i])%mod;</pre>
83
       return solve(n,c,vector<11>(a.begin(),a.begin()+c.size()));
84
```

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```
85|};
```

```
5.10 多项式-快速数论变换
```

```
1 const 11 mod=998244353,modg=3,modi=332748118;
   int a[maxn],b[maxn];
 3 <u>int</u> pow_mod(<u>int</u> a, <u>int</u> b){
     int ans=1;
     while(b){
 6
       \underline{if}(b&1) ans=(ll)ans*a\mod;
       a=(ll)a*a\%mod,b>>=1;
 8
 9
     return ans;
10 }
11 int add(int a, int b){
12
     a+=b; if (a>mod) return a-mod;
13
     return a;
14 }
15 | \underline{int}  sub(\underline{int}  a,\underline{int}  b){
16
     a-=b; \underline{if}(a<0) \underline{return} a+mod;
17
     return a;
18|}
19 <u>#define</u> arr <u>vector</u><<u>int</u>>
20 <u>const</u> <u>int</u> maxl=2e4+10;//卷积单个数组的最大长度
21 arr operator*(arr&,arr&b){
     int len=a.size();
23
     arr c(len));
24
     rep(i,0,len-1) c[i]=a[i]*b[i];
     return c;
25
26 }
27 class nubmer{public:
28
     int rev[max1<<2],len,pw;</pre>
29
     <u>int</u> wt[2][max1<<2];
30
     void init_0(){
       <u>int</u> len=1;//最开始的初始化,整个程序一次就够
31
32
       while(len<=maxl) len<<=1;</pre>
33
       for(int mid=1;mid<len;mid<<=1){</pre>
34
          11 wn1=pow_mod(modg,(mod-1)/(mid<<1));</pre>
35
         11 wn2=pow_mod(modi,(mod-1)/(mid<<1));</pre>
36
         wt[0][mid]=wt[1][mid]=1;
37
         11 wt1=wn1,wt2=wn2;
38
         rep(j,1,mid){
39
            wt[0][mid+j]=wt1;wt1=wt1*wn1%mod;
40
            wt[1][mid+j]=wt2;wt2=wt2*wn2%mod;
41
       }
42
     }
43
|44|
     void init(int n){
45
       len=1,pw=0;
46
       \underline{\text{while}}(\text{len} \leq n) \text{ len} \leq 1, ++pw; --pw;
       rep(i,0,len-1) rev[i]=rev[i>>1]>>1|(i&1)<<pw;
47
48
49
     void transform(arr &a,int flag){
50
       ll* f=flag==1?wt[0]:wt[1];
       if(a.size()!=len) a.resize(len);
51
       rep(i,0,len-1) <u>if</u>(i<rev[i]) swap(a[i],a[rev[i]]);
52
53
       for(int mid=1;mid<len;mid<<=1){</pre>
54
         <u>for</u>(<u>int</u> r=mid<<1,j=0;j<len;j+=r){
55
            11 *p=f+mid;
56
            for(int k=0;k<mid;++k,++p){</pre>
57
              <u>int</u> x=a[j+k],y=(*p)*a[j+k+mid]%mod;
              a[j+k+mid]=sub(a[j+k],y),a[j+k]+=y;
58
59
              \underline{if}(a[j+k]>mod) a[j+k]-=mod;
60
         }
61
62
63
       <u>if</u>(flag==-1) {
         11 inv=pow_mod(len,mod-2);
64
65
         rep(i,0,len-1){
66
            a[i]=a[i]*inv%mod;
67
            <u>if</u>(a[i]<0)a[i]+=mod;
68
69
       }
```

```
70
71
    void fix(arr &a){//分治ntt优化
72
      while(!a.empty()&&!a.back()) a.pop_back();
73
    }//会破坏掉a,b数组,视情况可以去掉引用
74
    arr mul(arr &a,arr &b){
75
      int la=a.size(),lb=b.size();
      <u>if</u>(la*lb<=1000){
76
77
        arr c(la+lb,0);
        rep(i,0,la-1) rep(j,0,lb-1)
78
79
          c[i+j]=(c[i+j]+(l1)a[i]*b[j])%mod;
80
81
82
      int n=la+lb;
83
      init(n);
84
      transform(a,1); transform(b,1);
85
      arr c=a*b;
86
      rep(i,0,len-1)c[i]=(ll)a[i]*b[i]%mod;
87
      transform(c,-1);
88
      return c;
89
91| }ntt;
```

5.11 多项式-快速沃尔什变换

```
\frac{\text{#define}}{\text{define}} add(a,b) ((a+=b)>=mod?a-=mod:a)
   class walsh{public:
     void transform_or(ll *a, int len, int flag){
       \underline{for}(\underline{int} i=1;i<len;i<<=1)
5
         for(int p=i<<1,j=0;j<len;j+=p)</pre>
            for(int k=0;k<i;++k)</pre>
 6
7
              add(a[i+j+k],flag==1?a[j+k]:mod-a[j+k]);
8
     void transform_and(ll *a, int len, int flag){
       for(int i=1;i<len;i<<=1)</pre>
10
         for(int p=i<<1,j=0;j<len;j+=p)</pre>
11
12
            for(int k=0;k<i;++k)</pre>
13
              add(a[j+k],flag==1?a[i+j+k]:mod-a[i+j+k]);
14
15
     void transform_xor(ll *a, int len, int flag){
16
       <u>for(int</u> i=1;i<len;i<<=1)
17
         for(int p=i<<1,j=0;j<len;j+=p)</pre>
18
            for(int k=0;k<i;++k){</pre>
19
              <u>int</u> x=a[j+k],y=a[i+j+k];
20
              a[j+k]=(x+y) \mod , a[i+j+k]=(x+mod-y) \mod ;
21
              if(flag==-1)
22
                a[j+k]=a[j+k]*inv2\mod, a[i+j+k]=111*a[i+j+k]*inv2\mod;
           }
23
24
25
   }fwt;
26
   int main(){
27
     \underline{\operatorname{cin}} >> n; n = (1 << n);
28
     rep(i,0,n-1) \underline{cin} >> a[i];
     rep(i,0,n-1) <u>cin</u>>>b[i];
29
30
     fwt.transform_or(a,n,1),fwt.transform_or(b,n,1);
     rep(i,0,n-1) ans[i]=a[i]*b[i]%mod;
31
32
     fwt.transform_or(a,n,-1),fwt.transform_or(b,n,-1);
33
     fwt.transform_or(ans,n,-1);
34
     rep(i,0,n-1) <u>cout</u><<ans[i]<<' ';<u>cout</u><<<u>endl</u>;
35
     fwt.transform_and(a,n,1),fwt.transform_and(b,n,1);
36
     rep(i,0,n-1) ans[i]=a[i]*b[i]%mod;
37
     fwt.transform_and(a,n,-1),fwt.transform_and(b,n,-1);
38
     fwt.transform_and(ans,n,-1);
39
     rep(i,0,n-1) cout << ans[i] << ' '; cout << endl;</pre>
     fwt.transform_xor(a,n,1),fwt.transform_xor(b,n,1);
40
41
     rep(i,0,n-1) ans[i]=a[i]*b[i]%mod;
42
     fwt.transform_xor(a,n,-1),fwt.transform_xor(b,n,-1);
43
     fwt.transform_xor(ans,n,-1);
44
     rep(i,0,n-1) <u>cout</u><<ans[i]<<' ';<u>cout</u><<<u>endl</u>;
|45| }
```

5.12 线性代数-异或线性基

1 template < typename T, const int len> class lbass {public:

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```
T d[len+1]; int cnt; bool flag;
    void init(){flag=0;memset(d,0,sizeof d);}
3
    bool insert(T x){
|4|
5
      for(int i=len;x&&i>=0;--i)
6
         <u>if</u>((T)1<<i&x){
           if(!d[i]) {d[i]=x;return true;}
           else x^=d[i];
9
10
      flag=1;
11
      <u>return</u> <u>false</u>;
12
     //线性基和x异或的最值
13
    T querymax(T x=0){
14
      per(i,0,len)x=max(x,x^d[i]);
15
16
      return x;
17
18
    T querymin(T x=0){
19
      per(i,0,len)x=min(x,x^d[i]);
20
       <u>return</u> x;
21
22
     //求所有异或值去重后的第k大,需要预处理
23
    T p[len+1];
24
    void makekth(){
25
      cnt=0;
26
      memset(p,0,sizeof p);
27
      T ans=0;
28
      per(i,0,len)per(j,0,i-1)
        <u>if</u>((T)1<<j&d[i]) d[i]^=d[j];
29
30
      rep(i,0,len) <u>if</u>(d[i]) p[cnt++]=d[i];
31
32
    T querykth(T k){
      <u>if</u>(flag)--k;//包含零
33
       if(!k) return 0;
34
35
      T res=0;
36
       if(k>=(T)1<<cnt) return -1;
37
      per(i,0,len) <u>if</u>((T)1<<i&k)res^=p[i];
38
       return res;
39
40|};
41 template typename T, const int len lbass T, len merge (const lbass T, len &a, const lbass T, len &b) {
42
    lbass<T,len>res=a;
43
    rep(i,0,len) res.insert(b.d[i]);
44
    return res;
45| \}
```

5.13 线性代数-线段树维护区间线性基

```
class segtree{public:
  #define nd node[now]
  #define ndl node[now<<1]</pre>
4
  #define ndr node[now<<1|1]</pre>
5
     struct segnode{
       int l,r,flag,val;
6
       <u>int</u> d[32];
       inline void init(){val=flag=0;memset(d,0,sizeof d);}
 9
       inline void insert(ll x){
10
         for(register int i=30;x&&i>=0;--i)
11
             if(x&(111<<i)){
|12|
               <u>if</u>(!d[i]) {d[i]=x;<u>return</u>;}
13
               else x^=d[i];
14
15
       int count(){int ans=0;per(i,0,30) if(d[i])ans++; return ans;}
16
       void update(int x){val^=x;flag^=x;}
17
18
     node[maxn << 2|3];
19
     <u>inline</u> segnode marge(segnode &a, segnode b) <u>const</u> {
20
       segnode ans;ans.init();
21
       per(i,0,30) ans.insert(a.d[i]),ans.insert(b.d[i]);
|22|
       ans.insert(a.val^b.val);
23
       ans.val=a.val;
24
       ans.l=a.l,ans.r=b.r;
25
       return ans;
```

```
<u>inline</u> <u>void</u> down(<u>int</u> now){
28
       if(nd.flag){
29
         ndl.update(nd.flag);ndr.update(nd.flag);
30
         nd.flag=0;
31
     }
32
33
     void maketree(int s,int t,int now=1){
34
       nd.l=s,nd.r=t;nd.init();
35
       if(s==t) {cin>>nd.val;return ;}
36
       maketree(s,(s+t)/2,now<<1);
37
       maketree((s+t)/2+1,t,now << 1|1);
38
       nd=marge(ndl,ndr);
39
40
     void update(int s,int t,int x,int now=1){
41
       if(s<=nd.l&&t>=nd.r) {nd.update(x); return;}
42
       down(now);
43
       if(s<=ndl.r) update(s,t,x,now<<1);</pre>
44
       if(t>ndl.r) update(s,t,x,now<<1|1);</pre>
45
       nd=marge(ndl,ndr);
46
47
     segnode query(int s,int t,int now=1){
48
       \underline{if}(s \le nd.1\&\&t \ge nd.r) {
49
         if(s==nd.1) {
50
           segnode x;x.init();
51
            return marge(x,nd);
52
         }else return nd;
53
54
       down(now);
55
       segnode ans;ans.init();
56
       if(s<=ndl.r) ans=marge(ans,query(s,t,now<<1));</pre>
57
       if(t>ndl.r) ans=marge(ans,query(s,t,now<<1|1));</pre>
58
       nd=marge(ndl,ndr);
59
       <u>return</u> ans;
     }
60
61| }tree;
62
   int main() {
63
       \underline{cin} >> n >> m;
64
       register int a,b,c,d;
65
       tree.maketree(1,n);
66
       \underline{\text{while}}(m--) {
67
            <u>cin</u>>>a>>b>>c;
68
            <u>if</u>(a==1)<u>cin</u>>>d;tree.update(b,c,d);
69
            else cout<<(1<<tree.query(b,c).count())<<endl;</pre>
70
71 }
```

5.14 杂项-高精度整数类

```
1 namespace bignumbers
  const int maxd=9999,dlen=4;
  class bignumber {public:
    int len,a[maxn];
5
    bignumber(){len=1;memset(a,0,sizeof(a));}
6
    bignumber(const 11);
    bignumber(<u>const</u> <u>char</u>*);
8
    bignumber(const bignumber &);
9
    bignumber & operator = (const bignumber &);
10
    friend istream& operator>>(istream&,bignumber&);
11
    bignumber operator+(const bignumber &)const;
12
    bignumber operator - (const bignumber &) const;
13
    bignumber operator*(const bignumber &)const;
    bignumber operator/(const int &)const;
14
    bignumber operator^(const int &)const;
15
    11 operator%(const 11 &)const;
16
17
    bool operator > (const bignumber &T) const;
18
    bool operator>(const int &t)const;
19
    void print();
20 };
21 bignumber::bignumber(const 11 b) {
22
    int c,d=b;len=0;
23
    memset(a,0,sizeof(a));
24
    while (d>maxd) {
25
      c=d-(d/(maxd+1))*(maxd+1);
```

```
d=d/(maxd+1);
27
       a[len++]=c;
     }
28
    a[len++]=d;
29
30|}
31 bignumber::bignumber(const char *s) {
32
     int t,k,index,L;
33
     memset(a,0,sizeof(a));
34
     L=strlen(s);
35
     len=L/dlen;
36
     if(L%dlen)len++;
37
     index=0;
38
     for(int i=L-1;i>=0;i-=dlen) {
39
       t=0,k=i-dlen+1;
40
       \underline{if}(k<0)k=0;
41
       rep(j,k,i)t=t*10+s[j]-'0';
42
       a[index++]=t;
43
44|}
45| bignumber::bignumber(const bignumber &T):len(T.len) {
46
     memset(a,0,\underline{sizeof}(a));
     rep(i,0,len-1)a[i]=T.a[i];
47
48|}
49 bignumber &bignumber::operator = (const bignumber &n) {
    memset(a,0,sizeof(a));
50
     rep(i,0,n.len-1)a[i]=n.a[i];
51
52
     len=n.len;
53
     return *this;
54 }
55 char ch[maxn*dlen];
56
  istream& operator>>(istream &in,bignumber &b) {
57
     in>>ch;
58
     int L=strlen(ch);
59
     int count=0,sum=0;
60
     for(int i=L-1;i>=0;) {
61
       <u>int</u> t=1;sum=0;
       for(int j=0;j<4&&i>=0;j++,i--,t*=10)
sum+=(ch[i] - '0')*t;
62
63
64
       b.a[count++]=sum;
65
66
     b.len=count++;
67
     return in;
68|}
69 bignumber bignumber::operator+(const bignumber &T)const {
70
     bignumber t(*this);
71
     int big=T.len>len?T.len:len;
72
     rep(i,0,big-1){
73
       t.a[i]+=T.a[i];
74
       <u>if</u>(t.a[i]>maxd)t.a[i+1]++,t.a[i]-=maxd+1;
75
76
     <u>if</u>(t.a[big]!=0) t.len=big+1;
77
     else t.len=big;
78
     return t;
79
80 bignumber bignumber::operator-(const bignumber &T)const {
81
     bool flag;
82
     bignumber t1,t2;
     if(*this>T) {
83
84
       t1=*this;t2=T;flag=0;
85
86
       t1=T;t2=*this;flag=1;
87
88
     int j,big=t1.len;
89
     rep(i,0,big-1){
90
       <u>if</u>(t1.a[i]<t2.a[i]) {
91
         j=i+1;
92
         \underline{\text{while}}(\text{t1.a[j]} == 0) \text{ j++;}
93
         t1.a[j--]--;
         while (j>i) t1.a[j--]+=maxd;
94
         t1.a[i]+=maxd+1-t2.a[i];
95
       } else t1.a[i]-=t2.a[i];
96
97
98
     t1.len=big;
```

```
while(t1.a[t1.len-1] == 0&&t1.len>1)
100
       t1.len--,big--;
     <u>if</u>(flag) t1.a[big-1]=0-t1.a[big-1];
101
102
     <u>return</u> t1;
103 }
104 bignumber bignumber::operator*(const bignumber &T)const {
105
     bignumber ret;
106
     int up,temp,temp1;
107|
     rep(i,0,len-1){
108
       up=0;
109
       rep(j,0,T.len-1){
110
         temp=a[i]*T.a[j]+ret.a[i+j]+up;
111
         if(temp>maxd) {
112
           temp1=temp-temp/(maxd+1)*(maxd+1);
113
           up=temp/(maxd+1),ret.a[i+j]=temp1;
114
         } else up=0,ret.a[i+j]=temp;
115
       if(up!=0)ret.a[i+T.len]=up;
116
     }
117
118
     ret.len=len+T.len;
     while(ret.a[ret.len-1] == 0&&ret.len>1)ret.len--;
119
120|
     return ret;
121|}
122 bignumber bignumber::operator/(const int &b)const {
123
     bignumber ret;
124
     <u>int</u> down=0;
125
     per(i,0,len-1) {
       ret.a[i]=(a[i]+down*(maxd+1))/b;
126
       down=a[i]+down*(maxd+1) - ret.a[i]*b;
127
128
     ret.len=len;
129
130
     while(ret.a[ret.len-1] == 0&&ret.len>1) ret.len--;
131
     <u>return</u> ret;
132|}
133 ll bignumber::<u>operator</u>%(<u>const</u> ll &b)<u>const</u> {
134
     int d=0;
135
     per(i,0,len-1)
136
       d=((d*(maxd+1))%b+a[i])%b;
137
     return d;
138|}
139 bignumber bignumber::<u>operator</u>^(<u>const</u> <u>int</u> &n)<u>const</u> {
140
     if(n==0)return 1;
|141|
     if(n==1)return *this;
142
     bignumber t, ret(1);
143
     int m=n,i;
144
     while(m>1) {
       t=*this;
145
       for(i=1;(i<<1)<=m;i<<=1)t=t*t;</pre>
146
147
       m-=i,ret=ret*t;
148
       if(m==1)ret=ret*(*this);
     }
149
150
     return ret;
151|}
152 bool bignumber::operator > (const bignumber &T) const {
     if(len>T.len)return true;
153
154
     else if(len==T.len) {
155
       int ln=len-1;
       while (a[ln] == T.a[ln] \&\& ln >= 0) ln --;
156
157
       if(ln>=0&&a[ln]>T.a[ln]) return true;
158
       else return false;
159
     }else return false;
160|}
161 bool bignumber::operator (const int &t)const {
162
     bignumber b(t);
163
     return *this>b;
164| \}
165 <u>void</u> bignumber::print() {
166
     printf("%d",a[len-1]);
     per(i,0,len-2)printf("%04d",a[i]);
167
    printf("\n");
168
170 } using bignumbers::bignumber;
```

5.15 杂项-分数类

```
1 //num是分子,den是分母,分母始终保持为正
  template<typename T>class farction{public:
    T num, den;
    farction(T num=0,T den=1) {
      if (den<0) {
6
        num=-num;
7
        den=-den;
8
9
      T g=_gcd(abs(num),den);
10
      this->num=num/g;
11
      this->den=den/g;
12
13
    farction operator +(const farction &o) const {
      return farction(num*o.den+den*o.num,den*o.den);
14
15
16
    farction operator -(const farction &o) const {
17
      return farction(num*o.den-den*o.num,den*o.den);
18
19
    farction operator *(const farction &o) const {
20
      return farction(num*o.num,den*o.den);
21
22
    farction operator /(const farction &o) const {
23
      return farction(num*o.den,den*o.num);
24
25
    bool operator <(const farction &o) const {</pre>
26
      return num*o.den<den*o.num;</pre>
27
28
    bool operator >(const farction &o) const {
29
      return num*o.den>den*o.num;
30
    bool operator ==(const farction &o) const {
31
32
      return num*o.den==den*o.num;
33
34 };
```

5.16 杂项-N进制快速幂优化

```
1|//底数不变,求值次数1e5以上的情况,可以加速10-100倍
 2|//1e9: maxp 1e3,maxv 1e9,100倍
3 //1e18: maxp 32000, maxv 1e18,14倍左右
  const 11 maxp=32000,maxv=1e18;
  const ll maxw=log(maxv)/log(maxp)+1;
  class basepow{public:
    11 pw[maxw][maxp];
    void init(ll base){
9
      base%=mod;
10
      rep(i,0,maxw-1)pw[i][0]=1;
      rep(i,1,maxp-1) pw[0][i]=(pw[0][i-1]*base)%mod;
11
12
      rep(i,1,maxw-1){
13
        pw[i][1]=pw[i-1][maxp-1]*pw[i-1][1]%mod;
|14|
        rep(j,2,maxp-1) pw[i][j]=pw[i][j-1]*pw[i][1]%mod;
15
    }
16
    inline 11 getpow(11 b,11 res=1,int cnt=0){
17
18
      while(b){
19
        res=res*pw[cnt++][b%maxp]%mod;
20
        b/=maxp;
21
22
      return res;
23
    }
24| }p;
```

6 String

6.1 KMP算法

```
1 //输入的数组从0开始,next函数为p,从1开始
2 template<typename T>class prefix{public:
3 int p[maxn],lens;
4 T *s;
5 void init(T *_s,int _lens){
6 s=_s,lens=_lens;
7 rep(i,0,lens-1) p[i]=0;
```

```
p[0] = -1;
9
       int now=0,pos=-1;
10
       while(now<lens)</pre>
11
         \underline{if}(pos==-1||s[now]==s[pos]) p[++now]=++pos;
12
         else pos=p[pos];
13
     vector<int> find(T *t,int lent){
14
15
       int now,pos=0;
16
       vector<int> ans;
17
       while(now<lent) {</pre>
         \underline{if}(pos==-1||t[now]==s[pos]) pos++,now++;
18
19
         else pos=p[pos];
20
         if(pos==lens) pos=p[pos],ans.push_back(now-lens);
21
22
       return ans:
23
24 };
25 prefix<char> kmp;
```

6.2 manacher算法

```
1 class manacher{
     char ma[maxn<<1]; int lenp[maxn<<1];</pre>
     void getp(char *s,int len){
       int p=0;
5
      ma[p++] = '$', ma[p++] = '#';
6
       forn(i,len) ma[p++]=s[i],ma[p++]='#';
7
       int r=0, mid=0;
8
      forn(i,(len+1)<<1){
9
         lenp[i]=r>i?min(lenp[(mid<<1)-i],r-i):1;</pre>
10
         while (ma[i+lenp[i]] == ma[i-lenp[i]]) lenp[i]++;
11
         if(i+lenp[i]>r)r=i+lenp[i],mid=i;
12
13
     }
14 }
```

6.3 后缀数组-倍增

```
1 namespace suffix{
     int tr[maxn],rank[maxn],sa[maxn],h[maxn],has[maxn],n;
3
     \underline{int} cmp(\underline{int} x,\underline{int} y,\underline{int} k){
       if(x+k>n||y+k>n) return 0;
4
5
       return rank[x] == rank[y] &&rank[x+k] == rank[y+k];
6
     void getsa(char *s,int _n,int m=233){
8
       int i,cnt;n=_n;
9
       <u>for</u>(i=1;i<=n;i++)has[s[i]]=0;
10
       for(i=0;i<=m;i++)has[i]=0;
11
       for(i=1;i<=n;i++)has[s[i]]++;
12
       <u>for</u>(i=1,cnt=0;i<=m;i++)<u>if</u>(has[i])tr[i]=++cnt;
13
       for(i=1;i<=m;i++)has[i]+=has[i-1];</pre>
14
       for(i=1;i<=n;i++)rank[i]=tr[s[i]],sa[has[s[i]]--]=i;</pre>
15
       for(int k=1;cnt!=n;k<<=1){</pre>
16
         for(i=1;i<=n;i++)has[i]=0;
17
         for(i=1;i<=n;i++)has[rank[i]]++;</pre>
18
         <u>for</u>(i=1;i<=n;i++)has[i]+=has[i-1];
         for(i=n;i>=1;i--)if(sa[i]>k)tr[sa[i]-k]=has[rank[sa[i]-k]]--;
19
         for(i=1;i<=k;i++)tr[n-i+1]=has[rank[n-i+1]]--;</pre>
20
21
         <u>for</u>(i=1;i<=n;i++)sa[tr[i]]=i;
         for(i=1,cnt=0;i<=n;i++)tr[sa[i]]=cmp(sa[i],sa[i-1],k) ? cnt:++cnt;</pre>
22
23
         for(i=1;i<=n;i++)rank[i]=tr[i];</pre>
24
       fill_n(h,n+2,0);
25
26
       for(int i=1;i<=n;i++){</pre>
27
          if(rank[i]==1)continue;
28
         for(int j=max(1,h[rank[i-1]]-1);;j++){
29
            <u>if</u>(s[i+j-1]==s[sa[rank[i]-1]+j-1])h[rank[i]]=j;
30
           <u>else</u> <u>break</u>;
31
32
33
     }
34 }
```

6.4 最小表示算法

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```
1 int minrep(int *s,int n){
      <u>int</u> i=0, j=1, k=0, t;
 3
      \underline{\text{while}}(i < n \& \& j < n \& \& k < n) \{
        t=s[i+k>=n?i+k-n:i+k]-s[j+k>=n?j+k-n:j+k];
        <u>if</u>(!t) ++k;
 5
 6
        else {
           if(t>0) i+=k+1;
           else j+=k+1;
 9
           <u>if</u>(i==j) ++j;
10
           k=0;
|11|
      }
12
13
      return min(i,j);
14|}
```

6.5 字符串哈希

```
1 const 11 decm=13331, modh=(11)1e17+13;
  inline ll mulh(ll a,ll b){
    ll tmp=(a*b-(ll)((long double)a/modh*b+1.0e-8)*modh);
    return tmp<0?tmp+modh:tmp;</pre>
4
5|}
6
  inline 11 addh(11 a,11 b){11 x=a+b; if(x<0)return a+b+modh; else if(x>=modh)x-modh; else return x;}
7 | 11 pw [maxn];
8 void init(int n=maxn-5) {pw[0]=1; rep(i,1,n) pw[i]=mulh(pw[i-1], decm);
9 <u>class</u> strhash{<u>public</u>:
    ll hs[maxn],len;
    void calhs(string &s,int n){len=n;rep(i,0,len-1)hs[i+1]=addh(mulh(hs[i],decm),s[i]);}
12
    inline ll geths(int a,int b){return addh(hs[b],-mulh(hs[a-1],pw[b-a+1]));}
13 }table;
```

6.6 自动机-tire树

```
1 const int csize=26,minc='a';
  class trie{public:
  #define nd node[now]
    struct tnode{
      int cnt,son[csize];
6
      tnode(){
7
        cnt=0:
8
        memset(son,0,sizeof son);
9
      }
10
    }node[maxn];
11
    int sz=0;
12
    void insert(char *s,int len){
13
      int now=0;
|14|
      rep(i,0,len-1){
15
        int ch=s[i]-minc;
16
        if(!nd.son[ch]) nd.son[ch]=++sz;
17
        now=nd.son[ch];
18
      node[now].cnt++;
19
20
21
    int find(char *s, int len){
22
      int ch,now=0;
23
      rep(i,0,len-1){
24
        ch=s[i]-minc;
        if(!nd.son[ch]) return -1;
25
26
        now=nd.son[ch];
27
28
      if(!nd.cnt) return -1;
29
      return 1;
    }
30
31 }tree;
```

6.7 自动机-tire图

```
1 #include<iostream>
2 #include<cstdio>
3 #include<cstring>
4 #include <vector>
5 #include <queue>
using namespace std;
8 const int N = 1000000 + 10, INF = 0x3f3f3f3f;
```

```
10 struct node {
    node *next[26];
|11|
12
    node *suff; //指向后缀节点
    bool flag;
13
14 } trie[N], *root;
|15| <u>int</u> tot;
16 char ori[N];
17 node* node_init() {
    trie[tot].flag = false;
18
    trie[tot].suff = NULL;
19
20
    memset(trie[tot].next, 0, sizeof trie[tot].next);
    return trie + tot++;
21
22 }
23 void trie_insert(char *s) {
24
    node *p = root;
25
    for(int i = 0; s[i]; i++) {
26
      \underline{int} j = s[i] - 'a';
      if(p->next[j] == NULL)
27
28
        p->next[j] = node_init();
29
      p = p-next[j];
30
31
    p->flag = <u>true</u>;
32 }
33 void trie_graph() {
34
    //trie[0]为虚拟节点,root为trie[1],trie[0]的所有边均指向root,方便以后操作
    for(int i = 0; i < 26; i++)
35
      trie[0].next[i] = root;
36
37
    root->suff = trie + 0;
    trie[0].suff = NULL;
38
39
    queue<node*> que;
    que.push(root);
40
41
    while(! que.empty()) {
42
      node *p = que.front();
      que.pop();
43
44
      for(int i = 0; i < 26; i++)
        if(p-\text{next}[i]) {
45
          //查看父亲节点的后缀节点是否存在编号为i的边,若没有,就一直往上找,直到根节点
46
47
          node *ptr = p->suff;
          while(ptr && ! ptr->next[i])
48
49
            ptr = ptr->suff;
          //要么找到了,要么循环到了虚拟节点后,后缀节点被置为root
51
          p->next[i]->suff = ptr->next[i];
52
          if(ptr->next[i]->flag)
53
            p->next[i]->flag = <u>true</u>;
54
          que.push(p->next[i]);
55
56
    }
57 }
58
  bool trie_query(char *s) {
59
    node *p = root;
    for(int i = 0; s[i]; i++) {
  int j = s[i] - 'a';
60
61
      while(true) {
62
        if(p->next[j] != NULL) {
63
64
          p = p-\text{next}[j];
65
          <u>if</u>(p->flag == <u>true</u>)
66
            return true;
67
          break;
68
        } else
          p = p \rightarrow suff;
69
70
    }
71
72
    <u>return</u> <u>false</u>;
73 }
74 <u>int</u> main() {
75
    int n;
    scanf("%d", &n);
76
    tot = 2;//tot+0是虚拟节点, tot+1是根节点, 故从2开始
77
78
    root = trie + 1;
79
    memset(trie[root].next, 0, sizeof trie[root].next);
80
    for(int i = 1; i <= n; i++) {</pre>
      scanf("%s", ori);
81
```

```
6.8 自动机-AC自动机
```

```
1 const int csize=128,minc=0;
  class autom{public:
3 #define nd node[now]
    struct acnode{int id,son[csize],fail;}node[maxn];
    int sz=0;
6
    queue<<u>int</u>> que;
    void clear(int n=maxn-5){
      memset(node,0,sizeof(acnode)*n);
      sz=0;
    }
10
11
    void insert(int id, char *s, int len=0){
12
      if(!len)len=strlen(s);
13
      int now=0;
      rep(i,0,len-1){
14
15
        int ch=s[i]-minc;
        if(!nd.son[ch]) nd.son[ch]=++sz;
16
17
        now=nd.son[ch];
18
19
      nd.id=id;
20
21
    void init(){
|22|
      int now=0;
23
      rep(i,0,csize-1) <u>if</u>(nd.son[i])
24
        que.push(nd.son[i]);
25
       while(!que.empty()){
26
        now=que.front();que.pop();
27
        rep(i,0,csize-1){
28
          if(nd.son[i]) {
            node[nd.son[i]].fail=node[nd.fail].son[i];
29
30
            que.push(nd.son[i]);
31
           }<u>else</u> nd.son[i]=node[nd.fail].son[i];
32
33
34
    }
35
    vector<int> query(char *t,int len=0){
36
       if(!len)len=strlen(t);
37
      int now=0;
38
      vector<int> ans;
39
      rep(i,0,len-1) {
        now=nd.son[t[i]-minc];
40
        for(int j=now; j; j=node[j].fail)
41
42
          if(node[j].id)ans.push_back(node[j].id);
43
44
       return ans;
45
46 }acam;
```

6.9 自动机-后缀自动机

```
const int maxn=1e5+10,maxsz=26+1,ch0='a';
  class saffixam {public:
     int fa[maxn<<1],son[maxn<<1][maxsz];</pre>
    int last,cnt,len[maxn<<1];</pre>
    void insert(int ch) {
5
      int pre=last,now=++cnt;
      last=now,len[now]=len[pre]+1;
      for(; pre&&!son[pre][ch]; pre=fa[pre])son[pre][ch]=now;
9
      if(!pre)fa[now]=1;
10
11
        int q=son[pre][ch];
12
        if(len[pre]+1==len[q])fa[now]=q;
13
        <u>else</u> {
14
          int nq=++cnt;
15
          memcpy(son[nq],son[q],sizeof(son[0]));
```

```
fa[nq]=fa[q];
17
          len[nq]=len[pre]+1;
18
          fa[q]=fa[now]=nq;
19
          for(; son[pre][ch]==q; pre=fa[pre])son[pre][ch]=nq;
20
      }
21
    }
|22|
23
    void init() {
24
      rep(i,0,cnt) {
25
        memset(son[i],0,sizeof son[0]);
26
        fa[i]=len[i]=0;
27
28
      last=cnt=1;
29
    }
30
    void insert(string &s ) {for(auto i:s)insert(i-ch0);}
31
    bool find(string &s) {
      int now=1;
32
33
      for(auto i:s) if(!(now=son[now][i-ch0]))return 0;
34
    }
35
36 } sam;
```

自动机-回文自动机 6.10

```
struct PAT {
     struct node {
       int len,num,fail,son[26];
     } t[maxn];
     int last,n,tot,s[maxn];
     void init() {
       memset(t,0,sizeof(t));
8
       tot=last=1;
9
       n=0;
10
       t[0].len=0;
11
       t[1].len=-1;
12
       t[0].fail=t[1].fail=1;
13
       s[0]=-1;
14
15
     int add(int c) {
16
       int p=last;
       s[++n]=c;
17
       \underline{\text{while}}(s[n]!=s[n-1-t[p].len])
18
19
         p=t[p].fail;
20
       <u>if</u>(!t[p].son[c]) {
21
         int v=++tot,k=t[p].fail;
22
         \underline{\text{while}}(s[n]!=s[n-t[k].len-1])
23
           k=t[k].fail;
24
         t[v].fail=t[k].son[c];
25
         t[v].len=t[p].len+2;
26
         t[v].num=t[t[v].fail].num+1;
27
         t[p].son[c]=v;
28
29
       last=t[p].son[c];
30
       return t[last].num;
31
32| } T;
```

Others

常用头文件

```
1 #include <bits/stdc++.h>
2 #define 11 long long
3 #define rep(ii,a,b) for(int ii=a;ii<=b;++ii)
4 #define per(ii,a,b) for(int ii=b;ii>=a;--ii)
5 #define forn(i,x,g,e) for(int i=g[x];i;i=e[i].next)
 6 #define IO ios::sync_with_stdio(false);cin.tie(0);cout.tie(0)
  #define ull unsigned long long
  #define fi first
  #define se second
10 #define mp make_pair
11 #define pii pair<11,11>
12 #define all(x) x.begin(),x.end()
13 <u>#define</u> show(x) <u>cout</u> << #x << "=" << x << <u>endl</u>
14 <u>#define</u> showa(a,b) <u>cout</u><<#a<<'<u>[</u>'<<b<<"]="<<a[b]<<<u>endl</u>
```

```
15 #define show2(x,y) cout << #x << "=" << x << " " << #y << "=" << y << end]
16 #define show3(x,y,z) cout<<#x<<"="<<x<" "<<#y<<"="<<y<" "<<#z<<"="<<z<<endl
  #define show4(w,x,y,z) cout<<#w<<"="<<w<" "<*#x<<"="<<x<" "<*#y<<"="<<y<" "<*#z<<"="<<z<<endl
17
  #define show5(v,w,x,y,z) cout<<#v<<"="<<v<" "<<#w<<"="<<x<<" "<<#x<<"="<<x<<" "<<#y<<"="<<y<<" "<<#z<<" "
181
       ="<<z<<end1
19 \frac{\text{#define}}{\text{mean}} showa2(x,a,b) \frac{\text{cout}}{\text{cout}} showa2(x,a,b) \frac{\text{cout}}{\text{cout}} showa2(x,a,b) \frac{\text{cout}}{\text{cout}}
20 using namespace std;//head
21 const int maxn=1e6+10, maxm=2e6+10;
22 const 11 INF=0x3f3f3f3f, mod=1e9+7;
23 int casn,n,m,k;
24 | \underline{int}  main() {
29 //struct node{
30 // int x,y;
31 // friend bool operator<(const node&a,const node&b){
32 //
         return a.x==b.x?a.y<b.y:a.x<b.x;
33 // }
34 //}a[maxn];
35 // auto _start=chrono::high_resolution_clock::now();
36 // auto _end=chrono::high_resolution_clock::now();
37 // cerr<<"elapsed time: "<<chrono::duration<double,milli>(_end-_start).count()<<" ms\n";
38 //int size=(64)<<20; //64MB
       __asm__("movq %0, %%rsp\n"::"r"((char*)malloc(size)+size));
```

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快速读写

```
1 namespace io{
     const int L=(1<<21)+1;</pre>
3
     char ibuf[L],*iS,*iT,obuf[L],*oS=obuf,*oT=obuf+L-1,c,st[55];int f,tp;
4
     #define gc() (iS==iT?(iT=(iS=ibuf)+fread(ibuf,1,L,stdin),(iS==iT?EOF:*iS++)):*iS++)
5
     inline void flush(){fwrite(obuf,1,oS-obuf,stdout);oS=obuf;}
     inline void putc(char x){*oS++=x; if(oS==oT)flush();}
     template<class I>
8
     inline void gi(I&x){
       for(f=1,c=gc();c<'0'||c>'9';c=gc())if(c=='-')f=-1;
9
10
       for(x=0; c <= '9' \&\&c >= '0'; c=gc())x=x*10+(c&15);x*=f;
11
     template<class I>
12
13
     inline void print(I x){
       <u>if</u>(!x)putc('0'); <u>if</u>(x<0)putc('-'), x=-x;
|14|
       while(x)st[++tp]=x%10+'0',x/=10;
15
16
       while(tp)putc(st[tp--]);
17
18
     <u>inline</u> <u>void</u> gs(<u>char</u>*s, <u>int</u>&l){
       for(c=gc();c!='_'\&\(c<'a'||c>'z');c=gc());
19
       for(1=0;c=='_'||c<='z'&&c>='a';c=gc())s[1++]=c;
20
21
22 };
```

快速读写加强版

```
1 namespace fastio{//支持读取整数,字符串,输出整数
  bool isdigit(char c){return c>=48&&c<=57;}</pre>
  const int maxsz=1e7;
  class fast_iostream{public:
    char ch=get_char();
6
    bool endf=1,flag;
    char get_char(){
8
      static char buffer[maxsz],*a=buffer,*b=buffer;
9
      return b==a&&(b=(a=buffer)+fread(buffer,1,maxsz, stdin),b==a)?EOF:*a++;
10
11
    template<typename type>bool get_int(type& tmp){
12
      flag=tmp=0;
      while(!isdigit(ch)&&ch!=EOF){flag=ch=='-';ch=get_char();};
13
      if(ch==EOF)return endf=0;
14
15
      do{tmp=ch-48+tmp*10;}while(isdigit(ch=get_char()));
16
      if(flag)tmp=-tmp;
17
      return 1;
18
19
    int get_str(char* str){
20
      char* tmp=str;
```

```
while(ch=='\r'||ch=='\n'||ch==' ')ch=get_char();
|22|
       if(ch==EOF)return(endf=0),*tmp=0;
       \underline{do}\{*(tmp++)=ch; ch=get\_char();\}\underline{while}(ch!='\r'\&\&ch!='\n'\&\&ch!=' '\&\&ch!=EOF);
23
       *(tmp++\bar)=0;
24
25
       return(int)(tmp-str-1);
26
27
    fast_iostream& operator>>(char* tmp){get_str(tmp); return *this;}
28
     template<typename type>fast_iostream& operator>>(type& tmp){get_int(tmp);return *this;}
|29|
    operator bool() const {return endf;}
30| };
31 template < typename type > void put(type tmp) {
32
    if (tmp==0){putchar(48); return;}
33
    static int top,stk[21];
     if (tmp<0){tmp=-tmp;putchar('-');}</pre>
34
     while(tmp)stk[++top]=tmp%10,tmp/=10;
35
36
     while(top)putchar(stk[top--]+48);
37
38 } fastio::fast_iostream io;
```

LIS

```
int num[maxn],dp[maxn],len;
   int lwb(int now){
     int l=0,r=len,ans=0;
     while(l<r){
       int mid=(l+r)>>1;
 6
       if(dp[mid]>=now) l=mid+1;
       else r=mid;
 8
 9
     return 1;
10|}
11
   int lwb2(int now){
12
     int l=0,r=len;
13
     \underline{\text{while}}(1 < r) \{
14
       <u>int</u> mid=(l+r)>>1;
       if(dp[mid]>=now) r=mid;
15
16
       else l=mid+1;
17
18
     return 1;
19|}
20 <u>int</u> main(){
21
     \underline{cin} >> n;
22
     <u>for(int</u> i=1;i<=n;i++) <u>cin</u>>>num[i];
23
     for(int i=1;i<=n;i++){</pre>
24
       if(dp[len]>=num[i]) dp[++len]=num[i];
25
       else dp[lwb(num[i])]=num[i];
26
27
     cout<<len<<endl;</pre>
28
     len=0;
29
     memset(dp,0,sizeof dp);
30
     for(int i=1;i<=n;i++){</pre>
31
       if(dp[len] < num[i]) dp[++len] = num[i];</pre>
32
       else dp[lwb2(num[i])]=num[i];
33
     cout << len << endl;
34
35|
```

矩阵运算

```
class matrix{public://mod
       int a,b;
 3
       vector<vector<11> > x;
 4
       matrix(<u>int</u> a=1,<u>int</u> b=1){
          this->a=a,this->b=b;x.resize(a);
 5
 6
          for(auto &i:x){i.resize(b);std::fill(all(i),0);}
       \underline{\text{void}} \ \ e(\underline{\text{int}} \ \ n)\{a=b=n; x=\text{matrix}(n,n).x; \underline{\text{for}}(\underline{\text{int}} \ \ i=0; i < n; i++)x[i][i]=1; \}
 8
 9
       \underline{\text{void}} \text{ fill(ll } xx=0) \{ \underline{\text{for}}(\underline{\text{int}} \text{ i=0;i<a;i++}) \underline{\text{for}}(\underline{\text{int}} \text{ j=0;j<b;j++}) x[i] [j] = xx; \}
10
       void fill(vector<vector<11>> &y){x=y;a=y.size();b=y[0].size();}
11
       matrix operator *(matrix &m){
12
          matrix ans(a,m.b);
13
          \frac{\text{for}(\text{int } \text{i=0;i<a;i++})}{\text{for}(\text{int } \text{j=0;j<m.b;j++})} \frac{\text{for}(\text{int } \text{k=0;k<b;k++})}{\text{if}} (\text{x[i][k]\&\&m.x[k][j])}
             ans.x[i][j]=(mod+ans.x[i][j]+(x[i][k]*m.x[k][j]+mod)%mod)%mod;
|14|
15
          return ans;
16
```

```
matrix operator +(matrix &m){
18
       matrix ans(a,m.b);
        for(int i=0;i<a;i++)for(int j=0;j<b;j++)ans.x[i][j]=(x[i][j]+m.x[i][j]+mod)%mod;</pre>
19
20
        <u>return</u> ans;
21
     matrix operator -(matrix &m){
22
23
       matrix ans(a,m.b);
24
        for(int i=0;i<a;i++)for(int j=0;j<b;j++)ans.x[i][j]=(x[i][j]-m.x[i][j]+mod)%mod;</pre>
25
       return ans;
26
27
     matrix pow(ll p){
28
       matrix ans;ans.e(a);matrix t;t.fill(x);
29
        \underline{\text{while}}(p)\{\underline{\text{if}}(p\&1) \text{ ans=t*ans;t=t*t;p>>=1;}\underline{\text{return}} \text{ ans;}
30
31 };
```

7.6 莫队算法

```
1 ll a [maxn], ans [maxn], sz, sum;
 2 class block{public:
     11 cnt[maxn];
     struct node{ll l,r,id;
 4
 5
       <u>bool</u> <u>operator</u> <(<u>const</u> node &b)<u>const</u> {
 6
          <u>if</u>(1/sz!=b.1/sz) <u>return</u> 1<b.1;
 7
          \underline{if}((1/sz)\&1) \underline{return} r < b.r;
 8
         return r>b.r;
 9
10
     };
|11|
     void update(int pos,ll flag=1){
12
       sum+=flag*a[pos]*(cnt[a[pos]]*211+flag);
13
       cnt[a[pos]]+=flag;
14
15 }ask;
   vector<block::node> tab;
16
17
   int main() {
18
     \underline{cin}>>n>>m;sz=sqrt(n);
19
     rep(i,1,n) <u>cin</u>>>a[i];
20
     rep(i,1,m) {
21
       int a,b;cin>>a>>b;
|22|
       tab.push_back({a,b,i});
     }
23
     sort(all(tab));
24
25
     <u>int</u> l=tab[0].1,r=tab[0].1-1;
26
     for(auto now:tab){
27
       while(l>now.l) ask.update(--1,1);
       while(l<now.l) ask.update(l++,-1);</pre>
28
29
       while(r<now.r) ask.update(++r,1);</pre>
30
       while(r>now.r) ask.update(r--,-1);
31
       ans[now.id] = sum;
32
33|}
```

7.7 大随机数

```
1| template<typename T>class random_xor{public:
    T s1,s2,s3,s4;
3 //
       random_xor(){
4 //
       srand(time(0)+rand());
5 //
       s1=(T)rand()*rand()*rand();
6 //
       s2=(T)rand()*rand()*rand();
7 //
8 //
     Τ
       getint(){
9 //
       T s3=s1,s4=s2;
       s1=s4;s3^=s3<<23;
10 //
       s2=s3^s4^(s3>>17)^(s4>>26);
11 //
12 //
       return s2+s4;
13 // }
14
   random_xor(){
15
      srand(time(0)+rand());
16
      s1=(T)rand()*rand()*rand();
17
      s2=(T)rand()*rand()*rand();
18
      s3=(T)rand()*rand()*rand();
19
      s4=(T)rand()*rand()*rand();
20
```

```
21
    T getint() {
       T t=s1^(s1<<11);
22
23
      s1=s2;s2=s3;s3=s4;
24
       return s4=s4^(s4>>19)^t^(t>>8);
25
26
    T get(T 1,T r){\underline{\text{return}} getint()%(r-l+1)+l;}
27
    T operator ()(T 1,T r)\{return getint()\%(r-1+1)+1;\}
28
      operator ()(){return getint();}
29|};
30 random_xor<<u>unsigned</u> __int128> a;
31 srand(time(0));
32| //mt19937::result_type seed=time(0);
33 auto randint=bind(uniform_int_distribution<int>(1,r),mt19937(rand()));
34 auto randfloat=bind(uniform_real_distribution<double>(1,r),mt19937(rand()));
```

7.8 大质数

```
1 1e9+7,1e9+9,1e9+21,1e9+33,

2 1e9+87,1e9+93,1e9+97,1e9+103,

3 1e9+123,1e9+181,1e9+207,1e9+223,

4 1e9+241,1e9+271,1e9+289,1e9+297,

5 1e9+321,1e9+349,1e9+363,1e9+403,
```

7.9 编译器位操作

```
#pragma GCC target ("popcnt")
2 __builtin_parity(n);//1的个数的奇偶性
3 __builtin_popcount(n);//1个数
4 __builtin_ctz(n);//末尾0个数,n!=0
5 __builtin_clz(n);//前导0的个数,n!=0
6 __builtin_ffs(n);//最后1的位置,从1开始
7 __builtin_parity(n)//x中1的奇偶性
8 uint32_t __builtin_bswap32(uint32_t x)//按字节翻转
```

7.10 最大子矩阵算法

```
1 /*
 2 把二维看成一维
4 再把起点行和终点行间每一列的数值压缩到每一个点上
   然后求一个最长连续子段和
6
  复杂度O(n^3)
7
  */
  #include <string.h>
9
  #include <stdio.h>
10 #include <iostream>
11 #include <algorithm>
12 using namespace std;
13 const int maxn=1e3+10;
14 const int maxm=1e6+10;
15 const int INF=0x3f3f3f3f;
16 #define 11 long long
17 | \underline{int} | casn,n,m,k;
18 <u>int</u> smax(<u>int</u> a[], <u>int</u> len){
19
    int mx=0,sub=0;
    for(int i=1;i<=len;i++){</pre>
20
21
       sub=max(a[i],sub+a[i]);
|22|
       mx=max(sub,mx);
23
|24|
    return mx;
25|}
26
  int arr[maxn];
27 int dp[maxn] [maxn];
28 int main(){
29 #define test
30 #ifdef test
    freopen("in.txt", "r", stdin); freopen("out.txt", "w", stdout);
31
32 #endif
34
     while (\operatorname{scanf}("%d", \&n)) {
35
       for(int i=1;i<=n;i++){</pre>
         for(int j=1;j<=n;j++){
  scanf("%d",&dp[i][j]);</pre>
36
37
38
       }
39
```

```
<u>int</u> ans=-INF;
       for(int i=1;i<=n;i++){</pre>
41
42
         memset(arr,0,sizeof arr);
43
         for(int j=i;j<=n;j++){</pre>
           for(int k=1;k<=n;k++){</pre>
44
45
             arr[k]+=dp[j][k];
46
47
           ans=max(ans,smax(arr,n));
48
49|
       printf("%d\n",ans);
50
51
  #ifdef test
52
     fclose(stdin);fclose(stdout);system("out.txt");
53
54 #endif
55
     return 0;
56 }
```

7.11 约瑟夫环算法

```
1 #include <bits/stdc++.h>
   using namespace std;
3 typedef long long 11;
6| 11 ysf(11 n,11 k,11 num) //n为人数 k为每次点名的数第num个出列的人的序号 0- n-1
   {
7
8
       \underline{if}(k==1)
9
           return num-1;
10
       11 \text{ ans}=0;
11
       if(num<k)</pre>
12
13
           ans=(k-1)%(n-num+1);
14
           ll p=n-num+1;
15
           for(11 i=2; i<=num; i++)</pre>
16
17
               ans=(ans+k)%(++p);
18
19
       }
20
                                               //num很大时用乘法加速
       else
21
|22|
           ans=-1;
23
           for(ll i=n-num+1; i<=n; i++)</pre>
24
25
               ll j=min(n,(i-1)+((i-1)-ans+(k-2))/(k-1));
26
               ans=(ans+k*(j-i+1))%j,i=j;
27
           }
28
       }
       return ans;
30
31 }
32 <u>int</u> main()
33| {
34
       11 T,T2;
       cin >> T;
35
       T2=T;
36
37
       while (T--)
38
39
           ll n, k, num;
40
           \underline{cin}>>n>>num>>k;
           cout<<"Case #"<<T2-T<<": "<<ysf(n,k,num)+1<<<u>endl</u>;
41
42
45| \}
```

7.12 java高精度

```
package acm;
import java.math.BigInteger;
import java.util.Scanner;

public class Main{
   public static final int maxn=1000,maxm=200000;
   public static BigInteger gcd(BigInteger a,BigInteger b) {
   if(b.compareTo(BigInteger.ZERO)==0)return a;
}
```

```
else return gcd(b, a.remainder(b));
9
10
     public static void main(String[] argc){
11
       int[]a=new int[maxn];
       for(int i=2;i<maxn;i++) a[i]=i;</pre>
12
       \underline{\text{for}}(\underline{\text{int}} \text{ i=2;i<maxn;i++}) \underline{\text{if}}(a[i]!=0)
13
         for(int j=i*2;j<maxn;j=j+i) a[j]=0;</pre>
14
15
       Scanner cin=new Scanner(System.in);
16
       int casn=cin.nextInt();
       <u>int</u>[] prime=<u>new</u> <u>int</u>[10000];
17
       for(int ii=1;ii<=casn;++ii) {</pre>
18
19
         BigInteger x=cin.nextBigInteger();
20
         BigInteger res=BigInteger.ONE;
21
         BigInteger now=BigInteger.ONE;
22
         <u>int</u> i=1;
23
         while(true) {
24
           i++;
25
           if(a[i]==0) continue;
26
           if(res.multiply(BigInteger.valueOf(a[i])).compareTo(x)<=0) {</pre>
27
             res=res.multiply(BigInteger.valueOf(a[i]));
28
             now=now.multiply(BigInteger.valueOf(a[i]+1));
29
           }else break;
30
         BigInteger g=gcd(res,now);
31
32
         res=res.divide(g);
33
         now=now.divide(g);
34
         System.out.println(res+"/"+now);
35
     }
36
37 }
```